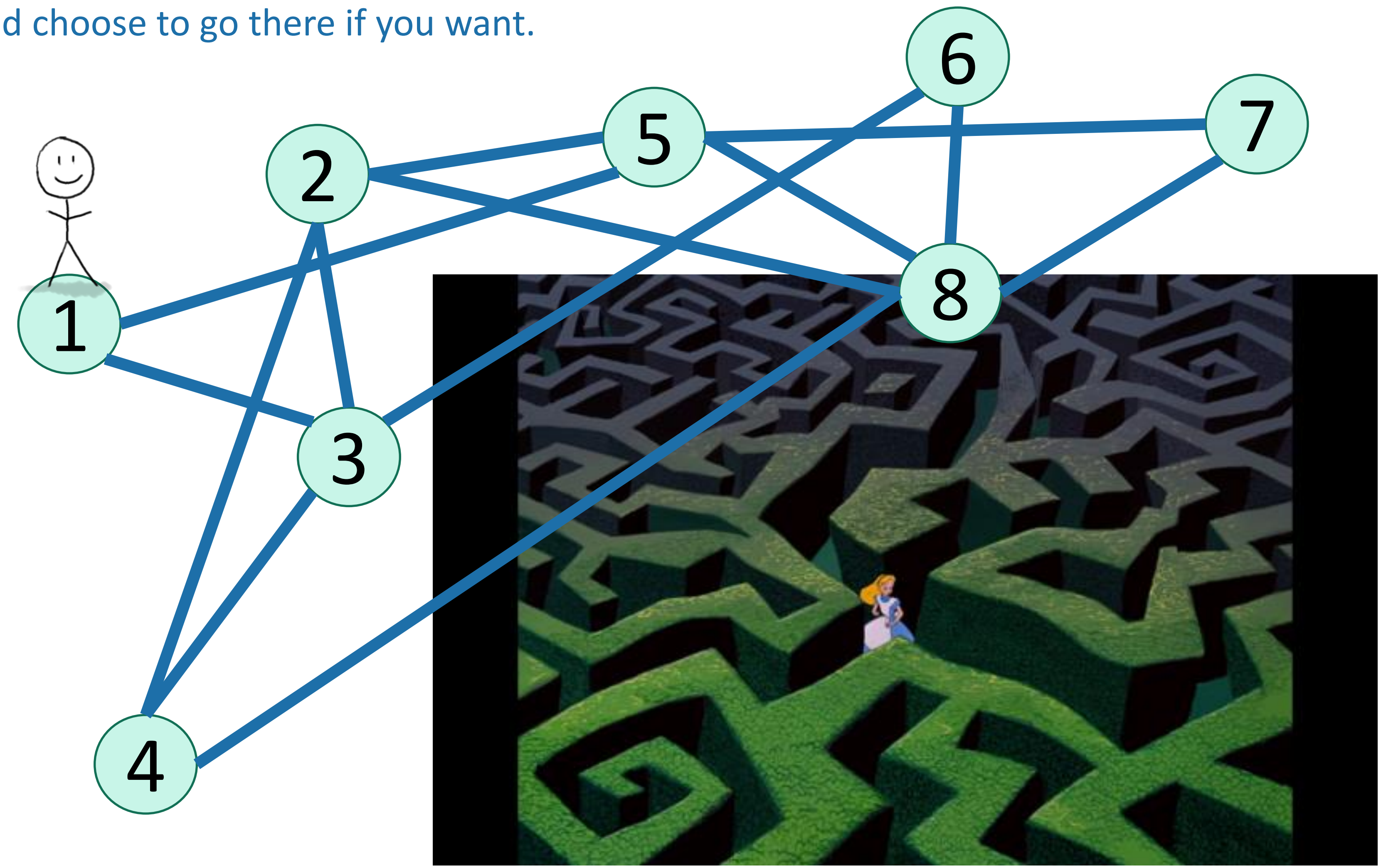
Advanced Data Structures and Algorithms

Depth First Search (DFS)

Depth-first search

How do we explore a graph?

At each node, you can get a list of neighbors, and choose to go there if you want.



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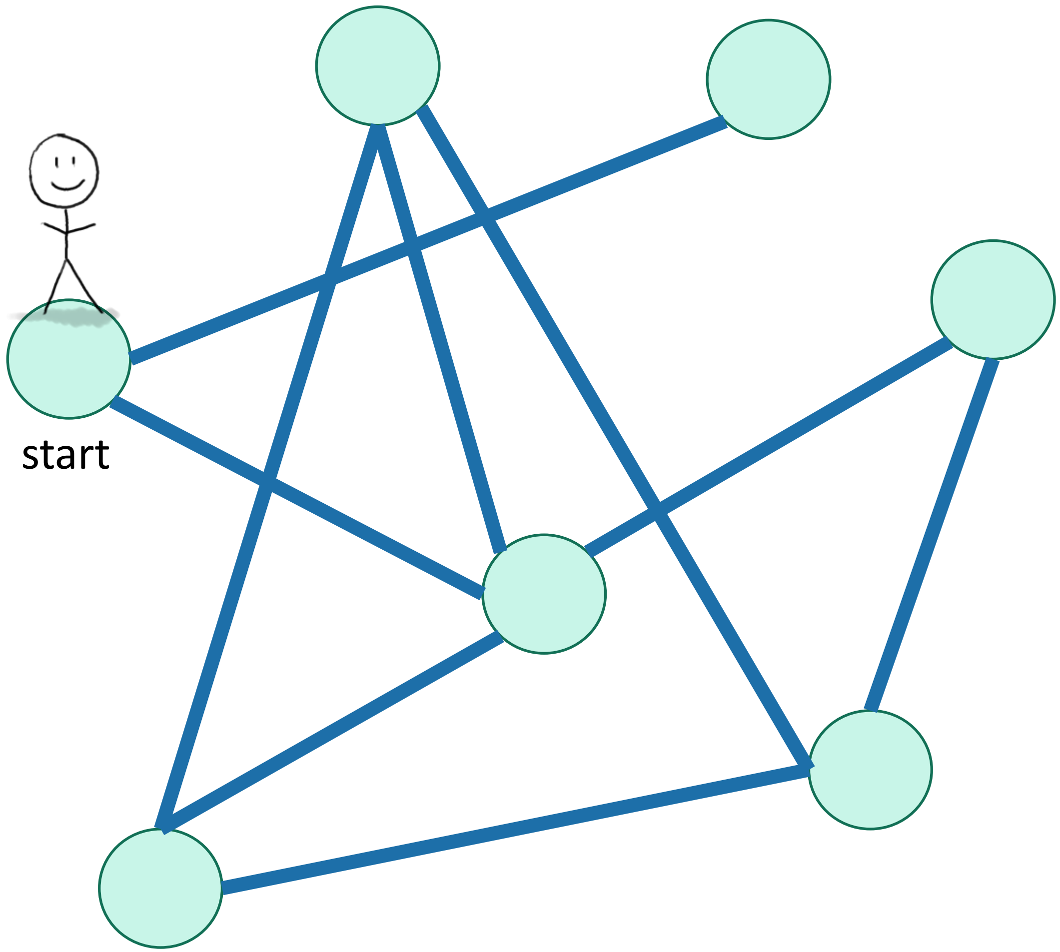

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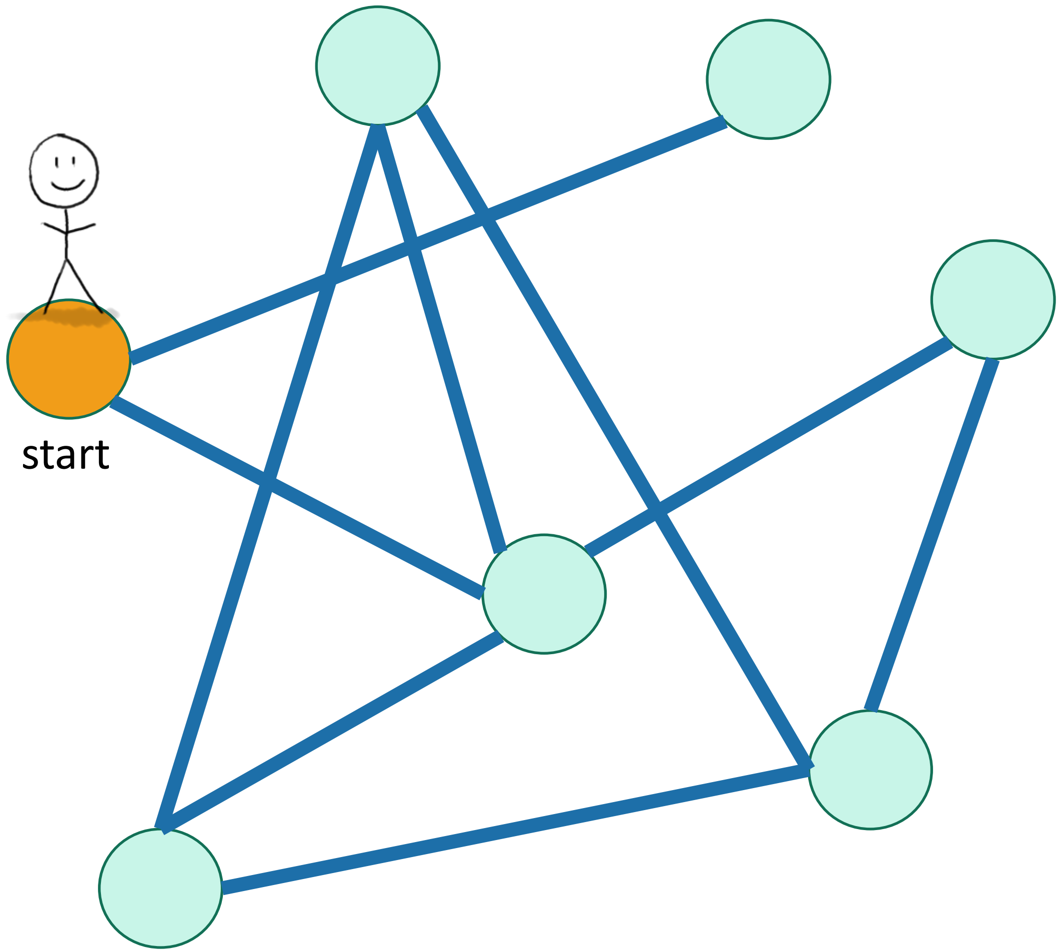

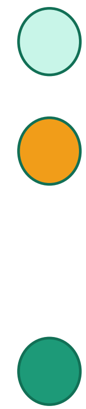

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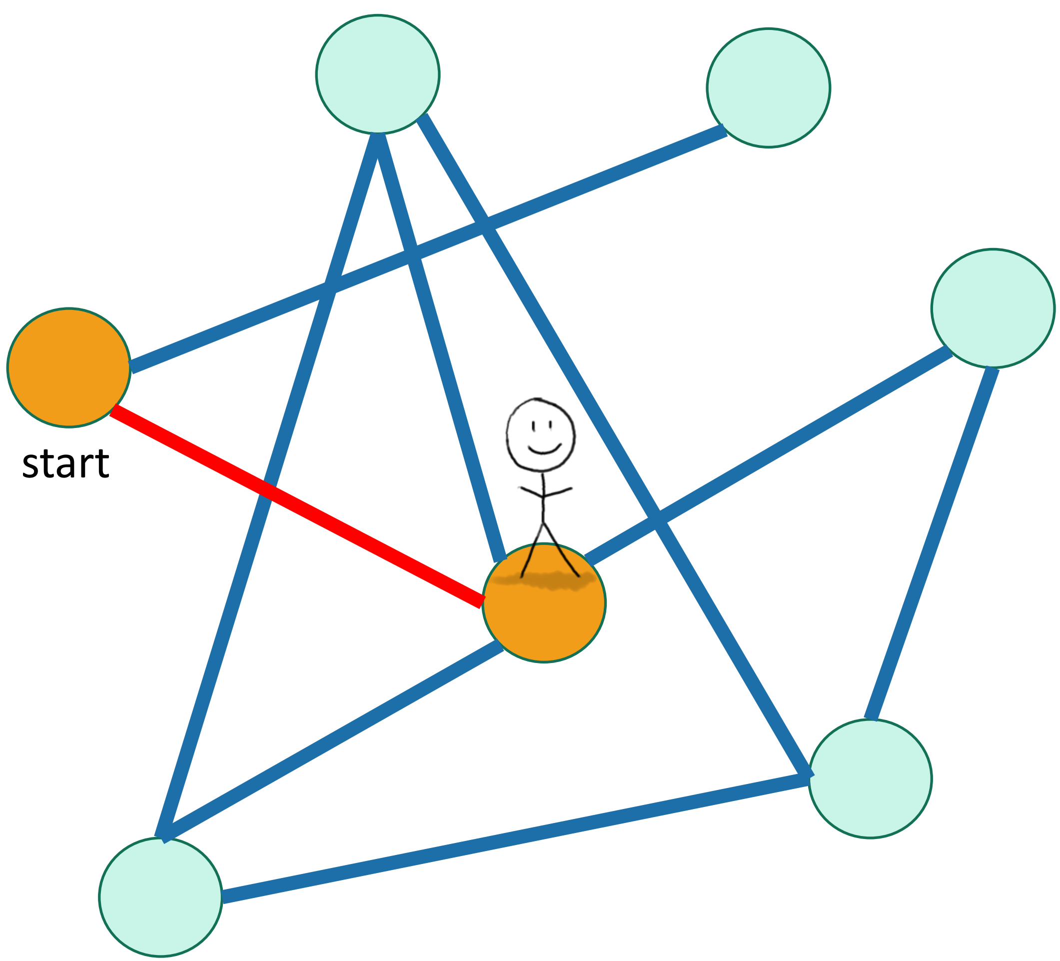

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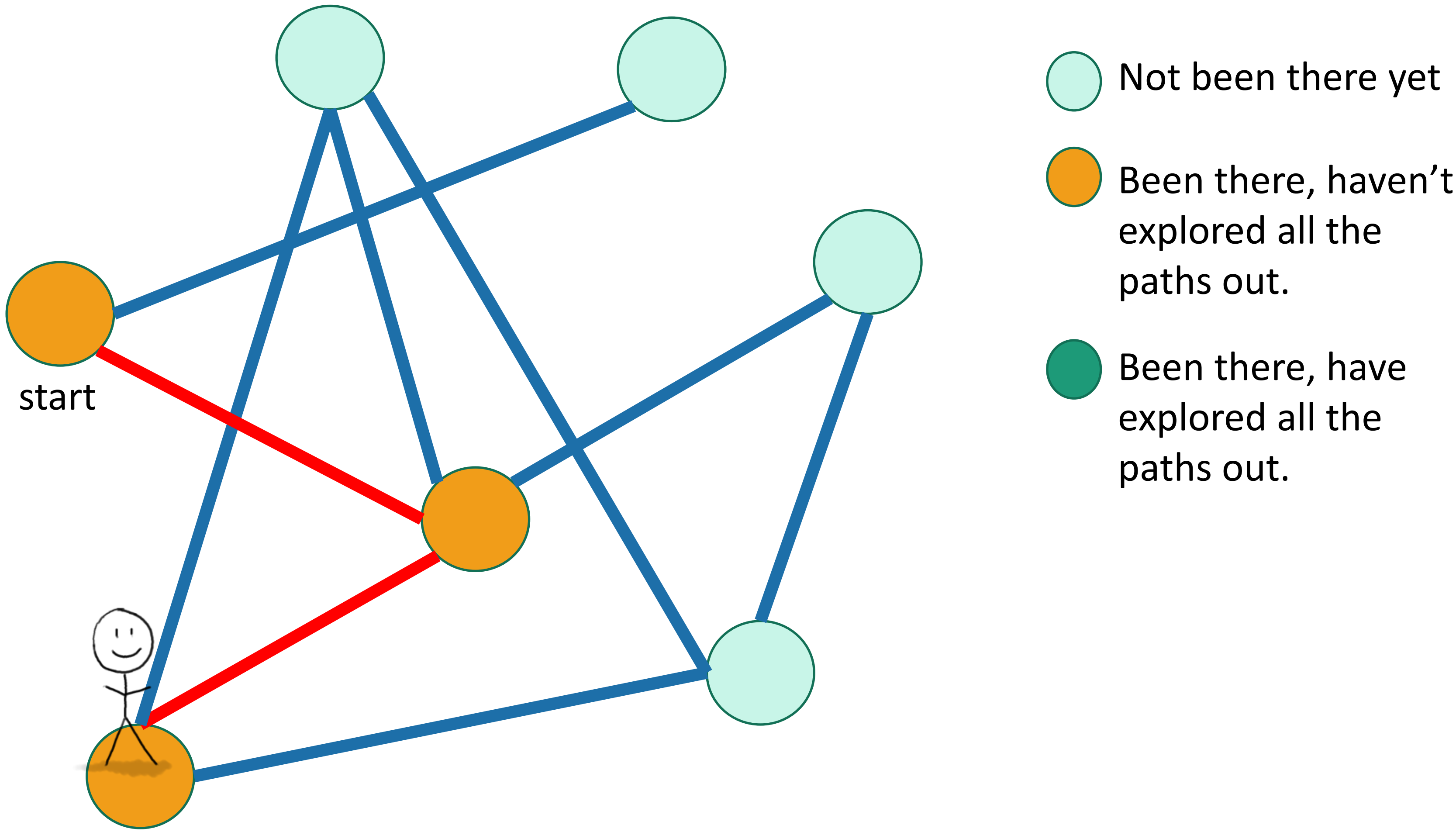

