

# Course: Algorithms

## Autumn 2018

# Single Source Shortest Path Algorithm - Dijkstra

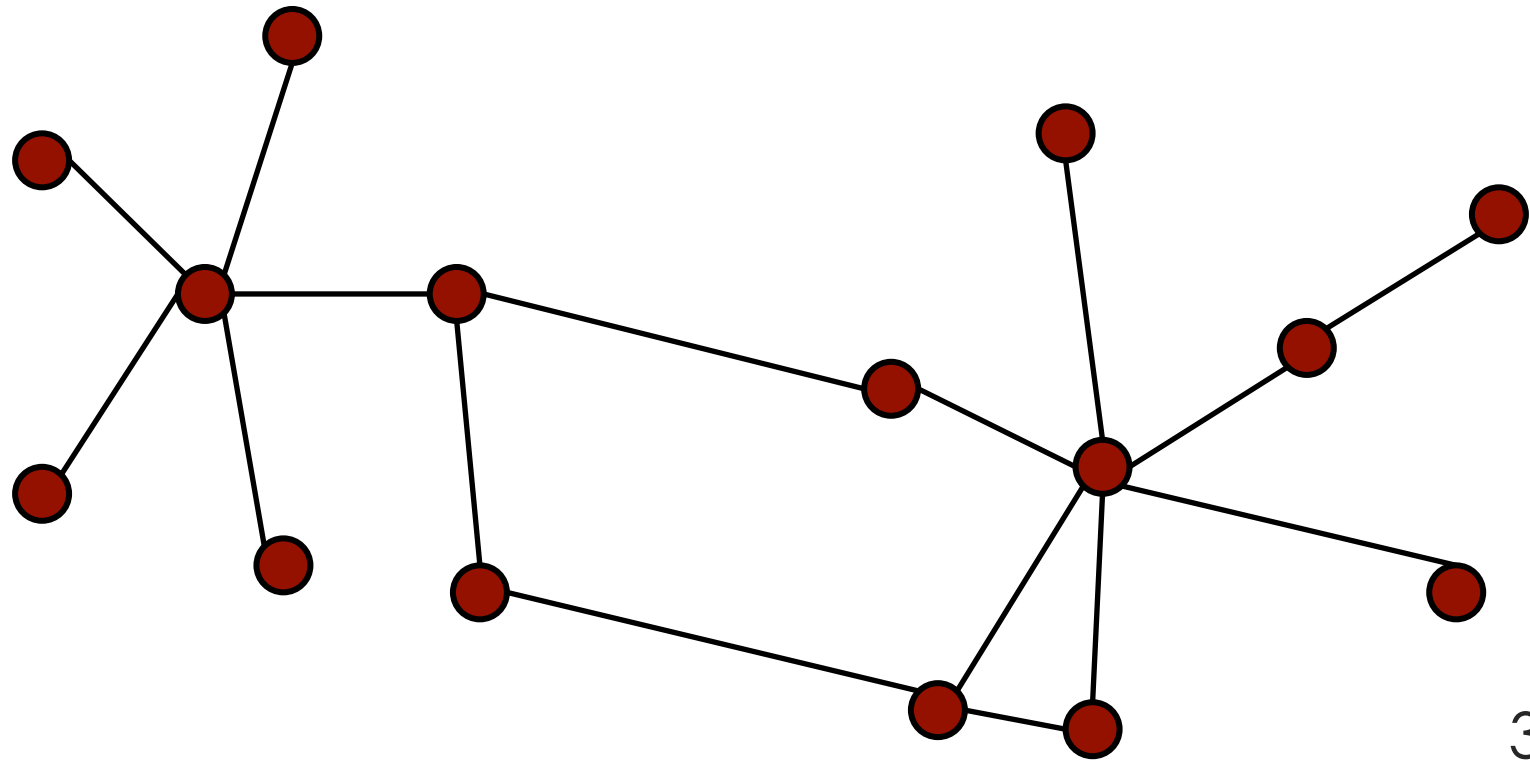
This lecture covers the interesting aspects of the single source shortest path algorithm – The popular Dijkstra's Algorithm. We will also look at its computation complexity with the limitations.

