



INDIAN INSTITUTE OF  
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TECHNOLOGY

# Breadth-First Search and Depth-First Search

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PDPM

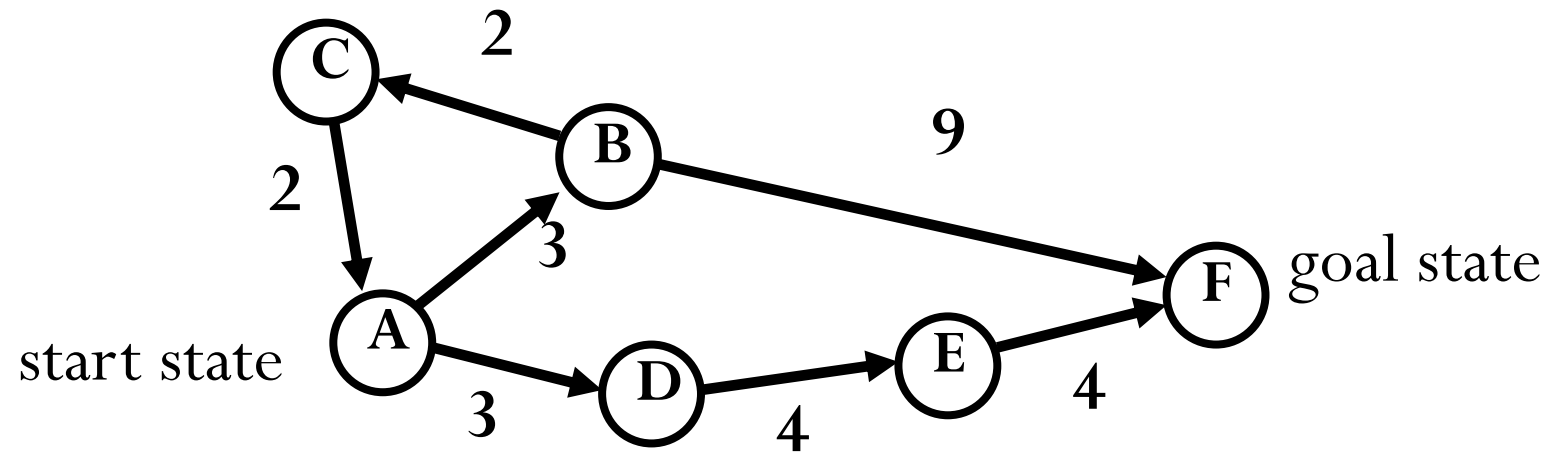
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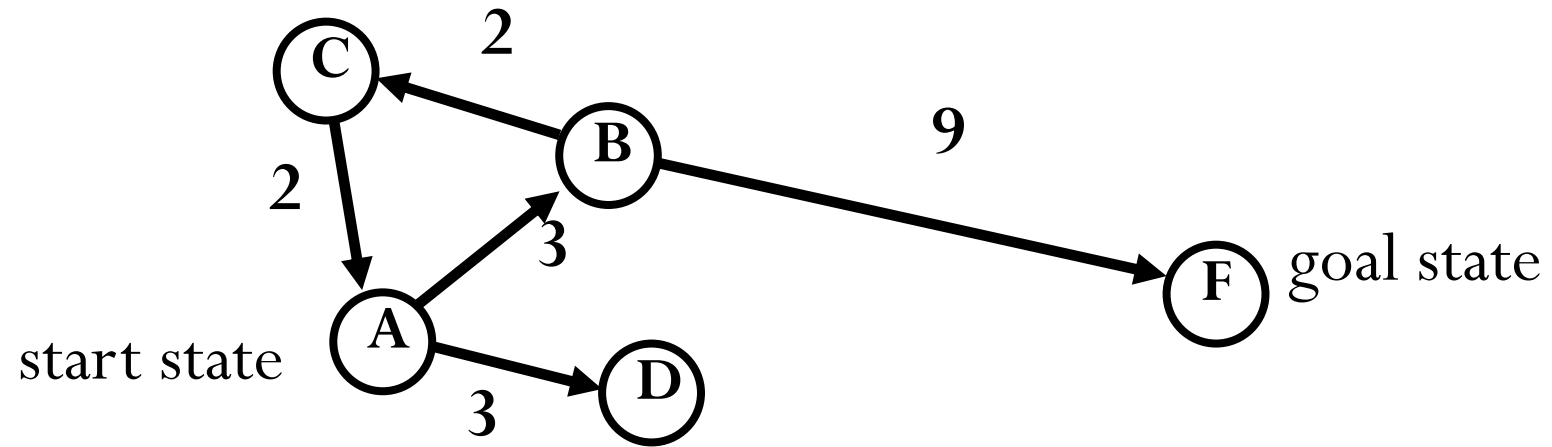
# Search

- We have some actions that can change the state of the world
  - Change induced by an action perfectly predictable
- Try to come up with a sequence of actions that will lead us to a goal state
  - May want to minimize number of actions
  - More generally, may want to minimize total cost of actions
- Do not need to execute actions in real life while searching for solution!
  - Everything perfectly predictable anyway

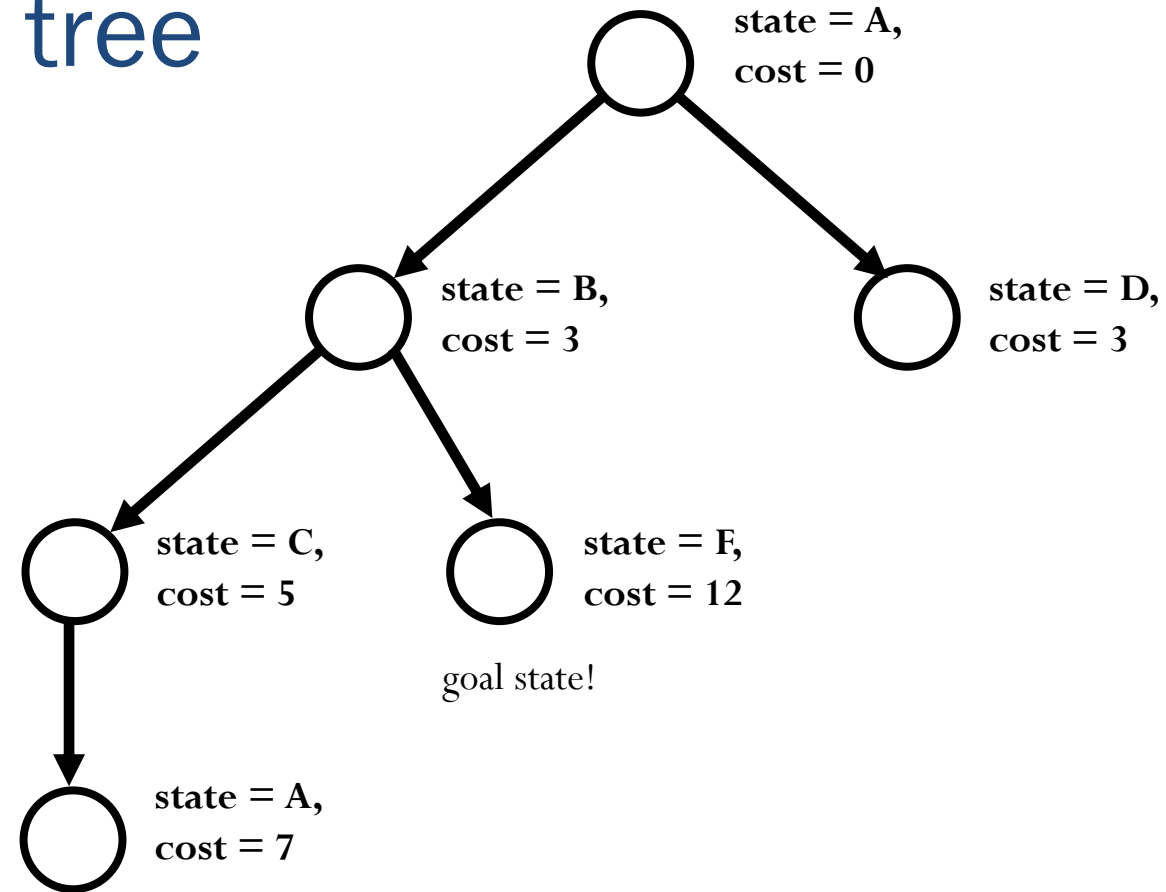
# A simple example: traveling on a graph