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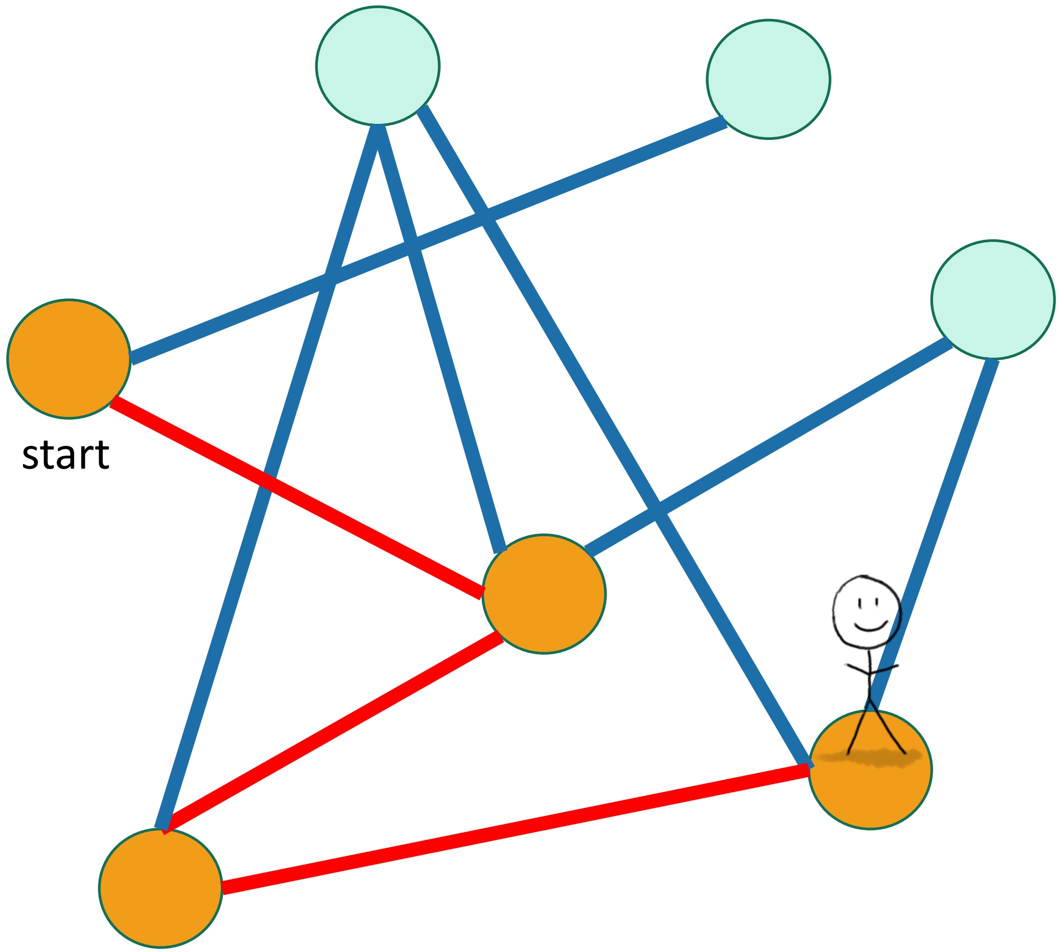

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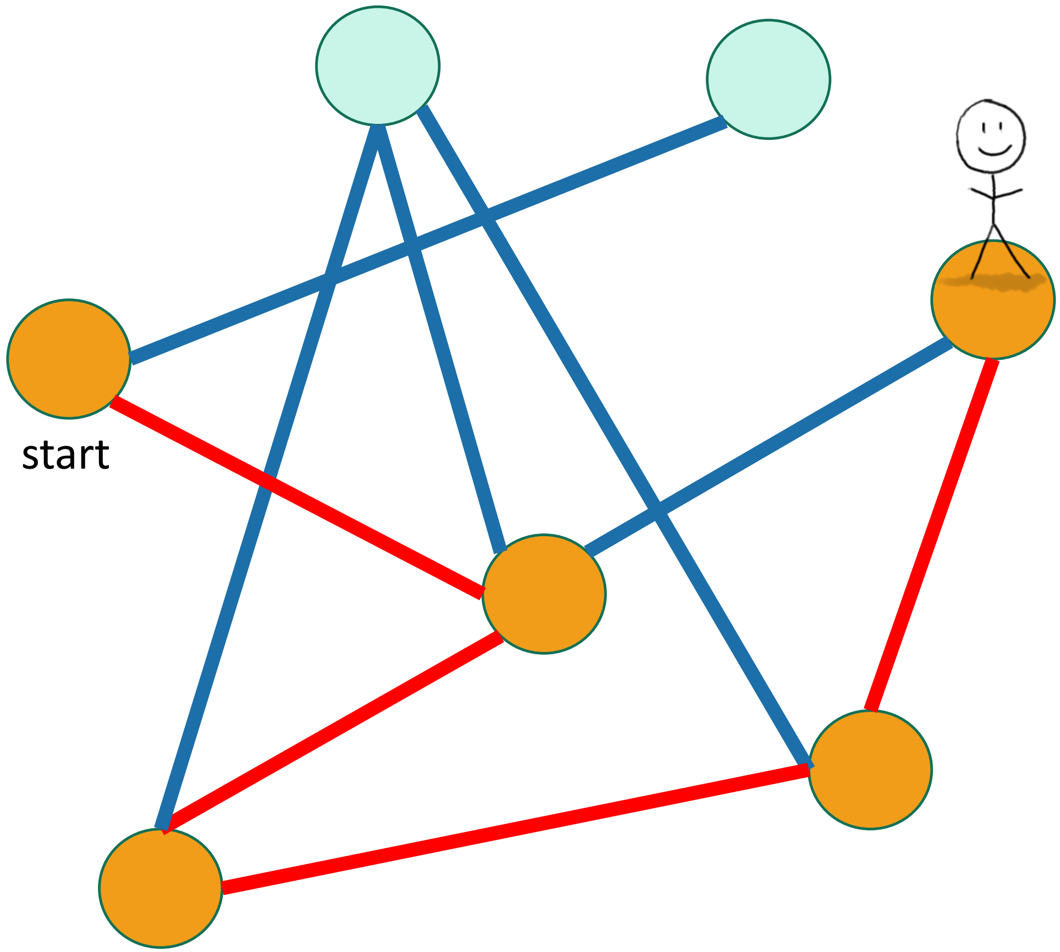

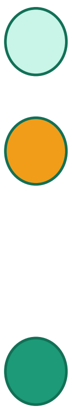

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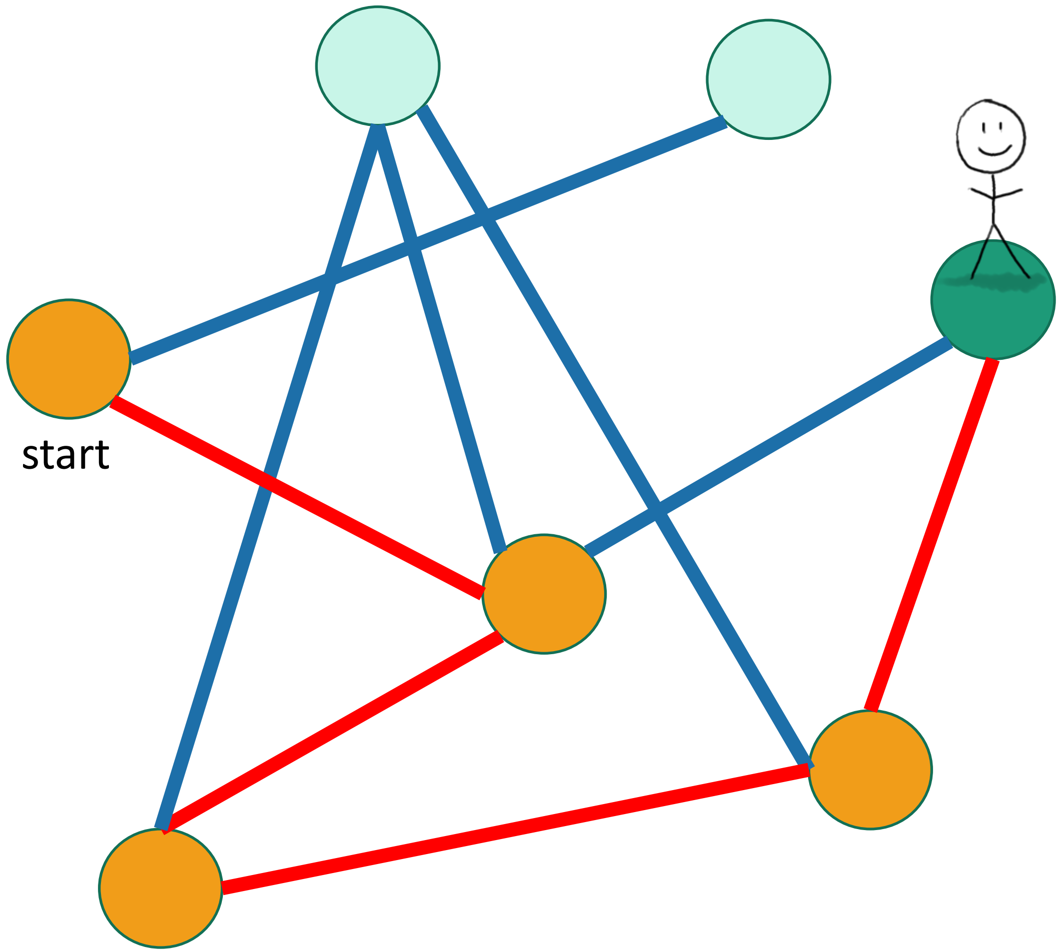

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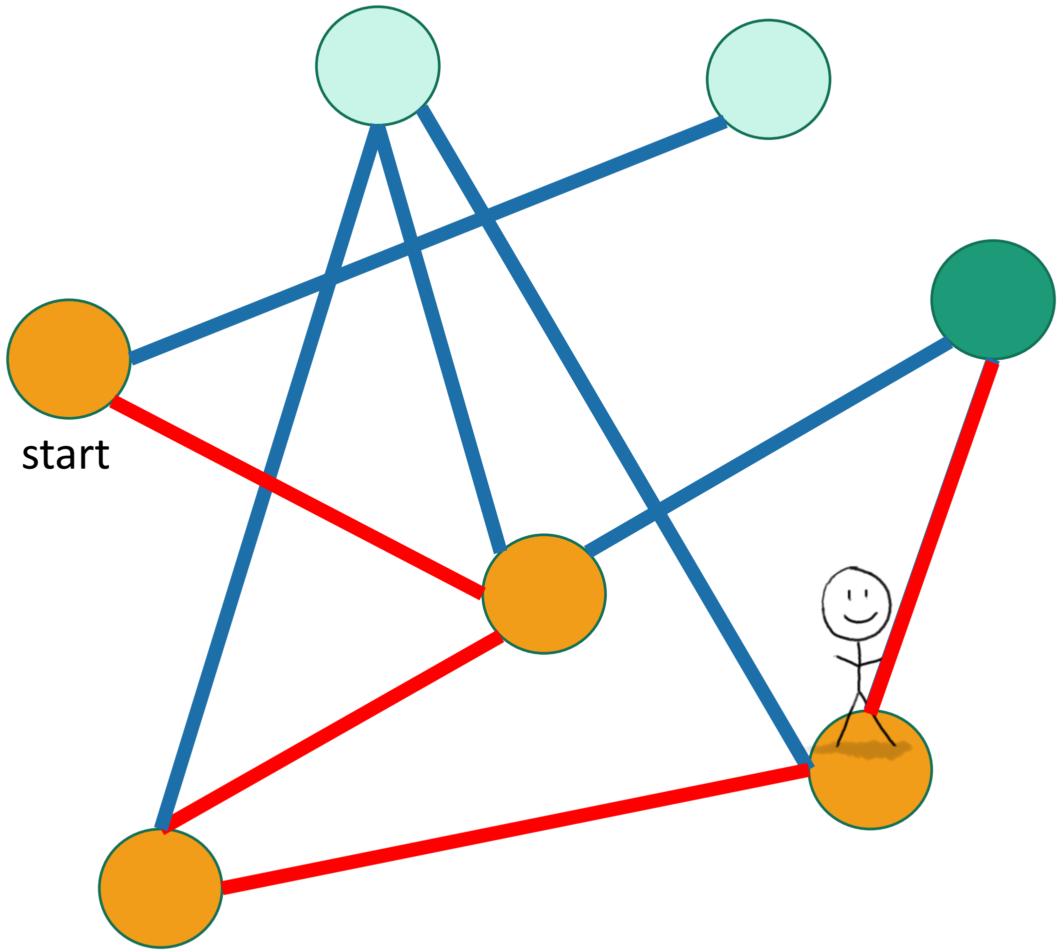

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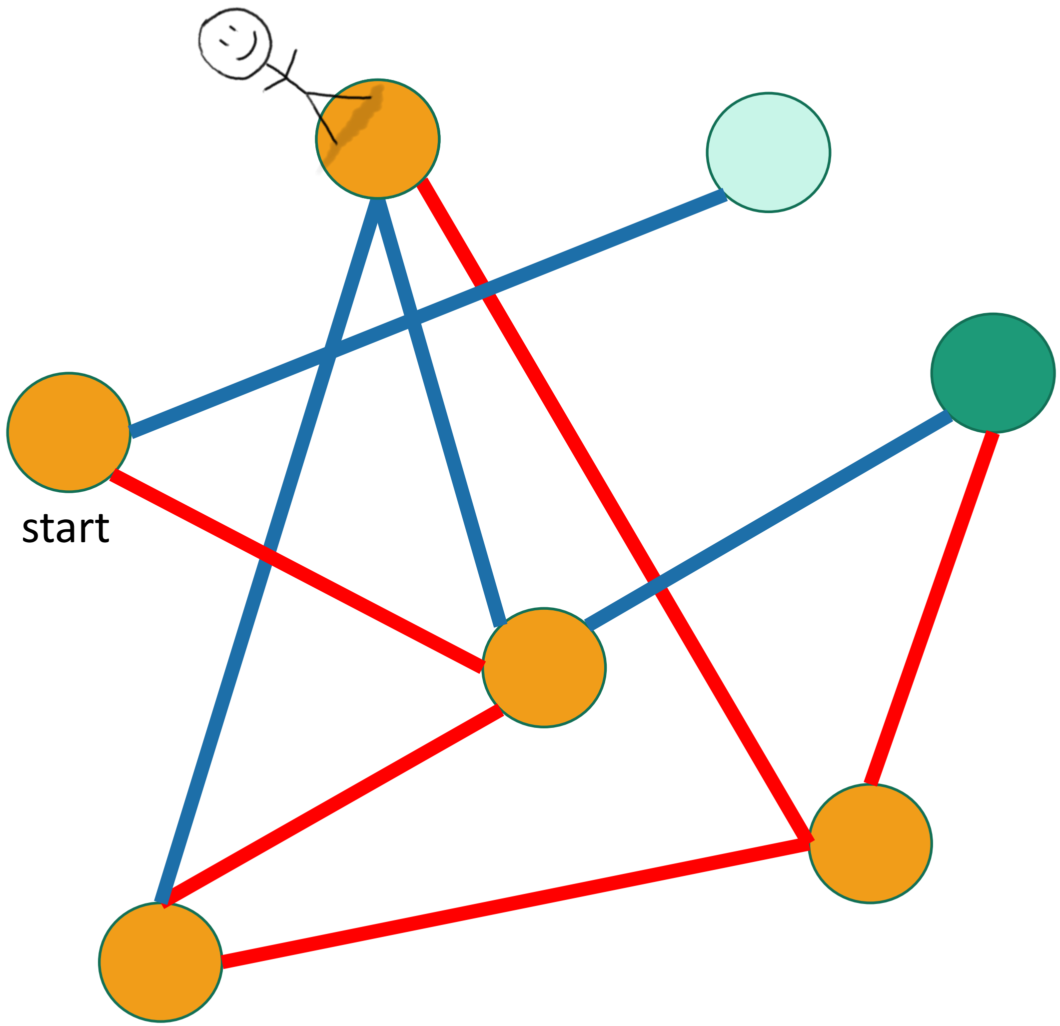

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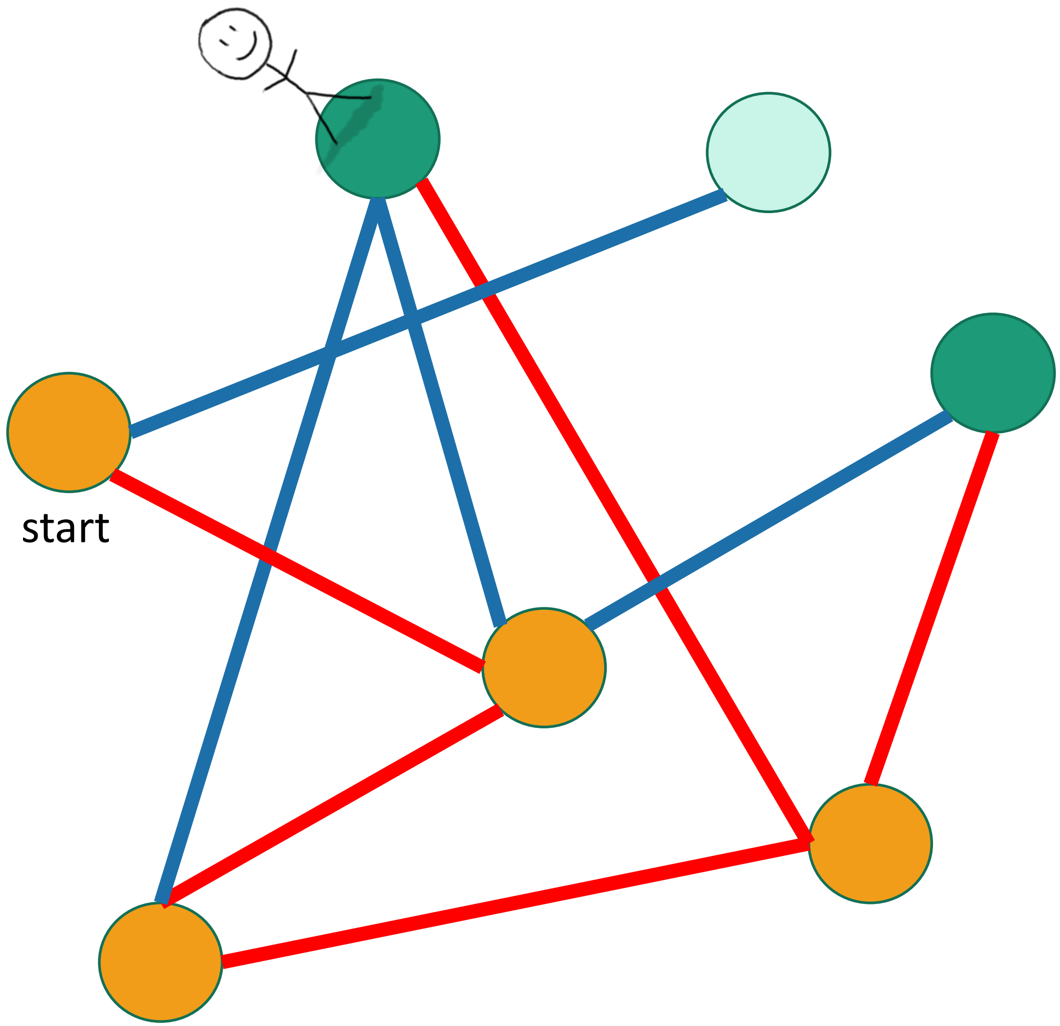