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# Overlay Construction

## Divide-and-Conquer Algorithm

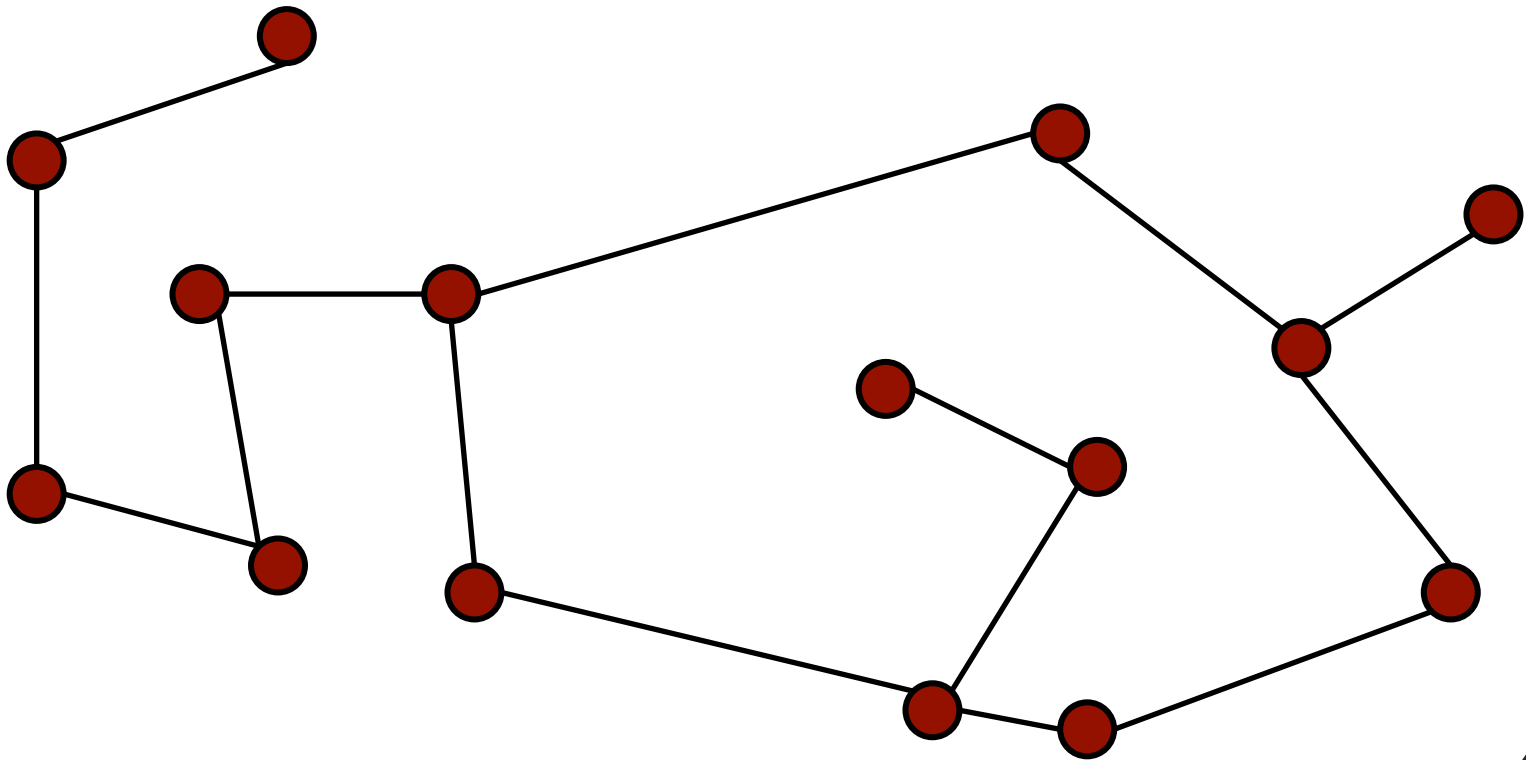
Initially: connected graph



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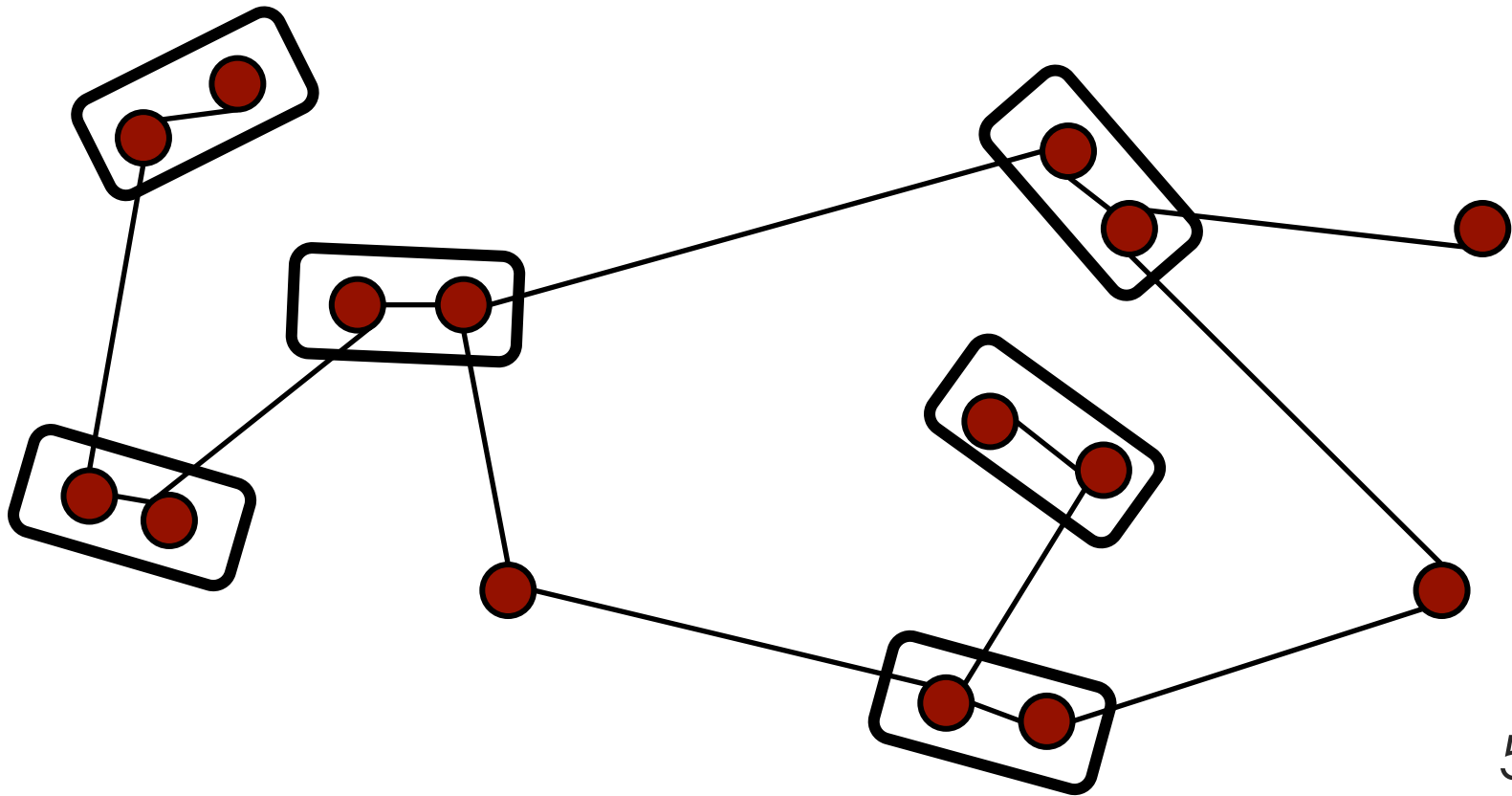
# Divide-and-Conquer Algorithm

