

Distance Vector Algorithm – Problems, Solutions and a Standard

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Recap

- Nodes exchange with their neighbors their current routing table information (destination, estimated cost)
- On receipt of a message, nodes update cost to destination based on Bellman-ford equation
- Messages sent periodically as well as when table changes

Reference Node C

D	C	H
A	5	A
B	3	B
D	4	D

Routing Table of C

(1)

To	A
A	0
B	1
C	5

Message from A
C to A: C = 5

D	C	H
A	5	A
B	3	B
D	4	D

Routing Table of C

D	C	H
A	5	A
B	3	B
D	4	D

Routing Table of C

(2)

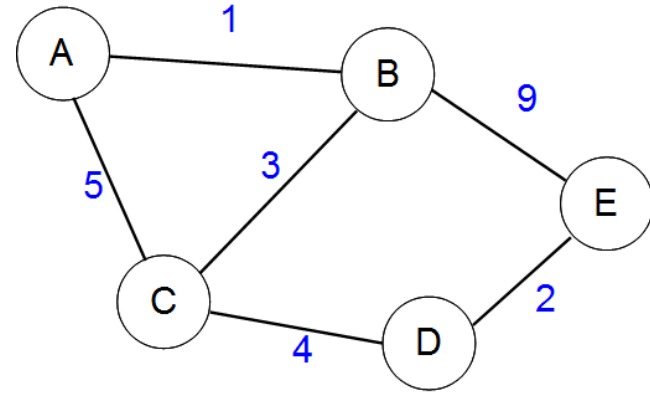
To	B
A	1
B	0
C	3
E	9

Message from B
C to B: C = 3

D	C	H
A	4	B
B	3	B
D	4	D
E	9	B

Routing Table of C

Example



D	C	H
A	4	B
B	3	B
D	4	D
E	9	B

Routing Table of C

(3)

To	D
C	4
D	0
E	2

Message from D
C to D: C = 4

D	C	H
A	4	B
B	3	B
D	4	D
E	6	D

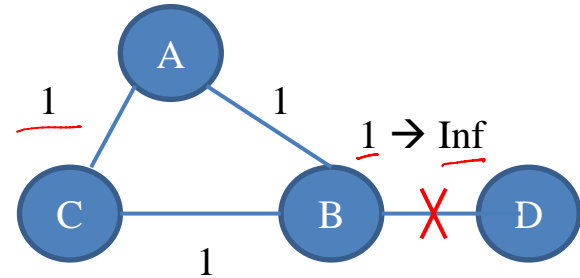
Routing Table of C

Every path has its puddle!

Counting to Infinity

Distance to Node D

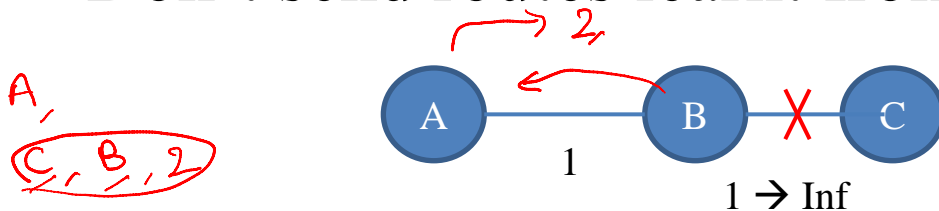
Mesg.	A	B	C
	2,B	∞ , -	2,B
$B \rightarrow A$	∞ , -	∞ , -	2,B
$C \rightarrow A$	3,C	∞ , -	2,B
$B \rightarrow C$	3,C	∞ , -	∞ , -
$A \rightarrow B$	3,C	4,A	∞ , -
$C \rightarrow A$	∞ , -	4,A	∞ , -
$B \rightarrow C$	∞ , -	4,A	5,B
$A \rightarrow B$	∞ , -	∞ , -	5,B



State maintained by nodes A, B and C

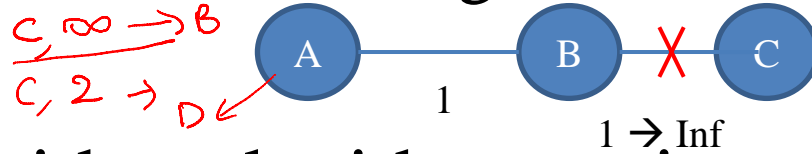
Partial Solutions

- Make infinity small
 - Use for example 16 to represent infinity (assumes max no of hops under 16)
↗ hop count
 - Bounds time it takes to count to infinity
- Split horizon
 - Don't send routes learnt from a neighbor back to it