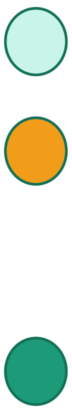

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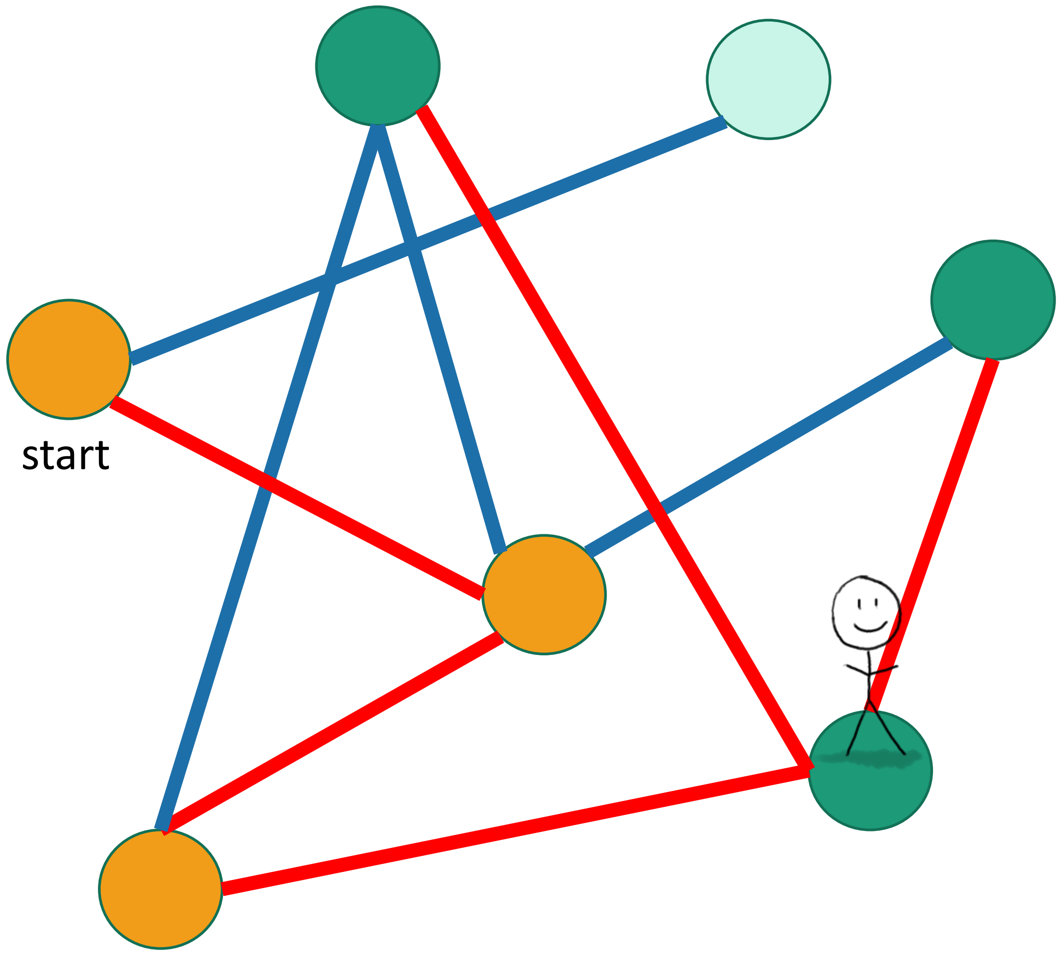

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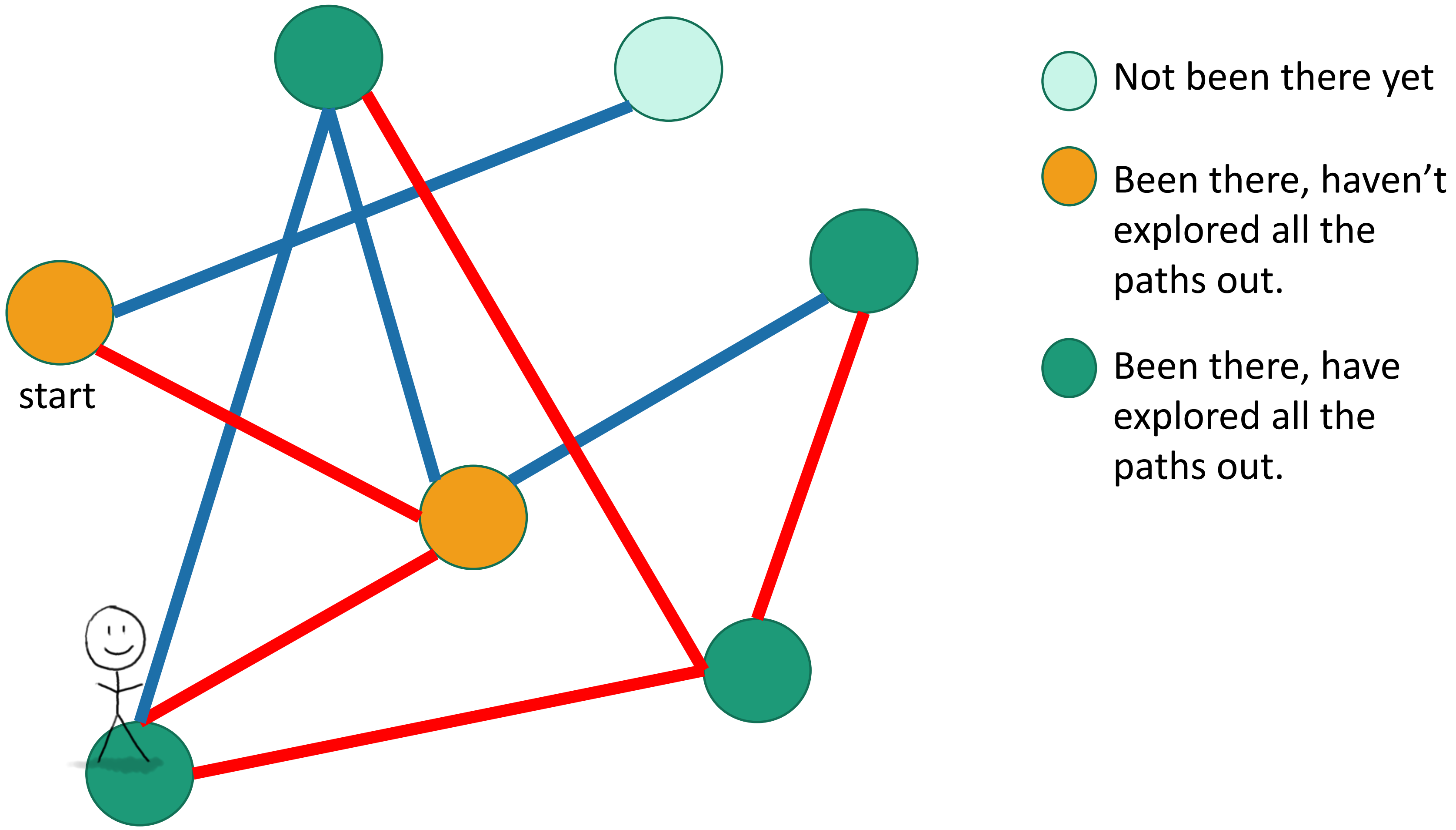

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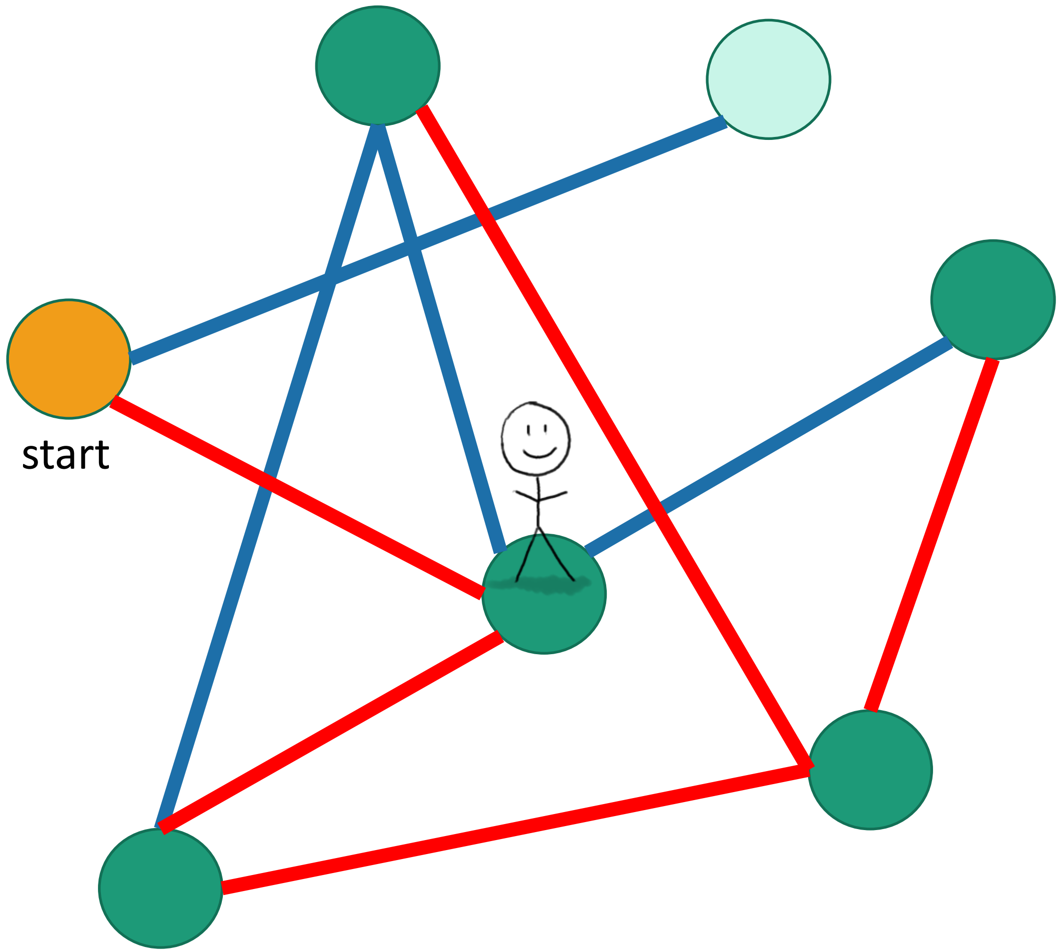

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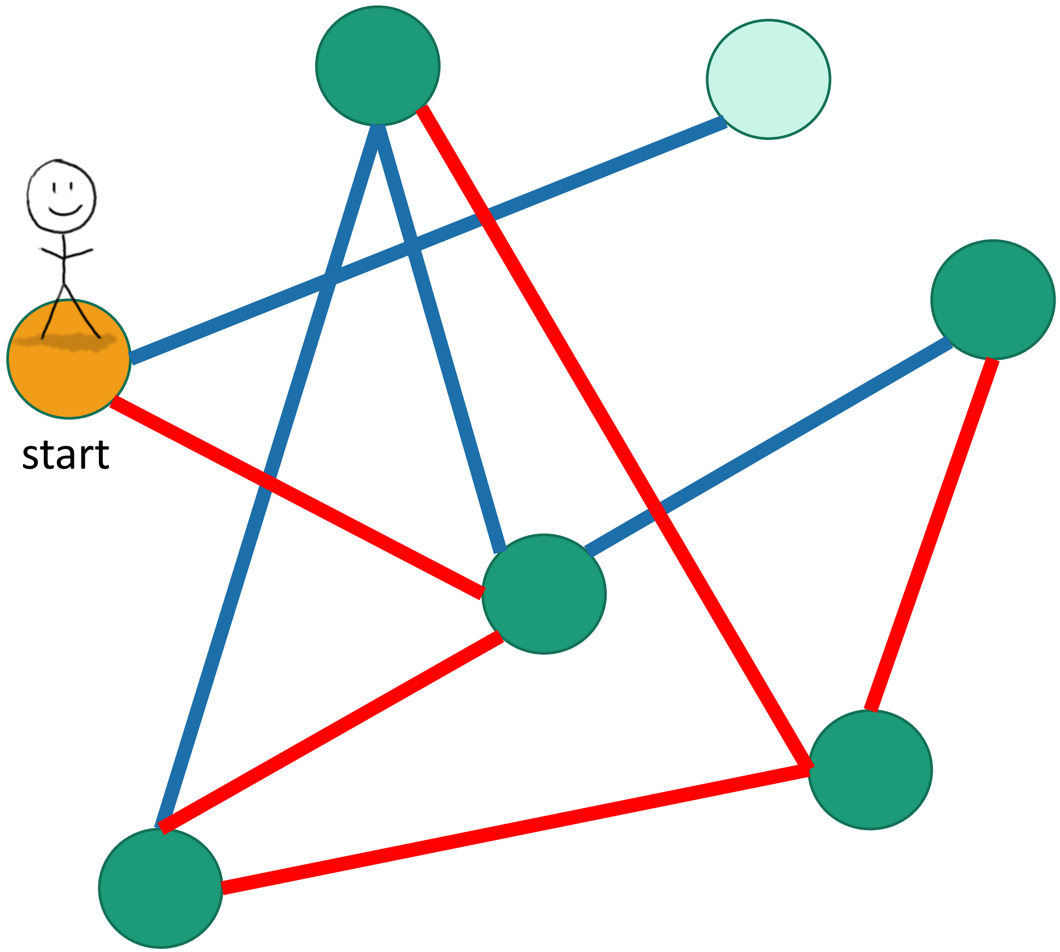

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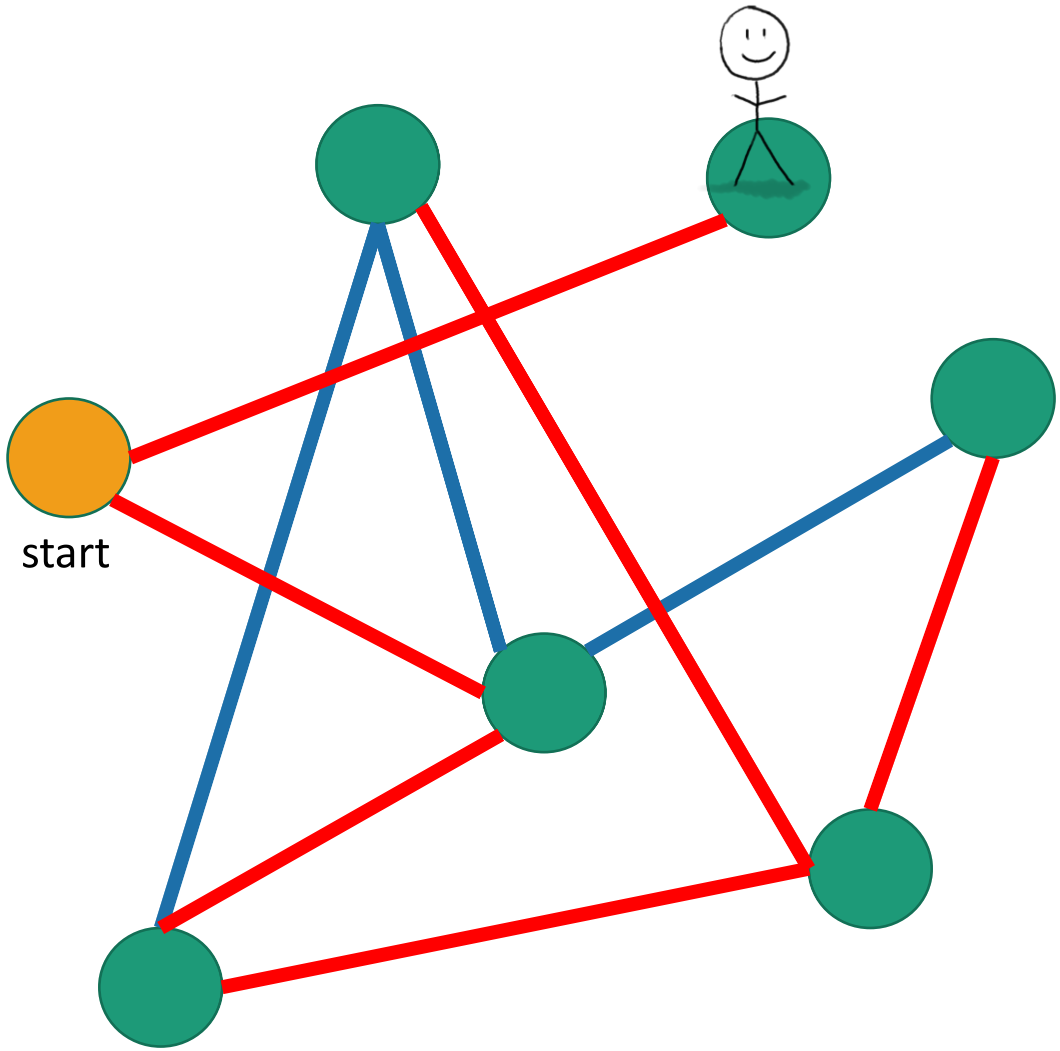