Sparsify: reduce the degree



# Divide-and-Conquer Algorithm

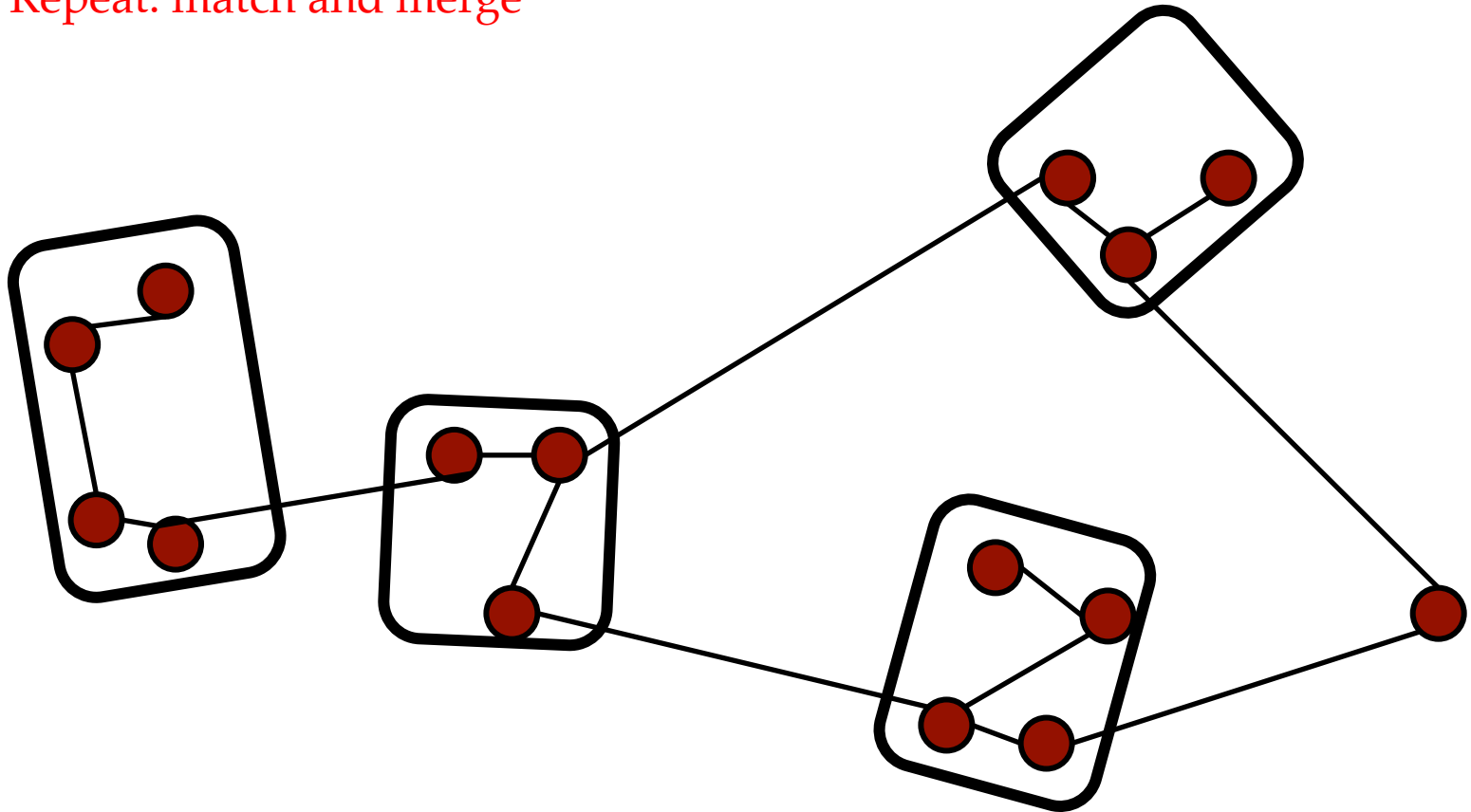
Repeat: match and merge



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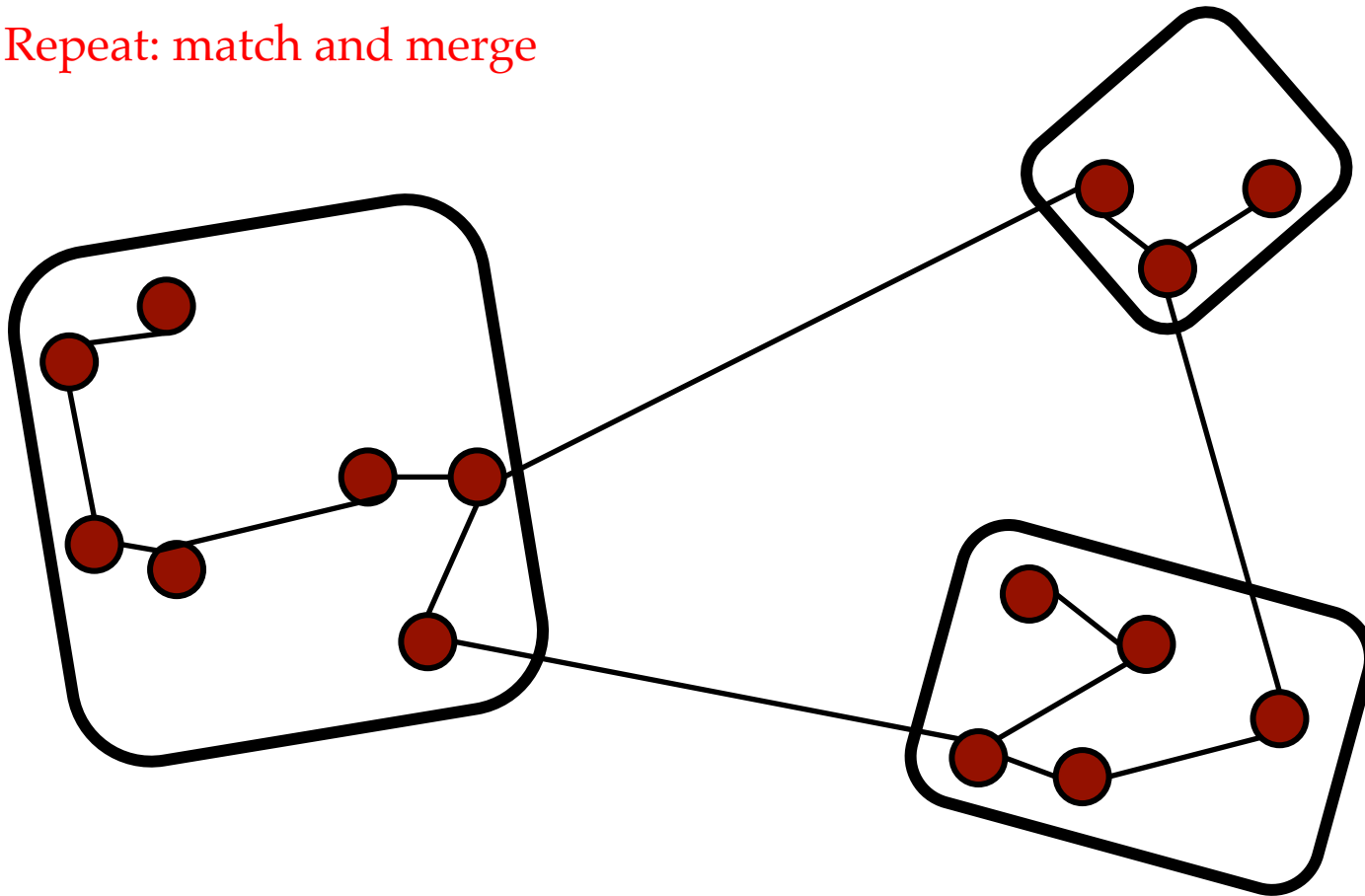
# Divide-and-Conquer Algorithm

