Project 2

Automated Reasoning

1. Basic Model Checking

Model Checking is a simple enumeration method of checking for entailment. In this KB the knowledge base entails α , the query. Model Checking is a method of inference by enumeration. The algorithm takes input clauses of Knowledge Base(KB) followed by query.

```
Input Format: Clause1, Clause2, Clause3, . . . . | = Query
```

The algorithm assigns all possible truth – table (model) assignments and checks whether the Query can be inferred.

```
Output: # When the Query can be inferred
Query can be inferred
Model 1
Model 2
......
# When the Query cannot be inferred
Query cannot be inferred
```

2. Advanced Propositional Inference

The WalkSAT algorithm for SAT (Satisfiability) checking is implemented as Part 2. WalkSAT is a local search algorithm to solve Boolean satisfiability problems. The input to the algorithm are Boolean clauses that are in conjunctive normal form. The clauses used in this project have been converted to CNF before giving them as input. The algorithm starts by assigning a random value to each variable in the formula. Then checks for satisfiability, if the assignment satisfies all clauses, the algorithm terminates, returning the model (assignment). Otherwise, a variable is flipped. The algorithm keeps on performing these operations until a satisfiable assignment is achieved, or the maximum number of flips are achieved.

The algorithm takes the inputs of clauses & query and requires the maximum number of allowable flips, and the flipping probability. Variable is flipped in the following way, first the WalkSAT algorithm picks an unsatisfied clause, then a variable in that clause is flipped. The clause is picked at random among unsatisfied clauses. The variable is picked that will result in the fewest previously satisfied

clauses becoming unsatisfied, with some probability of picking one of the variables at random.

The algorithm may restart with a new random assignment if no solution has been found for too long, as a way of getting out of local minima of numbers of unsatisfied clauses.

In this project the maximum allowable flips are set to 10000 and the probability of random flip is set to 0.5.

```
Input Format: Clause1, Clause2, Clause3, . . . . | = Query
```

The algorithm returns a model if input is satisfiable else returns that Failure

```
Output: # When the Query is satisfiable
Satisfiable
Model
```

When the Query cannot be inferred Failure

When the algorithm returns a "failure" it is either that the sentence is unsatisfiable or that the algorithm could not find a model for the given sentence with the set maximum number of flips.

3. Representation

In this project the propositional logic sentences are taken in the following format: Clause1, Clause2, Clause3,.... |= Query

Each variable, symbol are separated by a space

```
^ - Logical AND
V - Logical OR
-> - Logical Implication
<> - Logical Biconditional
|= is used to separate KB and Query
```

4.Problems

Both the algorithms take the input propositional sentence and display the output. Once an input is tested the programs ask if there are any more sentences to test, to which "y" has to be pressed to test more sentences and "n" to terminate.

Given below are the solutions obtained from running the programs. The screenshots of both the Enumeration Method and WalkSAT program are also shown.

Modus Ponens:

O can be inferred

Enumeration Method:

```
Enter the Proposition Sentence to be tested P,P->Q |= Q Q can be inferred

Model 1 : {'Q': True, 'P': True}

Time taken = 0.0004718019999998546 s

Enter y to check another sentance or n to exit y

Enter the Proposition Sentence to be tested P,P->Q |= ! Q

! Q cannot be inferred

Time taken = 0.000341333000001498 s
```

WalkSAT:

```
Enter the Proposition Sentence to be tested P, ! P v Q |= ! Q Unsatisfiable Time taken = 1.741035058 s

Do you want to check any more sentences (Y/N) y

Enter the Proposition Sentence to be tested P, ! P v Q |= Q Satisfiable

Model = \{'Q': True, 'P': True\}

Time taken = 0.0003572019999964482 s

Do you want to check any more sentences (Y/N) n
```

WUMPUS WORLD (SIMPLE):

There is no pit at location [1,2]

Enumeration Method:

WalkSAT:

HORN CLAUSES

Unicorn is Magical and Horned, but Mythical cannot be inferred.