# Searching for a solution

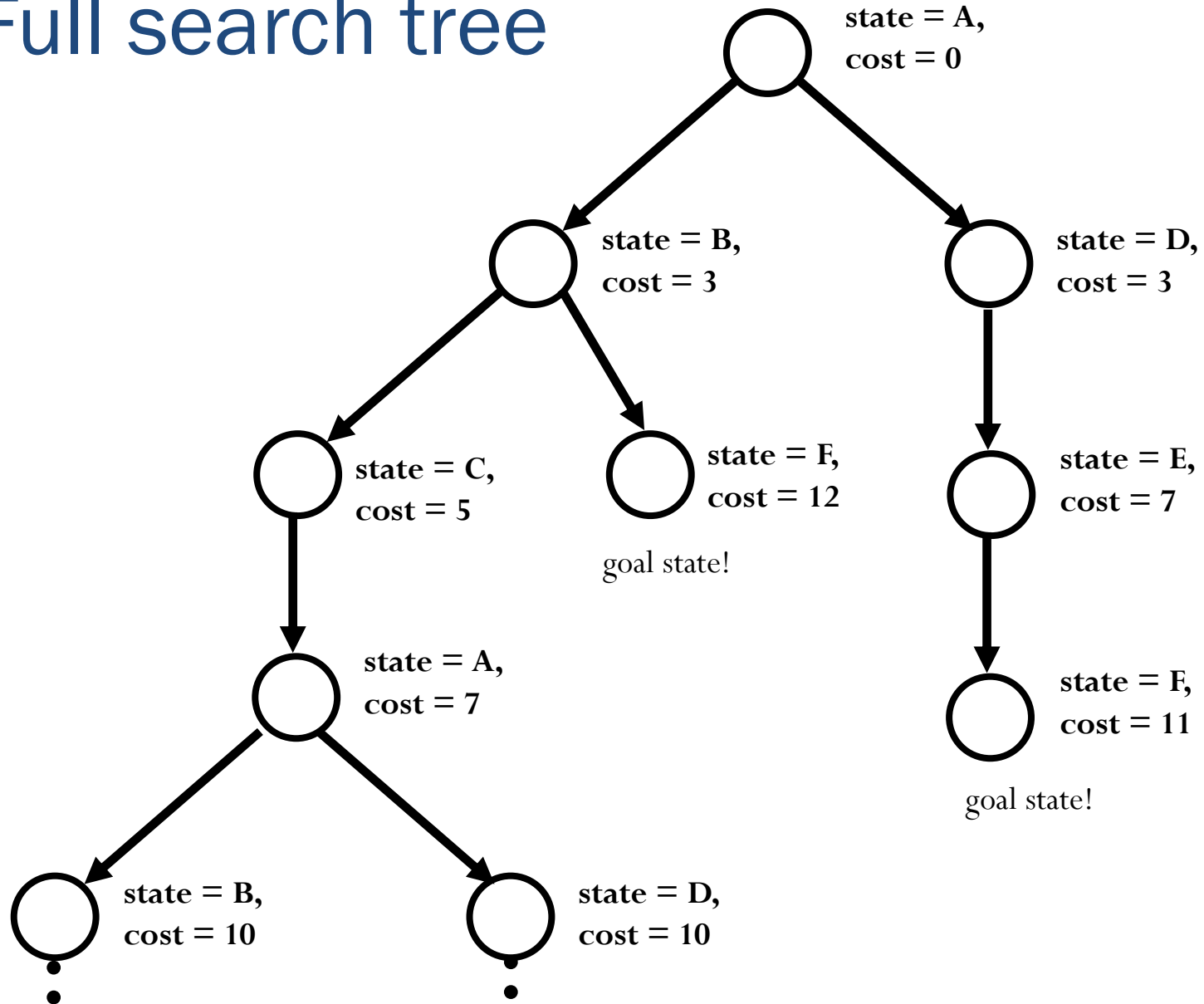


# Search tree

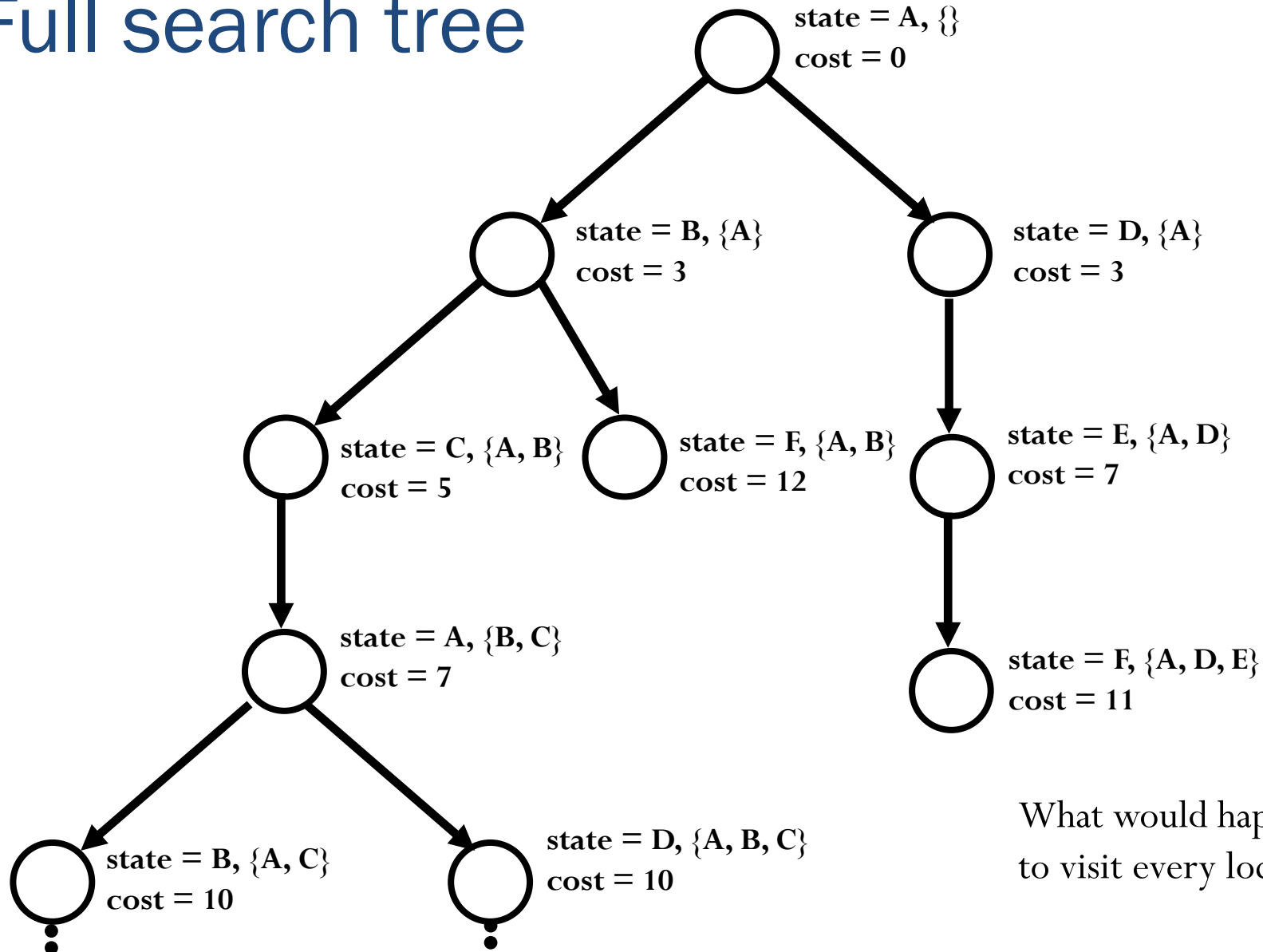


search tree nodes and states are not the same thing!

# Full search tree



# Full search tree



What would happen if the goal were to visit every location twice?

# Key concepts in search

- Set of **states** that we can be in
  - Including an **initial state**...
  - ... and **goal states** (equivalently, a **goal test**)
- For every state, a set of **actions** that we can take
  - Each action results in a new state
  - Typically defined by **successor function**
    - Given a state, produces all states that can be reached from it
- **Cost function** that determines the cost of each action (or **path** = sequence of actions)
- **Solution**: path from initial state to a goal state
  - **Optimal solution**: solution with minimal cost



# Uninformed search

- Given a state, we only know whether it is a goal state or not
- Cannot say one nongoal state looks better than another nongoal state
- Can only traverse state space blindly in hope of somehow hitting a goal state at some point
  - Also called **blind search**
  - Blind does **not** imply unsystematic!

# Searching Examples

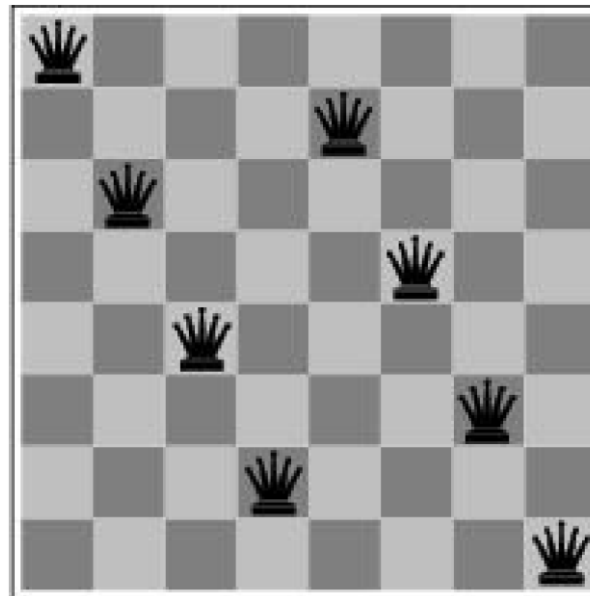
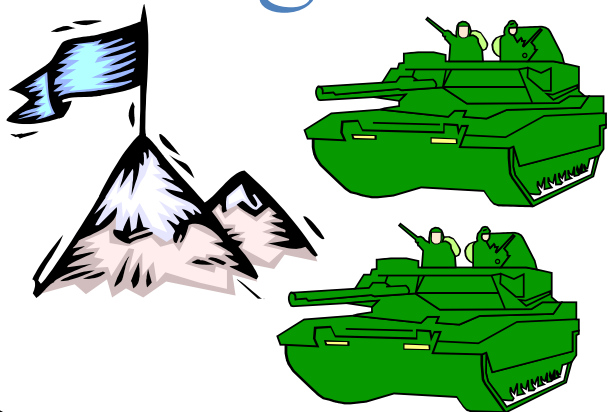


Rush Hour: Move cars forward and backward to “escape”

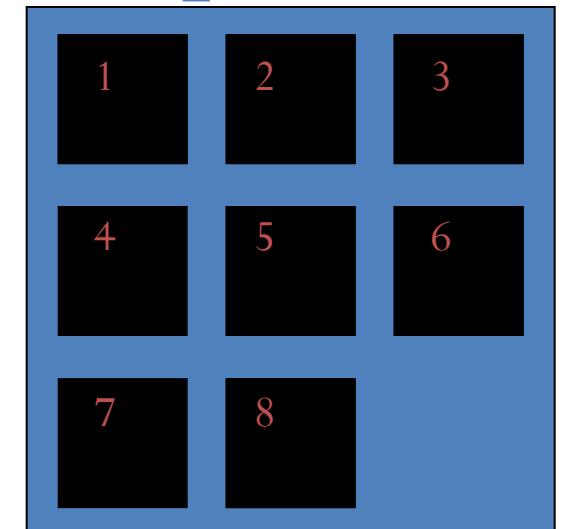


8-queens problem

Logistics



8-puzzle



# Generic search algorithm

- **Fringe** = set of nodes **generated** but not **expanded**
- $\text{fringe} := \{\text{initial state}\}$
- loop:
  - if fringe empty, declare failure
  - choose and remove a node  $v$  from fringe
  - check if  $v$ 's state  $s$  is a goal state; if so, declare success
  - if not, expand  $v$ , insert resulting nodes into fringe
- Key question in search: Which of the generated nodes do we expand next?