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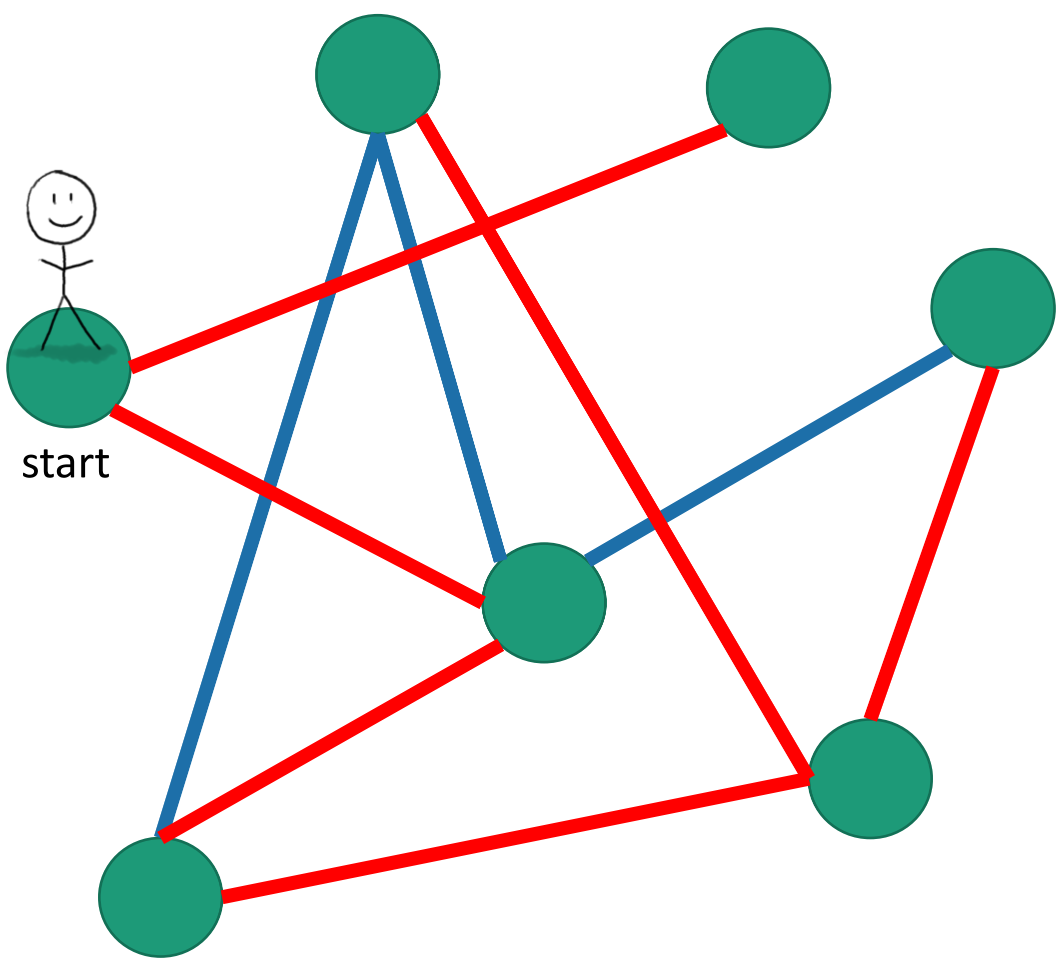

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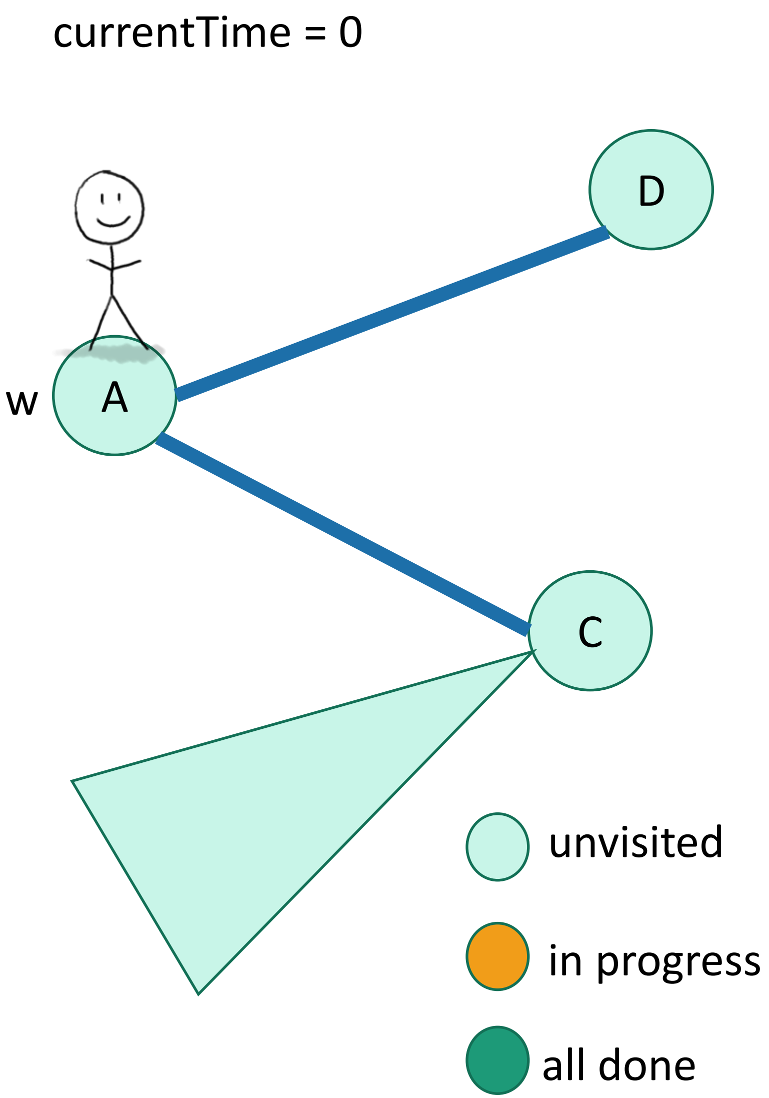
Depth First Search
Pseudocode

- Each vertex keeps track of whether it is:
  - Unvisited
  - In progress
  - All done

- Each vertex will also keep track of:
  - The time we **first enter it**.
  - The time we finish with it and mark it **all done**.

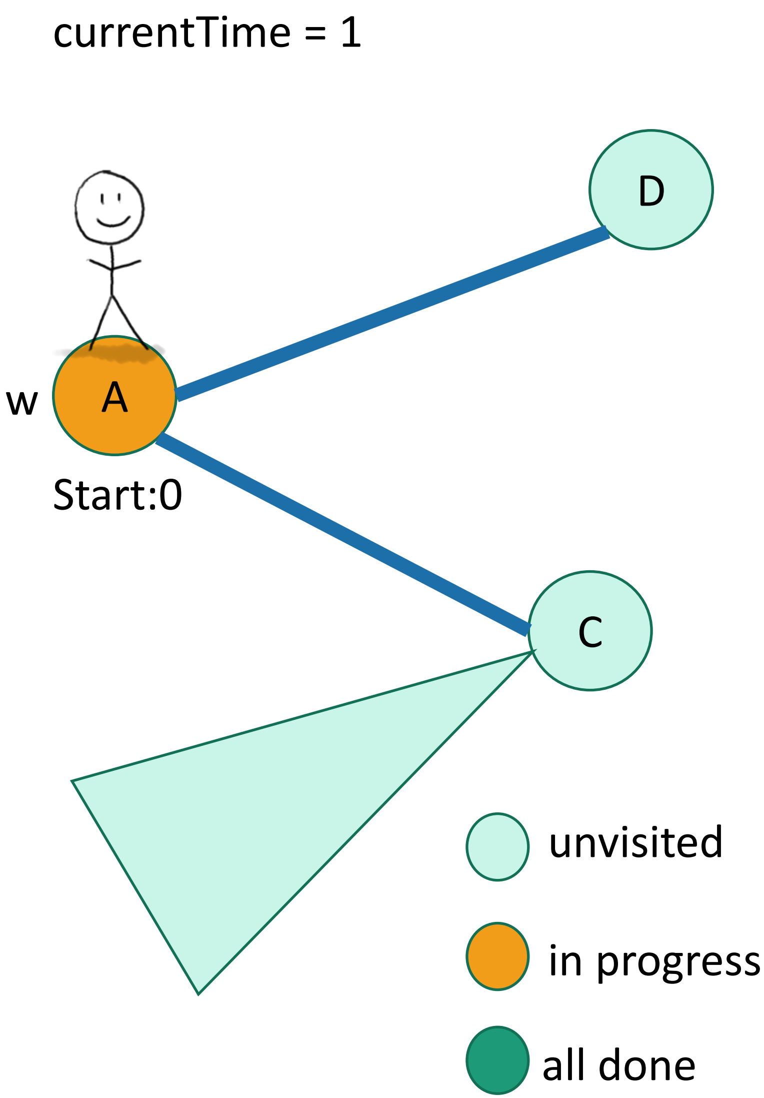
You might have seen other ways to implement DFS than what we are about to go through. This way has more bookkeeping – the bookkeeping will be useful later!

Depth First Search



- **DFS**(w, currentTime):
  - w.startTime = currentTime
  - currentTime ++
  - Mark w as **in progress**.
  - **for** v in w.neighbors:
    - **if** v is **unvisited**:
      - currentTime
        = **DFS**(v, currentTime)
      - currentTime ++
  - w.finishTime = currentTime
  - Mark w as **all done**
  - **return** currentTime

Depth First Search



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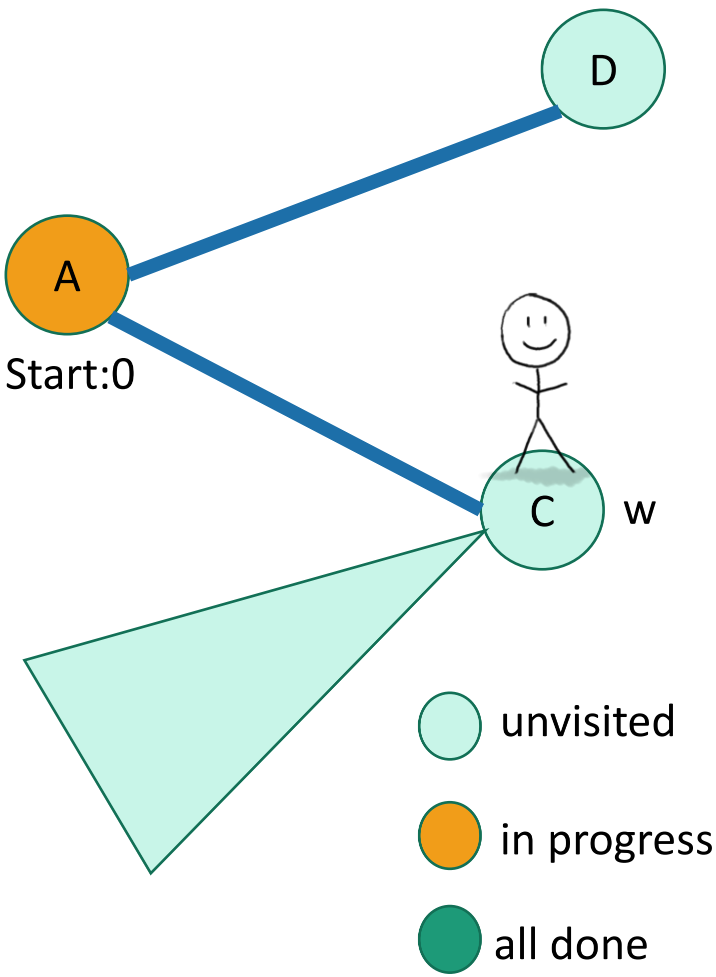
      - currentTime ++
  - w.finishTime = currentTime
  - Mark w as **all done**
  - **return** currentTime

Depth First Search

currentTime = 1



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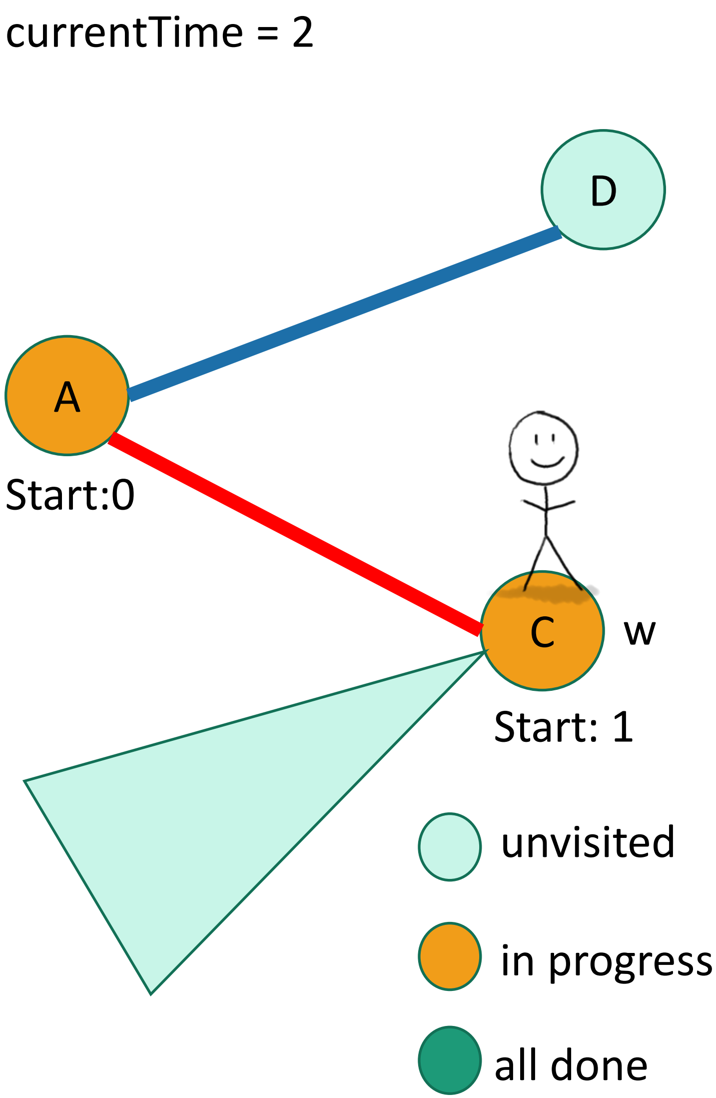

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      - currentTime ++
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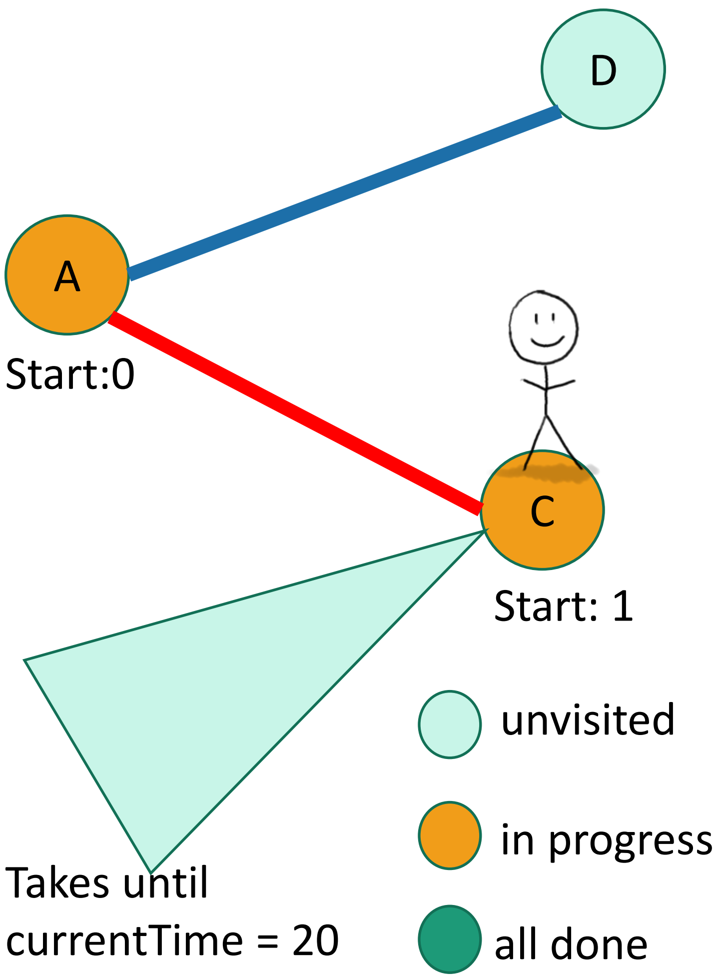
  - Mark w as **all done**
  - **return** currentTime

Depth First Search

currentTime = 20



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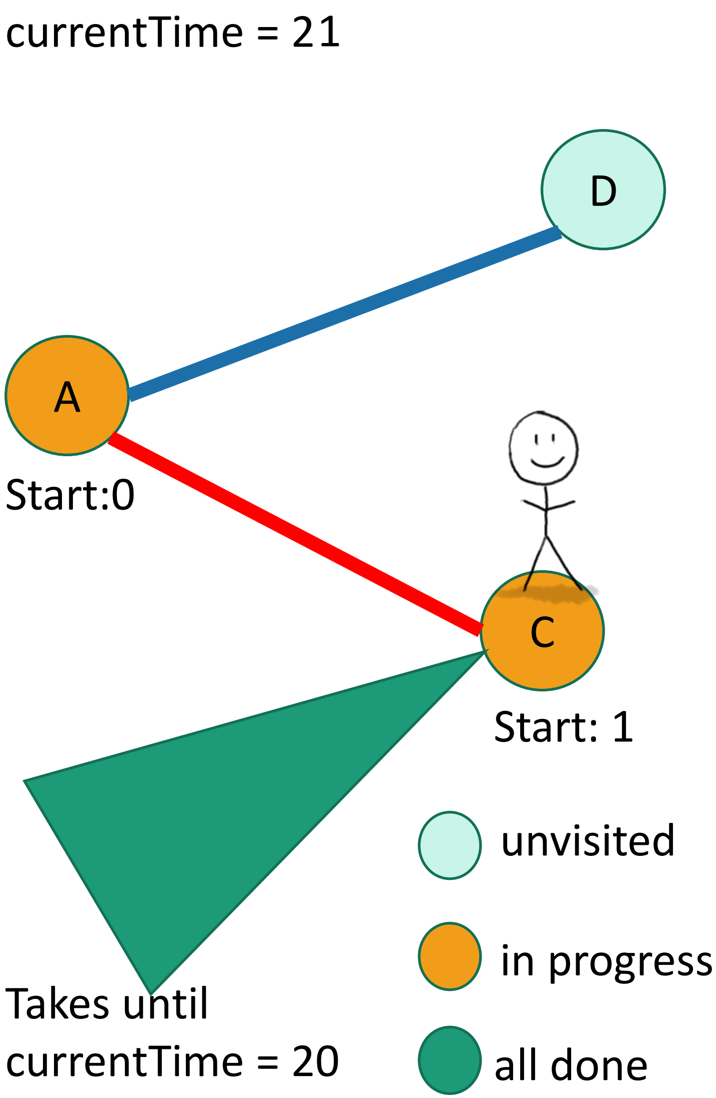
Depth First Search