Repeat: match and merge



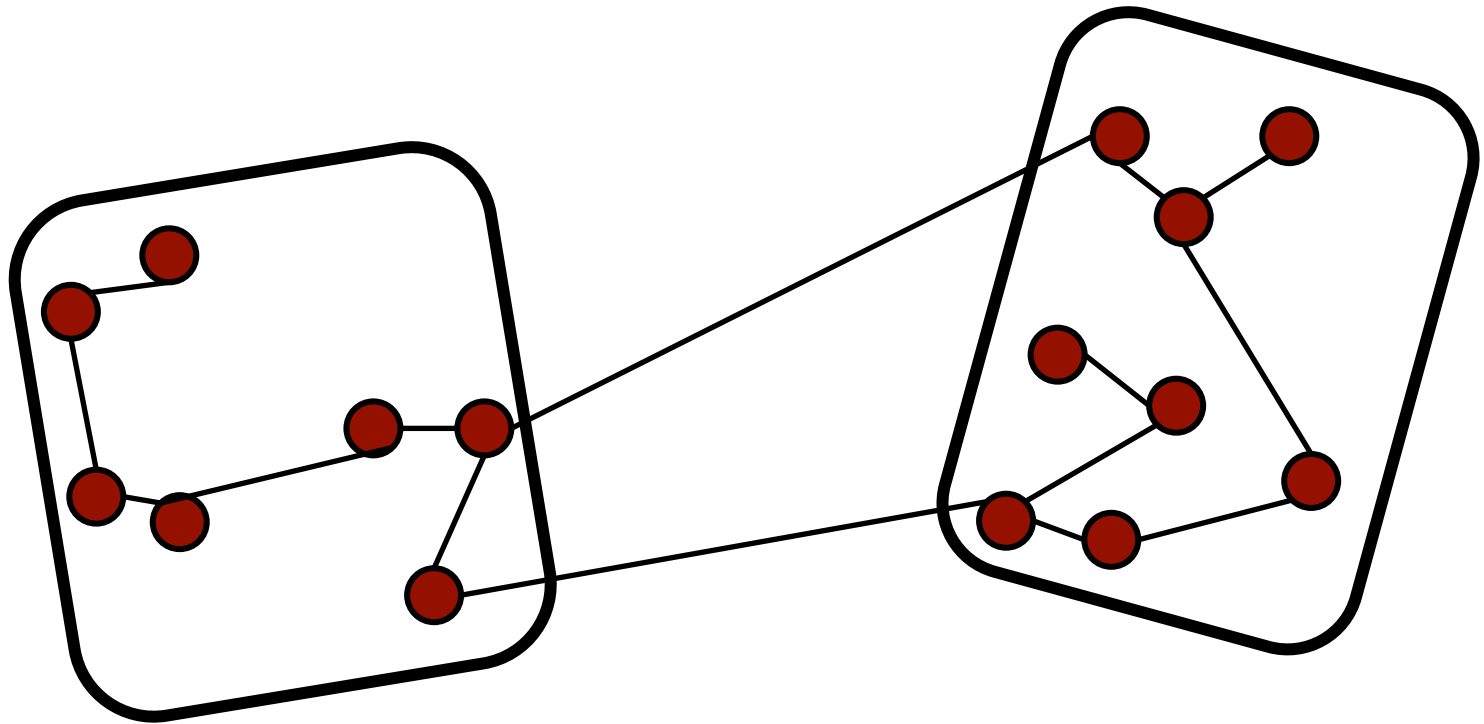
# Divide-and-Conquer Algorithm

Repeat: match and merge



# Divide-and-Conquer Algorithm

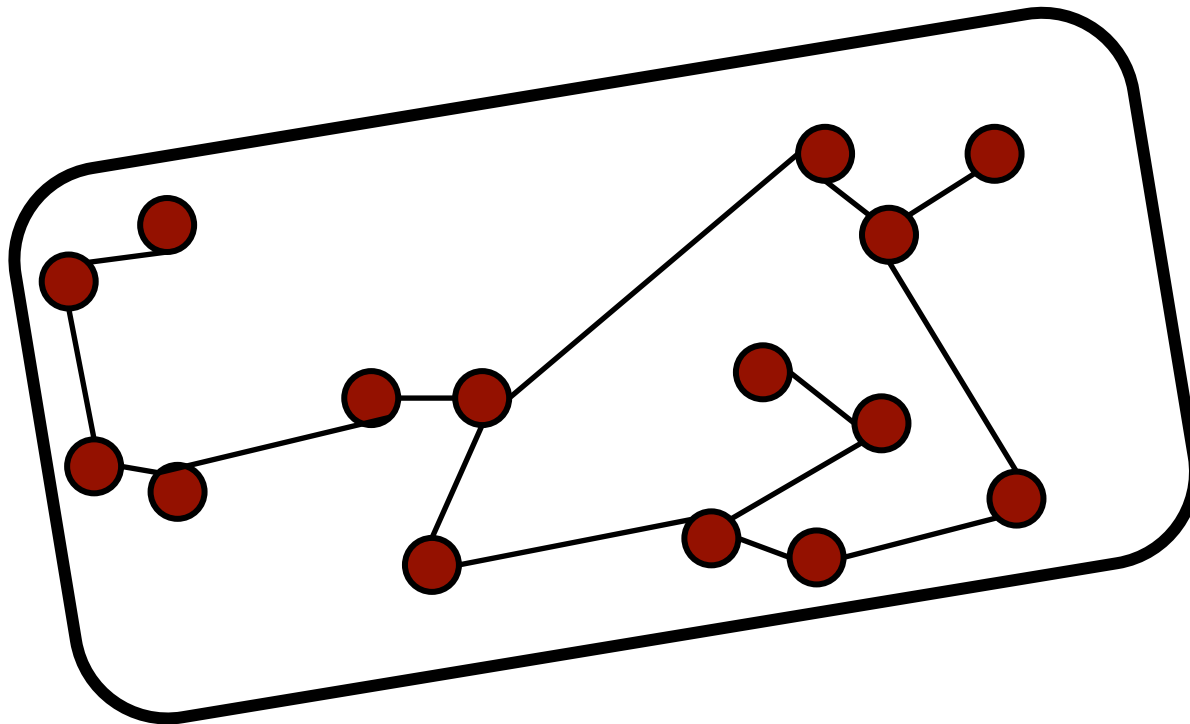
Repeat: match and merge





# Divide-and-Conquer Algorithm

Repeat: match and merge



# Efficiency Analysis

## Merging:

- Random walks:  $O(\log n)$  cost
- Bridge doubling:  $O(\log n)$  iterations
- Overall:  $O(\text{polylog } n)$  time to merge topologies.

## Divide-and-Conquer:

- Collection has small diameter  $\rightarrow O(\log n)$  cost to coordinate
- Collection merging:  $O(\text{polylog } n)$  cost per merge step.
- Number of iterations:  $O(\log n)$  matchings
- Overall:  $O(\text{polylog } n)$  time to form overlay.

# Applications of Overlays

- Peer-to-peer computing
- Multicast Routing
- Delay Tolerant Routing
  - Message will eventually reach the destination
- Social networks
- Small World Effects
- Community Detection Problems
- Graph Searches
- Stable Matching Problems
- Various interaction networks
  - Many more applications ...

# Dijkstra's Algorithm

- Single Source Shortest Path Algorithm proposed by Dijkstra in 1956 (Originally conceived)
- Basic Idea:
  - Two sets are maintained:
  - one set contains vertices included in shortest path tree and the other set includes vertices not yet included in shortest path tree
  - At every step of the algorithm, we find a vertex which is in the other set (set of not yet included) and has a minimum distance from the source
- Similar to Prim's algorithm for MST

# Pseudo Code

Procedure **Dijkstra**(Graph, source)  
begin

    create vertex set Q

    for each vertex v in Graph: // Initialization

        dist[v]  $\leftarrow$  INFINITY // Unknown distance from source to v

        prev[v]  $\leftarrow$  UNDEFINED // Prev. node in optimal path from source

        add v to Q // All nodes initially in Q (unvisited nodes)

    dist[source]  $\leftarrow$  0 // Distance from source to source

    while Q is not empty:

        u  $\leftarrow$  vertex in Q with min dist[u] // Select the node with least distance

        remove u from Q

        for each neighbor v of u: // where v is still in Q.

**weight  $\leftarrow$  dist[u] + length(u, v)**

**if weight < dist[v]: // A shorter path to v has been found**

**dist[v]  $\leftarrow$  weight; prev[v]  $\leftarrow$  u**

    return dist[], prev[]

