

Address Space Utilization

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Recap

- Hierarchical addressing based on classes (A,B,C) to handle scalability and different network sizes
 - IP address has a network part and a host part
 - Routers maintain entries corresponding to network portion
- What can potentially go wrong now?

Problem Statement

- Network part uniquely identifies a physical network
- Network with just 2 hosts needs Class C address
 - Efficiency: $\frac{2}{2^8} = 0.7\%$
- Network with 256 hosts needs a Class B address
 - Efficiency: $\frac{256}{2^{16}} = 0.4\%$

2 hosts

IP address - 32 bits
class B address → block of
IP addresses

190.23.*.*

16 bits 16

8 bits 2⁸

256 hosts


2³² ~ 4 billion hosts

Problem Statement

- Class B addresses in high demand (keeping future needs in mind)
 - Host addresses ~ 4 billion (2^{32}) ^{2^{24} hosts ~ very few}
 - Class B networks: 2^{14} (~16000)
 - Out of class B addresses, out of addresses to hand out
- Need to solve “Address assignment inefficiency”
 - Challenge: Within IP framework (using 32-bits)

“The gem cannot be polished without friction,
nor man perfected without trials.”

---Chinese proverb

“You see a mousetrap; I see free cheese and a
 ***** challenge!”

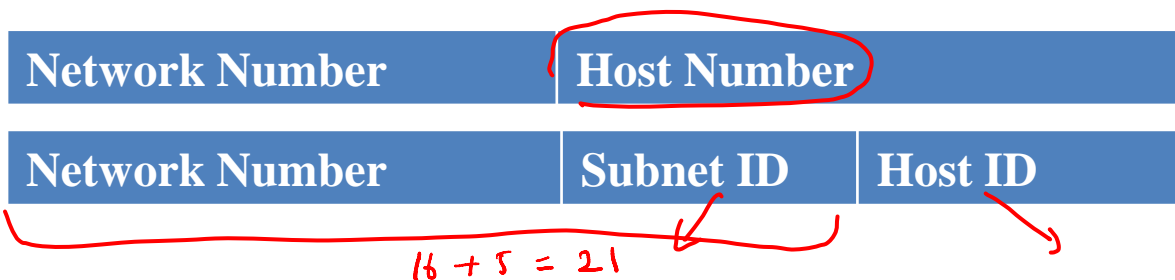
---Scroobius Pip

Specific Scenario –1

- An organization has 30 different physical networks, each network has about 2000 hosts
- Current Allocation: Allocate 30 class B addresses
 - Efficiency = 2000/2¹⁶ = 3%
- Will one class B address suffice?
 - Can support 2¹⁶ = 65536 hosts > 30*2000

Solution: Subnetting

- Introduce another level of hierarchy
 - Divide host part into subnet id and host id



- 16 host bits divided into 5 subnet bits (32 physical networks) and 11 host bits (2048 hosts)

- Address format: a.b.c.d/x, where x is # bits in net portion of address

– Example: 190.23.12.17, mask is 255.255.248.0 (/21)

- `ifconfig eth0 190.23.12.17 netmask 255.255.255.240`

- The bit-wise end of the IP address and the subnet mask give the subnet number of the host
- All hosts on a given physical network have the same subnet number and mask

Class B Address: 10111110.00010111.*.*; 190.23.0.0/16

$(2^{11} - 2)$

32 Subnets	Subnet ID <i>5 bits</i>	Subnet No <i>16 bits</i>	Hosts <i>11 bits</i>	<i>→ 0000 all zeros all ones broadcast</i>
	00000	<u>190.23.0.0</u> ✓	<u>190.23.0.1 – 190.23.7.254</u>	
	00001	190.23.8.0 ✓	190.23.8.1 – 190.23.15.254	
			
	11111	190.23.248.0 ✓	190.23.248.1 – 190.23.255.254	

- All routers outside organization have one entry (190.23.0.0/16)
- Routers within organization have more detailed entries corresponding to different subnets