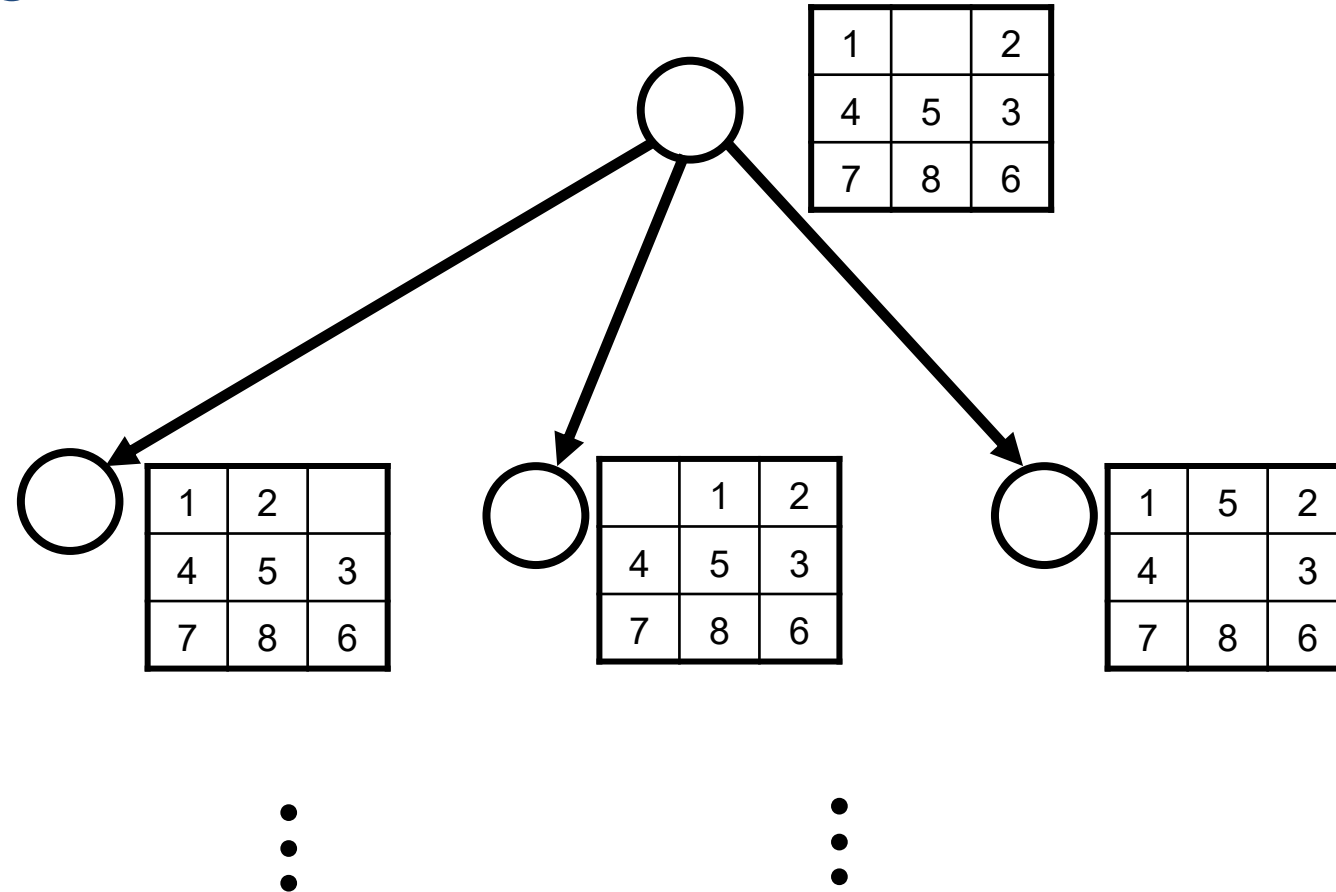
# 8-puzzle

1		2
4	5	3
7	8	6

1	2	3
4	5	6
7	8	

goal state

# 8-puzzle

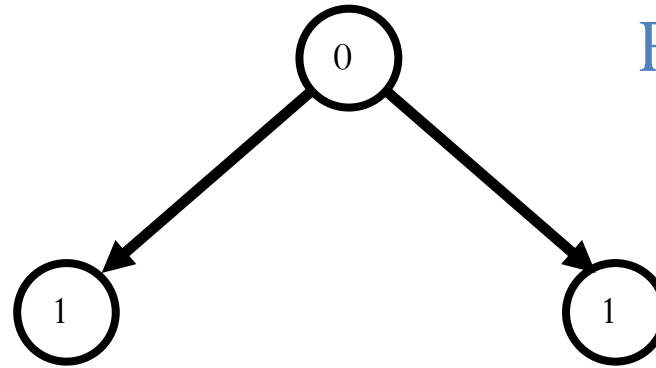


# Breadth-First Search

0

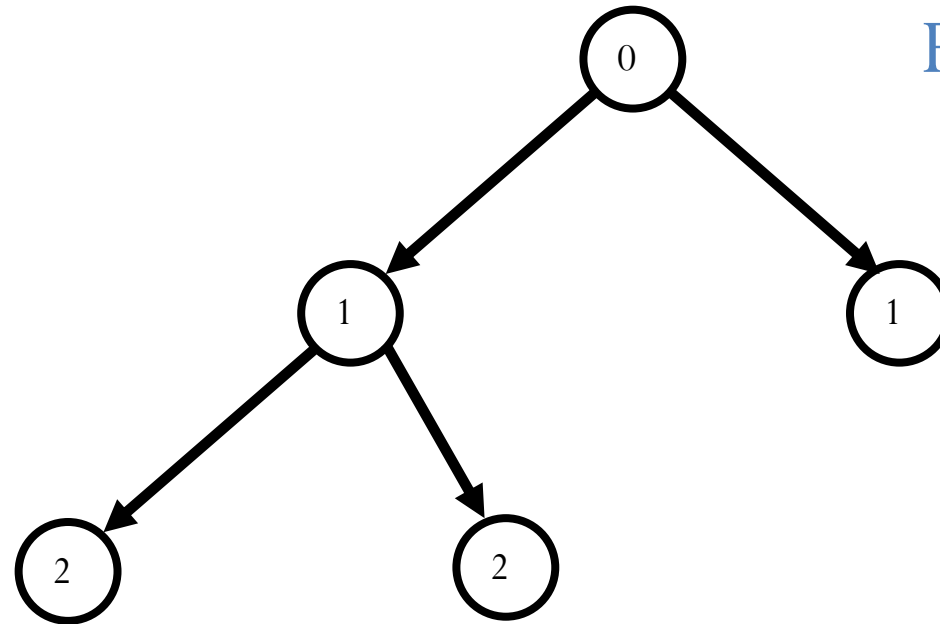
First-In-First-Out (FIFO)  
Queue

# Breadth-First Search



First-In-First-Out (FIFO)  
Queue

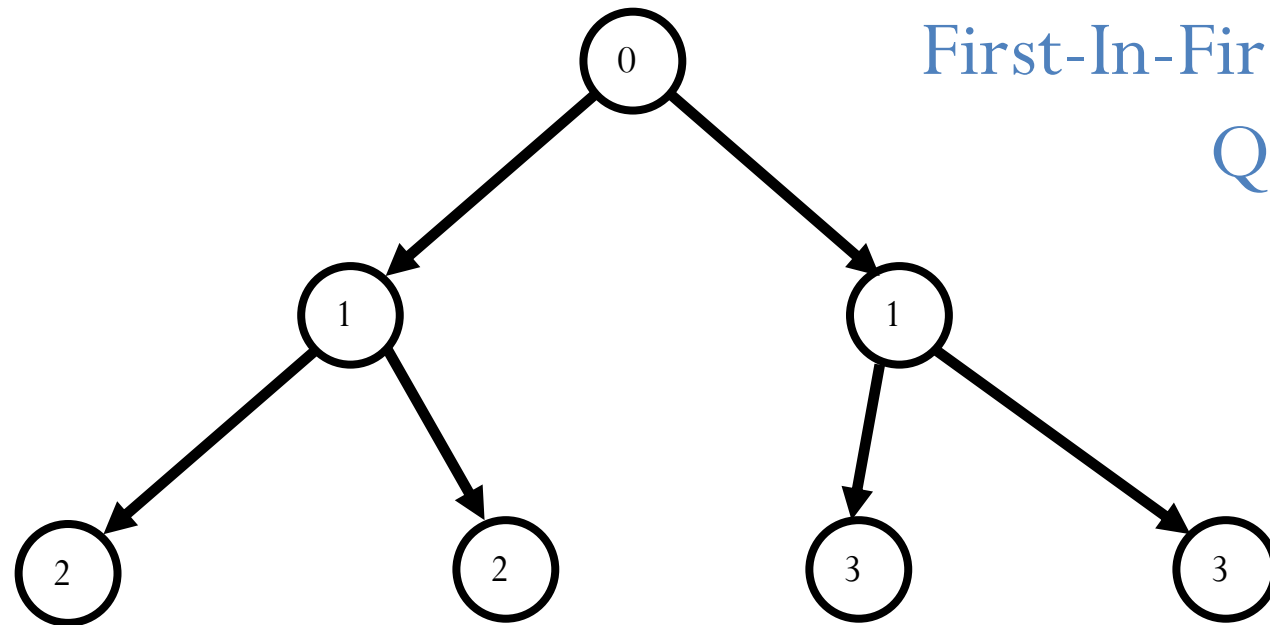
# Breadth-First Search



First-In-First-Out (FIFO)  
Queue