end

# Key Aspect

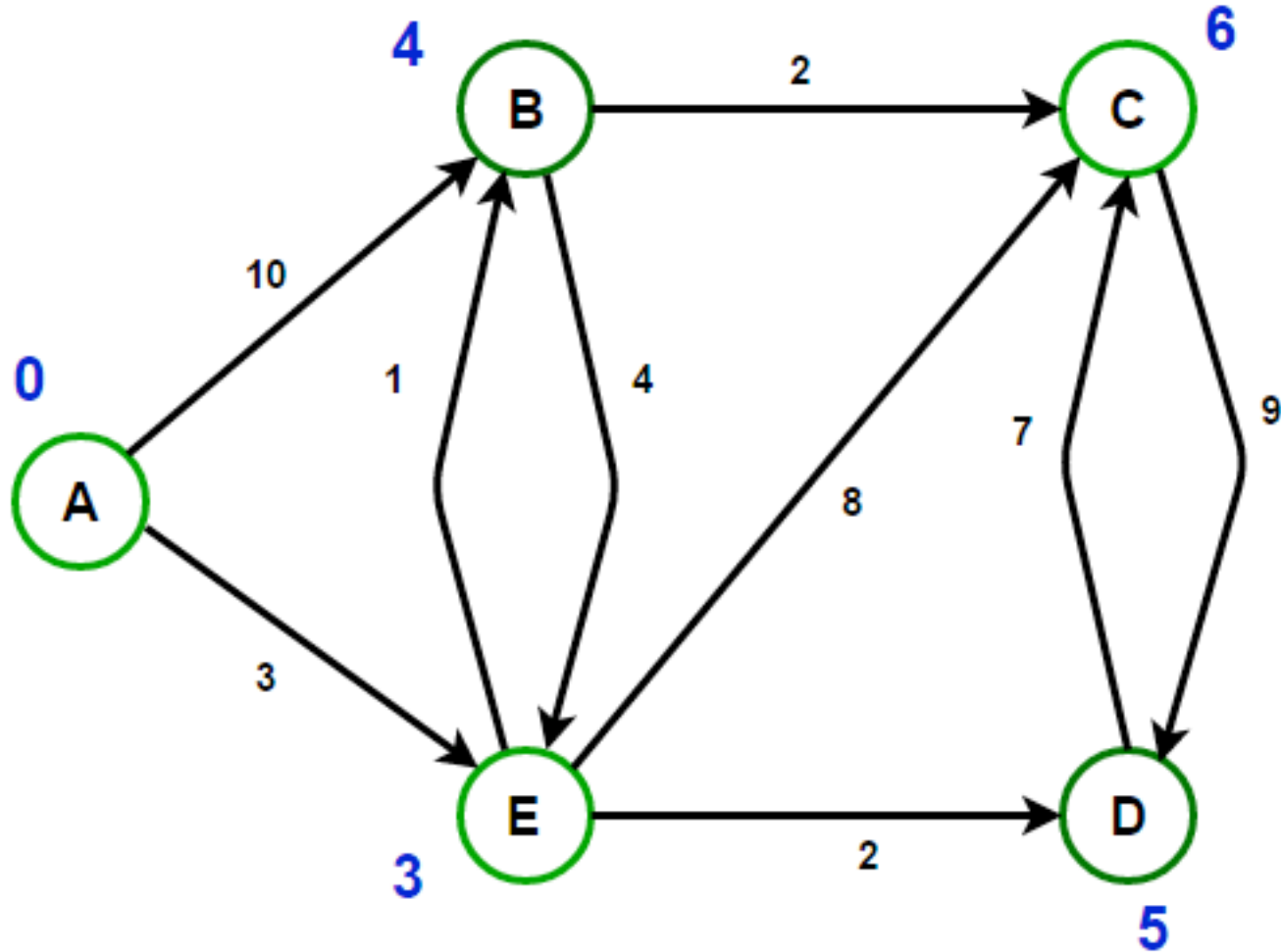
## Relaxation Principle

A shorter path to  $v$  has been found

$\text{weight} \leftarrow \text{dist}[u] + \text{length}(u, v)$

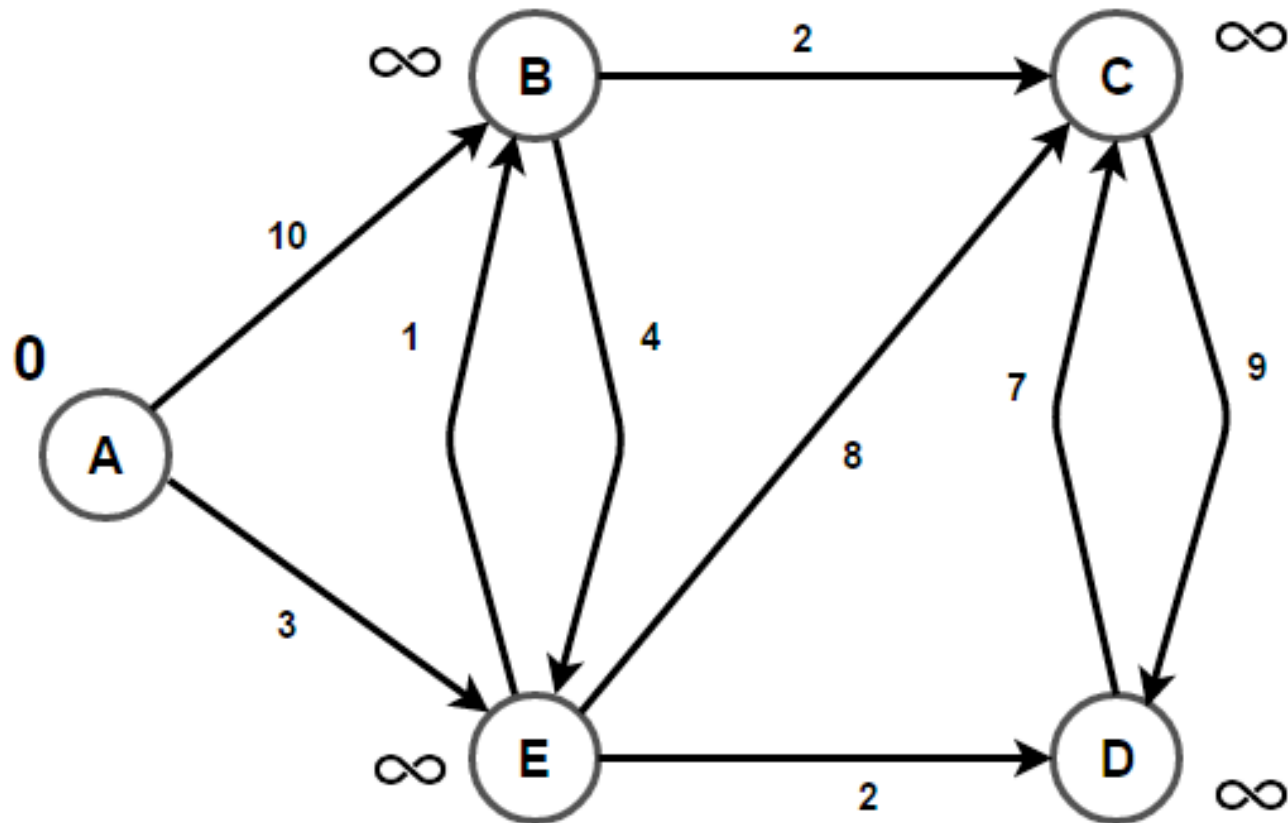
if (  $\text{weight} < \text{dist}[v]$  ) then  
     $\text{dist}[v] \leftarrow \text{weight}$

