

# Basic Programming and Data Structures

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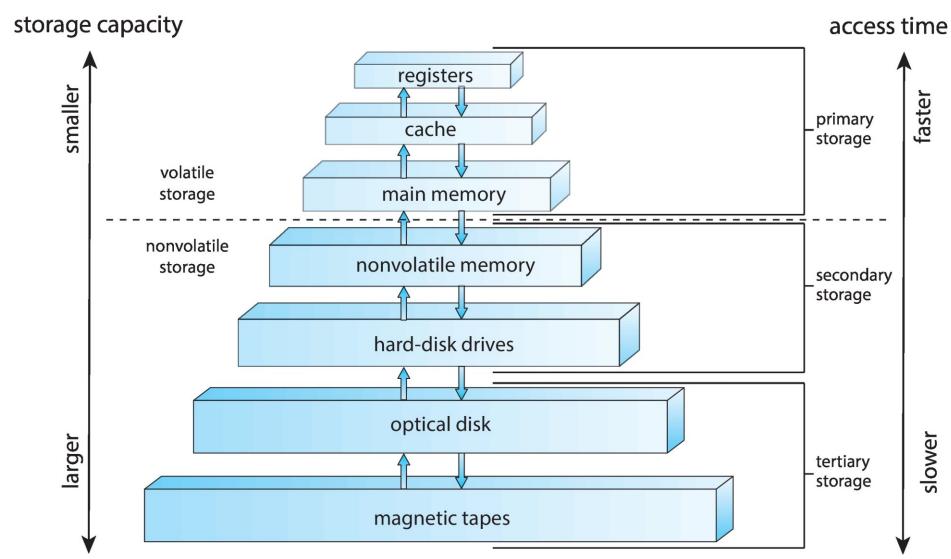
# Goals of the Course

- To become familiar with Data structures and Algorithms
- To develop ability of problem solving skill in programming
- To understand that Data structures are the building blocks of larger softwares
- To develop ability for analyzing existing Data structures and Algorithms
- To develop skills with the C, C++, Java for Data structures and Algorithms

- "Get your data structures correct first, and the rest of the program will write itself."
  - David Jones

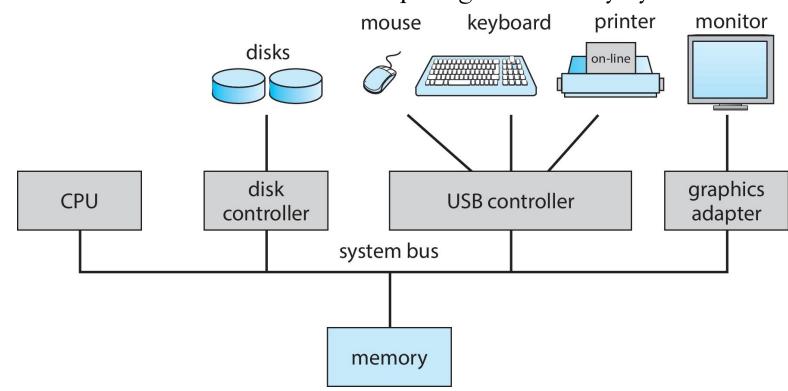
# Basic Program Execution on CPU and Memory

