### Extending SDN to the Data Plane

Anirudh Sivaraman, Keith Winstein, Suvinay Subramanian, Hari Balakrishnan

M.I.T.

http://web.mit.edu/anirudh/www/sdn-data-plane.html

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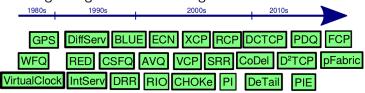
## Switch Data Planes today

Two key decisions on a per-packet basis:

- ► Scheduling: Which packet should be transmitted next?
- ► Queue Management: How long can queues grow? Which packet to drop?

### The Data Plane is continuously evolving

▶ The long lineage of in-network algorithms:



- Each scheme wins in its own evaluation.
- ► Some believe in a "silver bullet" knobless in-network method.

### We disagree: There is no silver bullet!

- ▶ Different applications care about different objectives.
- Applications use different transport protocols.
- Networks are heterogeneous.

### Early symptoms

- Hard to configure wired AQM for wireless links
- Several distinct point solutions for datacenters
  - ▶ DCTCP, HULL, D3, DeTail, PDQ, pFabric
- ▶ No consensus on the "right metric"
  - Minimizing missed deadlines
  - ► Flow Completion Time
  - Latency
  - Throughput
  - Tail Latency

## Quantifying "No Silver Bullet": Network Configurations

Configuration	Description
CoDel+FCFS	One shared FCFS queue with CoDel
CoDel+FQ	Per-flow fair queueing with CoDel on each queue
Bufferbloat+FQ	Per-flow fair queueing with deep buffers on each queue

# Quantifying "No Silver Bullet": Workloads and Objectives

Workload	Description	Objective
Bulk	Long-running TCP flow	Maximize throughput
Web	Switched TCP flow with ON and OFF periods	Minimize 99.9 %ile flow completion time
Interactive	Long-running real-time streaming app	Maximize throughput delay, i.e., "power"



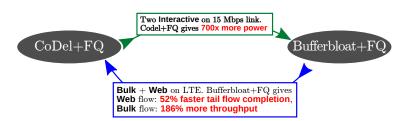


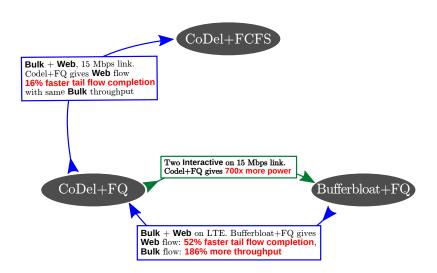
Bufferbloat+FQ

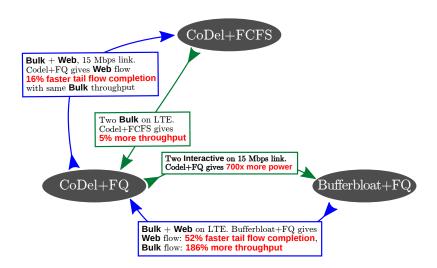


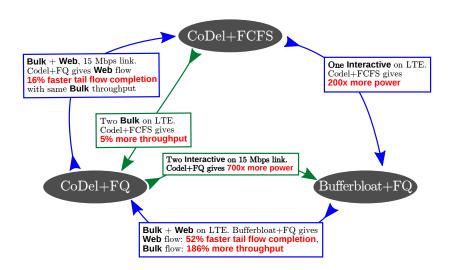


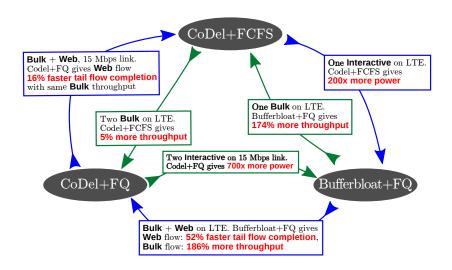












### Why is no single data plane configuration the best?

- Bufferbloat on variable-rate links helps throughput!
  - ▶ Variable-rate links have an inherent delay-throughput tradeoff
- ▶ FCFS is preferable to Fair Queuing in some cases
  - When equally aggressive flows compete, they don't need protection from each other
  - Helps reduce tail packet delay
- ► Fair Queuing is required in some cases
  - ▶ When competing flows aren't equally aggressive, isolation helps

### So what should the network designer do?

#### Architect a flexible data plane

- Programmable queue management and scheduling
- Not just for selecting among pre-built choices, but to change behavior in the field
- ▶ Because there is no silver bullet and innovation will continue!

### Controlled flexibility: Want performance, security

(Or, why this isn't the same as "active networks")

- Provide interfaces only to the head and tail of queues
- Operators specify only queue-management/scheduling logic
- ► No access to packet payloads (for now)

### Building such a data plane in four parts

- Hardware gadgets
  - Random number generators (RED, BLUE)
  - Binary tree of comparators (pFabric, SRPT)
- I/O interfaces
  - Drop/mark head/tail of queue
  - ▶ Interrupts for enqueue/dequeue
- State maintenance
  - Per-flow (WFQ, DRR)
  - ► Per-dst address (PF)
- A domain-specific instruction set
  - Expresses control flow
  - ▶ Implements new functions unavailable in hardware

### Feasibility study: CoDel

Synthesis numbers on Xilinx Kintex-7:

Resource	Usage	Fraction of FPGA
Slice logic	1,256	1%
Slice logic dist.	1,975	2%
IO/GTX ports	27	2%
DSP slices	0	0%
Maximum speed	$12.9 \times 10^6$ pkts/s ~10gbps	

- ▶ Small fraction of the FPGA's resources.
- ▶ Can be improved by pipelining or parallelizing.

#### Conclusion

- ► There is no silver bullet to in-network resource control because of application and network diversity
- ▶ Algorithms will continue to evolve: the data plane should help
- ► Directions to reproduce results: http://web.mit.edu/anirudh/www/sdn-data-plane.html

#### Limitations and Practical Considerations:

- ► Cannot express several network functions that need payloads.
- ▶ How do applications signal objectives to the network?
- ► Feasibility at 10G on high port-density switches.
- Mechanism to map flows onto per-port queues.
- Energy and Area overheads.