CS 535 Design and Analysis of Algorithms

Spring Semester, 2019

Homework 5

Assigned: April 2 Due: April 16

Problem 1 Let G = (V, E, c, s, t) be a flow network with integer capacities, and f be a feasible flow with integer values. Suppose there is a edge e with head s and such that f(e) = 1. Describe an O(|V| + |E|) algorithm (An English explanation may be enough) that produces another feasible flow f' with $|f| \le |f'|$ and such that f'(e) = 0.

Be precise though: if you use an algorithm from the texbook, explain which graph is the input of the algorithm. Justify the overall running time and correctness.

Problem 2 A multiple source-sink network is a tuple G = (V, E, c, S, T), where V is a set of vertices, E is a set of directed edges (parallel edges are allowed), $S \subset V$ is the set of sources, and $T \subset V$ is the set of sinks, c is a capacity function: $c: E \to Z_+$. Also, $S \cap T = \emptyset$. That is, sources are distinct from sinks.

A function $f: E \to R_+$ is called a *flow* if the following three conditions are satisfied:

1. conservation of flow at interior vertices: for all vertices u not in $S \cup T$,

$$\sum_{e \in \delta^-(u)} f(e) = \sum_{e \in \delta^+(u)} f(e) ;$$

2. capacity constraints: $f \leq c$ pointwise: i.e. for all $e \in E$,

$$f(e) \leq c(e)$$
.

Assume that non-negative quantities p_s , for $s \in S$, and q_t , for $t \in T$, are given. The goal of this problem is to determine if a valid flow exists such that for all $s \in S$:

$$\sum_{e \in \delta^+(s)} f(e) - \sum_{e \in \delta^-(s)} f(e) = p_s$$

and such that for all $t \in T$:

$$\sum_{e \in \delta^-(t)} f(e) - \sum_{e \in \delta^+(t)} f(e) = q_t.$$

Use Network Flows to give a polynomial-time algorithm for this **decision** problem (the answer is YES or NO). Hint: read chapter 26.1 of the textbook.

Problem 3 The edge connectivity of an undirected multigraph is the minimum number k of edges that must be removed to disconnect the graph. For example, the edge connectivity of a tree is 1, and the edge connectivity of a cyclic chain of vertices is 2. Show how to determine the edge connectivity of an undirected multigraph G = (V, E) by running a maximum-flow algorithm on at most |V| flow networks, each having O(|V|) vertices and O(|E|) edges.

Argue correctness.