# Data Storage Device Hierarchy



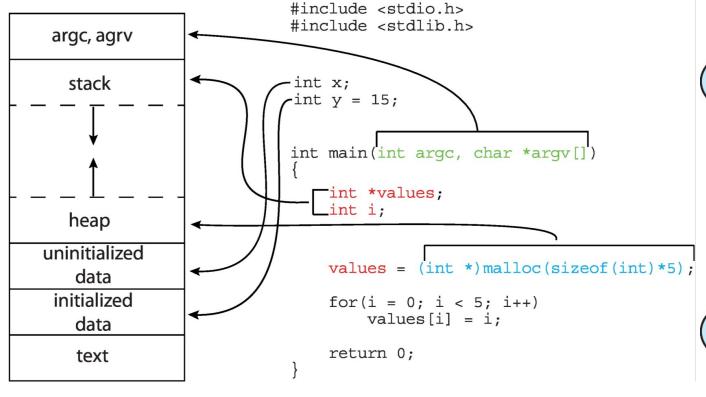
# Computer System Organization

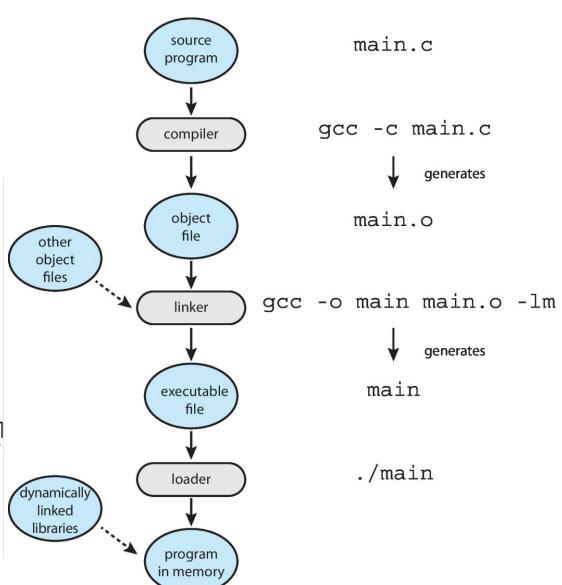
- Computer-system operation
  - One or more CPUs, device controllers connect through common **bus** providing access to shared memory
  - Concurrent execution of CPUs and devices competing for memory cycles



# Program to Process

• When you run an exe file, the OS creates a process = a running program

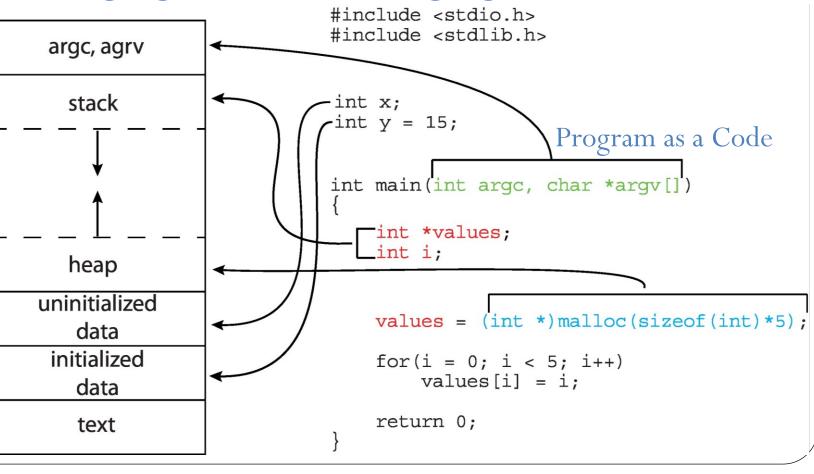




# Motivation: Program Code to Memory

- Abstraction of complex usage Program as a Memory (RAM or Cache).
- Conversion of High level language to Low level language
- Static/global variables are allocated in the executable
  - Local variables of a function on Stack
  - Dynamic allocation with malloc on the heap

Process as a Memory



# Program to Process

 Virtual address space is setup by OS during process creation

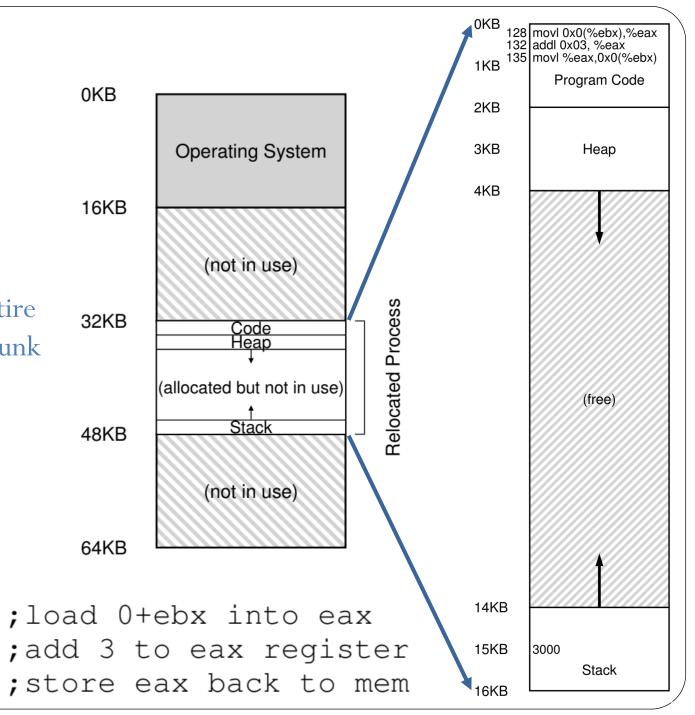
Simplified OS: places entire memory image in one chunk

```
void func() {
   int x = 3000;
   x = x + 3;
   ... Compiler
```