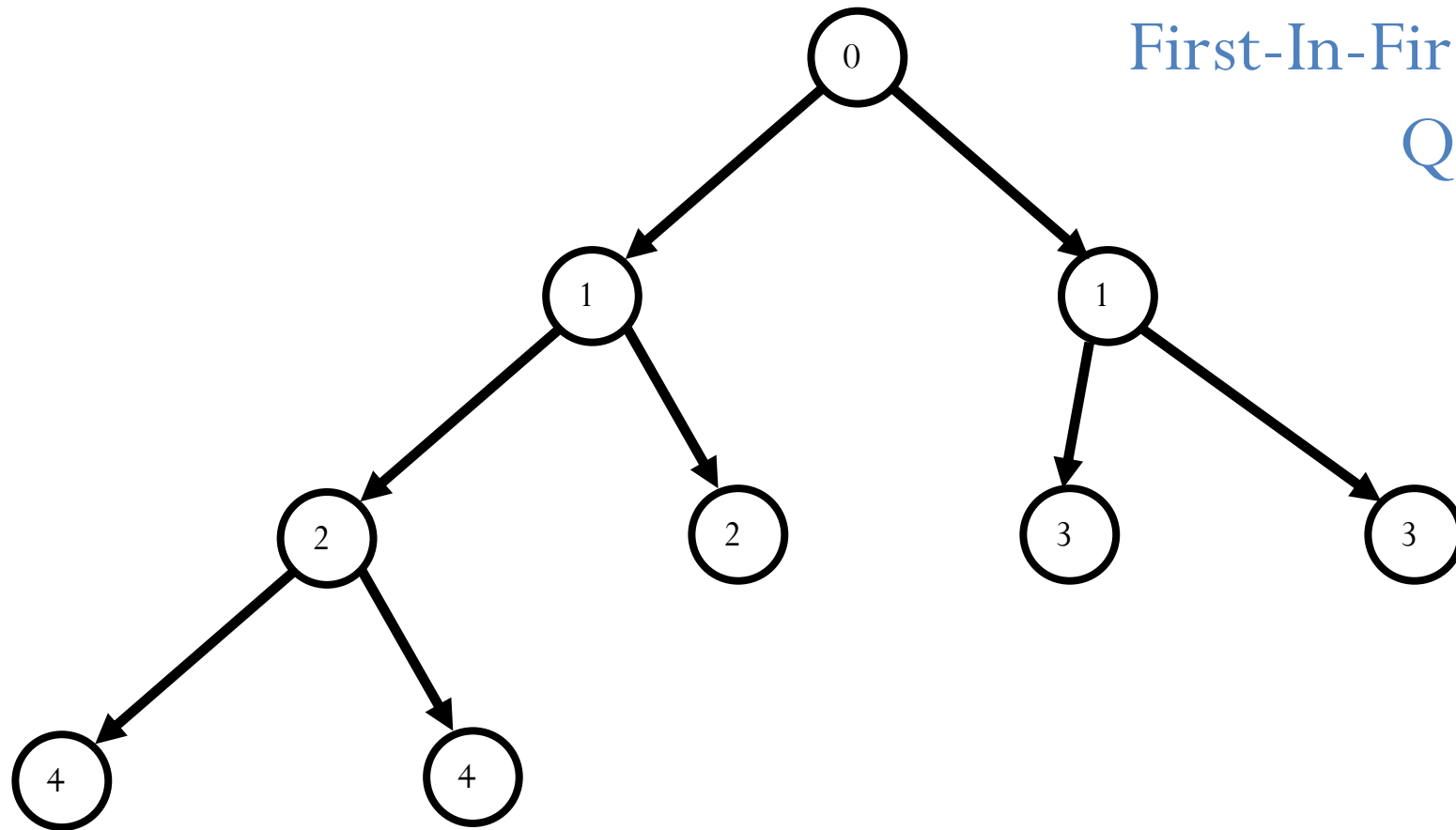


# Breadth-First Search



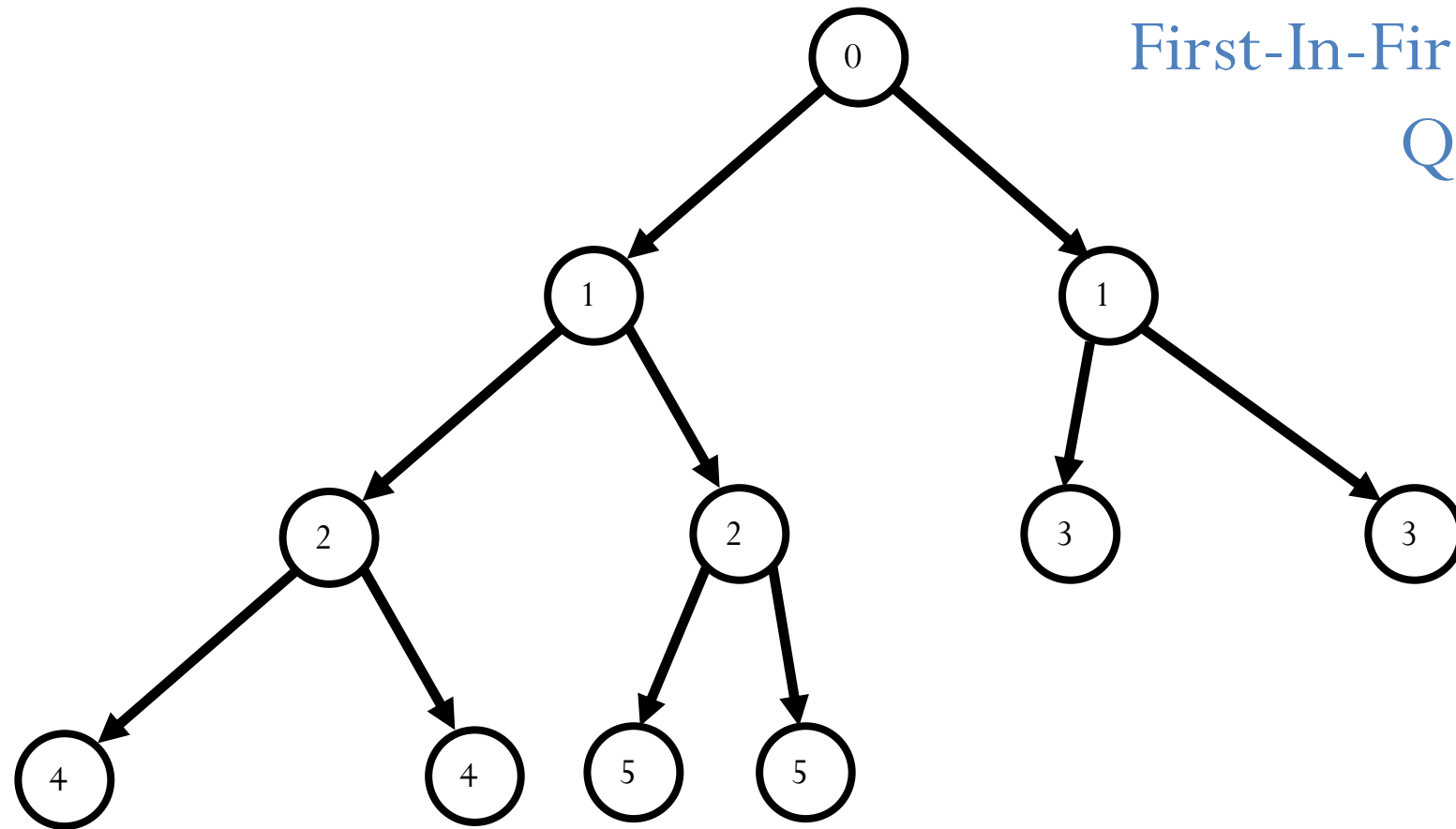
First-In-First-Out (FIFO)  
Queue

# Breadth-First Search



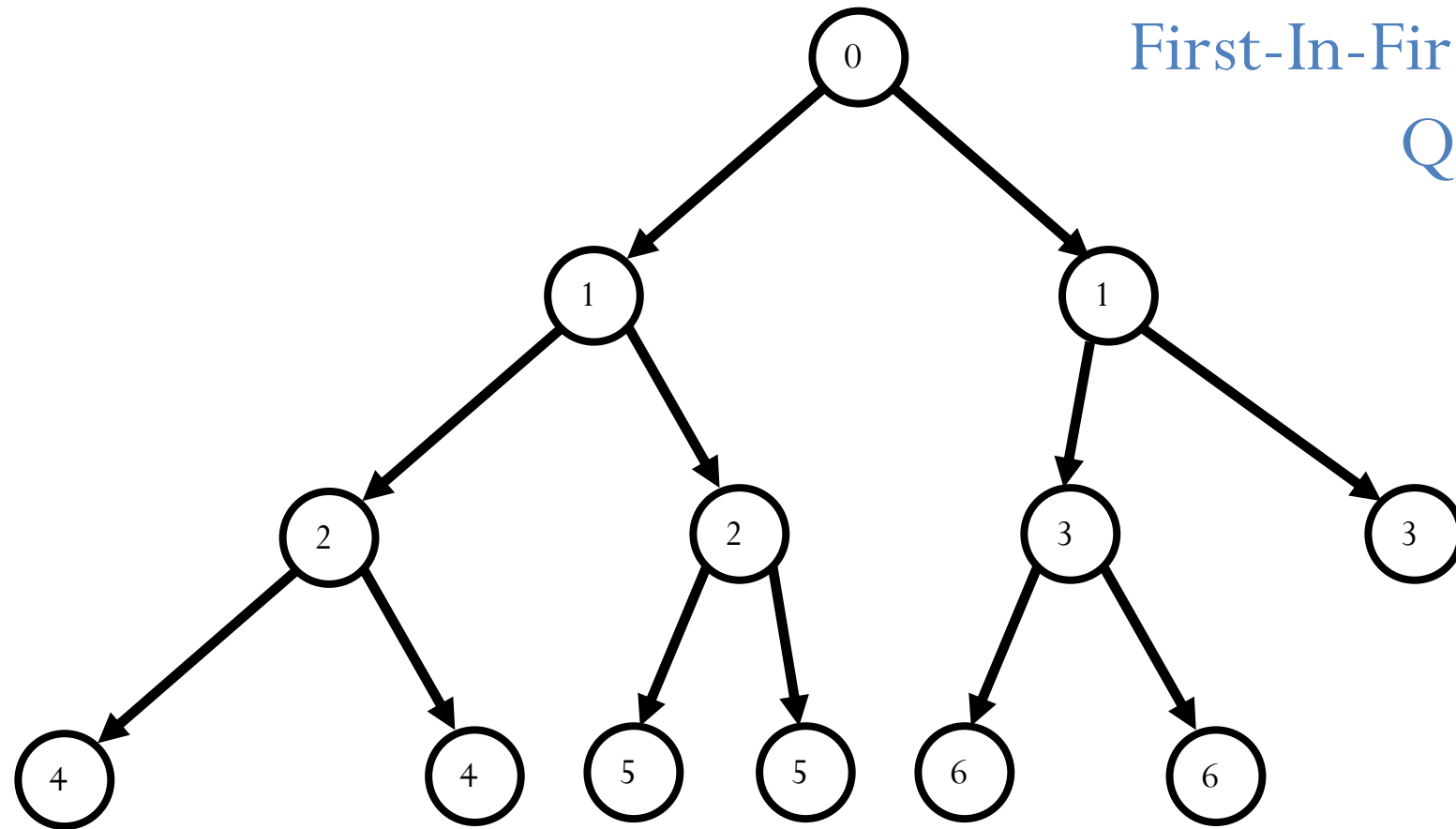
First-In-First-Out (FIFO)  
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# Breadth-First Search



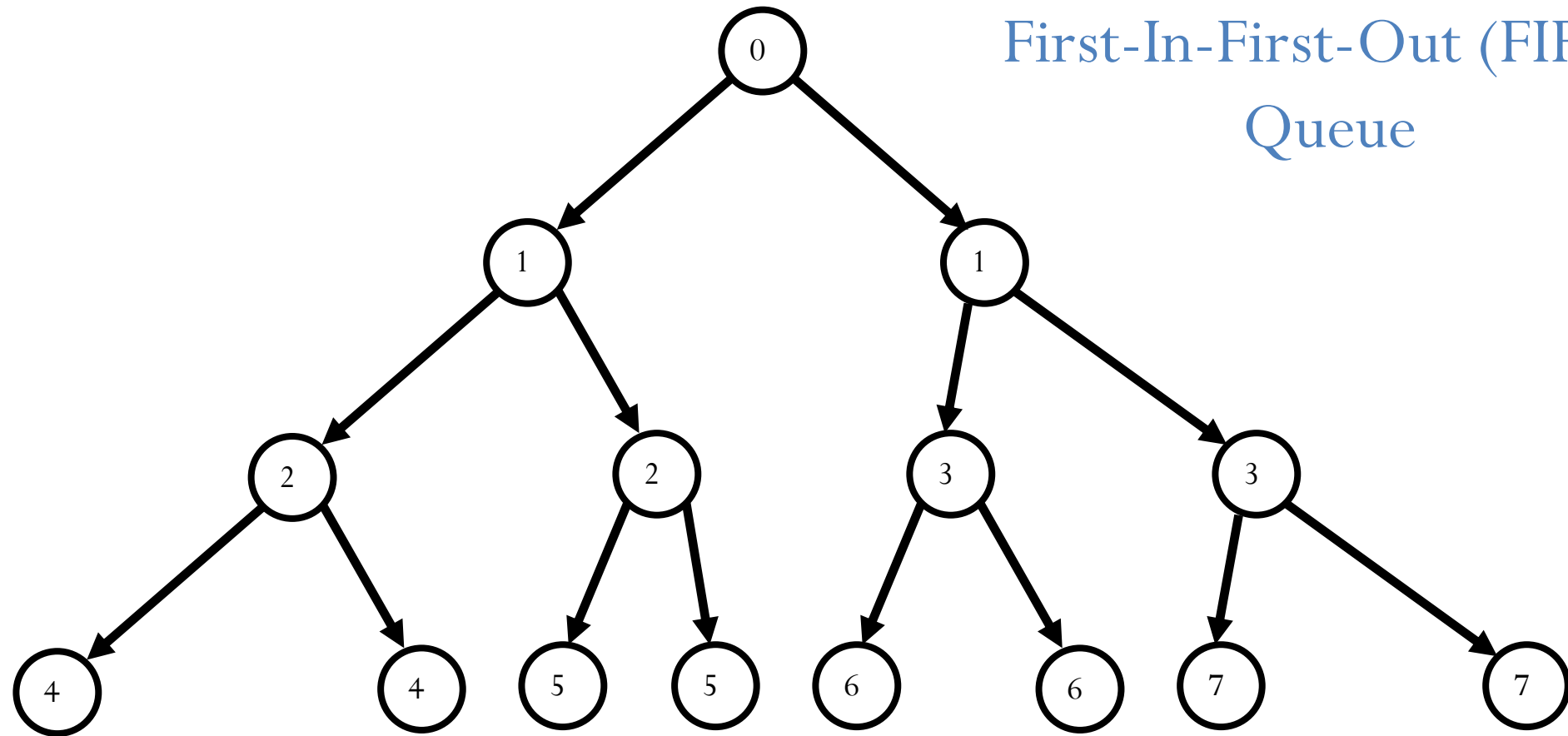
First-In-First-Out (FIFO)  
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# Breadth-First Search