Forwarding at a Router

D = Destination IP Address

For each forwarding table entry

$T = D \ \& \ \text{SubnetMask}$

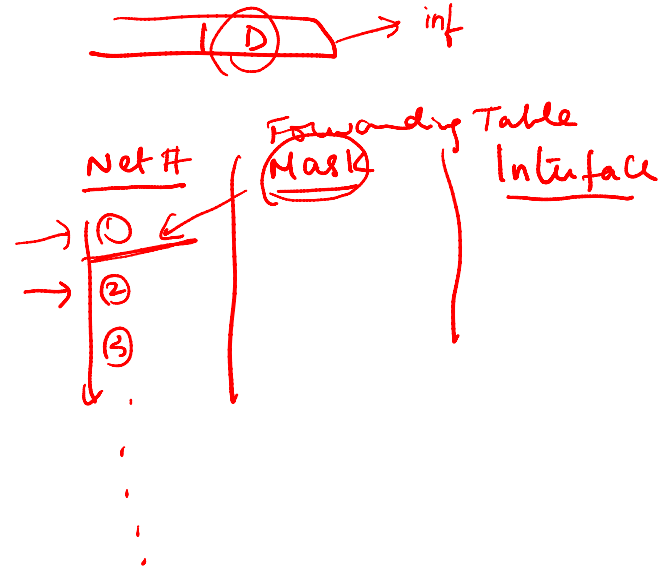
If $T == \text{SubnetNum}$

If Nexthop is an interface

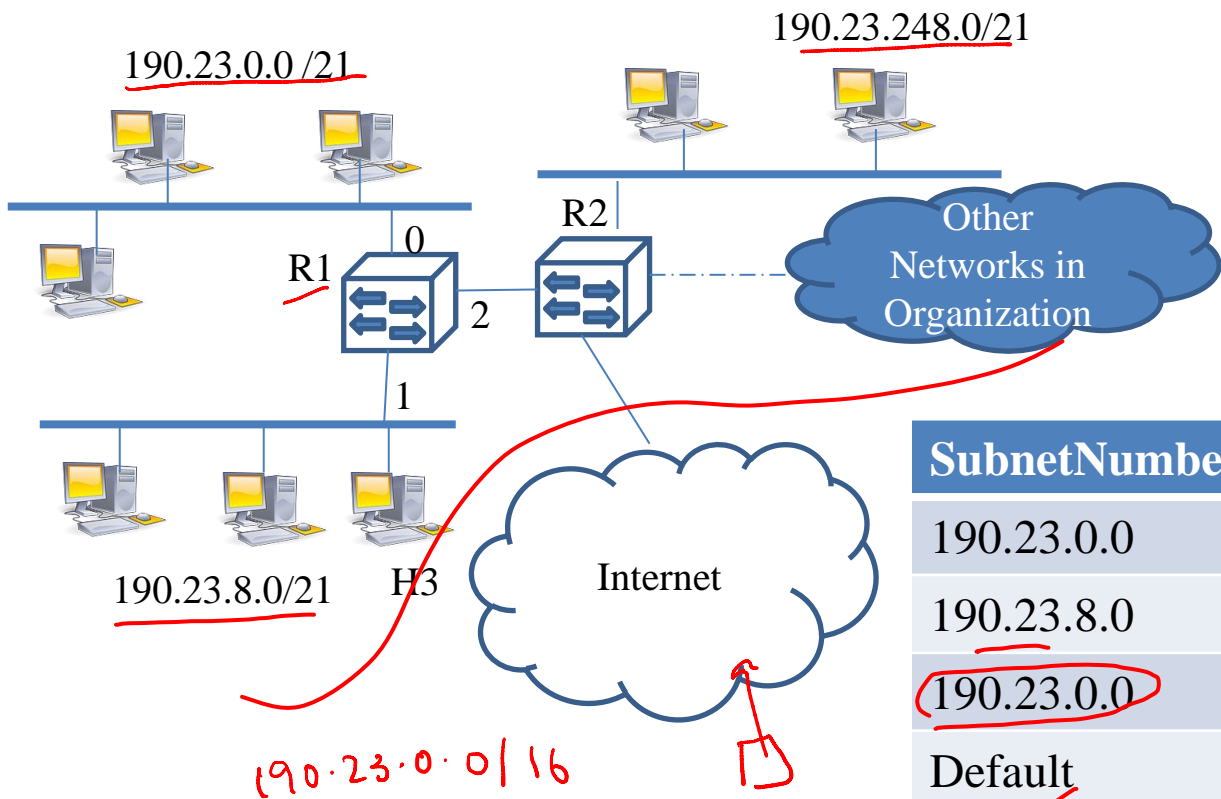
deliver datagram directly to destination