135: movl %eax, 0x0(%ebx)

128: movl 0x0(%ebx), %eax

132: addl \$0x03, %eax

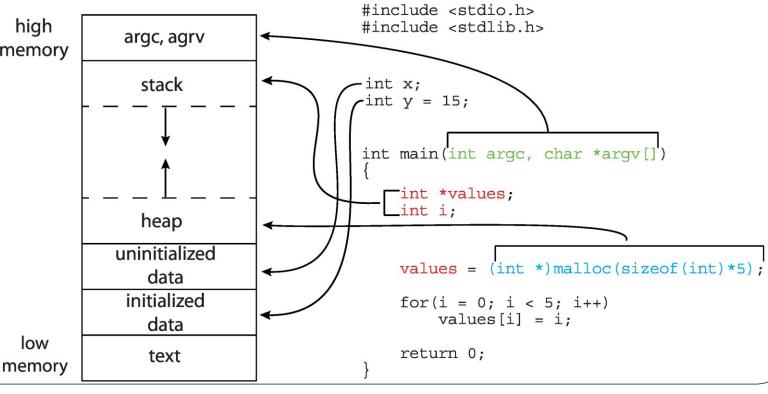


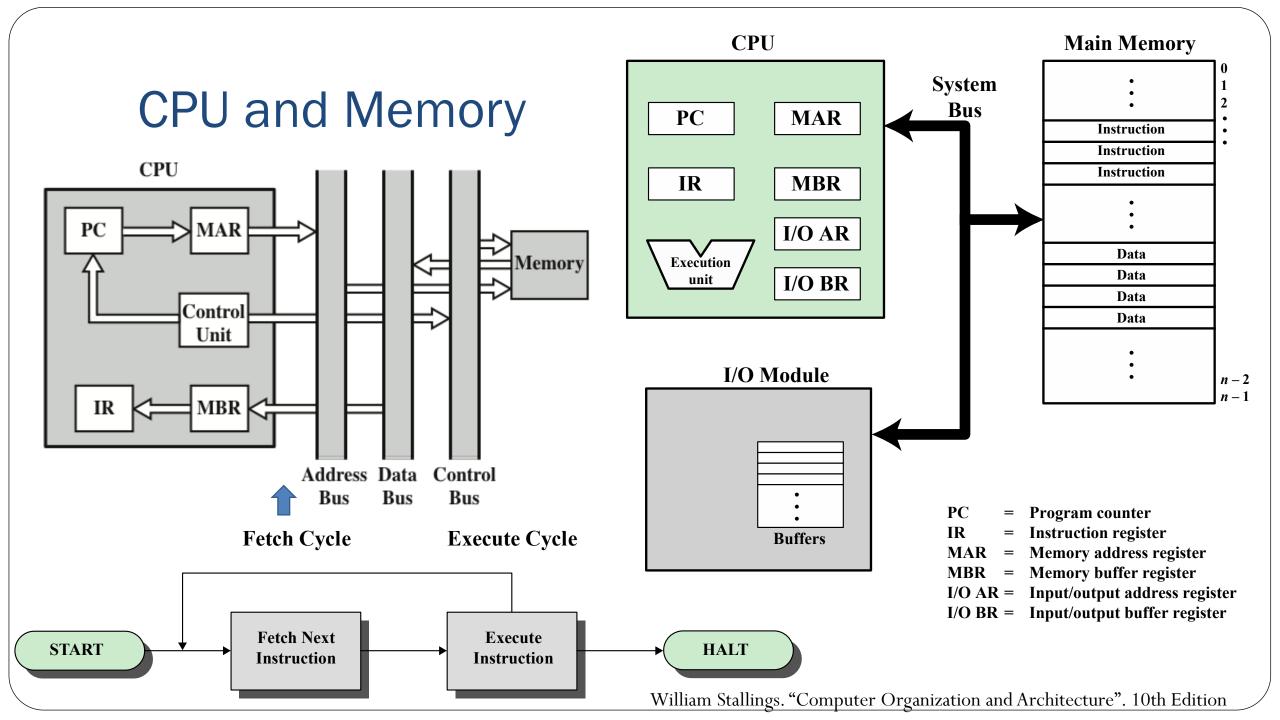
# **Program and Process**

- A unique identifier
- Points CPU program counter to current instruction — Other registers may store operands, return values etc. high memory

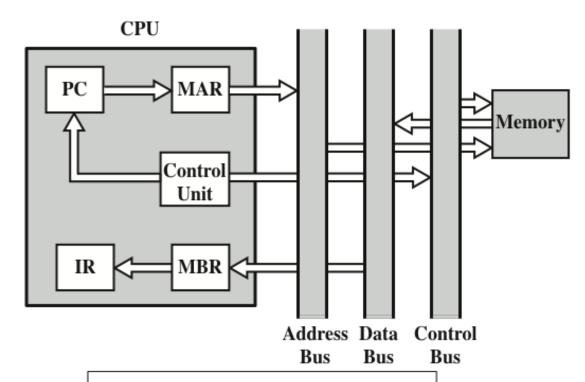
low

- CPU context: registers
  - Program counter
  - Current operands
  - Stack pointer





# **CPU** and **Memory**

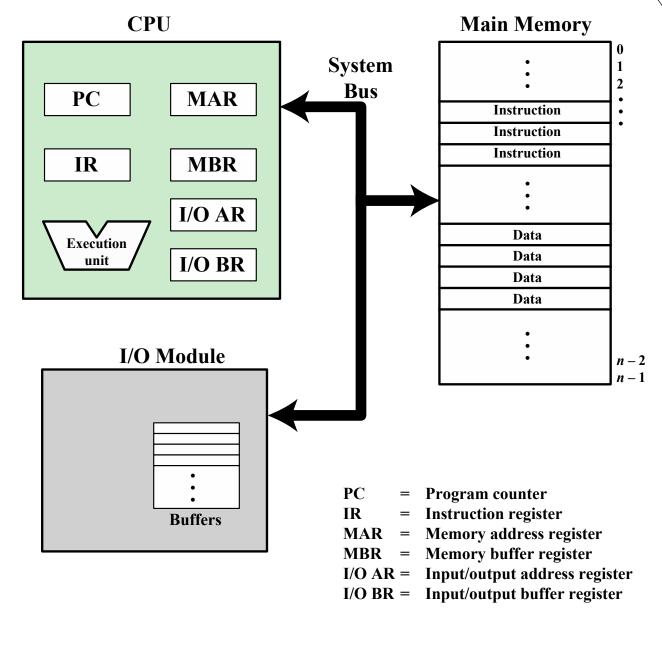


t1 : MAR ← PC

t2 : MBR ← MEMORY

PC ← (PC) + I

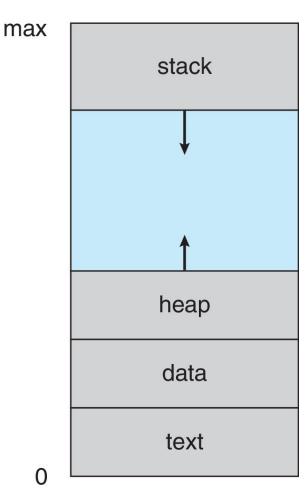
t3 : IR ← (MBR)



William Stallings. "Computer Organization and Architecture". 10th Edition

# **Program and Process**

- OS allocates memory and creates memory image
  - Loads code, data from disk exe
  - Creates runtime stack, heap
  - Opens basic files STD IN, OUT, ERR
  - Initializes CPU registers PC points to first instruction
- Memory image
  - Code & data (static)
  - Stack and heap (dynamic)