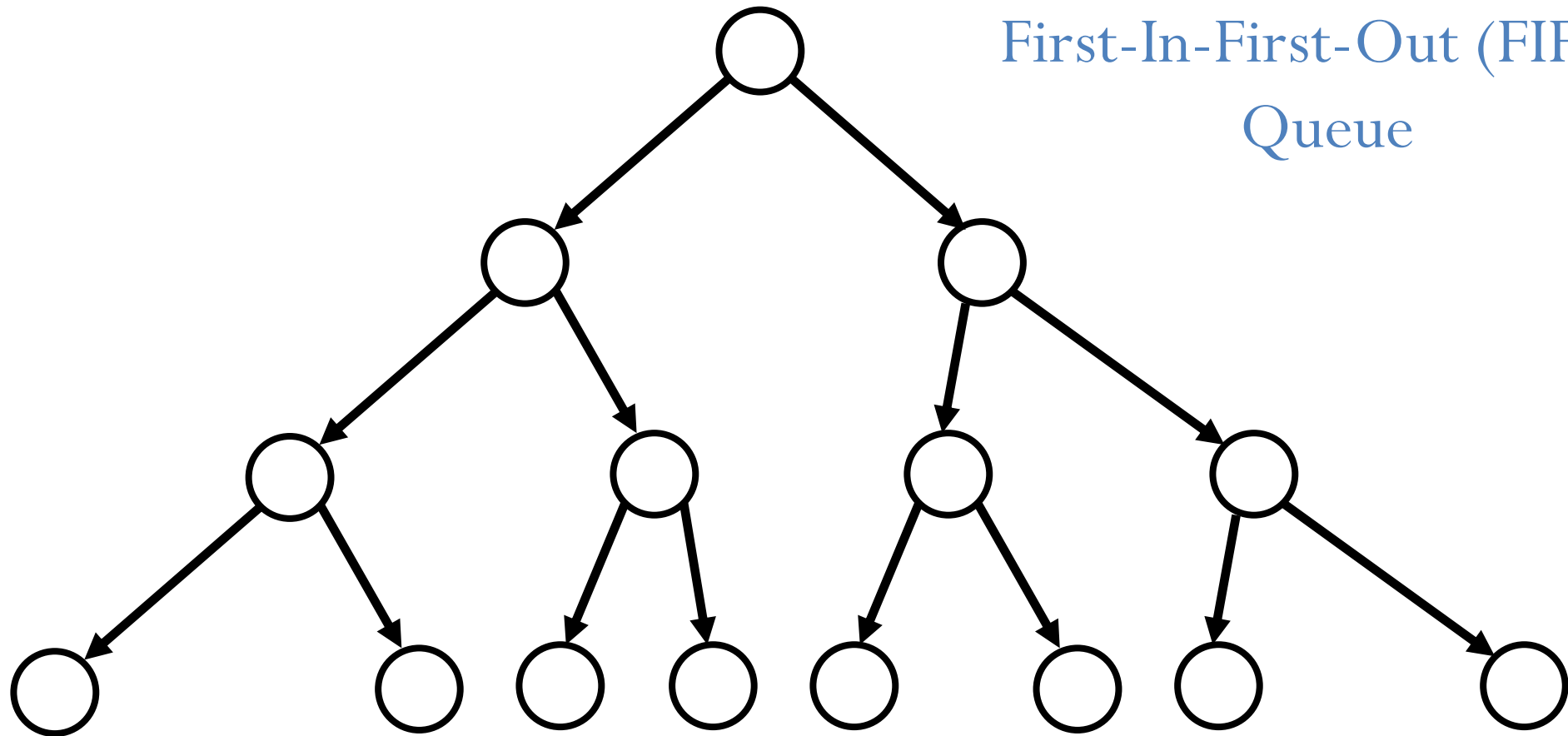


First-In-First-Out (FIFO)  
Queue

# Breadth-First Search



# Breadth-First Search

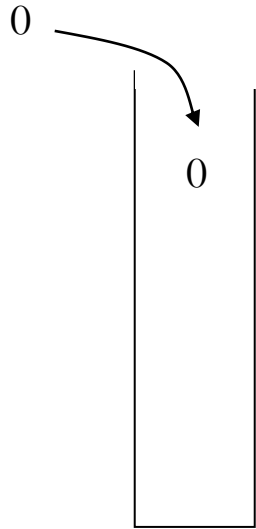


First-In-First-Out (FIFO)  
Queue

# Properties of Breadth-First Search

- Nodes are expanded in the same order in which they are generated
  - Fringe can be maintained as a First-In-First-Out (FIFO) queue
- BFS is **complete**: if a solution exists, one will be found
- BFS finds a **shallowest** solution
  - Not necessarily an optimal solution
- If every node has  $b$  successors (the **branching factor**), first solution is at depth  $d$ , then fringe size will be at least  $b^d$  at some point
  - This much space (and time) required ☹

# Depth-First Search

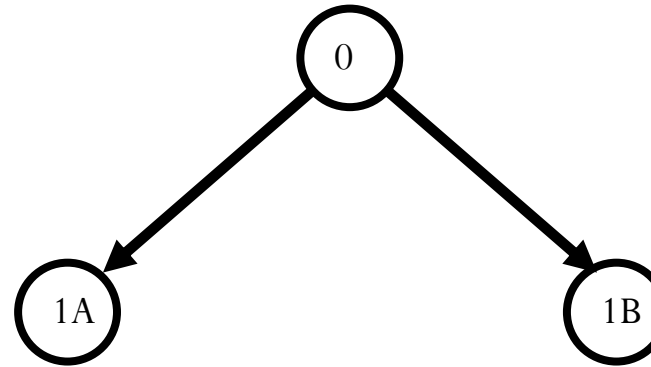
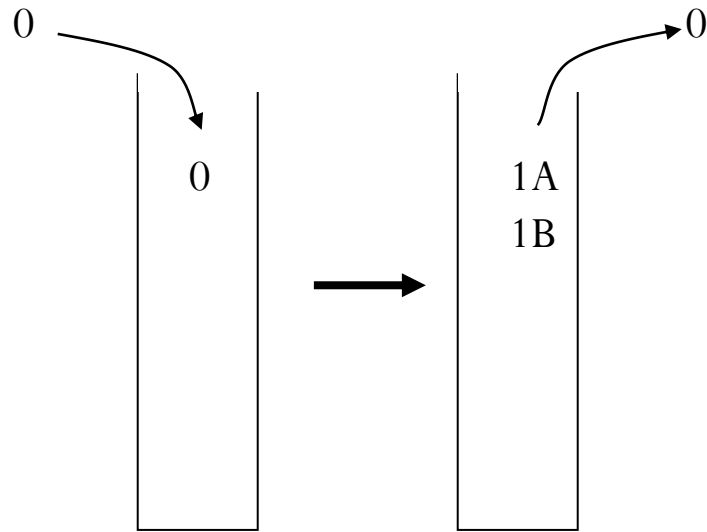


