

Welcome 😊

Agenda : 2D D.P

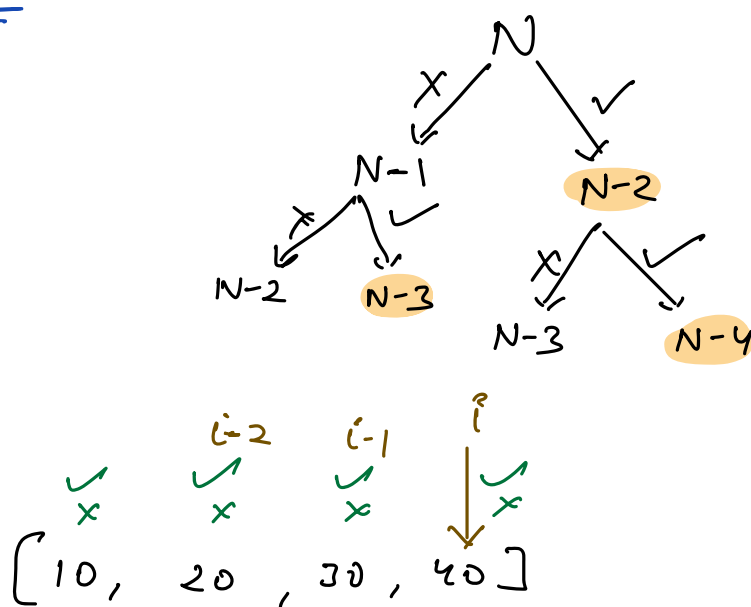
Q Find minimum subsequence sum from a given array, where selecting adjacent element is not allowed

eg: A : [ 5 , 4 , 8 ]  $\Rightarrow 5 + 8 \Rightarrow 13$

Quiz [ 10 , 20 , 30 , 40 ]  $10 + \cancel{30} + 40 = 50$

$20 + \cancel{30} + 40 = \underline{\underline{60}}$

Brute force



manSum [ i-1 ]

manSum [ i-2 ]

man ( manSum [ i-1 ] , manSum [ i-2 ] + arr [ i ] )

## Recursive DP

int dp[N] = {-1}

int mansum ( arr[], <sup>N-1</sup>inden

{ if ( inden < 0 ) return 0 ;

if ( dp[inden] != -1 ) return dp[inden]

f1 = mansum ( arr, inden - 1 )

f2 = arr[inden] + mansum ( arr, inden - 2 )

ans = max ( f1, f2 )

dp[inden] = ans

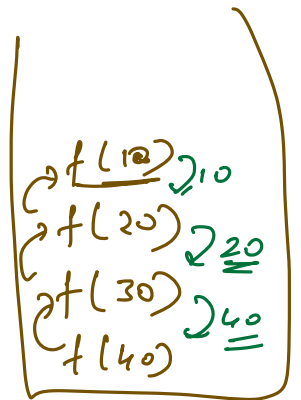
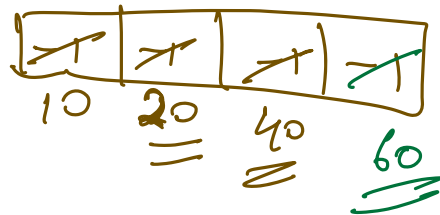
return ans

}

T.C  $\Rightarrow$   ~~$O(2^N)$~~   $\Rightarrow O(N)$

S.C =  $O(N+N)$

[10, 20, 30, 40]



[20, 10, 30, 40]

dp[0] = a[0]

dp[1] = max(a[0], a[1])

for ( i  $\rightarrow$  2 to N-1 )

{ dp[i] = max ( arr[i] + dp[i-2], dp[i-1] )

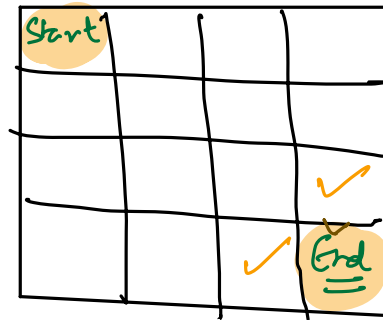
}  
return dp[N-1]

T.C  $\Rightarrow O(N)$

S.C  $\Rightarrow O(N)$

can reduce to  $O(1)$

Q Count Unique Paths.



Bruteforce

$$\text{ways}(i, j) = \text{ways}(i-1, j) + \text{ways}(i, j-1)$$

Recursive

int dp[N][m] = {-1}

int ways(i, j)

{

if (i == 0 || j == 0)  
return 1

if (dp[i][j] != -1) return dp[i][j]

dp[i][j] = ways(i-1, j) + ways(i, j-1)

return dp[i][j]

}

T.C =  $O(\frac{N \times m}{2}) \rightarrow O(N \times m)$

S.C =  $O(N \times m)$

Q No. of ways to reach bottom right with the given constraints.

	0	1	2
0	1	<del>0</del>	1
1	1	1	0
2	0	1	1

if (mat[i][j] != 0)

$$\text{ways}[i][j] = \text{ways}[i-1][j] + \text{ways}[i][j-1]$$

else

$$\text{ways}[i][j] = 0$$

Q Dungeons & Princes

-3	+2	+4	<del>-5</del>	
-6	5	-4	<del>6</del>	
-15	-7	<u>1</u> 9	<del>-2</del>	
<u>1</u> 2	<u>1</u> 10	<u>8</u> -3	<del>-4</del> 5	

Find min. health with which Prince should start so that he can reach to princess without dying.

Bruteforce

$$X + \text{arr}[i][j] = \min(\text{dp}[i+1][j], \text{dp}[i][j+1])$$

code

```
if ( arr[N-1][M-1] > 0 )
```

```
    dp[N-1][M-1] = 1
```

```
else {
```

```
    dp[N-1][M-1] = 1 + abs(arr[N-1][M-1])
```

```
}
```

// Fill last row & last column

```
for ( i = N-2 ; i >= 0 ; i-- )
```

```
{
```

```
    for ( j = M-2 ; j >= 0 ; j-- )
```

```
{
```

```
        dp[i][j] = max( 1, min( dp[i+1][j], dp[i][j+1] ) - arr[i][j] )
```

```
}
```

```
}
```

```
return dp[0][0]
```

T.C =  $O(N \times M)$

S.C =  $O(N \times M)$