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# Basic Data Structures

# Programing (Coding)

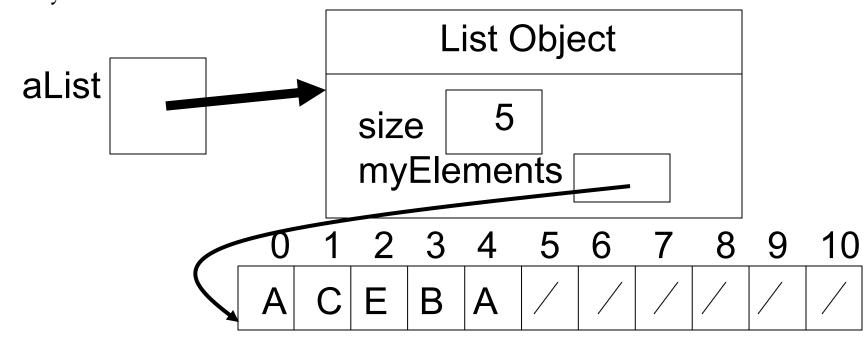
- All programs (or process) has data
  - To store, display, gather
  - In the form of information, numbers, images, sound
- Programmer decide to structure data
- Several option leads to program development
- Program properties
  - execution speed
  - memory requirements
  - maintenance (debugging, extending, etc.)

# **Code Implementation**

- Theoretically
  - abstract base class describes ADT
  - inherited implementations implement data structures
  - can change data structures transparently (to client code)
- Practice
  - different implementations sometimes suggest different interfaces (generality vs. simplicity)
  - performance of a data structure may influence form of client code (time vs. space, one operation vs. another)

# **Data Structures**

- A Data Structure is:
  - an implementation of an abstract data type and
  - "An organization of information, usually in computer memory", for better algorithm efficiency."



# Data structures

### Ideal data structure:

fast, elegant, memory efficient

### Generates tensions:

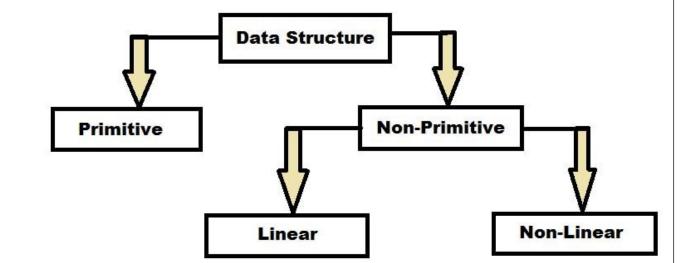
- time vs. space
- performance vs. elegance
- generality vs. simplicity
- one operation's performance vs. another's

## Dictionary ADT

- list
- binary search tree
- AVL tree
- Splay tree
- Red-Black tree
- hash table

## Data structures

- Data Structures are containers:
  - they hold other data
  - arrays are a data structure
  - ... so are lists
- Other types of data structures:
  - stack, queue, tree, binary search tree, hash table, dictionary or map, set, and on and on
  - www.nist.gov/dads/
  - en.wikipedia.org/wiki/List\_of\_data\_structures
- Different types of data structures are optimized for certain types of operations



# Data Structures Operations

- Data Structures will have 3 core operations
  - a way to add things
  - a way to remove things
  - a way to access things
- Details of these operations depend on the data structure
  - Example: List, add at the end, access by location, remove by location
- More operations added depending on what data structure is designed to do

# Built-in Data Type

• There are five basic data types associated with variables

```
- integer: a whole number
int
• float - floating point value: a number with a fractional part
• double - a double-precision floating point value
• char - a single character
• void - valueless special purpose type
• int a = 4000;
                              // positive integer data type
• float b = 5.2324;
                              // float data type
• char c = 'Z';
                              // char data type
• long d = 41657;
                             // long positive integer data type
• long e = -21556;
                              // long -ve integer data type
• int f = -185;
                              // -ve integer data type
• short g = 130;
                              // short +ve integer data type
• short h = -130;
                             // short -ve integer data type
• double i = 4.1234567890;
                              // double float data type
• float j = -3.55;
                              // float data type
```

# Built-in Data Type

### Advantage

- Simple: Really simple! Only FIVE types of data!!
- Easy to handle: Allocation of memory and operations are already defined
- Built-in support by programming language
  - The C library are there to deal with them

### Limitations

- There is a need for storing and handling variety of data types: Image, text, video, etc.
- Limited range
- Waste of memory
- No flexibility
- Error prone programming