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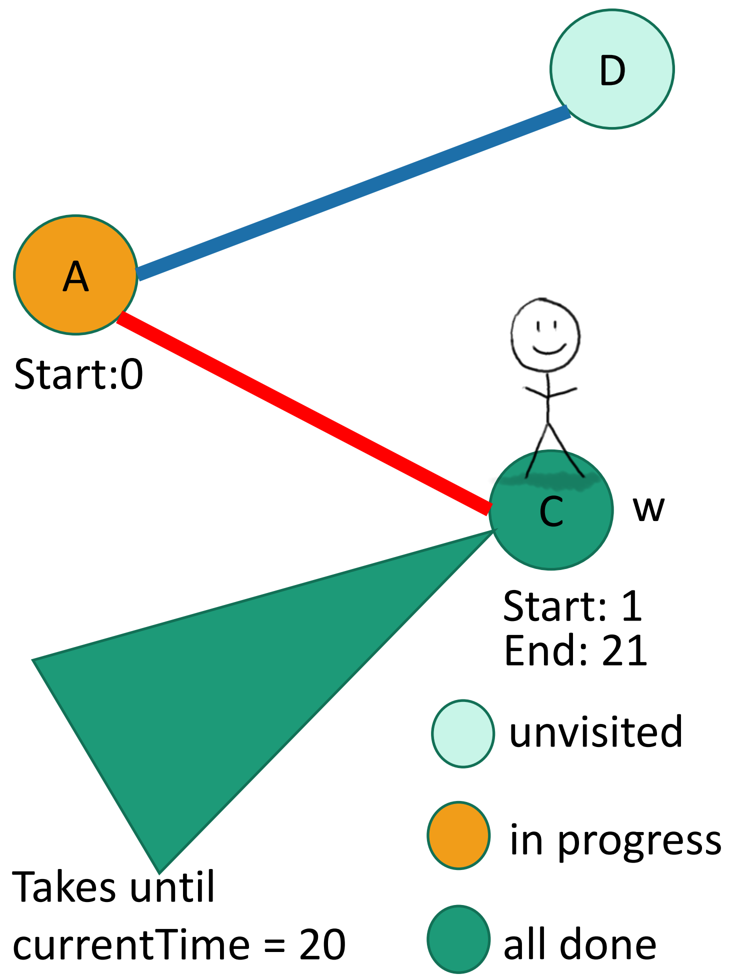
Depth First Search

currentTime = 21



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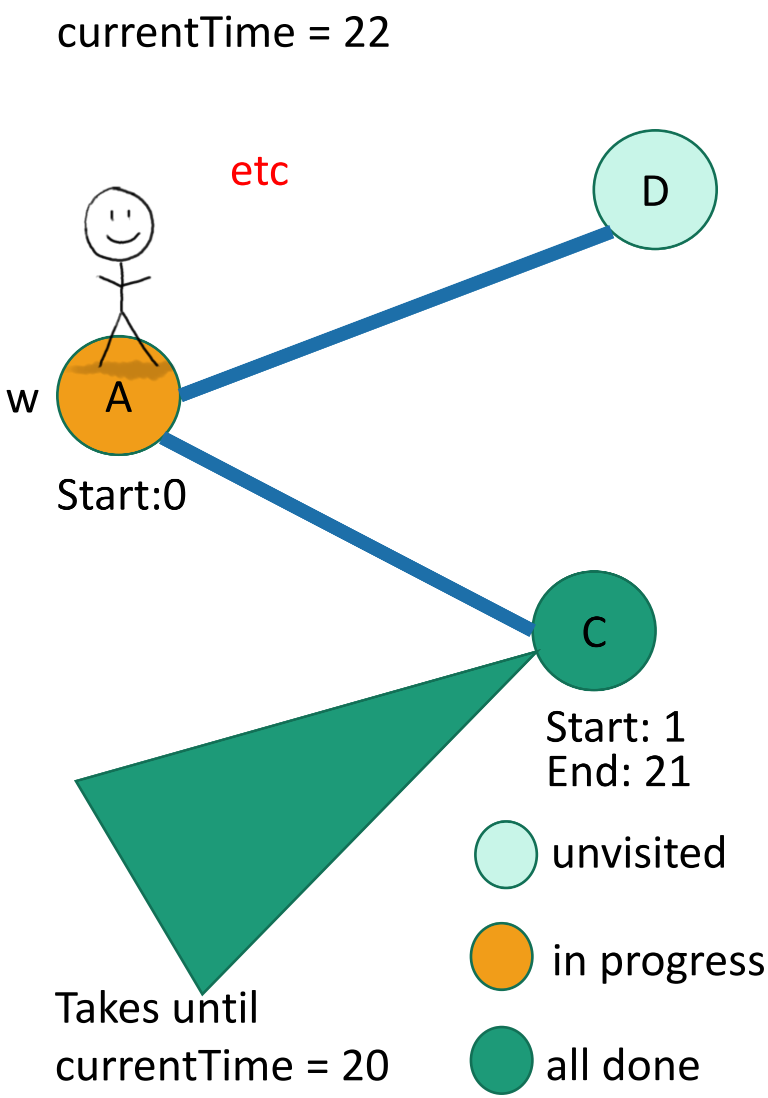

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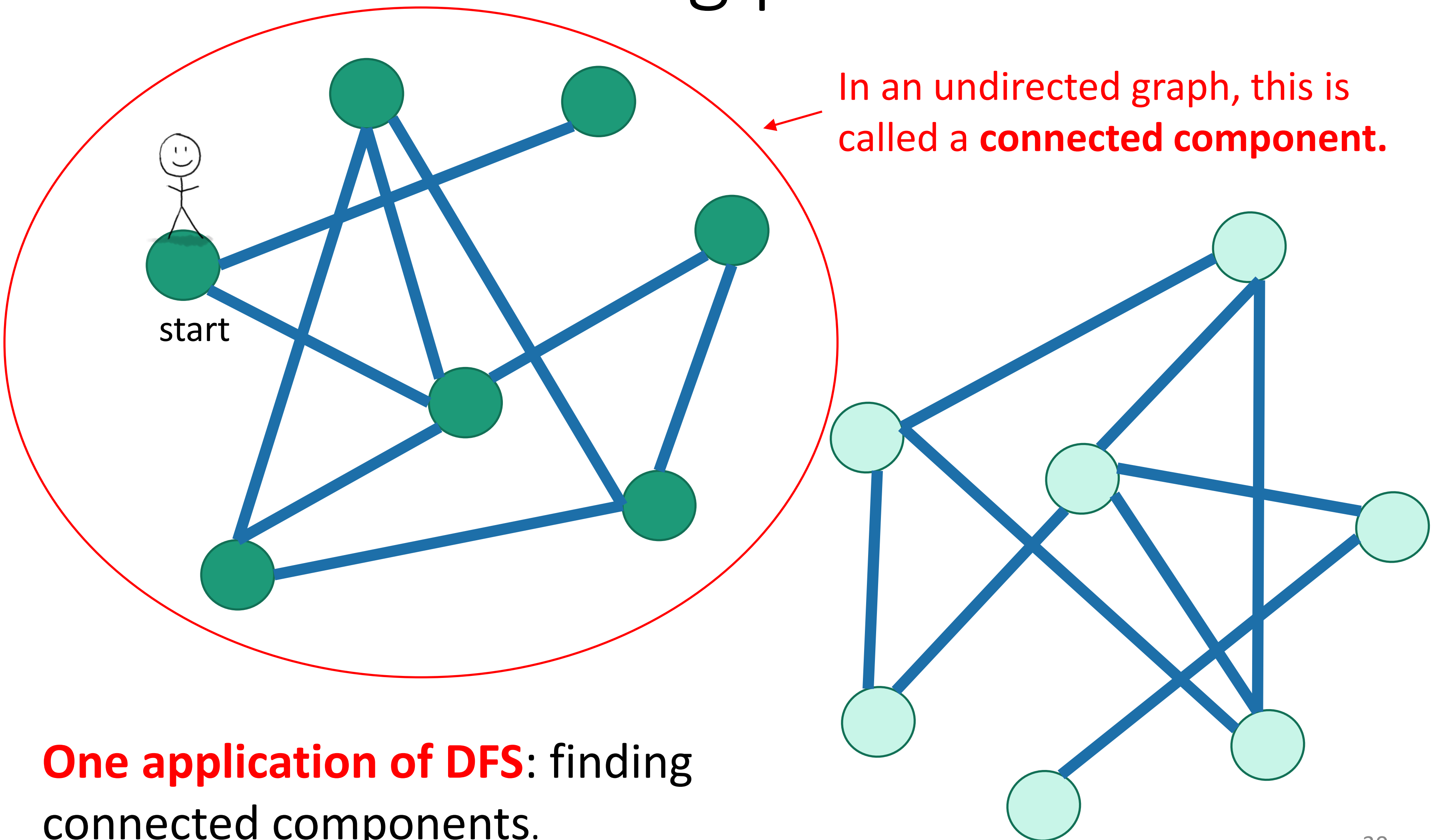
Fun exercise

- Write pseudocode for an iterative version of DFS.

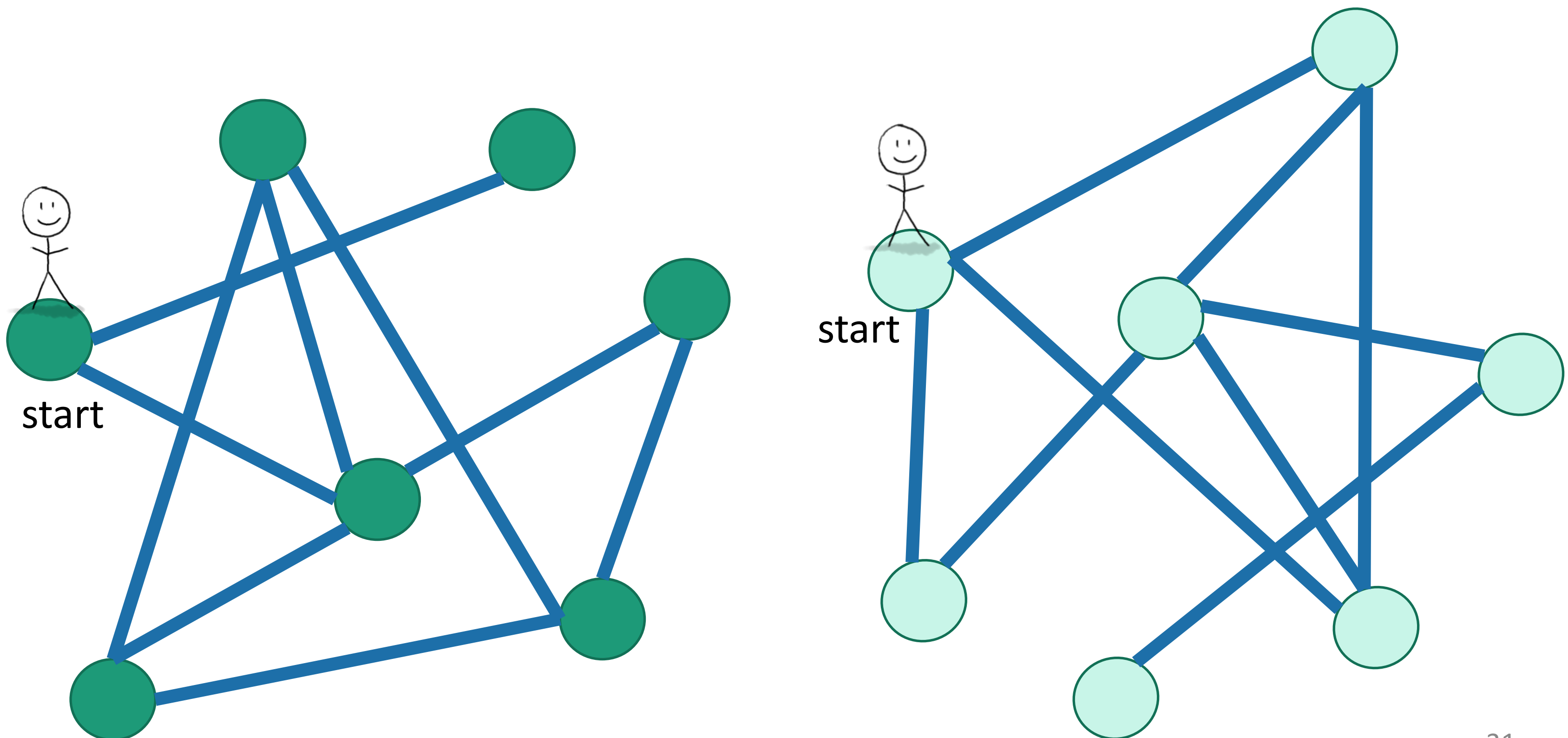
DFS finds all the nodes reachable from the starting point



One application of DFS: finding connected components.

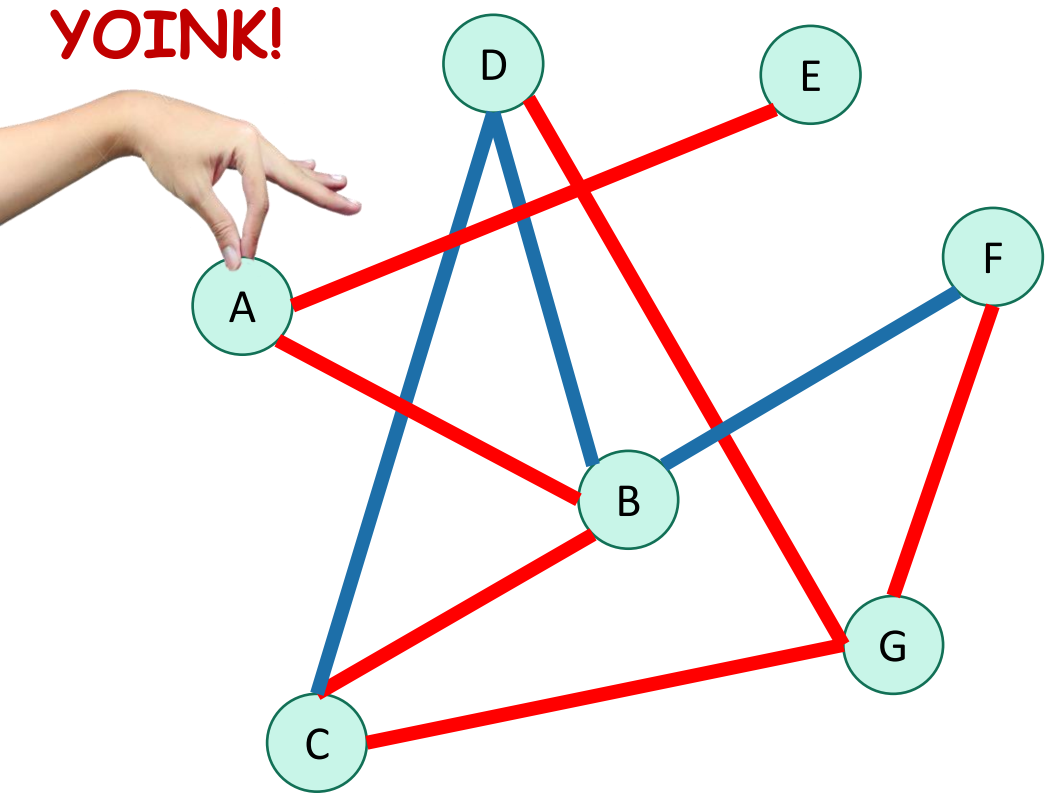
To explore the whole graph

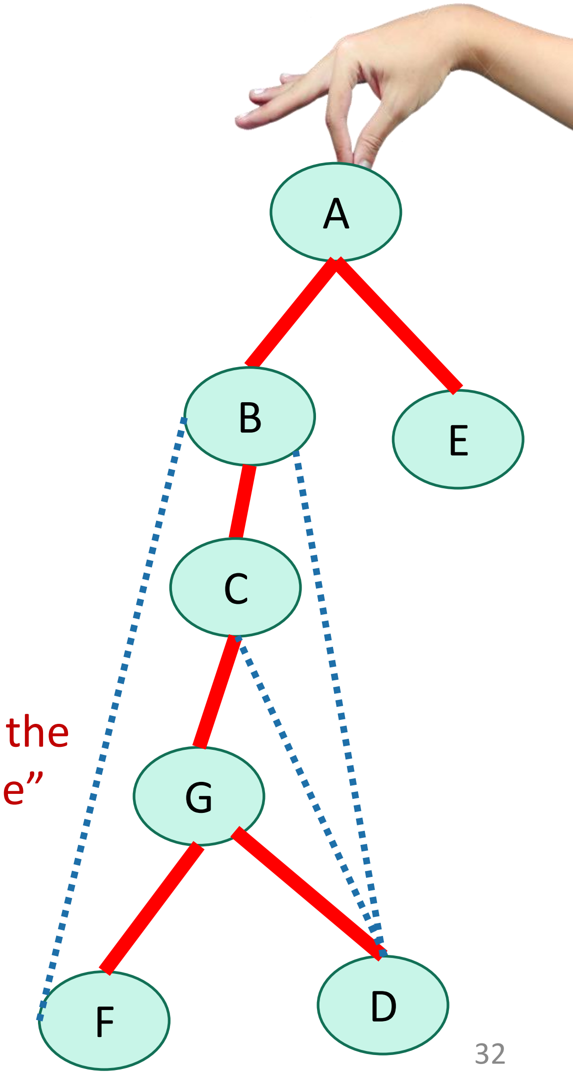
- Do it repeatedly!



Why is it called depth-first?

- We are implicitly building a tree:





Call this the
"DFS tree"

- First, we go as deep as we can.