0

Last-In-First-Out (LIFO)  
Stack

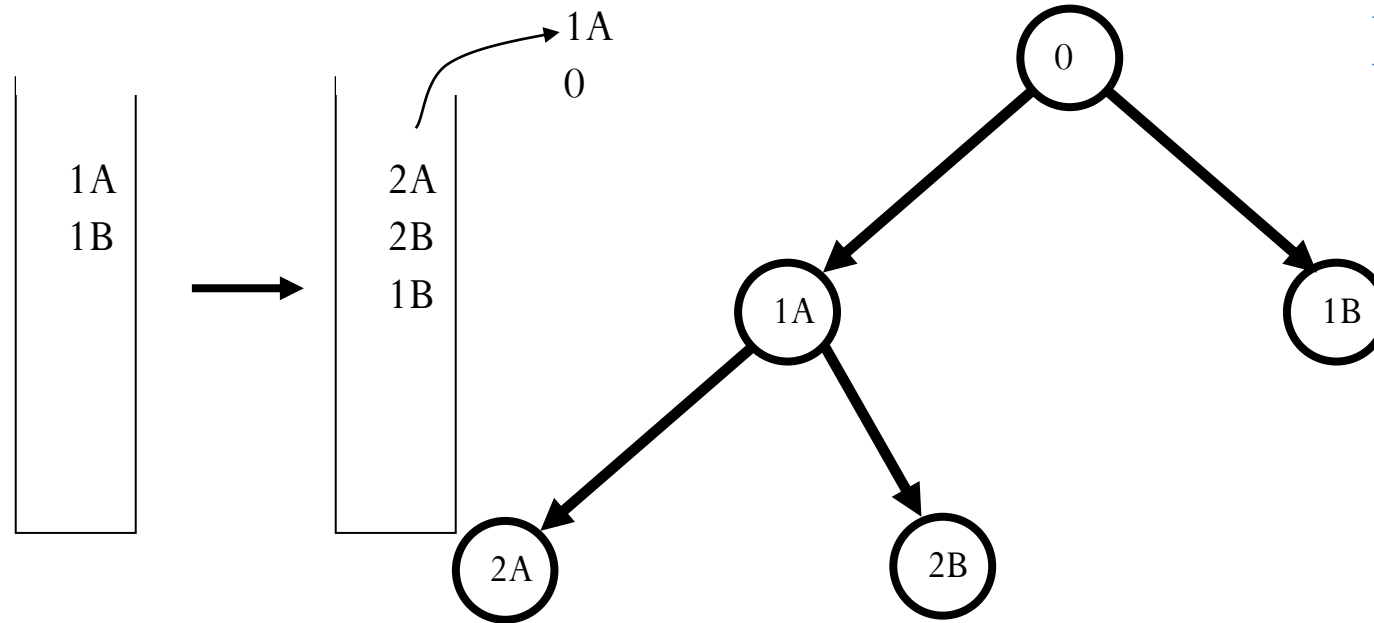


# Depth-First Search



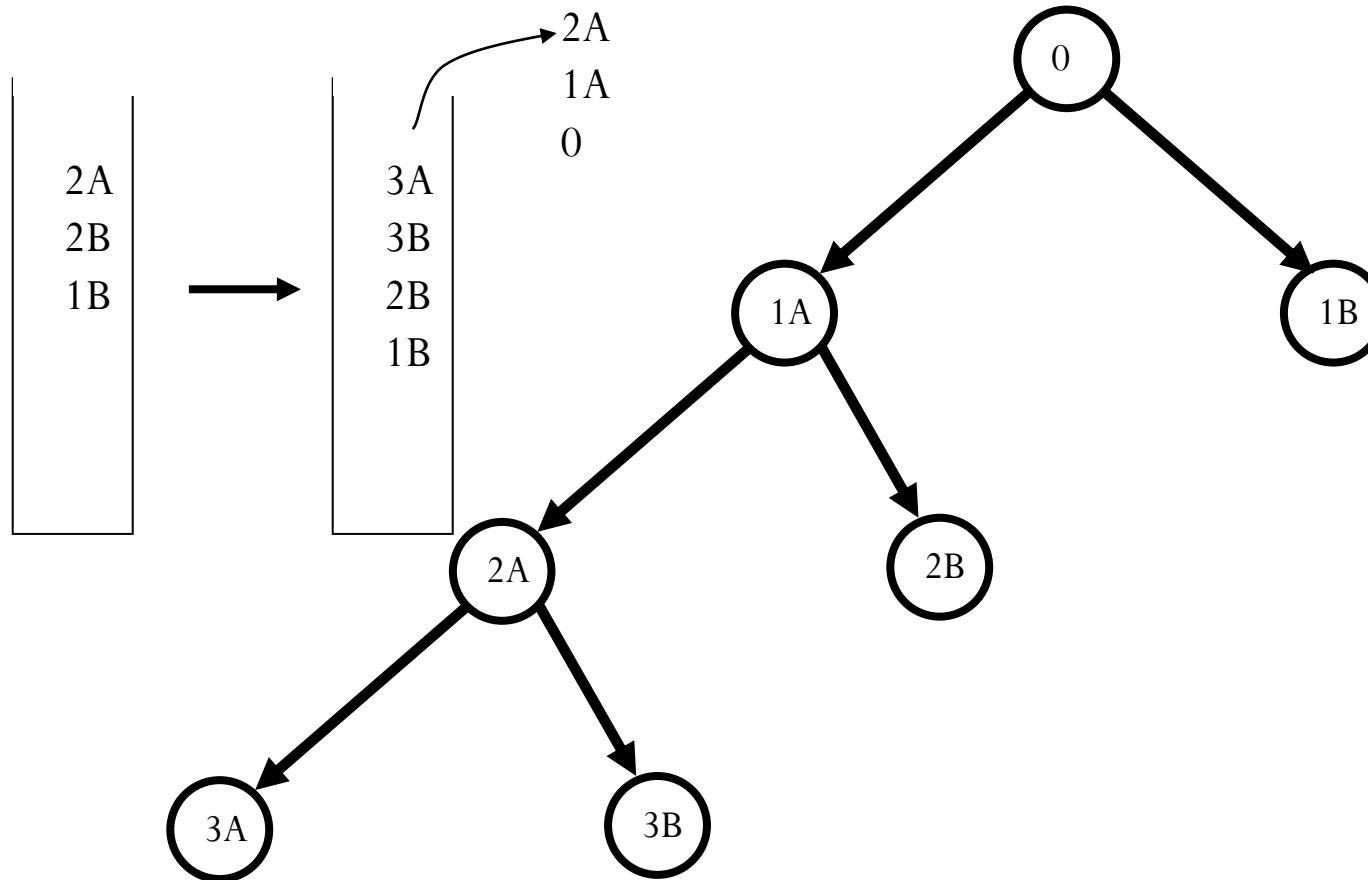
Last-In-First-Out (LIFO)  
Stack

# Depth-First Search



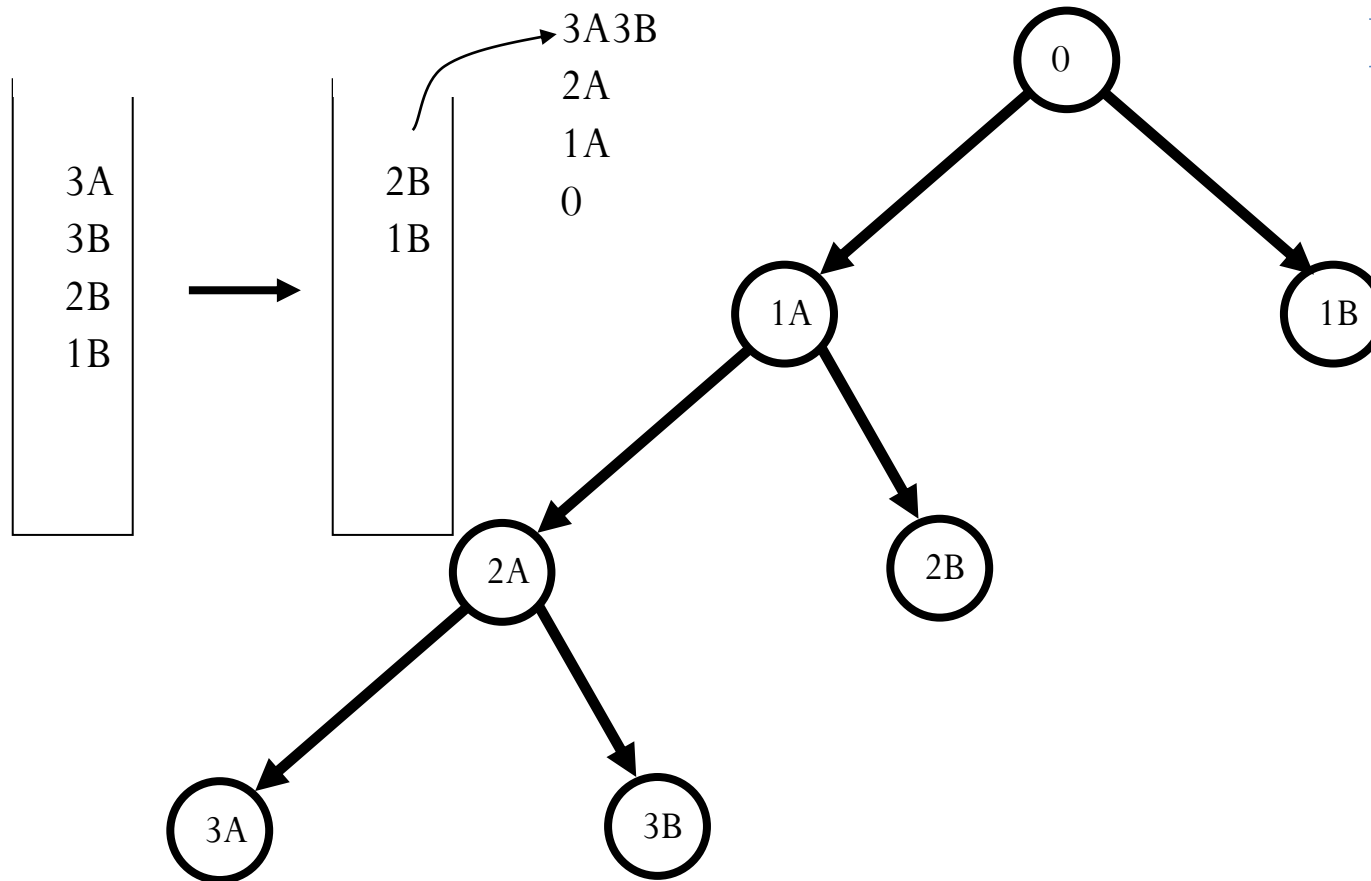
Last-In-First-Out (LIFO)  
Stack

# Depth-First Search



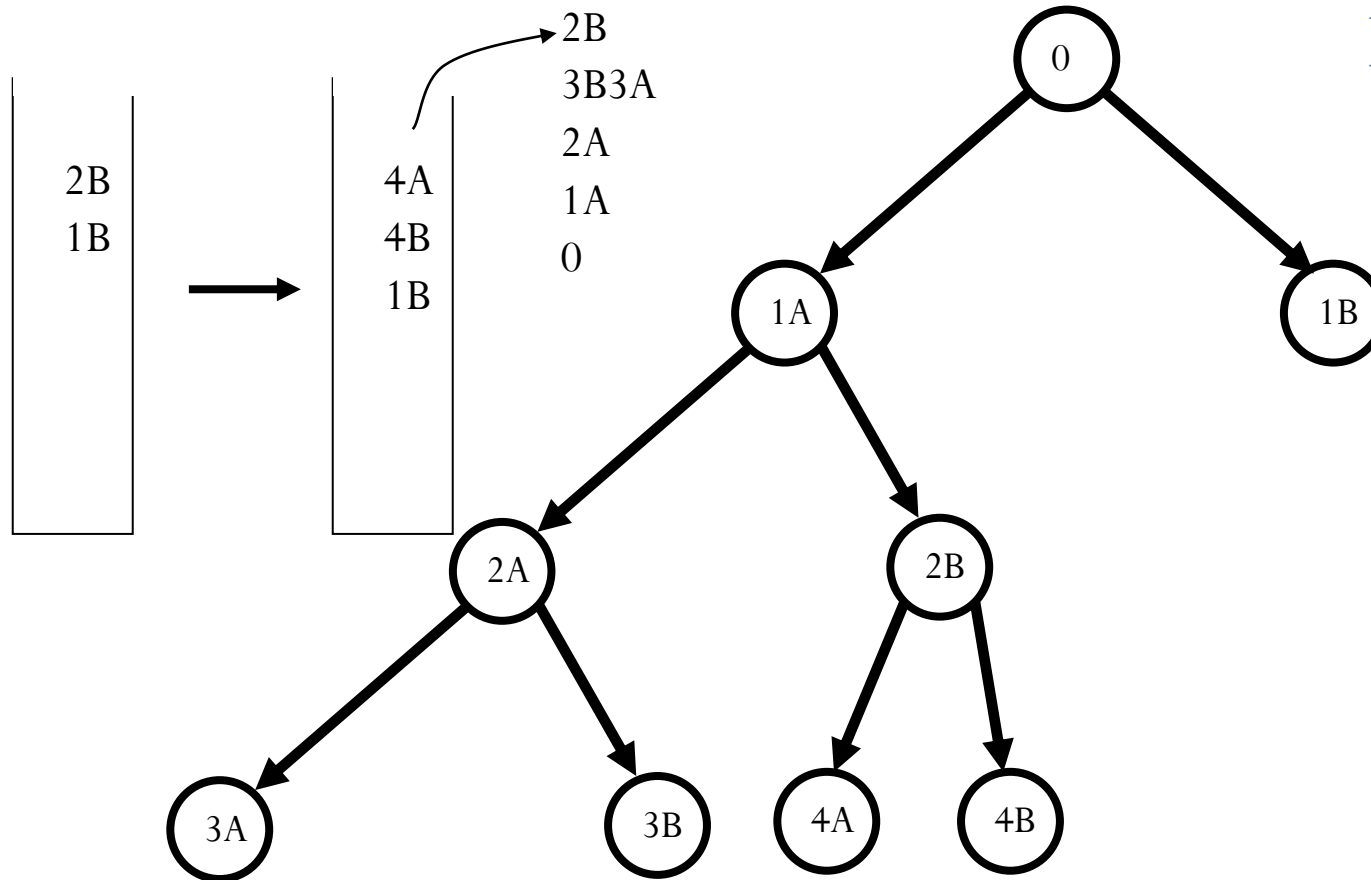
Last-In-First-Out (LIFO)  
Stack

# Depth-First Search



Last-In-First-Out (LIFO)  
Stack

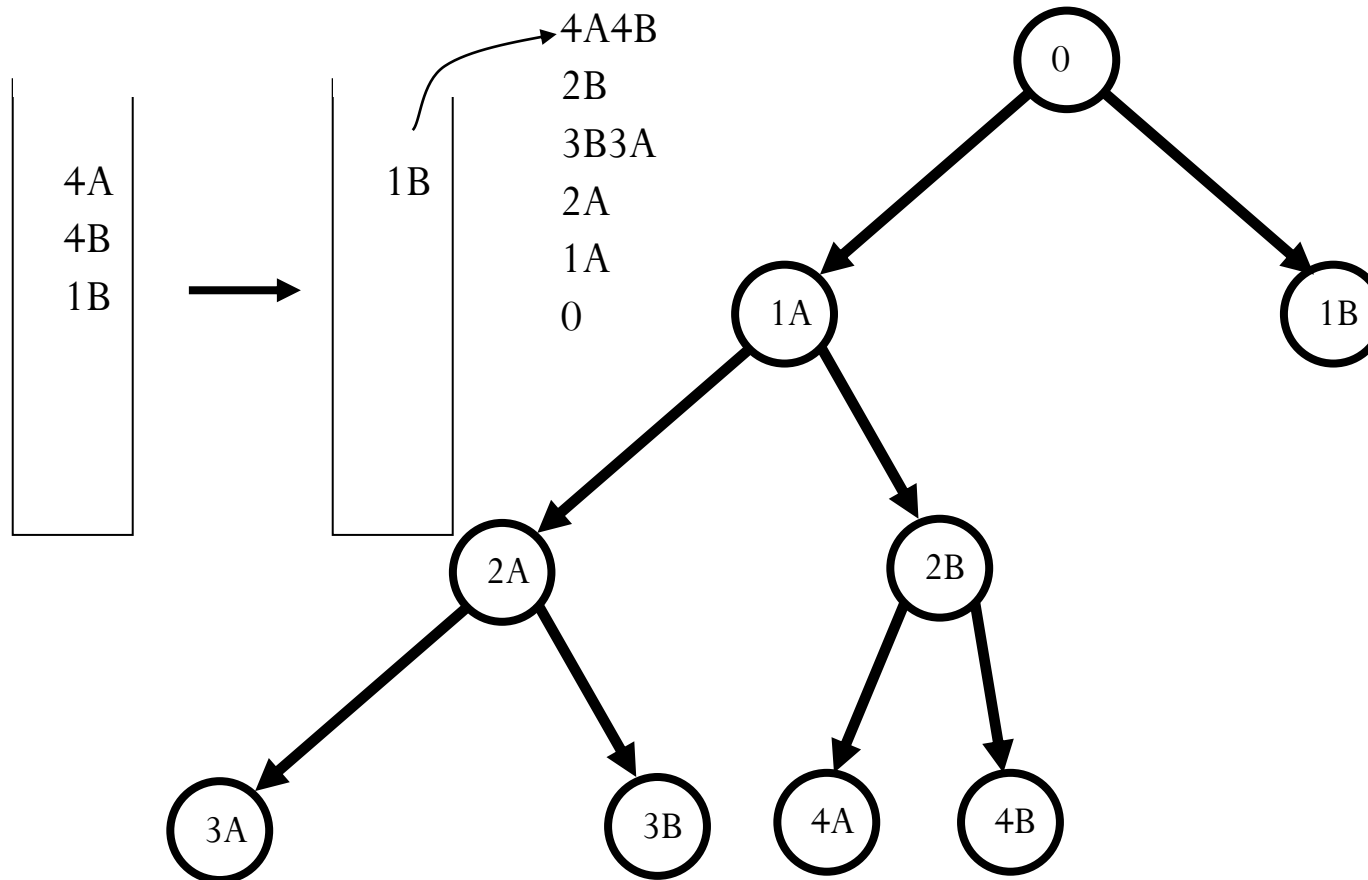
# Depth-First Search



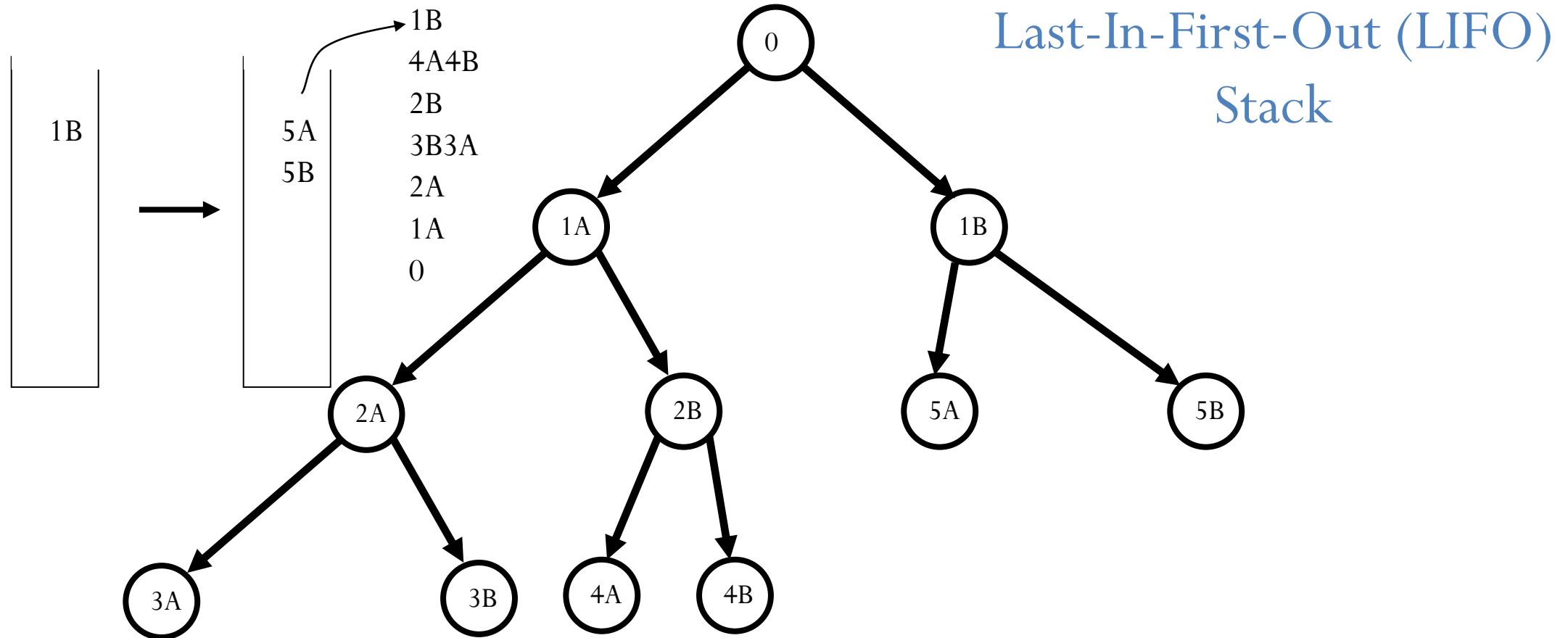
Last-In-First-Out (LIFO)  
Stack

# Depth-First Search

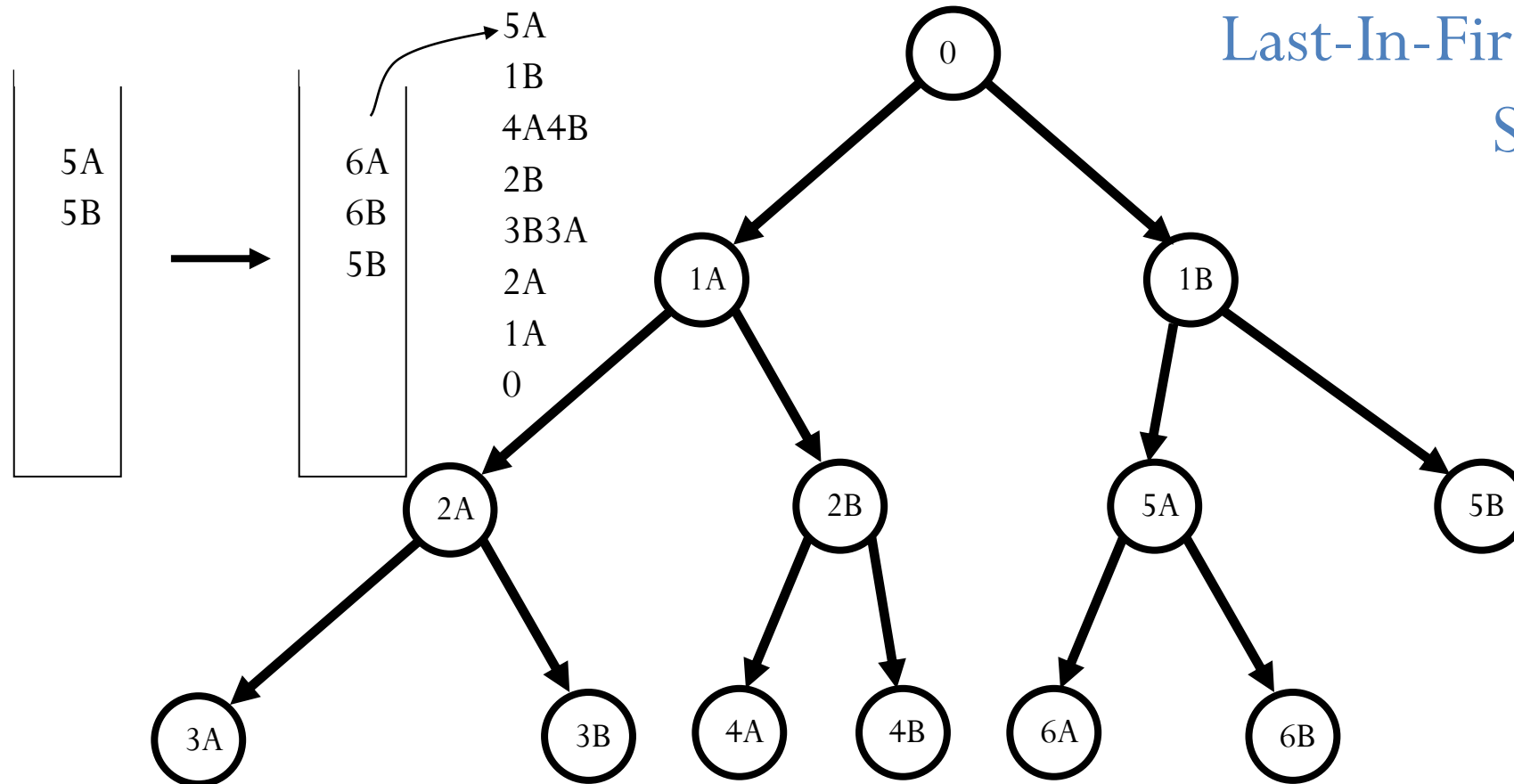
Last-In-First-Out (LIFO)  
Stack



# Depth-First Search

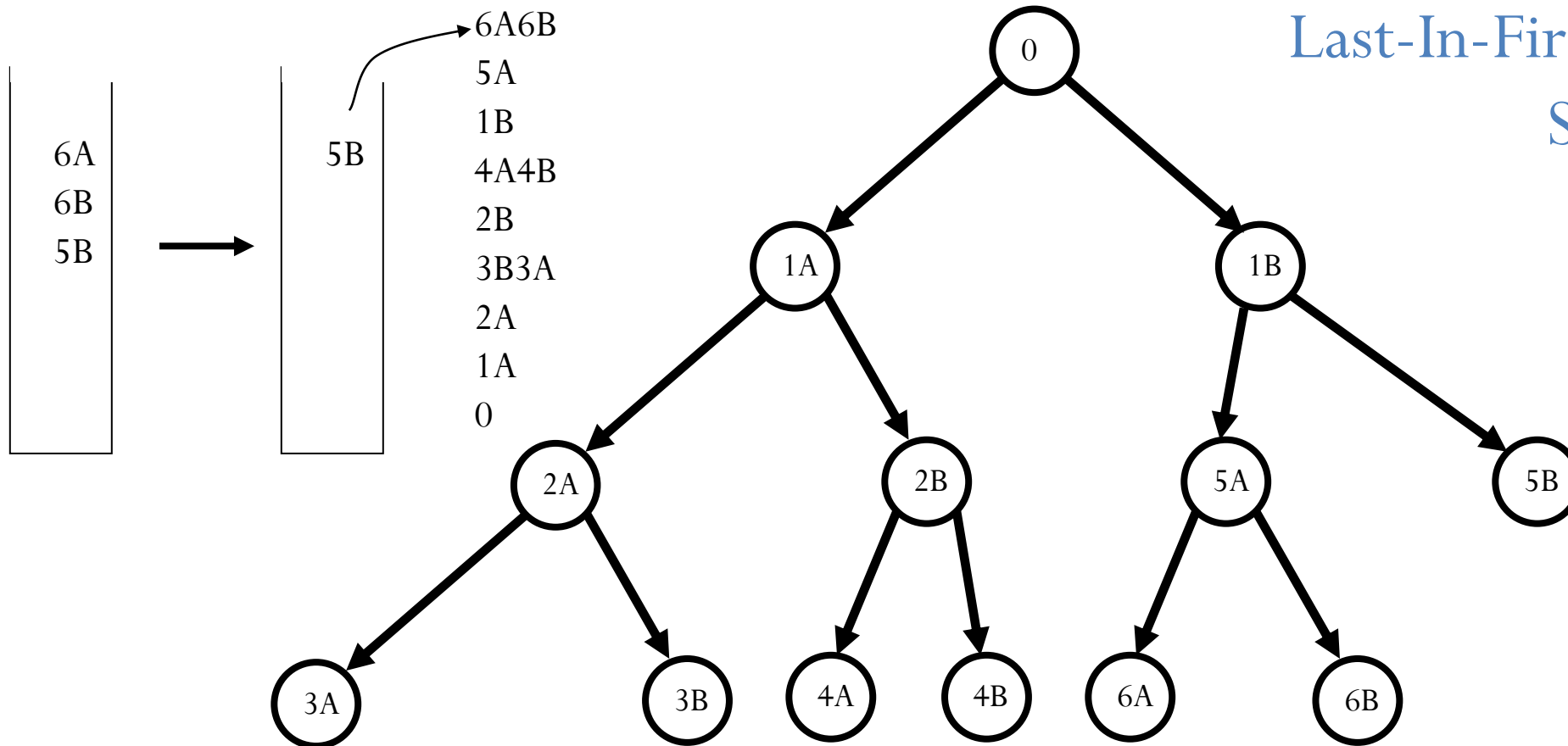


# Depth-First Search





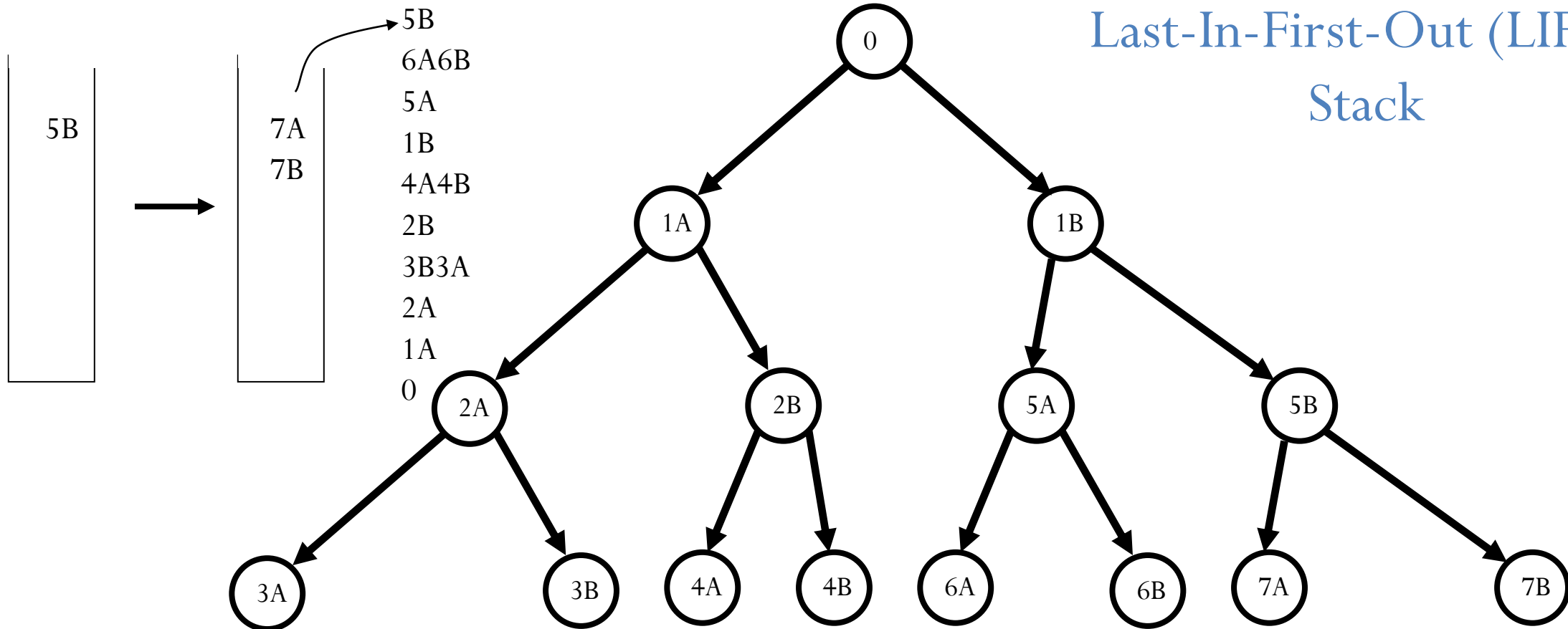
# Depth-First Search



Last-In-First-Out (LIFO)  
Stack

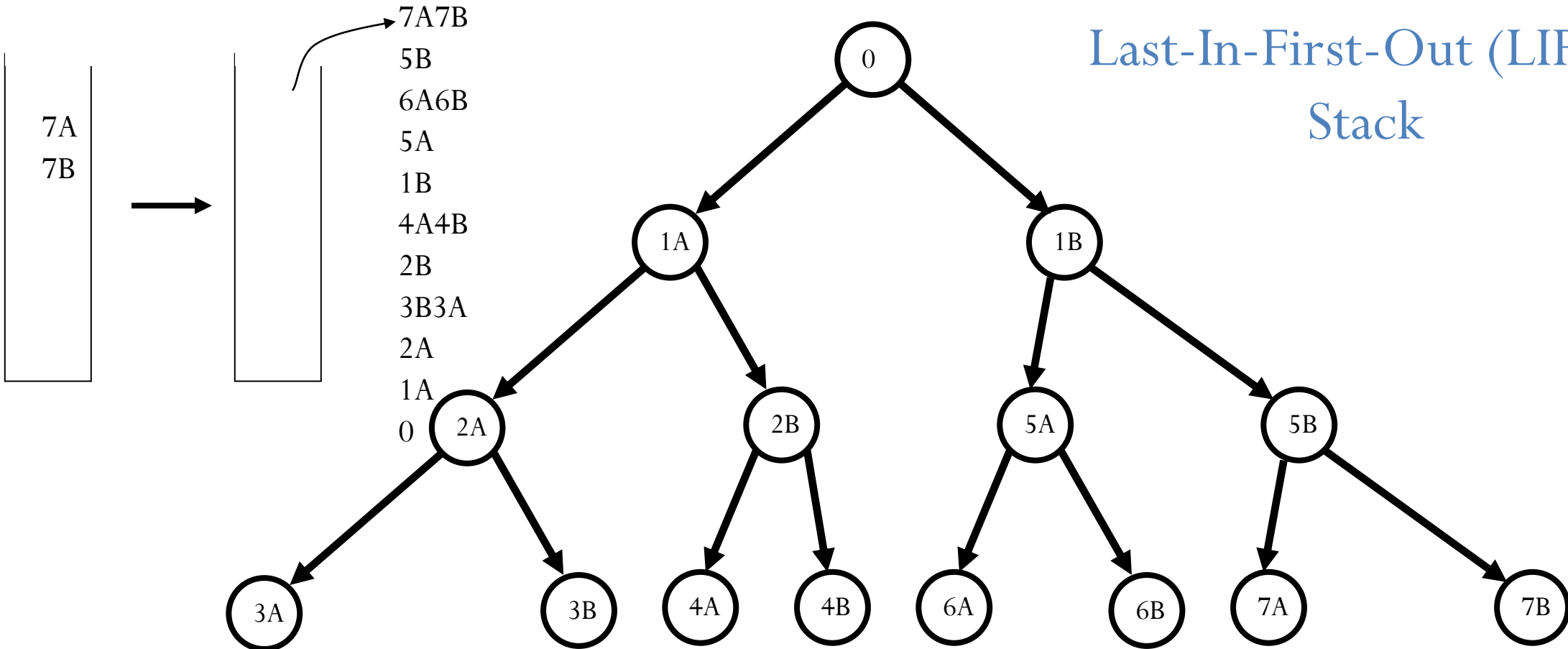
# Depth-First Search

Last-In-First-Out (LIFO)  
Stack

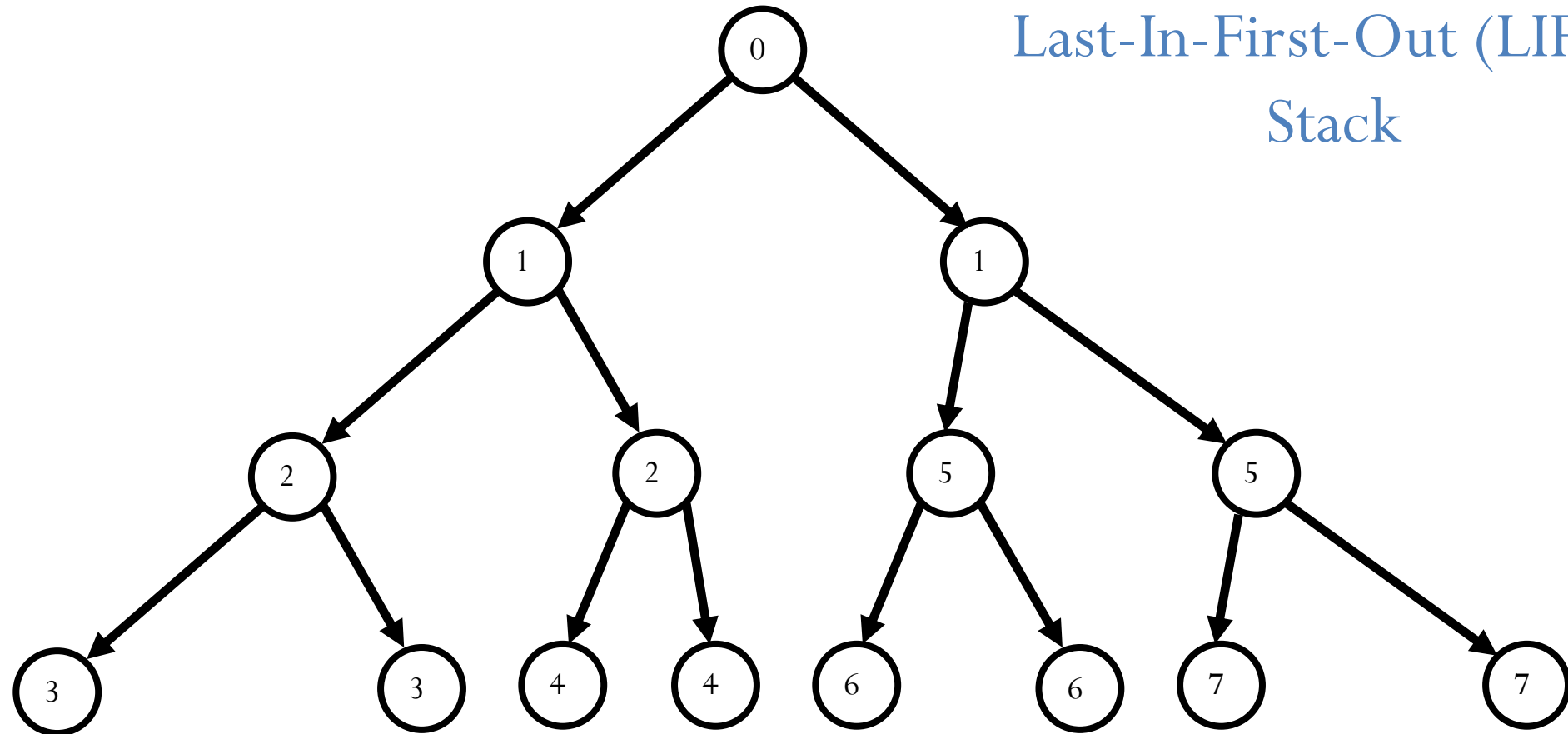