# **User-Defined Data Types**

- User can define their own data type
- Also, called Custom Data Type, Abstract Data Type (ADT), etc.
- All logically related data can be grouped into a form called structure
- Each member into the group may be a built-in data type or any other user defined data type
- No recursion, that is, a structure cannot include itself
- Examples:

```
Complex number: z = x + i y
```

Matrices:  $A_{\underline{}}(m \times n)$ 

Date: dd/mm/yy

Date: {int dd, int mm, int yy}

# **User-Defined Data Types**

### Advantage

- It is always convenient for handling a group of logically related data items.
- Examples: Student's record, name, roll number, and marks.
- Elements in a set: Used in relational algebra, database, etc.
- A non non-trivial data structure becomes a trivial.
  - Helps in organizing complex data in a more meaningful way

# Abstraction

- Because details of the implementation are hidden.
- When you do some operation on the list, say insert an element, you just call a function.
- Details of how the list is implemented or how the insert function is written is no longer required.

- Algorithm: Description of a step-by-step process to solve a problem
  - independent of High Level Language (HLL)
- Data Structure
  - A set of algorithms which implement an Abstract Data Type (ADT)
  - Data Structures: Arrays, Linked lists, Stacks, Queues, Matrices, Trees, Graphs
  - Usage: Searching, Sorting
- Abstract Data Type
  - An opportunity for an acronym
  - Mathematical description of an object and the set of operations on the object

- Present an ADT
- Motivate with some applications
- Repeat until browned entirely through
  - develop a data structure for the ADT
  - analyze its properties
    - efficiency
    - correctness
    - limitations
    - ease of programming
- Contrast data structure's strengths and weaknesses
  - understand when to use each one

- Abstract Data Types (aka ADTs) are descriptions of how a data type will work without implementation details
- Description can be a formal, mathematical description
- Java interfaces are a form of ADTs
  - some implementation details start to creep in

- Programming languages usually have a library of data structures
  - <u>Java collections framework</u>
  - C++ standard template library
  - .Net framework (small portion of VERY large library)
  - Python lists and tuples
  - Lisp lists

# Structure in C

# Defining a Structure

```
struct tag {
    member 1;
    member 2;
    :
    member m;
};
```

- **struct** is the required keyword
- tag is the name of the structure
- member 1, member 2,.. are individual member declarations

# Defining a Structure

```
char name[30];
int roll_number;
int total_marks;
char dob[10];

struct student s1, sList[100];

A new data-type
```

• Each member in a structure can be accessed with (.) operator called scope resolution operator

```
s1.name; sList[5].roll_number;
```

# Structures (records)

```
char name[10];
    int age;
    float salary;
    } person;

strcpy(person.name, "james");
person.age=10;
person.salary=35000;
```

```
typedef struct human_being {
        char name[10];
        int age;
        float salary;
or
typedef struct {
        char name[10];
        int age;
        float salary
        } human_being;
human_being person1, person2;
```

# Set Manipulation Structure

```
Set Manipulation
struct nodeS {
                   int element;
                                                Structure
                   struct nodeS *next;
                                                definition
typedef struct nodeS set;
set *union (set a, set b);
set *intersect (set a, set b);
set *minus (set a, set b);
                                   Function
void insert (set a, int x);
                                   prototypes
void delete (set a, int x);
int size (set a);
```

# Complex Numbers Structure

```
#include <stdio.h>
main()
   struct complex
                                     Structure definition
         float real;
                                     and
                                     Variable Declaration
         float complex;
                                                                               Scope
   } a, b, c;
                                                                               restricted
                                                                               within
                                                                               main()
   scanf ("%f %f", &a.real, &a.complex);
                                                  Reading a member variable
   scanf ("%f %f", &b.real, &b.complex);
   c.real = a.real + b.real;
                                                  Accessing members
   c.complex = a.complex + b.complex;
```

# Complex Numbers Structure

```
struct typeX {

float re;
float im;
};

Structure
definition
```

### typedef struct typeX complex;

```
complex *add (complex a, complex b);
complex *sub (complex a, complex b);
complex *mul (complex a, complex b);
complex *div (complex a, complex b);
complex *read();
void print (complex a);

Function
prototypes
```