Running time

To explore **just the connected component** we started in

- We look at each edge at most twice.
  - Once from each of its endpoints
- And basically we don't do anything else.
- So…

$$O(m)$$

Running time

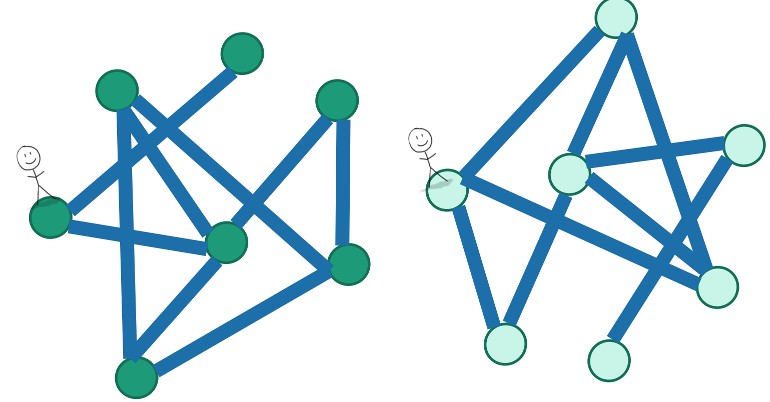
To explore just the connected component we started in

- Assume we are using the linked-list format for G.
- Say C = (V', E') is a connected component.
- We visit each vertex in C exactly once.
  - Here, "visit" means "call DFS on"
- At each vertex w, we:
  - Do some book-keeping: O(1)
  - Loop over w's neighbors and check if they are visited (and then potentially make a recursive call): O(1) per neighbor or O(deg(w)) total.
- Total time:
  - $\sum_{w \in V'} (O(\deg(w)) + O(1))$
  - $= O(|E'| + |V'|)$
  - $= O(|E'|)$

In a connected graph, $|V'| \leq |E'| + 1.$

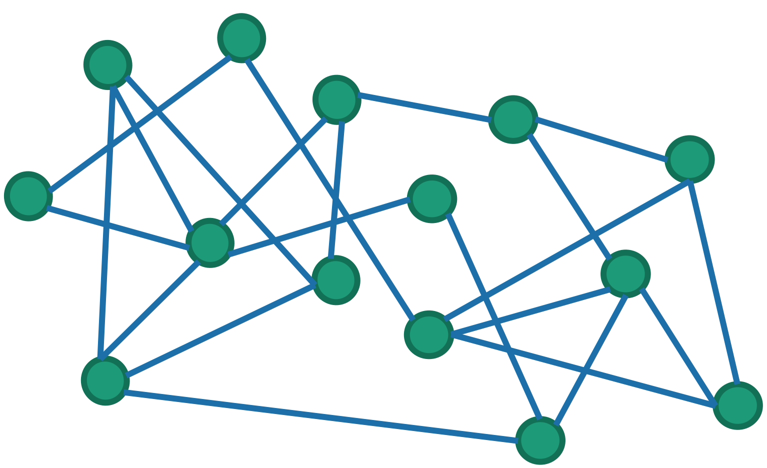
Running time

To explore the whole graph

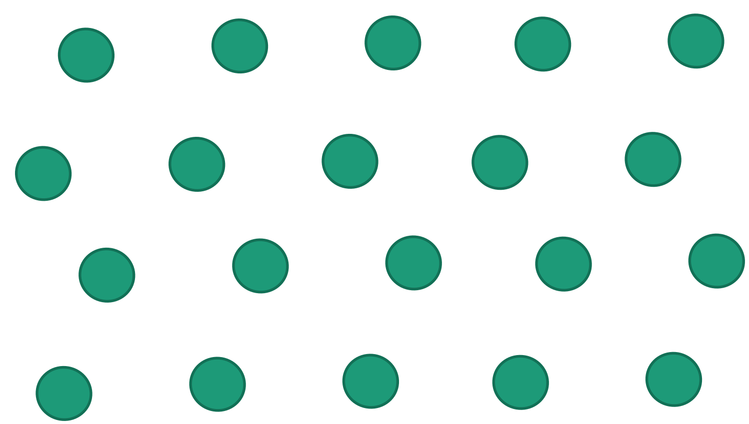


- Explore the connected components one-by-one.
- This takes time O(n + m)
  - Same computation as before:

$$\sum_{w \in V} (O(\deg(w)) + O(1)) = O(|E| + |V|) = O(n + m)$$



or

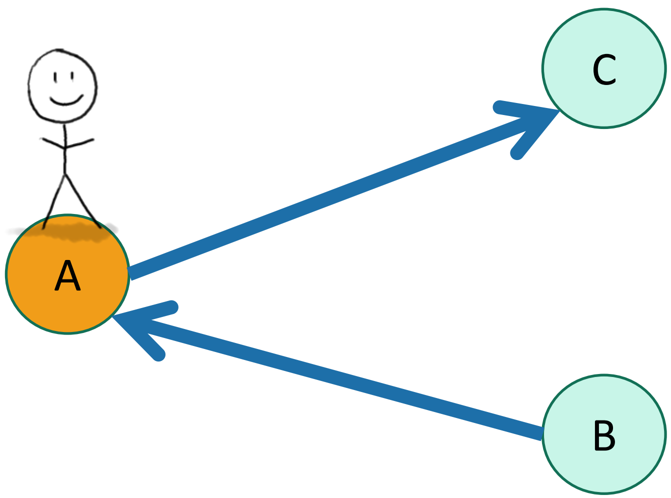


Here the running time is O(m) like before

Here m=0 but it still takes time O(n) to explore the graph.

You check:

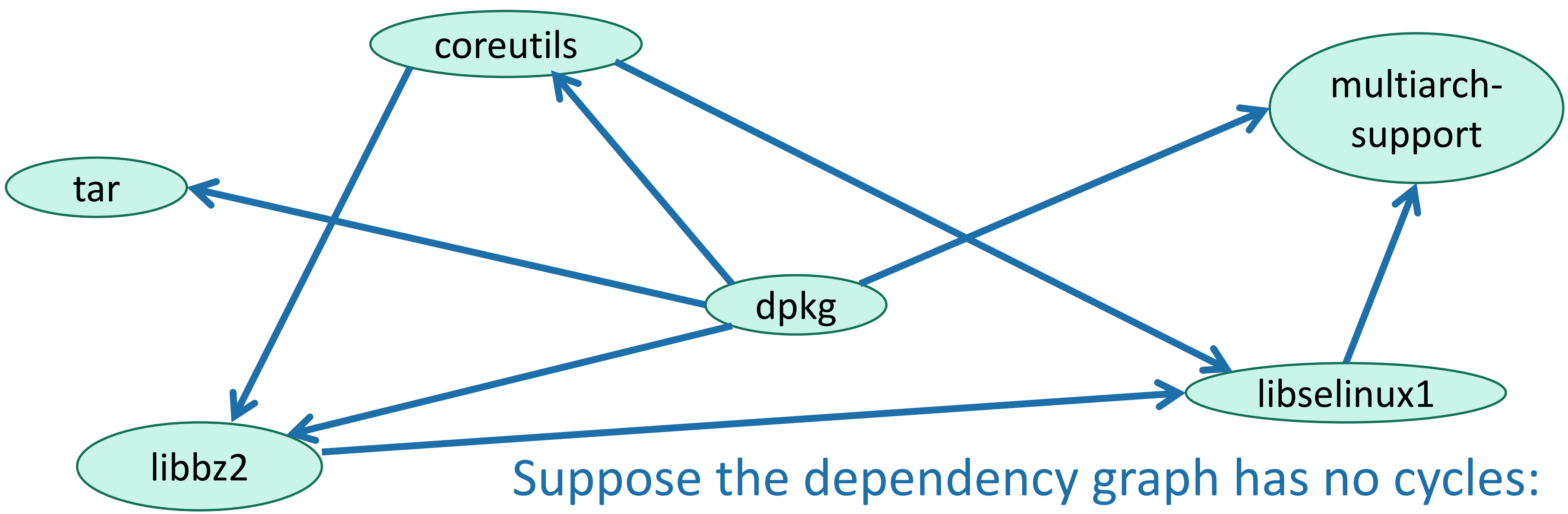
DFS works fine on directed graphs too!



Only walk to C, not to B.

Application of DFS: topological sorting

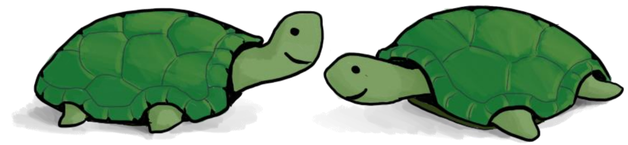
- Find an ordering of vertices so that all of the dependency requirements are met.
  - Aka, if v comes before w in the ordering, there is not an edge from w to v.

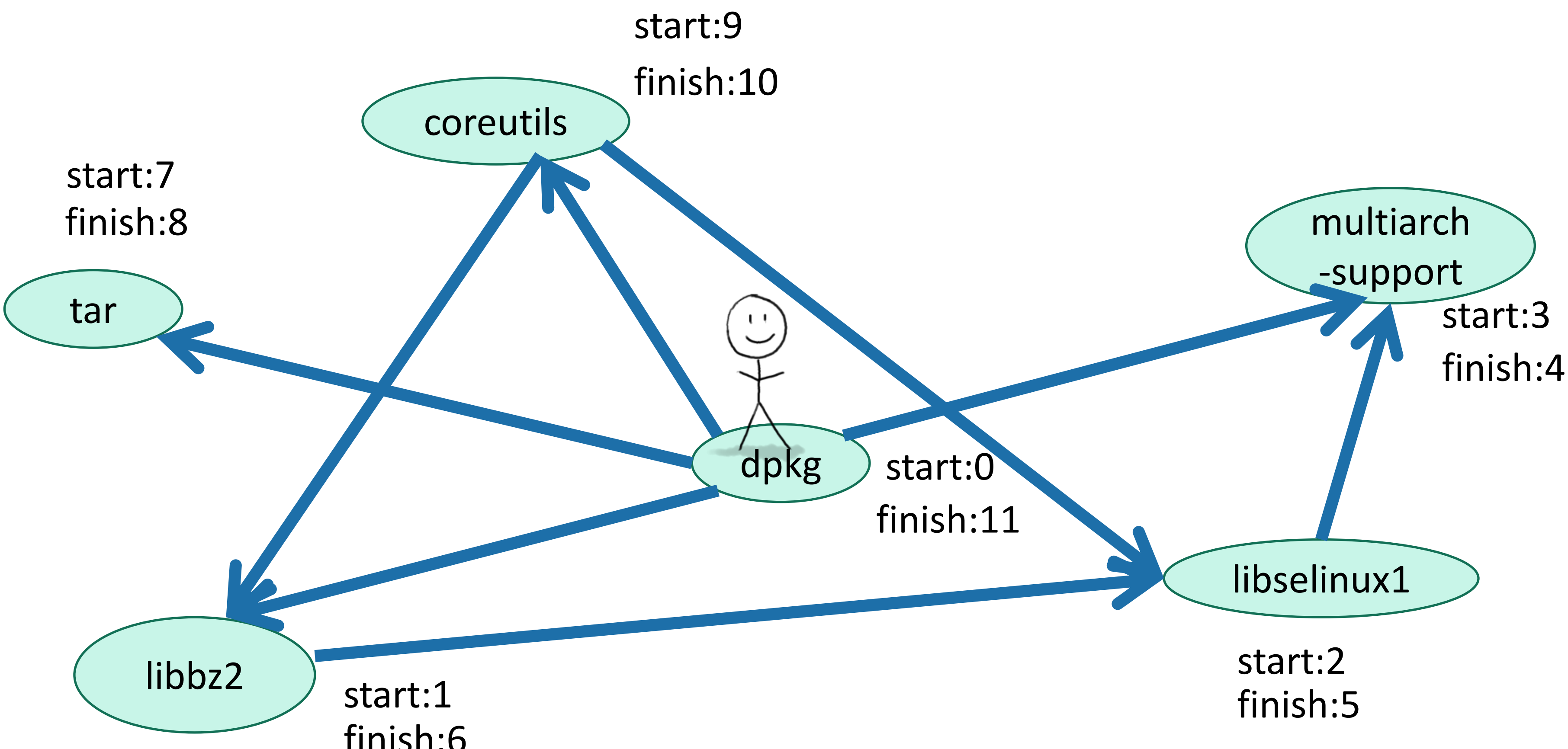


Suppose the dependency graph has no cycles:
it is a **Directed Acyclic Graph (DAG)**

Let’s do DFS

What do you notice about the finish times? Any ideas for how we should do topological sort?





Suppose the underlying graph has no cycles

Finish times seem useful

Claim: In general, we'll always have:

finish: [larger]

finish: [smaller]

To understand why, let's go back to that DFS tree.

A more general statement

(this holds even if there are cycles)

This is called the “parentheses theorem”

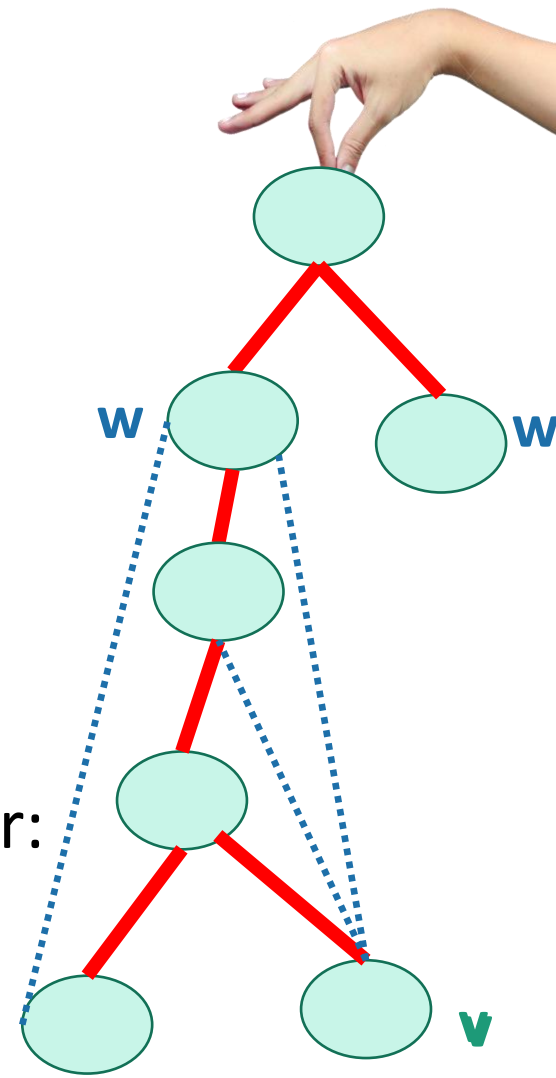
(check this statement carefully!)

- If v is a descendant of w in this tree:

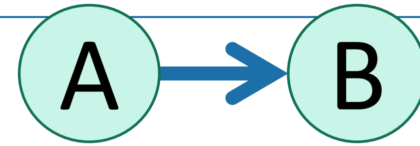
- If w is a descendant of v in this tree:

- If neither are descendants of each other:

(or the other way around)

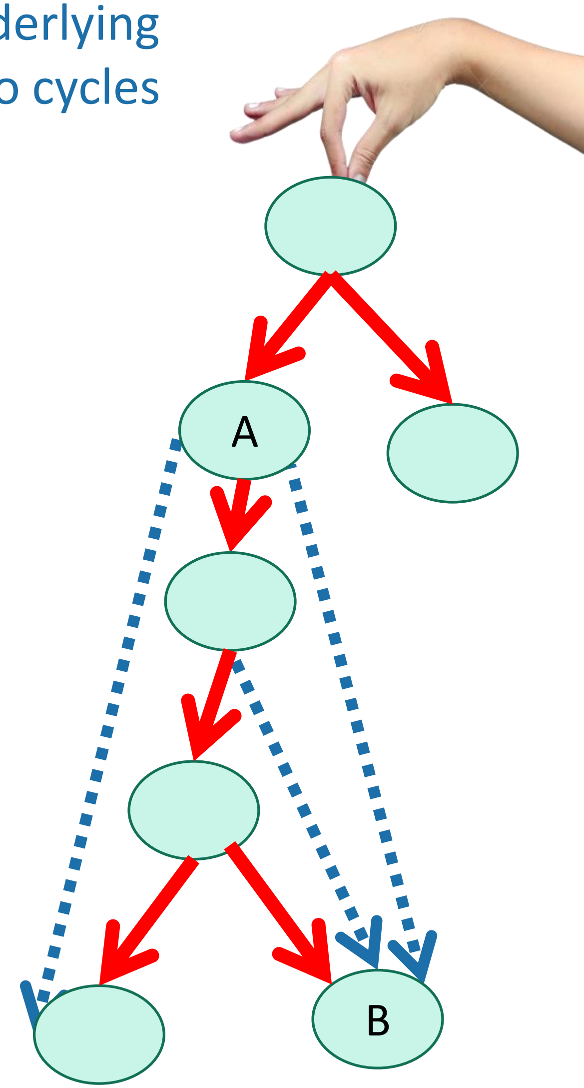


So to prove this →

If 

Then B.finishTime < A.finishTime

Suppose the underlying graph has no cycles

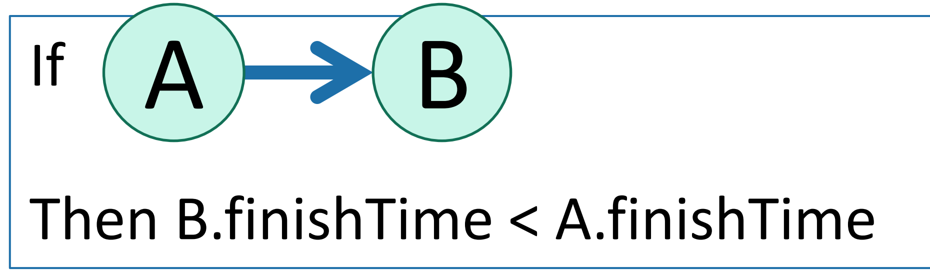


- **Case 1**: B is a descendant of A in the DFS tree.

- Then

- aka, B.finishTime < A.finishTime.