# Dijkstra's - Illustration



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# Dijkstra's - Illustration

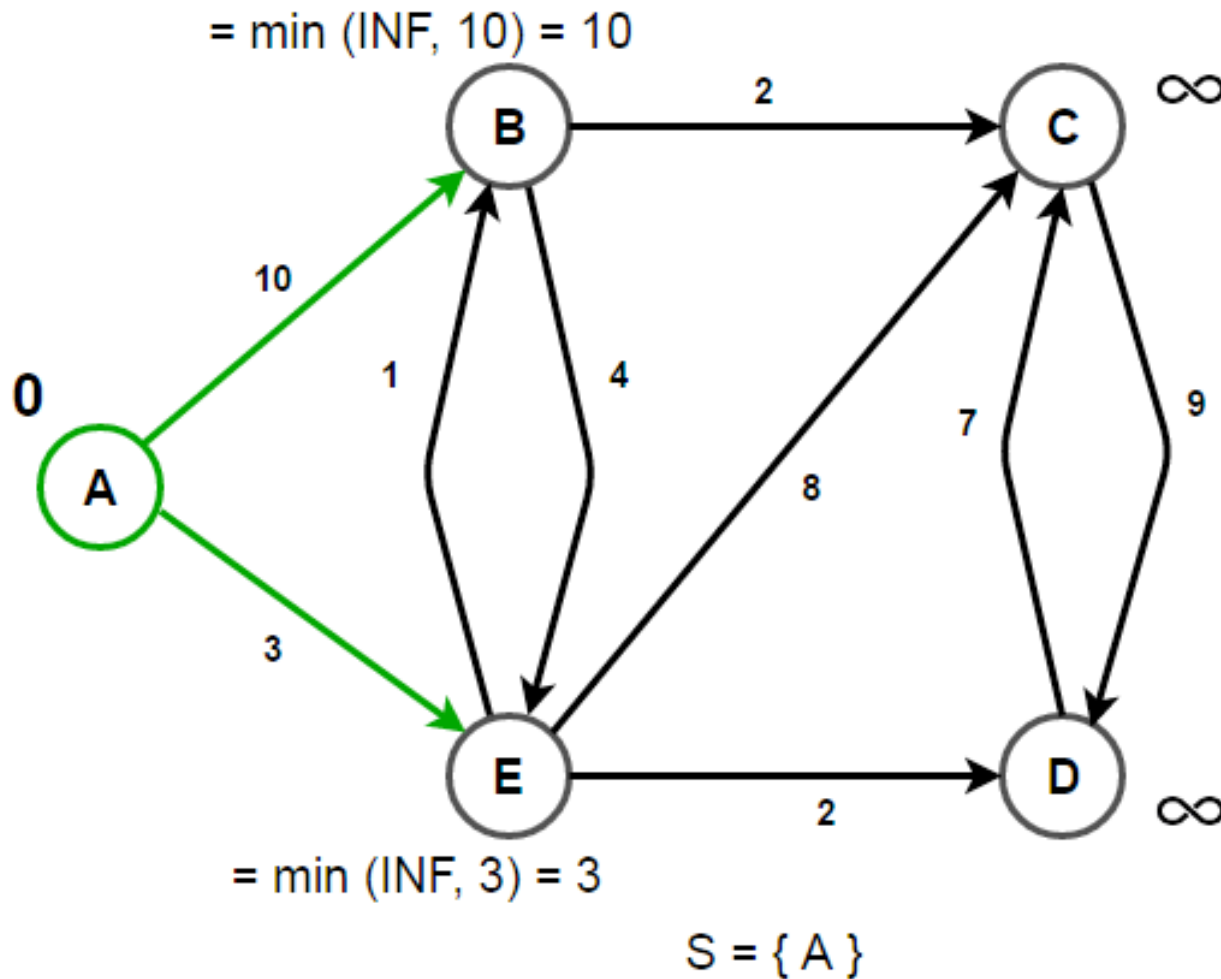


$S = \{\}$

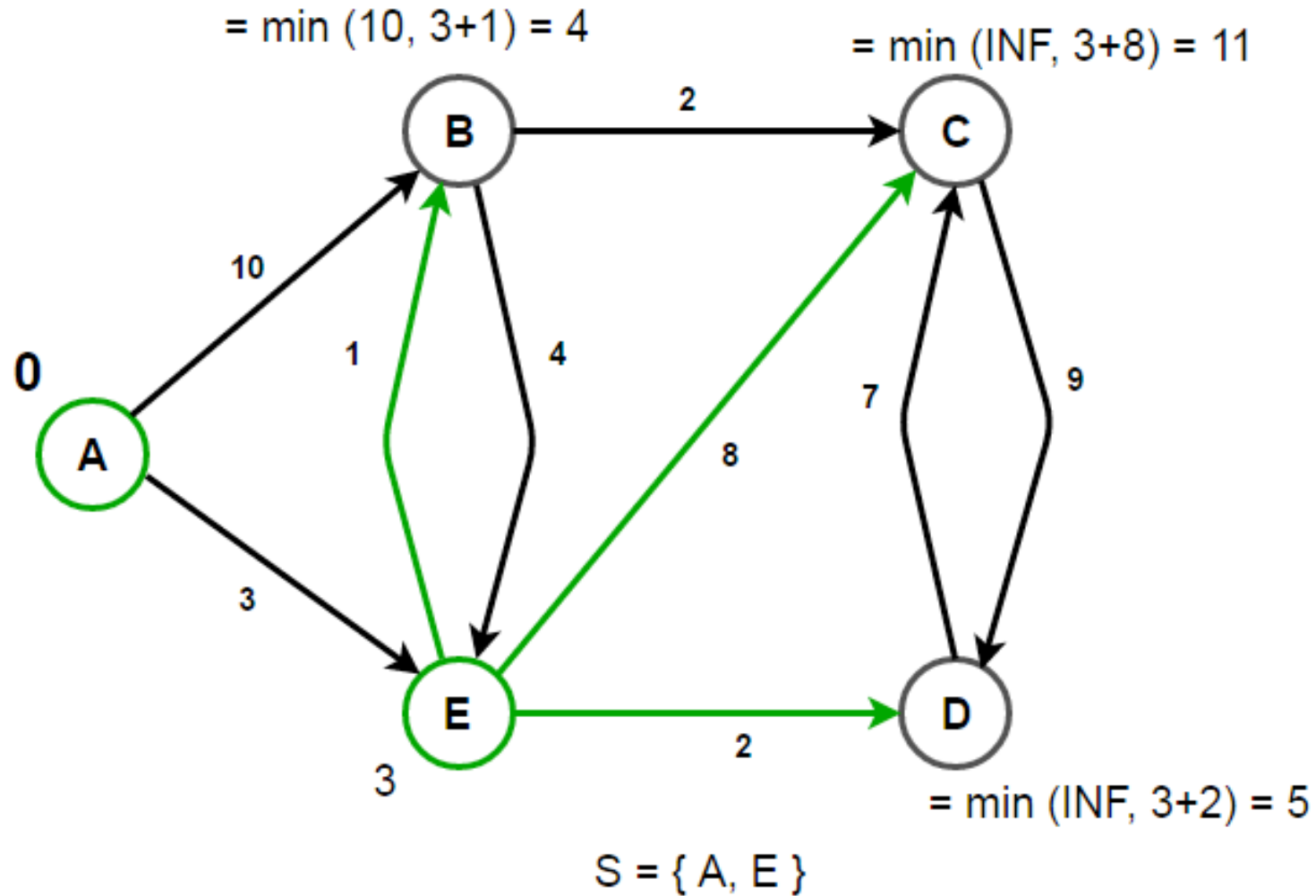
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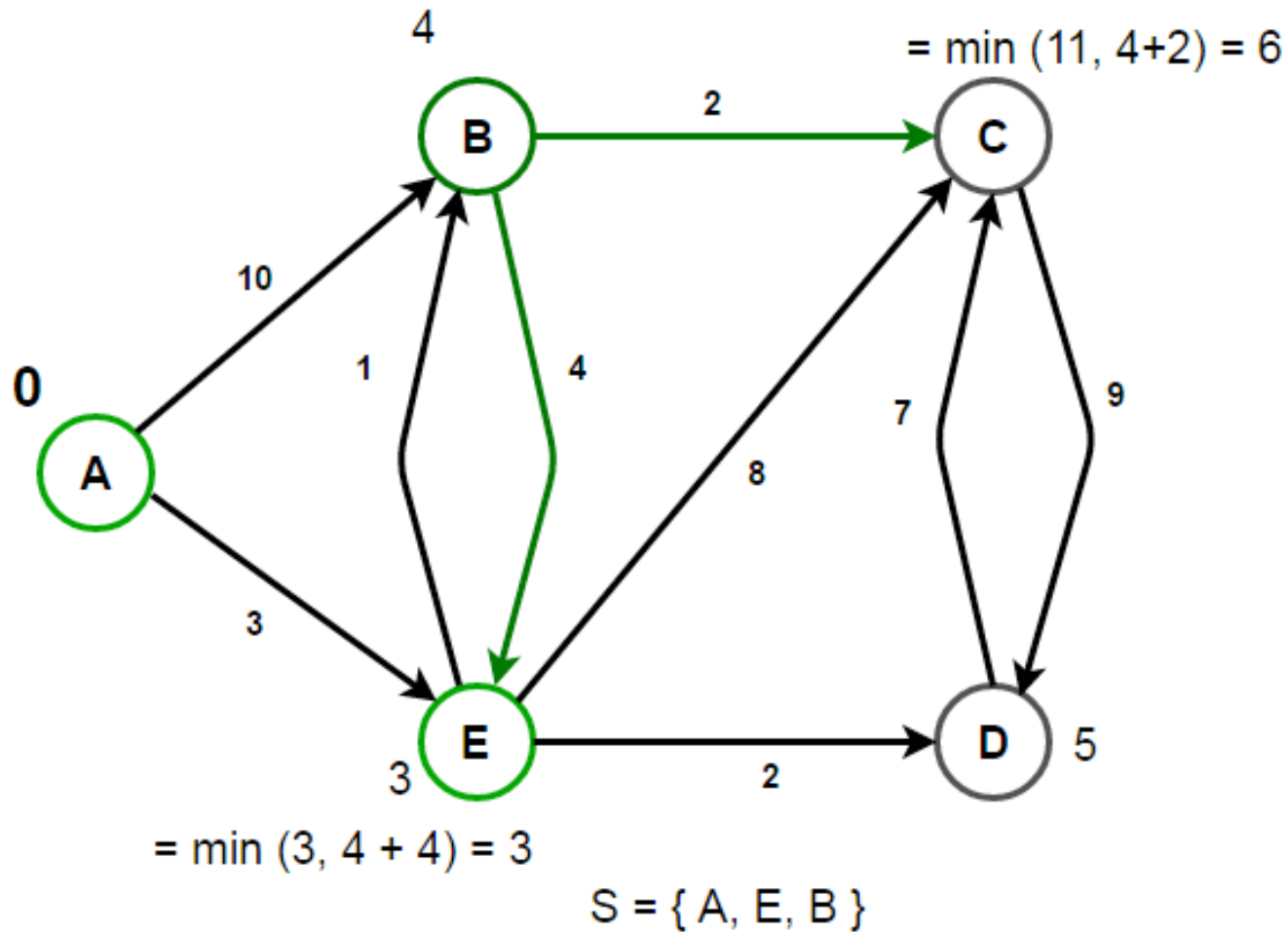
# Dijkstra's - Illustration