# Depth-First Search

Last-In-First-Out (LIFO)  
Stack



# Depth-First Search



Last-In-First-Out (LIFO)  
Stack

# Implementing Depth-First Search

- Fringe can be maintained as a Last-In-First-Out (LIFO) queue (aka. a stack)
- Also easy to implement recursively:
- DFS(node)
  - If goal(node) return solution(node);
  - For each successor of node
    - Return DFS(successor) unless it is *failure*;
  - Return *failure*;

# Properties of depth-first search

- Not complete (might cycle through nongoal states)
- If solution found, generally not optimal/shallowest
- If every node has  $b$  successors (the **branching factor**), and we search to at most depth  $m$ , fringe is at most  $b^m$ 
  - Much better space requirement 😊
  - Actually, generally don't even need to store all of fringe
- Time: still need to look at every node
  - $b^m + b^{m-1} + \dots + 1$  (for  $b > 1$ ,  $O(b^m)$ )
  - **Inevitable** for uninformed search methods...

# Combining good properties of BFS and DFS

- **Limited depth DFS:** just like DFS, except never go deeper than some depth  $d$
- **Iterative deepening DFS:**
  - Call limited depth DFS with depth 0;
  - If unsuccessful, call with depth 1;
  - If unsuccessful, call with depth 2;
  - Etc.
- Complete, finds shallowest solution
- Space requirements of DFS
- May seem wasteful timewise because replicating effort
  - Really not that wasteful because **almost all effort at deepest level**
  - $db + (d-1)b^2 + (d-2)b^3 + \dots + 1b^d$  is  $O(b^d)$  for  $b > 1$

# Searching solution evaluation

- Comparing multiple searching algorithm based on
  - Completeness: does it always find a solution if one exist?
  - Time complexity: How long depends on number of nodes
  - Space complexity: Memory depends on number of nodes
  - Optimality: Find shortest path (or least cost solution)?
  - Systematicity: does it visit each state at most once?



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