Enumeration Method:

```
Enter y to check another sentance or n to exit y
Enter the Proposition Sentence to be tested Mythical -> ! Mortal , ! Mythical -> Mortal ^ Mammal , ! Mortal v Mammal -> Horned , Horned -> Magical |= Mythical mythical cannot be inferred

Time taken = 0.00:20935710000033935 s
Enter y to check another sentance or n to exit y
Enter the Proposition Sentence to be tested Mythical -> ! Mortal , ! Mythical -> Mortal ^ Mammal , ! Mortal v Mammal -> Horned , Horned -> Magical |= ! Mythical interaction of the proposition sentence to be tested Mythical -> ! Mortal , ! Mythical -> Mortal ^ Mammal , ! Mortal v Mammal -> Horned , Horned -> Magical |= ! Mythical interaction of the proposition sentence to be tested Mythical -> ! Mortal , ! Mythical -> Mortal ^ Mammal , ! Mortal v Mammal -> Horned , Horned -> Magical |= Magical can be inferred

Model 1 : ('Magical': True, 'Horned': True, 'Mortal': True, 'Mythical': True, 'Mammal': True}

Model 2 : ('Magical': True, 'Horned': True, 'Mortal': False, 'Mythical': True, 'Mammal': True}

Model 3 : ('Magical': True, 'Horned': True, 'Mortal': False, 'Mythical': True, 'Mammal': False})

Enter y to check another sentance or n to exit y
Enter the Proposition Sentence to be tested Mythical -> ! Mortal , ! Mythical -> Mortal ^ Mammal , ! Mortal v Mammal -> Horned , Horned -> Magical |= Horned Horned can be inferred

Model 1 : ('Magical': True, 'Horned': True, 'Mortal': True, 'Mythical': False, 'Mammal': True})

Model 2 : ('Magical': True, 'Horned': True, 'Mortal': False, 'Mythical': True, 'Mammal': True})

Model 3 : ('Magical': True, 'Horned': True, 'Mortal': False, 'Mythical': True, 'Mammal': True})

Model 3 : ('Magical': True, 'Horned': True, 'Mortal': False, 'Mythical': True, 'Mammal': False})

Time taken = 0.01088670499999116 s
```

WalkSAT:

4. LIARS AND TRUTH-TELLERS

- a. Amy is a liar (Amy cannot be inferred and ! Amy can be inferred)
 Bob is a liar (Bob cannot be inferred and ! Bob can be inferred)
 Cal tells the truth (Cal cannot be inferred)
- b. Amy tells the truth (Amy can be inferred)Bob is a liar (Bob cannot be inferred and ! Bob can be inferred)Cal is a liar (Cal cannot be inferred and ! Cal can be inferred)

Enumeration Method:

a.

```
Enter the Proposition Sentence to be tested
                                                 Amy <> Cal ^ Amy , Bob <> ! Cal , Cal <> Bob v ! Amy |= Amy
Amy cannot be inferred
Time taken = 0.0006675050000000127 s
Enter y to check another sentance or n to exit
                                                 Amy <> Cal ^ Amy , Bob <> ! Cal , Cal <> Bob v ! Amy |= ! Amy
Enter the Proposition Sentence to be tested
! Amy can be inferred
Model 1 : {'Cal': True, 'Bob': False, 'Amy': False}
Time taken = 0.0016668220000006784 s
Enter y to check another sentance or n to exit
Enter the Proposition Sentence to be tested
                                                 Amy <> Cal ^ Amy , Bob <> ! Cal , Cal <> Bob v ! Amy |= Bob
Bob cannot be inferred
Time taken = 0.0007993830000003754 s
Enter y to check another sentance or n to exit
Enter the Proposition Sentence to be tested
                                                 Amy <> Cal ^ Amy , Bob <> ! Cal , Cal <> Bob v ! Amy |= ! Bob
! Bob can be inferred
Model 1 : {'Cal': True, 'Bob': False, 'Amy': False}
Time taken = 0.0013491129999998464 s
Enter y to check another sentance or n to exit
Enter the Proposition Sentence to be tested
                                                 Amy <> Cal ^ Amy , Bob <> ! Cal , Cal <> Bob v ! Amy |= Cal
Cal can be inferred
Model 1 : {'Cal': True, 'Bob': False, 'Amy': False}
Time taken = 0.001353344999998285 s
Enter y to check another sentance or n to exit n
```