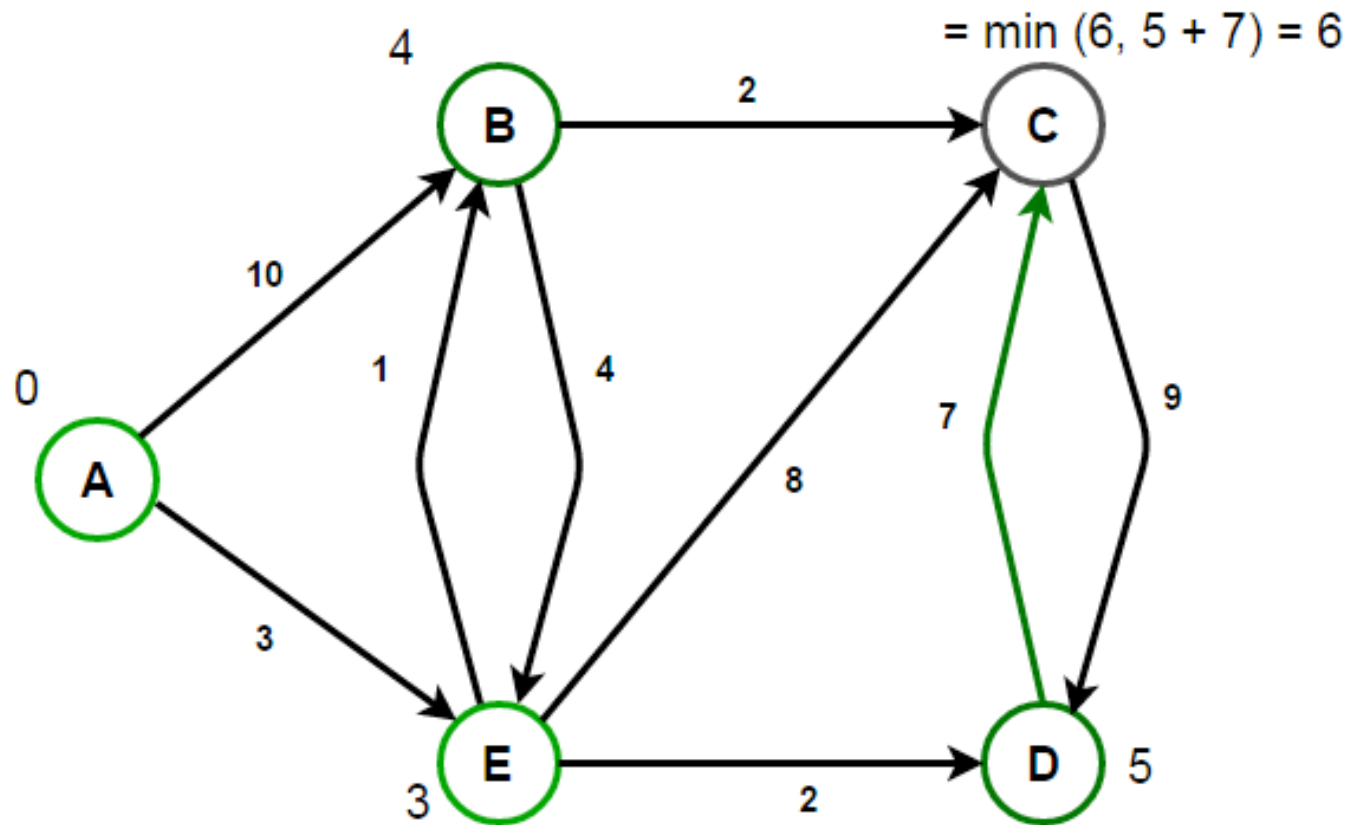
# Dijkstra's - Illustration



# Dijkstra's - Illustration



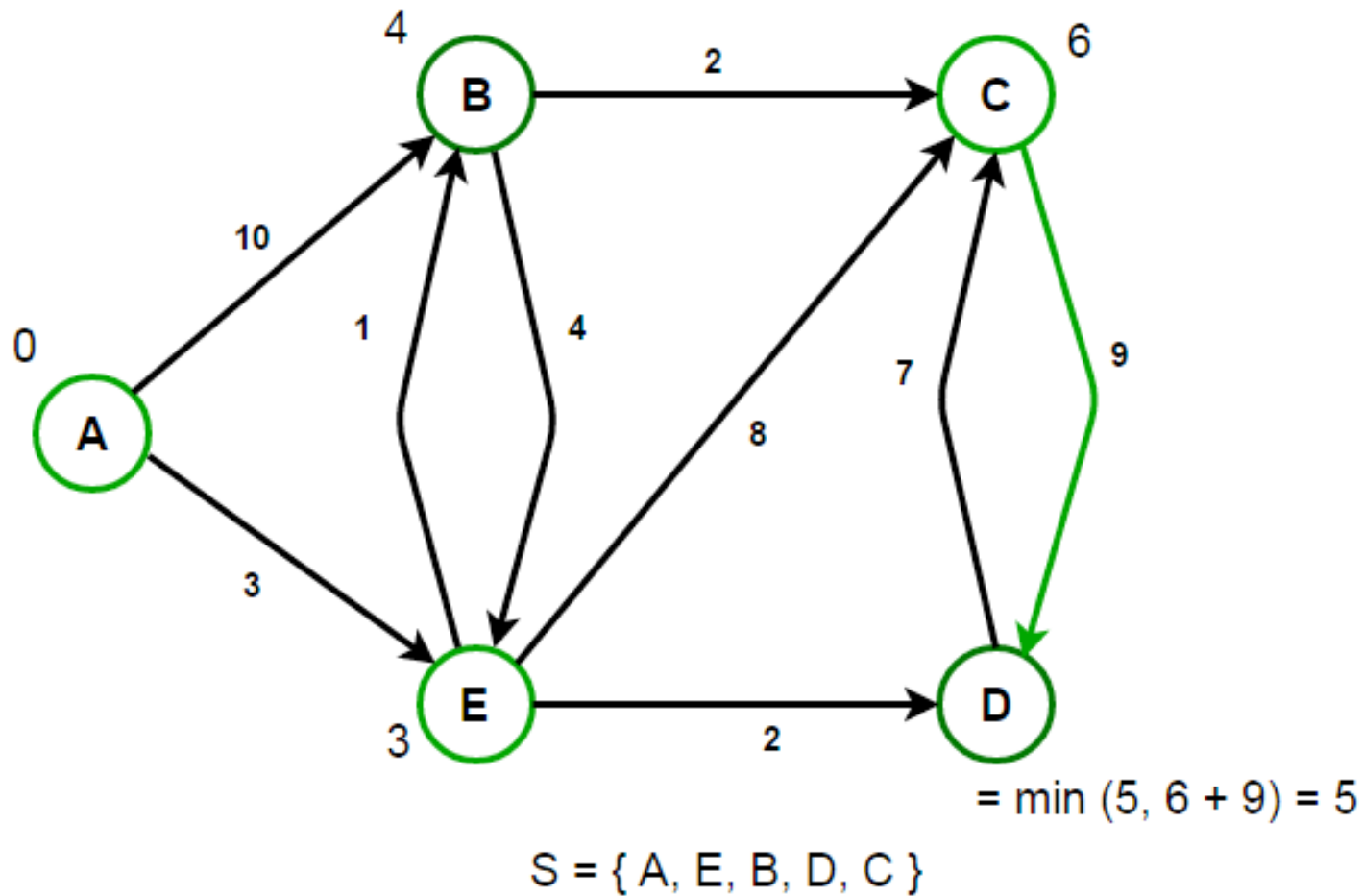
# Dijkstra's - Illustration



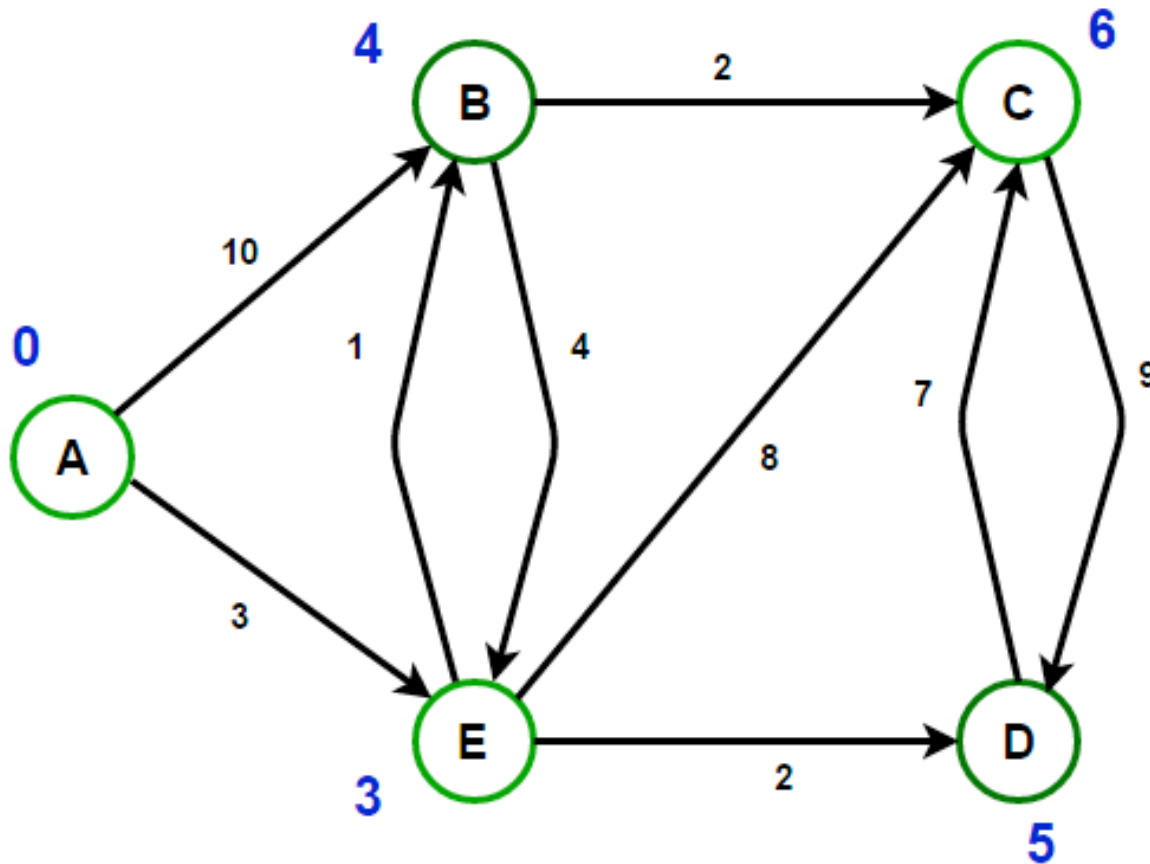
$S = \{A, E, B, D\}$

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# Dijkstra's - Illustration



# Dijkstra's - Illustration



# Important Points

- Dijkstra's algorithm can be used for directed graphs as well
- This algorithm finds shortest distances from source to all vertices
- Time Complexity of the implementation is  $O(V^2)$ 
  - If the input graph is represented using adjacency list, it can be reduced to  $O(E \log V)$  with the help of binary heap
- Dijkstra's algorithm doesn't work for graphs with negative weight edges
  - For graphs with negative weight edges:
    - Bellman-Ford algorithm