# Add Two Complex Numbers Structure

```
#include <stdio.h>
typedef struct complex
 float real;
 float imag;
 complex;
int main()
{ complex n1, n2, temp;
  printf("For 1st complex number \n");
  printf(" Enter re & im part respectively:\n");
  scanf("%f%f", &n1.real, &n1.imag);
  printf("\nFor 2nd complex number \n");
  printf("Enter re & im part respectively:\n");
  scanf("%f%f", &n2.real, &n2.imag);
 temp = add(n1, n2);
  printf("Sum = %.1f + %.1fi", temp.real, temp.imag);
return 0;
```

```
complex add(complex n1, complex n2)
  complex temp;
   temp.real = n1.real + n2.real;
   temp.imag = n1.imag + n2.imag;
  return(temp);
OUTPUT
For 1st complex number
Enter re & im part respectively: 2.3
4.5
For 2nd complex number
Enter re & im part respectively: 3.4
Sum = 5.7 + 9.5i
```

# **Comparing Dates Structure**

```
#include <stdio.h>
struct date { int dd, mm, yy;
int date_cmp(struct date d1, struct date d2);
void date_print(struct date d);
int main(){
  struct date d1 = \{7, 3, 2015\};
  struct date d2 = \{24, 10, 2015\};
  int cmp = date\_cmp(d1, d2);
  date_print(d1);
  if (cmp == 0)
   printf(" is equal to");
  else if (cmp > 0)
   printf(" is greater, i.e., later than ");
  else printf(" is smaller, i.e., earlier than");
  date_print(d2);
  return 0;
```

```
/* compare given dates d1 and d2 */
int date_cmp(struct date d1, struct date d2) {
 if (d1.dd == d2.dd & d1.mm == d2.mm & d1.yy == d2.yy)
          return 0;
 else if (d1.yy > d2.yy \mid | d1.yy == d2.yy & d1.mm > d2.mm \mid |
          d1.yy == d2.yy && d1.mm == d2.mm && d1.dd > d2.dd
          return 1;
 else return -1;
/* print a given date */
void date_print(struct date d) {
   printf("%d/%d/%d", d.dd, d.mm, d.yy);
OUTPUT:
7/3/2015 is smaller, i.e., earlier than 24/10/2015
```

# Sparse Matrices - Data Structure

# 1-D Array Representation

- Implementation of the abstract list data structure using programming language
  - "Backing" Data Structure
- Arrays are contiguous memory locations with fixed capacity
- Allow elements of same type to be present at specific positions in the array
- Index in a List can be mapped to a Position in the Array
  - Mapping function from list index to array position

# Matrix Multiplication

```
// Given 2-D arrays: a[n][n], b[n][n]
// Output 2-D array: c[n][n] initialized to 0
for (i = 0; i < N; i++)
  for (j = 0; j < N; j++)
    for (k = 0; k < N; k++)
    c[i][j] += a[i][k] * b[k][j];</pre>
```

# Symmetric Matrix

- An  $n \times n$  matrix can be represented using 1-D array of size  $\frac{n(n+1)/2}{2}$  by storing either the lower or upper triangle of the matrix
- Use one of the methods for a triangular matrix
- Optimization: The elements that are not explicitly stored may be computed from those that are stored.

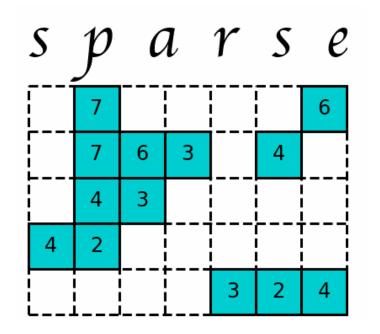
```
2 4 6 0
4 1 9 5
6 9 4 7
0 5 7 0
```

# **Sparse Matrices**

- Only a small subset of items are populated in matrix
  - Students and courses taken, faculty and courses taught
  - Adjacency matrix of social network graph
  - vertices are people, edges are "friends"
- Rows and columns are people, cell has 0/1 value
- Why not use regular 2-D matrix?
  - 1-D representation
  - Array of arrays representation