Else

deliver datagram to NextHop (router)



Example



Forwarding
Routing Table at R1

SubnetNumber	Mask	Nexthop
190.23.0.0	<u>255.255.248.0</u>	Inf0
<u>190.23.8.0</u>	255.255.248.0	Inf1
<u>190.23.0.0</u>	<u>255.255.0.0</u>	<u>R2</u>
<u>Default</u>	0.0.0.0	<u>R2</u>

```

kameswari@asterix:~$ route
Kernel IP routing table
Destination      Gateway          Genmask          Flags  Metric  Ref    Use  Iface
10.129.0.0        *                255.255.0.0      U        1        0        0  eth0
link-local        *                255.255.0.0      U       1000        0        0  eth0
default          router.it.iitb.  0.0.0.0          UG         0        0        0  eth0
kameswari@asterix:~$
kameswari@asterix:~$
kameswari@asterix:~$ route -n
Kernel IP routing table
Destination      Gateway          Genmask          Flags  Metric  Ref    Use  Iface
10.129.0.0        0.0.0.0          255.255.0.0      U        1        0        0  eth0
169.254.0.0        0.0.0.0          255.255.0.0      U       1000        0        0  eth0
0.0.0.0          10.129.250.1    0.0.0.0          UG         0        0        0  eth0
kameswari@asterix:~$ |

```

A close-up, full-frame image of a deep red velvet curtain. The fabric is heavily draped, creating a series of vertical, wavy folds that catch the light, giving it a rich, textured appearance. The color is a vibrant, slightly dark red.

INTERMISSION

Specific Scenario -- 2

- An organization has a physical network with 4000 hosts
- Current Solution: Give a class B address
 - Efficiency: $4000/2^{16} = \underline{6\%}$ ←
- How about assigning multiple class C addresses?
16 → 2⁸ hosts
- Problem: 16 entries for same organization in the routing table

Solution--2

- Assign multiple contiguous class C addresses & aggregate
- 222.7.16.* through 222.7.31.*, top 20 bits in this range are the same (0001 bits, 20-bit network number)
- Advertise 222.7.16/20 as the organization's network address
- Goes by the name supernetting

Conclusions

- Subnetting: One class address shared among many physical networks
- Supernetting: Multiple class addresses shared among one physical network (Autonomous system -- AS) *→ organization*
- Network portion can take on any length

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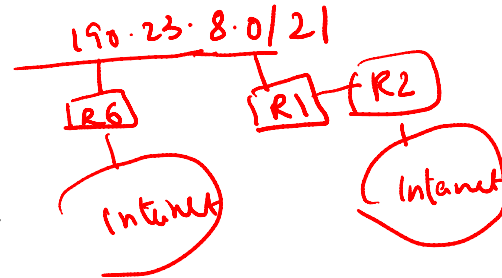
2 to (32)

Classless Interdomain Routing (CIDR)

- Use a new notation to represent network numbers (also called IP prefixes)
- Address block represented as A/X, where A is the address prefix and X is the prefix length
 - X can range from 2 till 32
 - X is represented as a network mask as well 255.255.248.0 ^{/21}
- E.g. 222.7.16/20 (Mask 255.255.240.0) represents addresses in the range (222.7.16.0 to 222.7.31.255)
entry → interna

Longest Prefix Match

- Routers do a prefix match.
 - Does destination address fall in the range of addresses captured by prefix?
- Prefix match works if Internet topology is a tree
 - Shortest path between networks is unique
- Internet is a graph
 - Many networks multi-home
 - Many matching prefixes



Example

- Two prefixes in a forwarding table 190.23.8.0/21 and 190.23.0.0/16
 ↑ R6
 → R2
 – Address 190.23.8.1 matches both
- Go with the longest prefix match (e.g. 190.23.8.0/21)
- Challenge: Longest match between destination IP address and variable length prefixes in forwarding table
 - Lot of research in this space

Summary

- Class based addressing was found not to be scalable
- Subnetting: Share a single class address among multiple networks
- Supernetting: Share multiple class addresses on a single network
- Lead to CIDR (classless addressing) and Longest prefix match --- widely used now