# Complexity

- Worst-case performance

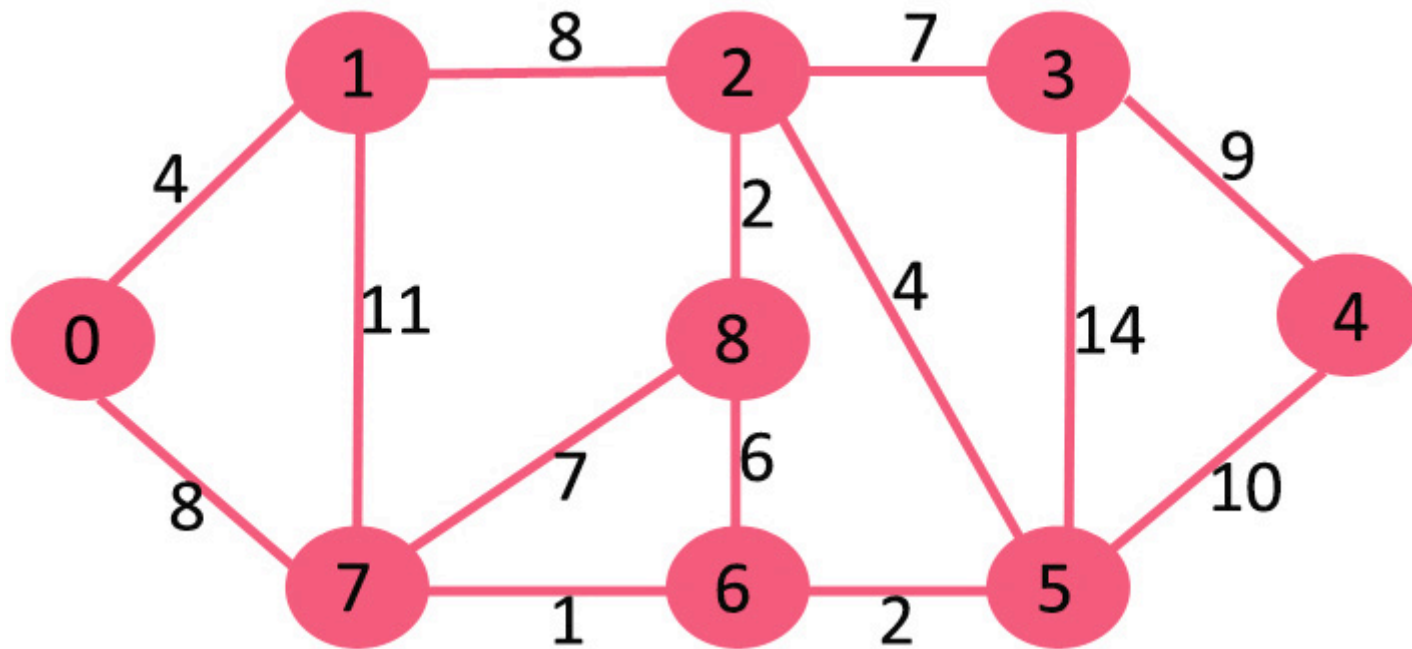
$$( |E| + |V| \log |V| )$$

- Why??
  - Visiting all nodes
  - In every chosen node, checking all nodes satisfying the relaxation principle.
- Try Dijkstra's algorithm  
with a Priority Queue

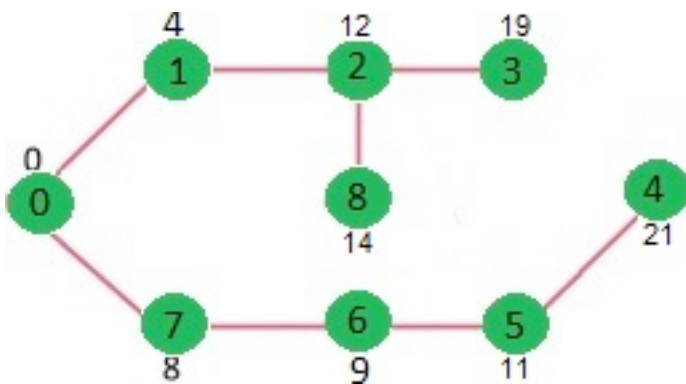
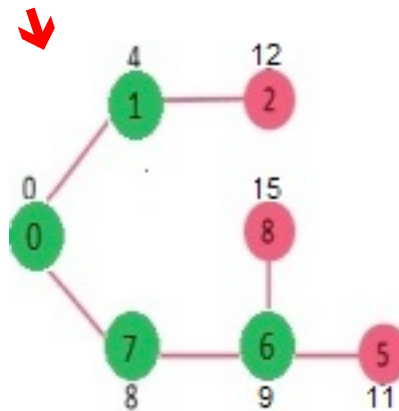
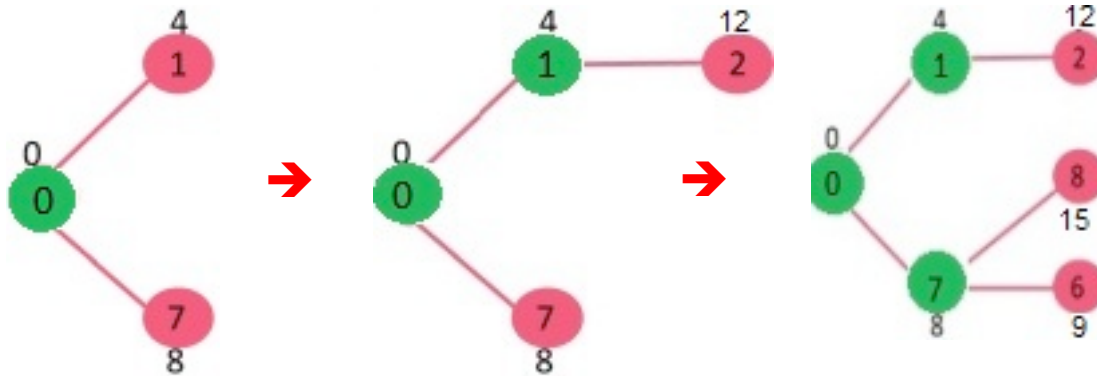
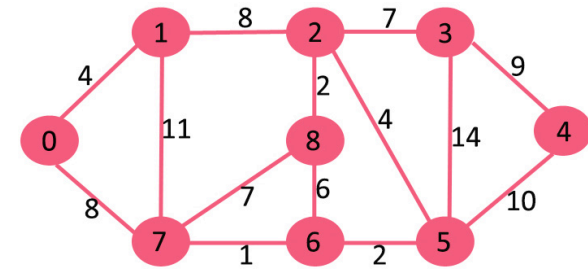


# Exercise

- Apply Dijkstra's Algorithm



# Dijkstra's - Exercise



# Help among Yourselves?

- **Perspective Students** (having CGPA above 8.5 and above)
- **Promising Students** (having CGPA above 6.5 and less than 8.5)
- **Needy Students** (having CGPA less than 6.5)
  - Can the above group help these students? (Your work will also be rewarded)
- You may grow a culture of **collaborative learning** by helping the needy students

# Assistance

- You may post your questions to me at any time
- You may meet me in person on available time or with an appointment
- TA s would assist you to clear your doubts.
- You may leave me an email any time (email is the best way to reach me faster)

# Thanks ...



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