# **Sparse Matrix**

- A matrix is sparse
  - if many of its elements are zero
- A matrix that is not sparse is dense
- The boundary is not precisely defined
  - Diagonal and tridiagonal matrices are sparse
  - We classify triangular matrices as dense
- Two possible representations
  - array
  - linked list



### **Dense Matrix**

1	2	31	2	9	7	34	22	11	5
11	92	4	3	2	2	3	3	2	1
3	9	13	8	21	17	4	2	1	4
8	32	1	2	34	18	7	78	10	7
9	22	3	9	8	71	12	22	17	3
13	21	21	9	2	47	1	81	21	9
21	12	53	12	91	24	81	8	91	2
61	8	33	82	19	87	16	3	1	55
54	4	78	24	18	11	4	2	99	5
13	22	32	42	9	15	9	22	1	21

### **Sparse Matrix**

1		3		9		3			
11		4						2	1
		1				4		1	
8				3	1				
			9			1		17	
13	21		9	2	47	1	81	21	9
				19	8	16			55
54	4				11				
		2					22		21

# Array Representation of Sparse Matrix

- The nonzero entries may be mapped into a 1D array in row-major order
- To reconstruct the matrix structure, need to record the row and column each nonzero comes from

0	0	8	0	0	0	
0	7	0	0	0	0	L
0	0	0	0	0	5	l
			0		0	
0	0	0	0	1	0	

Rows	Columns	values
5	6	5
0	2	8
1	1	7
2	5	5
3	1	3
4	4	1

	0						
0	0	0	9	0	8	0	0
(0)	4						

(a) A 
$$4 \times 8$$
 matrix

(b) Its representation

# Array Representation of Sparse Matrix

```
template < class T >
class Term {
private:
    int row, col;
    T value;
};
```

0 0 8 0 0 0 0 7 0 0 0 0 0 0 0 0 0 5 0 3 0 0 0 0 0 0 0 0 1 0



	Rows	Columns	values
	5	6	5
	0	2	8
	1	1	7
	2	5	5
	3	1	3
ſ	4	4	1

```
template < class T >
class sparseMatrix {
private:
       int rows, cols,
       int terms;
       Term<T> *a;
       int MaxTerms;
public:
```

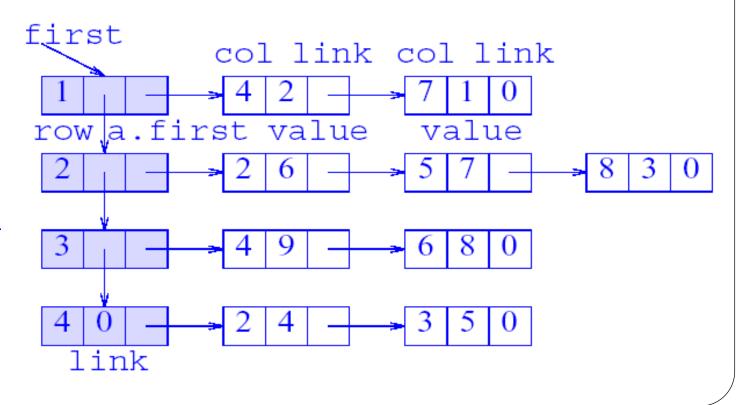
# Linked Representation of Sparse Matrix

- A shortcoming of the 1-D array of a sparse matrix is that we need to know the number of nonzero terms in each of the sparse matrices when the array is created
- A linked representation can overcome this shortcoming

(a) A  $4 \times 8$  matrix

a []	Ō	1	2	3	4	5	6	7	8
row	1		2	2	2	3	3	4	4
col	4	7	2	5	8	4	6	2	3
row col value	2	1	6	7	3	9	8	4	5

(b) Its representation



## References

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ขอบคุณ

תודה רבה Grazie Italian

Hebrew

Thai

ಧನ್ಯವಾದಗಳು

Kannada

Sanskrit

धन्यवादः

Ευχαριστώ

Greek

Thank You English

Gracias

Spanish

Спасибо

Russian

Obrigado

Portuguese

شكراً

https://sites.google.com/site/animeshchaturvedi07

Merci

French

Arabic

多謝

**Traditional** 

Chinese

धन्यवाद

Hindi

Danke

German



Simplified

Chinese

நன்றி

Tamil

**Tamil** 

ありがとうございました 감사합니다

Japanese

Korean