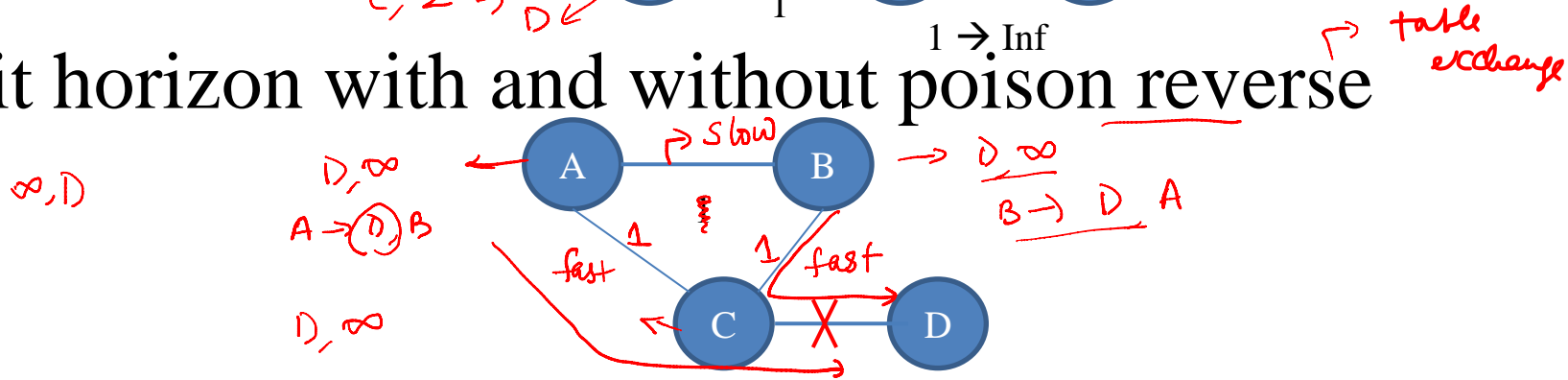


Partial Solutions

- Split horizon with poison reverse
 - Send routes learnt from a neighbor back to it but with infinite cost



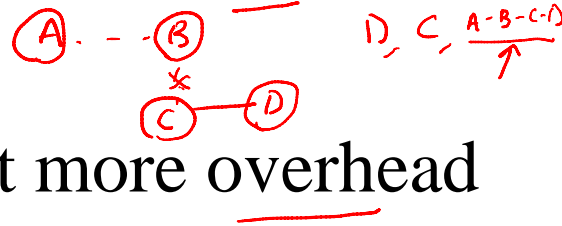
- Split horizon with and without poison reverse



- Both don't work for loops with more than 2 nodes

Partial Solutions

- Hold-Down Timer: Wait some time before propagating link failure
 - Slows down convergence
- Path-vector routing is a variation of distance-vector
 - Each node sends to its neighbors not just the cost, but the entire path to the destination
 - Avoids the looping problem of DV but more overhead



RIP

- Routing Information Protocol (RIP) is a standard that implements DV routing
- One of the oldest DV based protocol
 - Popular once, not used much due to convergence problems
↳ size

RIP Features

- Uses UDP and work over reserved port 520
- Period updates sent every 30 sec
- Supports multiple address families IP
- Cost of a link is 1 (finds minimum hop route)
- 16 represents infinity, *split horizon, hold-down timer*
- RIP can run only on very small networks

RIPv1 Packet Format

→ R₁ C₁ (NH₁)

0	Command	Version	Reserved	31
	Family of Net 1		Reserved	
	Address of Net 1 (IP Address)			
	Reserved			
	Reserved			
	Distance to Net 1			
	Family of Net 2		Reserved	
	Address of Net 2 (IP Address)			
	Reserved			
	Reserved			
	Distance to Net 1			

→ IP Prefix (NH₁)
Route

Reserved fields are set to all zero

$\left. \begin{matrix} IPF_1 \\ R \end{matrix} \right\} \left. \begin{matrix} IPF_2 \\ R \end{matrix} \right\} \left. \begin{matrix} IPF_3 \\ R \end{matrix} \right\}$
 $\begin{matrix} 1 & 1 & 1 \\ 0 & 0 & 0 \end{matrix}$
 $\begin{matrix} IPF_1 & 0 \\ IPF_2 & 0 \\ IPF_3 & 1 \end{matrix}$

1 to 25 sets of entries, each entry is 20 bytes

classless

Summary

- Distance vector is a distributed, dynamic algorithm
- Exchanges information locally to determine routes
- Suffers from poor convergence, routing loops
- RIP is a standard that implements the DV protocol
 - Handles above problems via (split horizon, hold-down timer and using a value of 16 to represent infinity)
- Better approach: Link-state routing