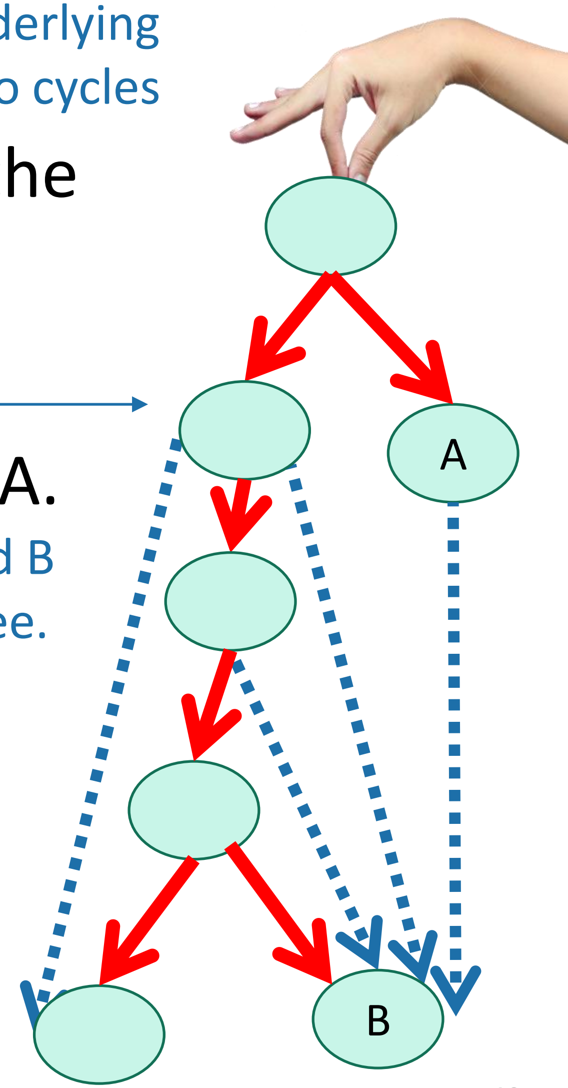
So to prove this →



Suppose the underlying graph has no cycles

- **Case 2**: B is a NOT descendant of A in the DFS tree.
  - Notice that A can't be a descendant of B or else there'd be a cycle; so it looks like this
- Then we must have explored B before A.
  - Otherwise we would have gotten to B from A, and B would have been a descendant of A in the DFS tree.
- Then

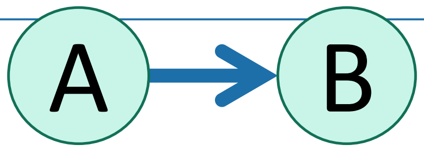
- aka, B.finishTime < A.finishTime.



Theorem

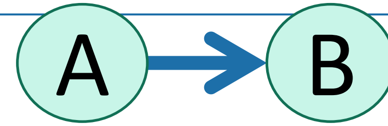
- If we run DFS on a directed acyclic graph,

If



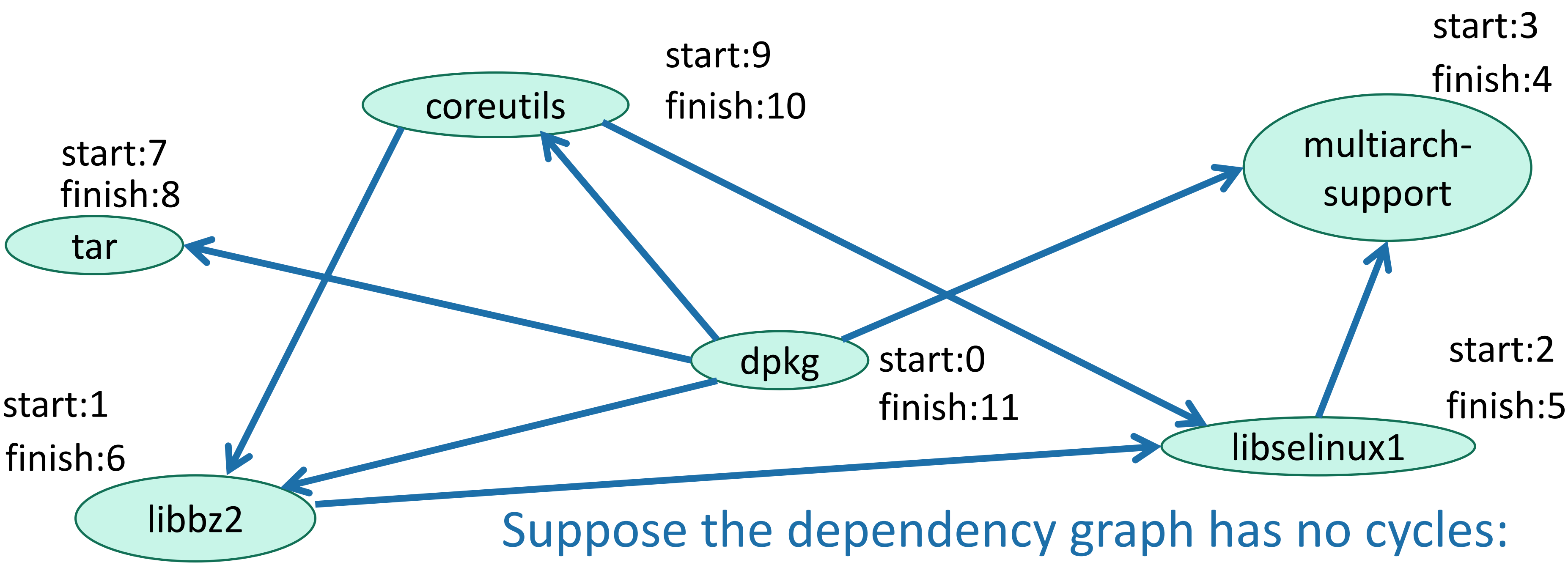
Then B.finishTime < A.finishTime

Back to topological sorting

If 

Then B.finishTime < A.finishTime

- In what order should I install packages?
- In reverse order of finishing time in DFS!

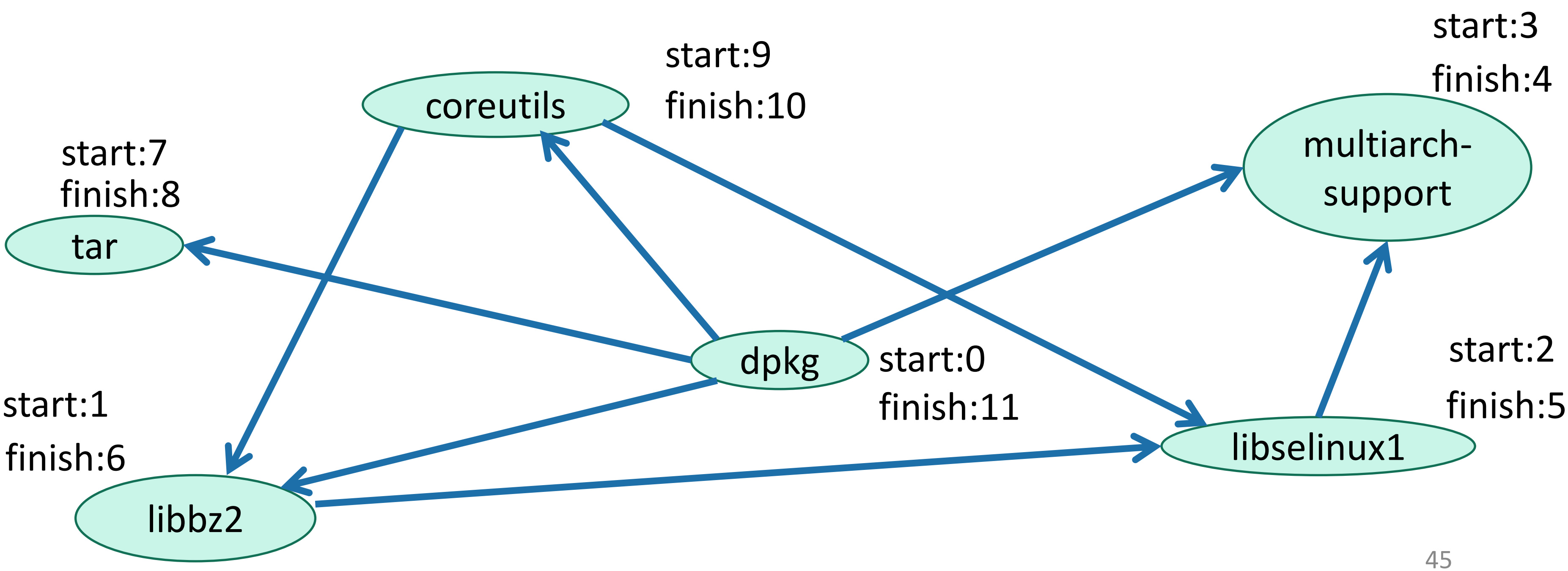


Suppose the dependency graph has no cycles: it is a **Directed Acyclic Graph (DAG)**

Topological Sorting (on a DAG)

- Do DFS
- When you mark a vertex as **all done**, put it at the **beginning** of the list.

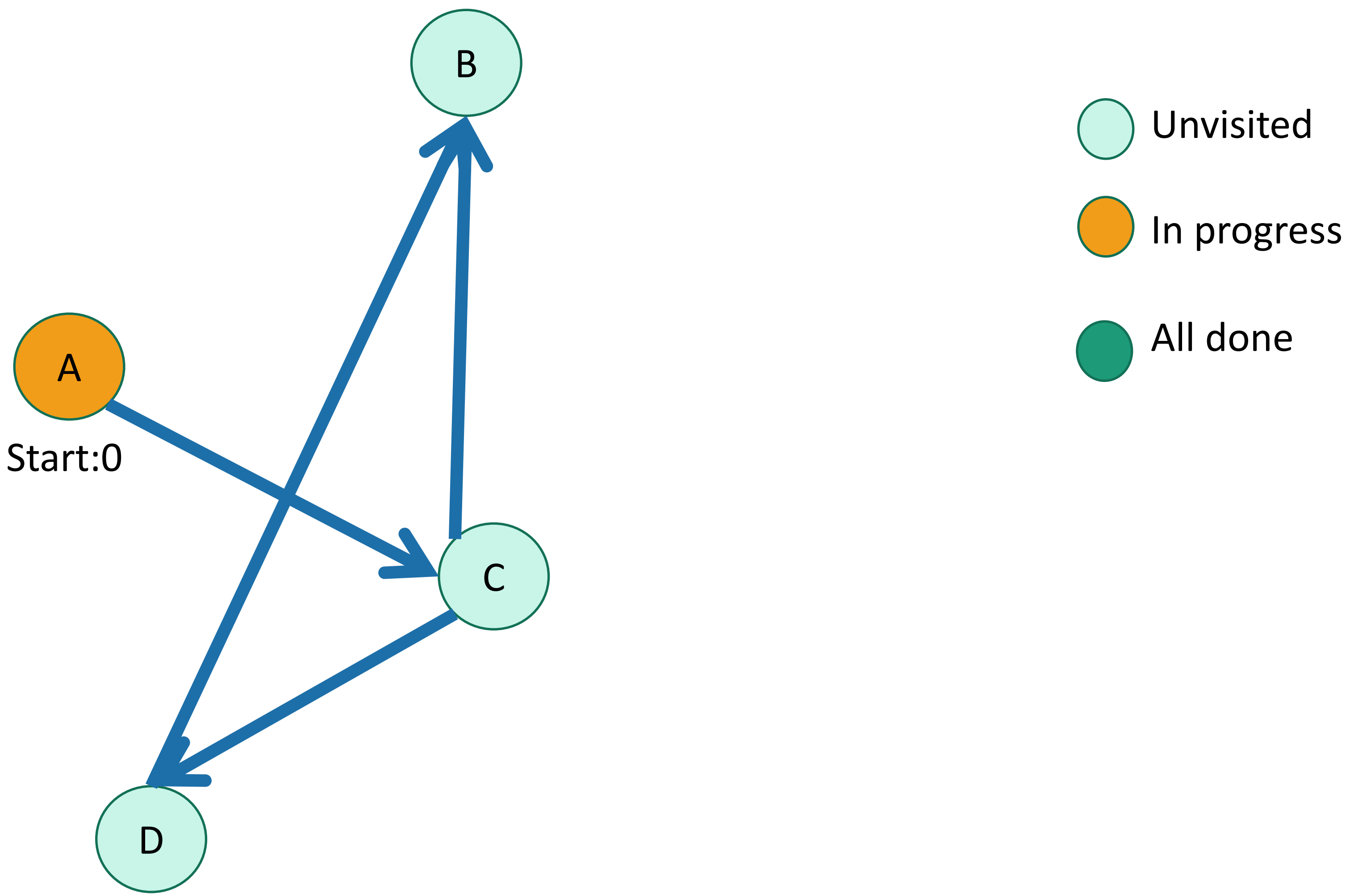
- dpkg
- coreutils
- tar
- libbz2
- libselinux1
- multiarch_support



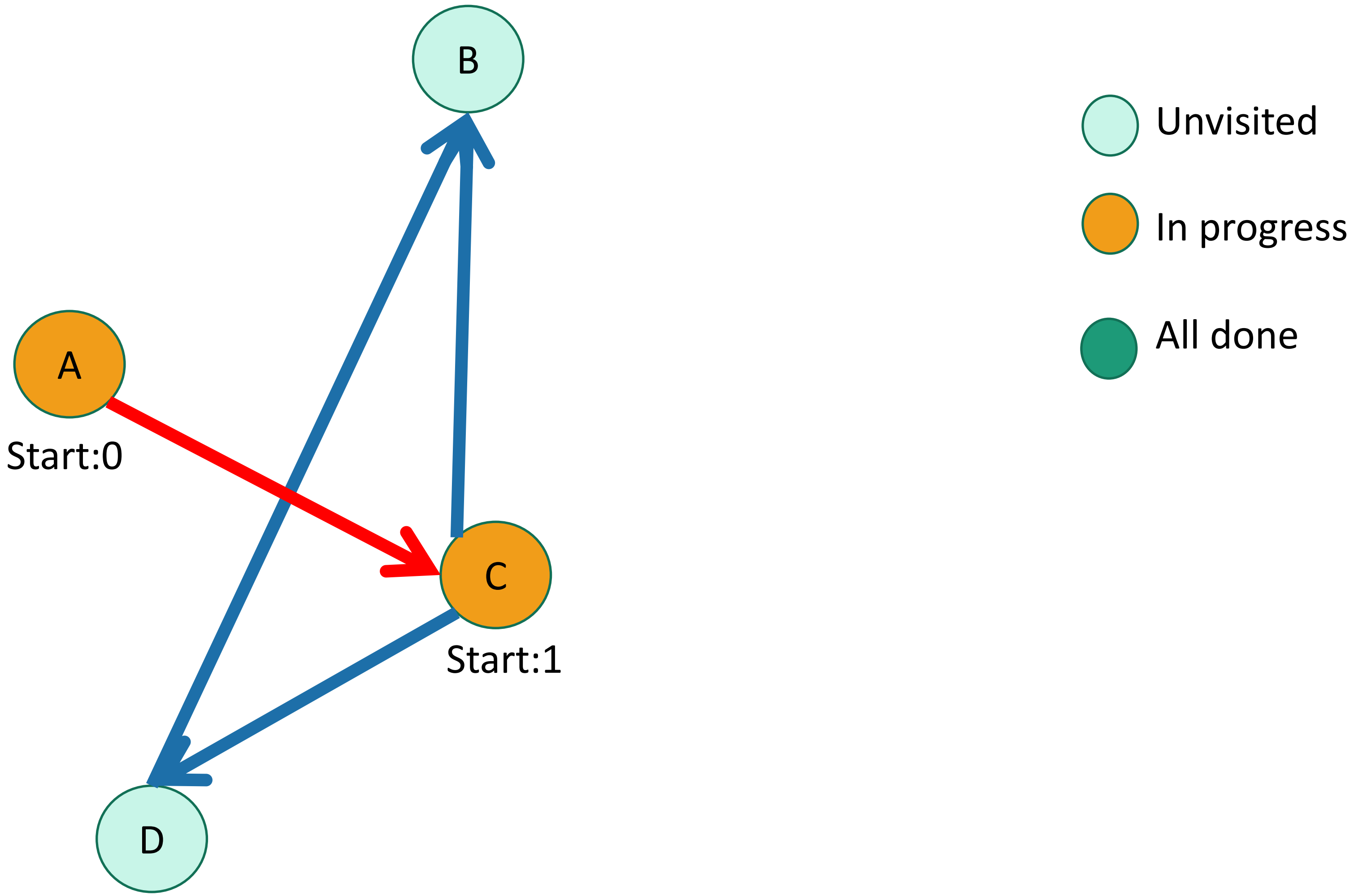
What did we just learn?

- DFS can help you solve the **topological sorting problem**
  - That’s the fancy name for the problem of finding an ordering that respects all the dependencies
- Thinking about the DFS tree is helpful.