b.

```
Enter the Proposition Sentence to be tested
                                                Amy <> ! Cal , Bob <> Amy ^ Cal , Cal <> Bob |= Amy
Amy can be inferred
Model 1 : {'Bob': False, 'Cal': False, 'Amy': True}
Time taken = 0.0011058080000001524 s
Enter y to check another sentance or n to exit
                                                Amy <> ! Cal , Bob <> Amy ^ Cal , Cal <> Bob |= Bob
Enter the Proposition Sentence to be tested
Bob cannot be inferred
Time taken = 0.0008360560000006956 s
Enter y to check another sentance or n to exit
                                                Amy <> ! Cal , Bob <> Amy ^ Cal , Cal <> Bob |= ! Bob
Enter the Proposition Sentence to be tested
! Bob can be inferred
Model 1 : {'Bob': False, 'Cal': False, 'Amy': True}
Time taken = 0.0010359899999983213 s
Enter y to check another sentance or n to exit
                                                Amy <> ! Cal , Bob <> Amy ^ Cal , Cal <> Bob |= Cal
Enter the Proposition Sentence to be tested
Cal cannot be inferred
Time taken = 0.0008466339999984029 s
Enter y to check another sentance or n to exit
                                                Amy <> ! Cal , Bob <> Amy ^ Cal , Cal <> Bob |= ! Cal
Enter the Proposition Sentence to be tested
! Cal can be inferred
Model 1 : {'Bob': False, 'Cal': False, 'Amy': True}
Time taken = 0.0010201209999962657 s
Enter y to check another sentance or n to exit n
```

WalkSAT:

a.

```
! Amy v ! Cal , ! Bob v ! Cal ^ Bob v Cal , ! Cal v Bob v ! Amy ^ Cal v ! Bob ^ Cal v Amy |= Amy
              oposition Sentence to be tested
 nsatisfiable<sup>'</sup>
Time taken = 3.8426244449999998 s
Do you want to check any more sentences (Y/N)
Enter the Proposition Sentence to be tested
                                                         ! Amy v ! Cal , ! Bob v ! Cal ^ Bob v Cal , ! Cal v Bob v ! Amy ^ Cal v ! Bob ^ Cal v Amy | = ! Amy
Model = {'Bob': False, 'Amy': False, 'Cal': True}
Time taken = 0.000639294999991515 s
Do you want to check any more sentences (Y/N)
                                                         ! Amy v ! Cal , ! Bob v ! Cal ^ Bob v Cal , ! Cal v Bob v ! Amy ^ Cal v ! Bob ^ Cal v Amy |= Bob
Enter the Proposition Sentence to be tested
Unsatisfiable
Time taken = 3.9557613430000025 s
Do \, you want to check any more sentences (Y/N)
                                                         ! Amy v ! Cal , ! Bob v ! Cal ^ Bob v Cal , ! Cal v Bob v ! Amy ^ Cal v ! Bob ^ Cal v Amy |= ! Bob
      the Proposition Sentence to be tested
Satisfiable
Model = {'Bob': False, 'Amy': False, 'Cal': True}
Time taken = 0.001334657000001016 s
Do you want to check any more sentences (Y/N) Enter the Proposition Sentence to be tested
                                                         ! Amy v ! Cal , ! Bob v ! Cal ^ Bob v Cal , ! Cal v Bob v ! Amy ^ Cal v ! Bob ^ Cal v Amy |= Cal
Satisfiable
Model = {'Bob': False, 'Amy': False, 'Cal': True}
Time taken = 0.000838524000023941 s
    you want to check any more sentences (Y/N)
                                                         ! Amy v ! Cal , ! Bob v ! Cal ^ Bob v Cal , ! Cal v Bob v ! Amy ^ Cal v ! Bob ^ Cal v Amy |= ! Cal
Enter the Proposition Sentence to be tested
 Jnsatisfiable
 ime taken = 4.9136013639999945 s
Do \, you want to check any more sentences (Y/N) \,
```

5. Conclusion

As one can see WalkSAT algorithm is less computationally costly than the enumeration method. But when a query is unsatisfiable the WalkSAT algorithm takes a lot of time when compared to enumeration method.

The ability of the WalkSAT algorithm to find a model depends on the maximum allowable flips and the probability of flipping (randomly), the chances of finding the model increase as we increase the maximum allowable flips and the probability of flipping.

6. REFERENCES:

- **1.** Russell, S. and Norvig, P., 2010. *Artificial Intelligence: A Modern Approach* 3rd ed.,
- 2. WalkSAT. En.wikipedia.org.

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