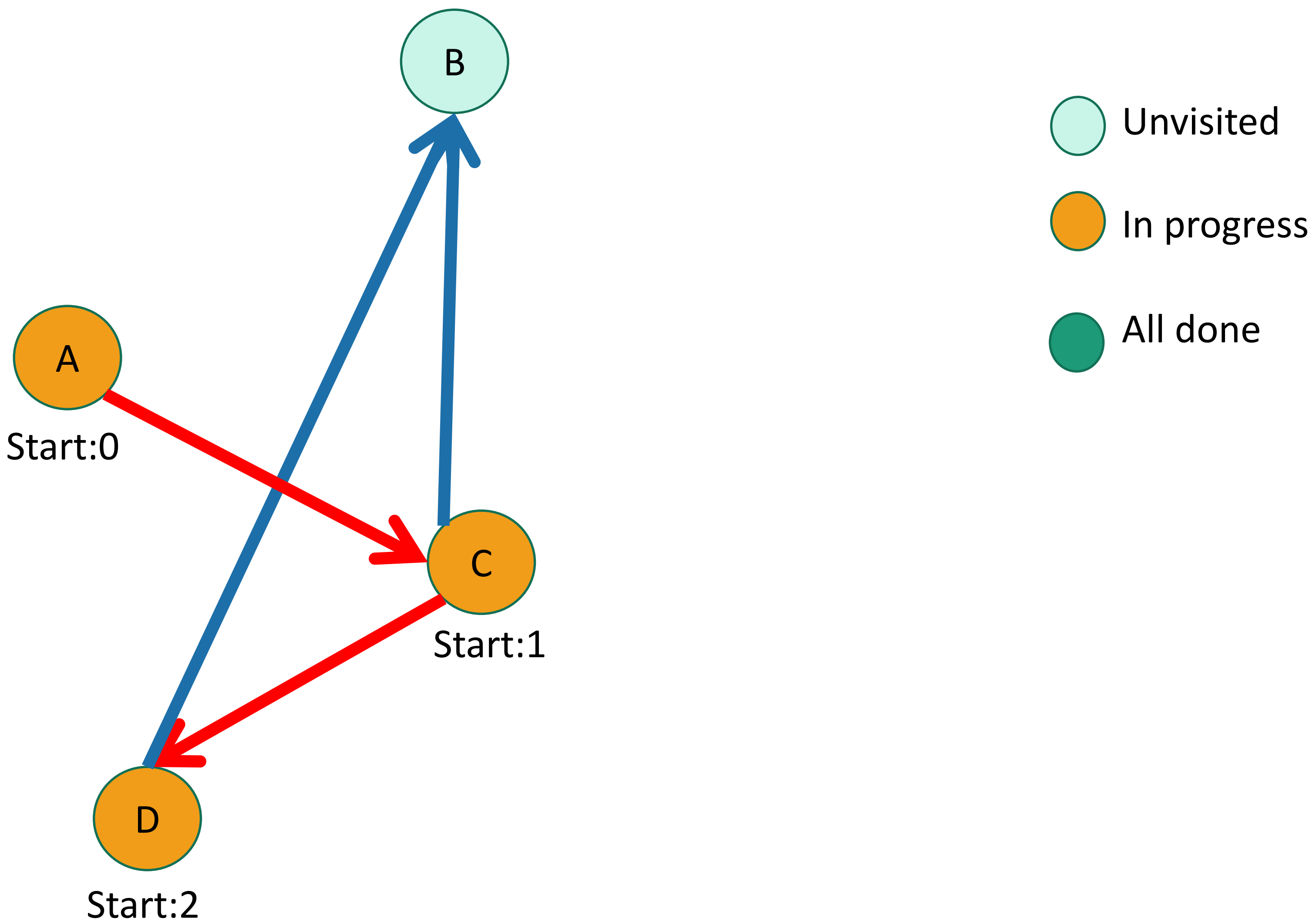
Example:



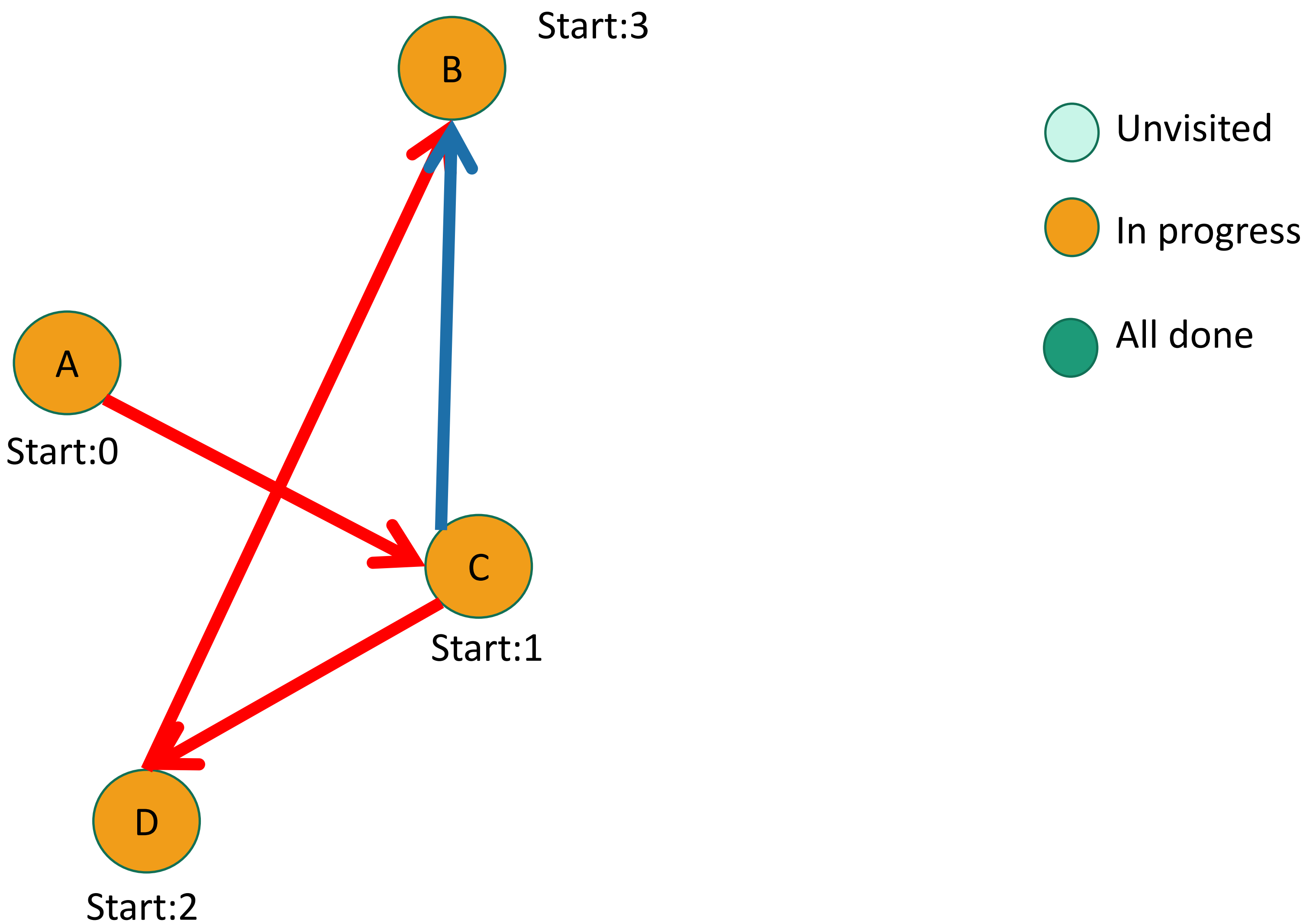
Example



Example

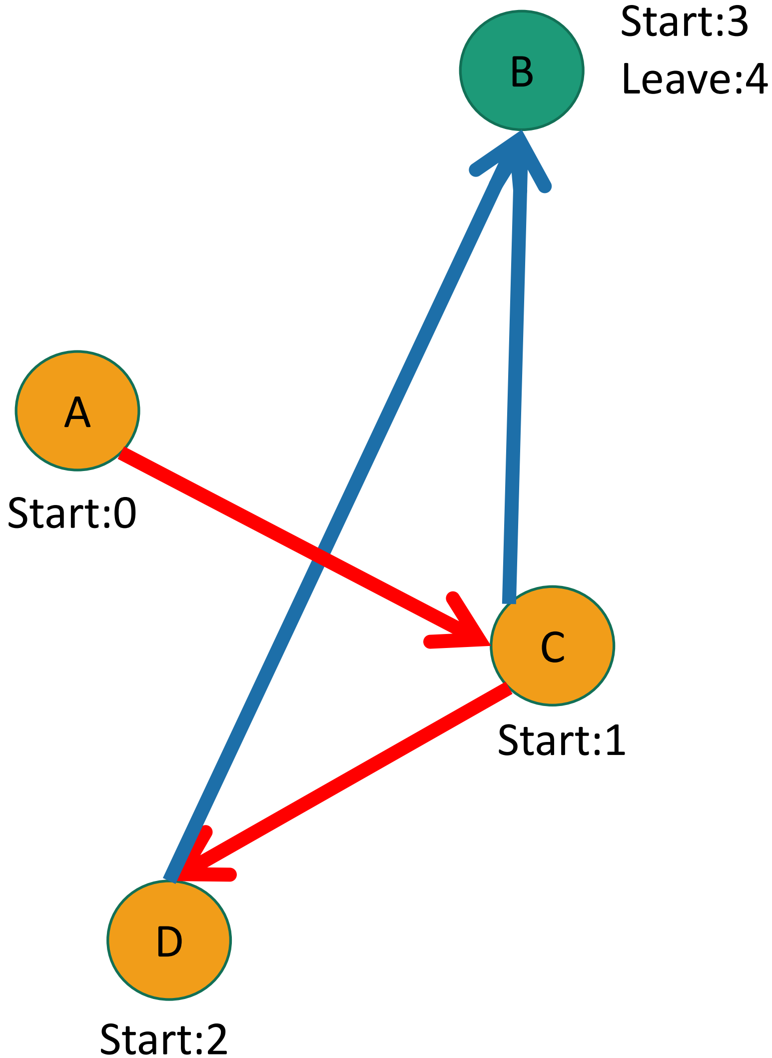


Example



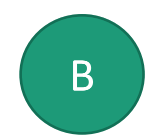
Start:3
B
A
Start:0
C
Start:1
D
Start:2
Unvisited
In progress
All done

Example



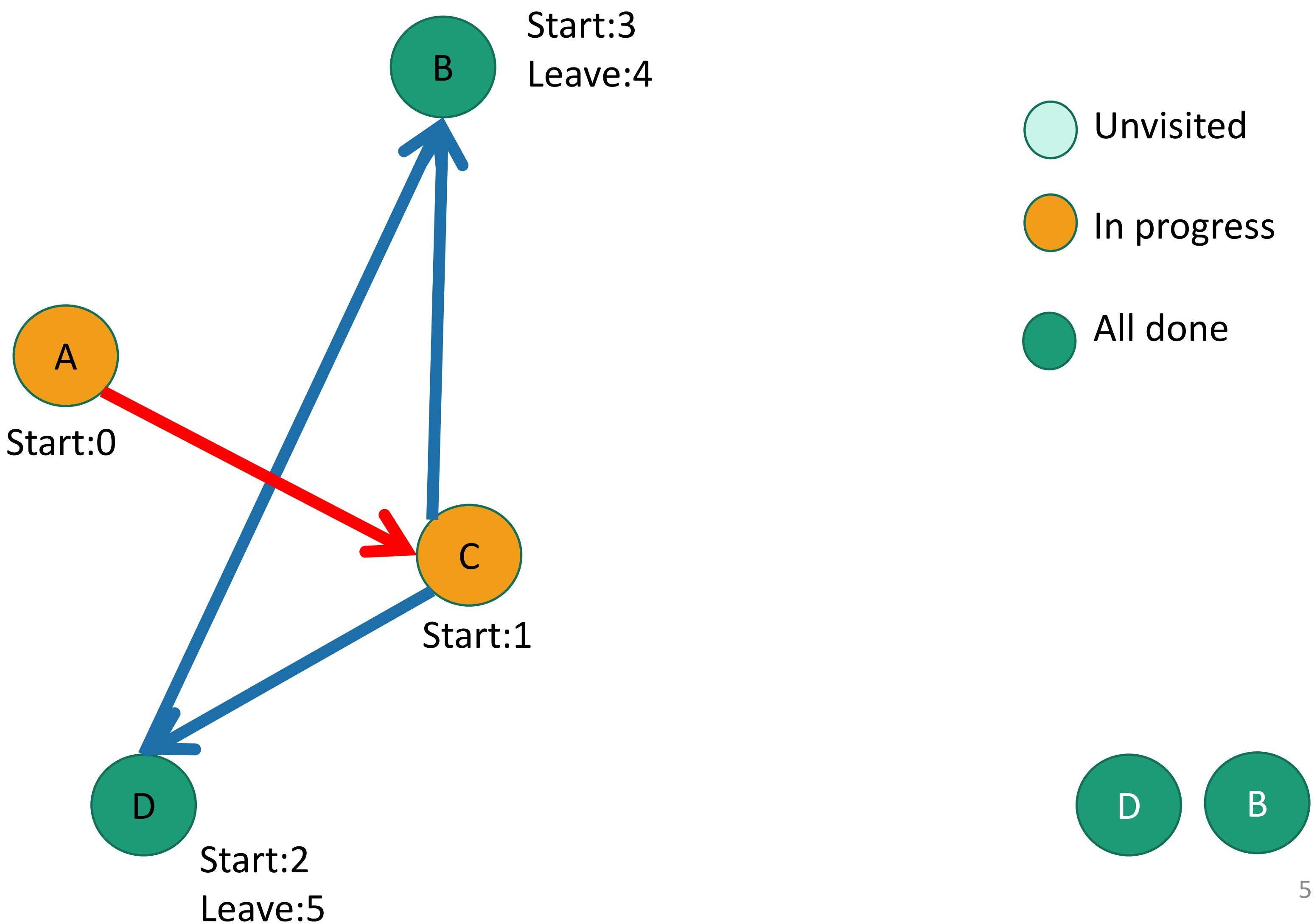
B
Start:3
Leave:4
A
Start:0
C
Start:1
D
Start:2

Unvisited
In progress
All done



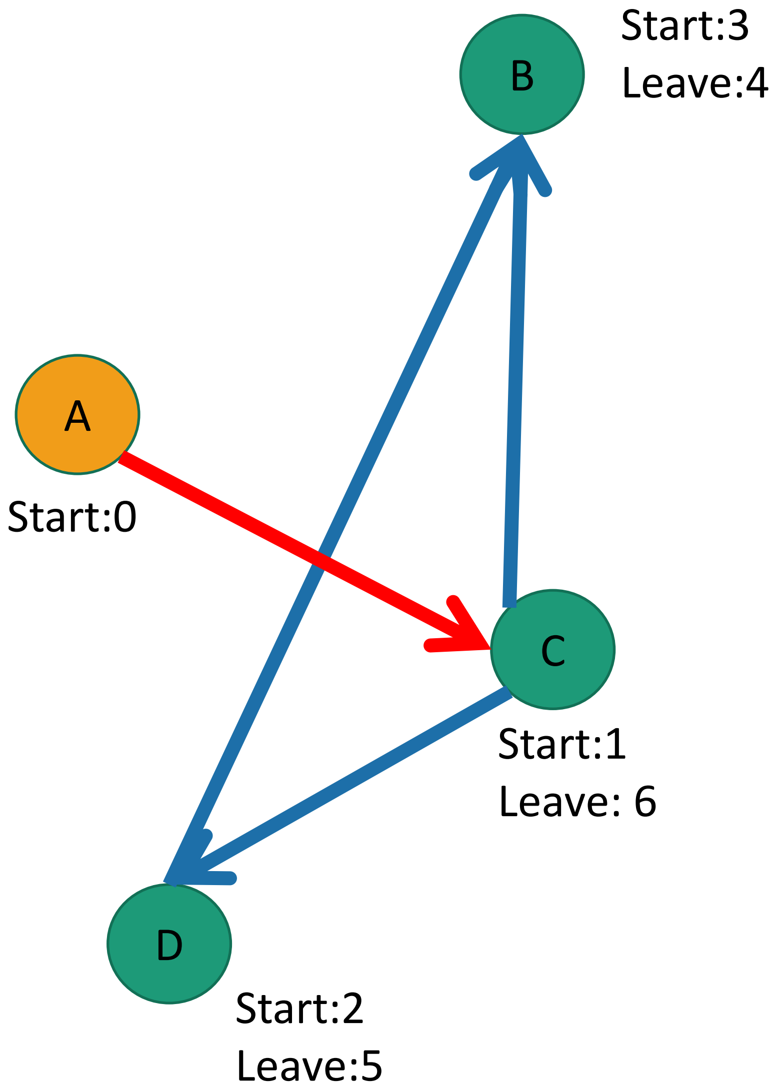
B

Example



B
Start:3
Leave:4
Unvisited
In progress
All done
A
Start:0
C
Start:1
D
Start:2
Leave:5
D
B

Example

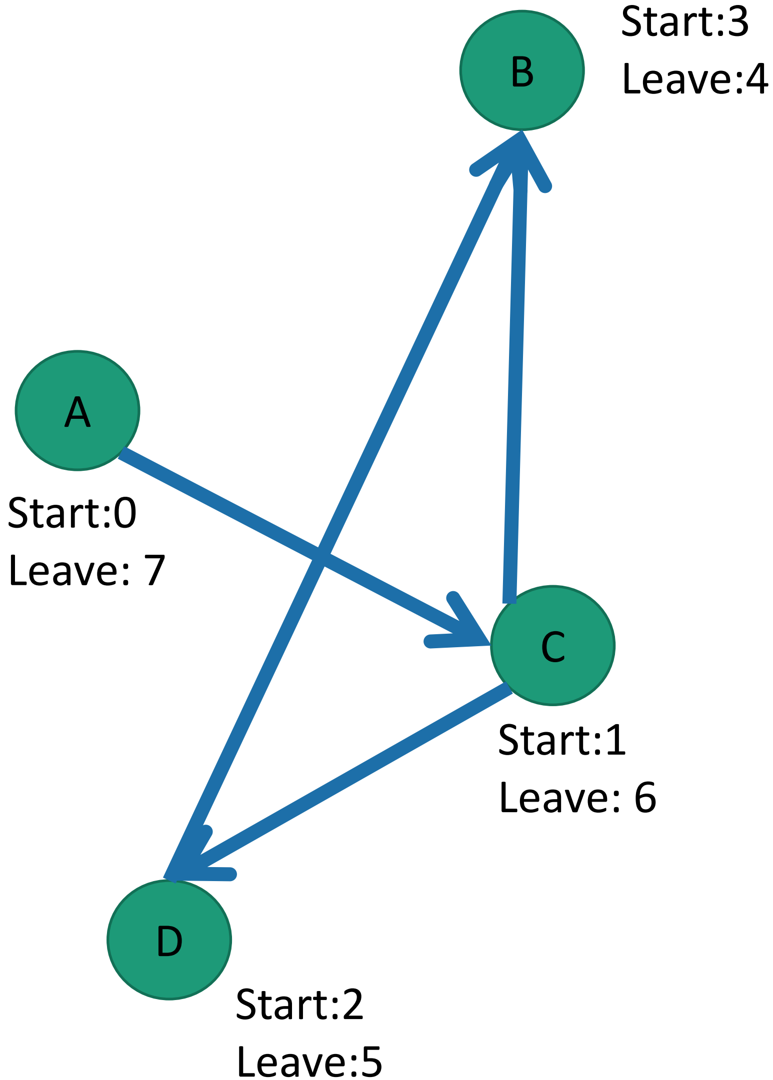


B
Start:3
Leave:4
A
Start:0
C
Start:1
Leave: 6
D
Start:2
Leave:5

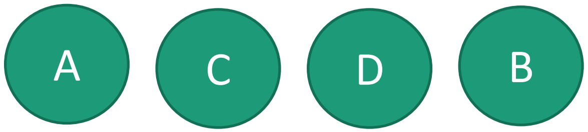
Unvisited
In progress
All done

C
D
B

Example

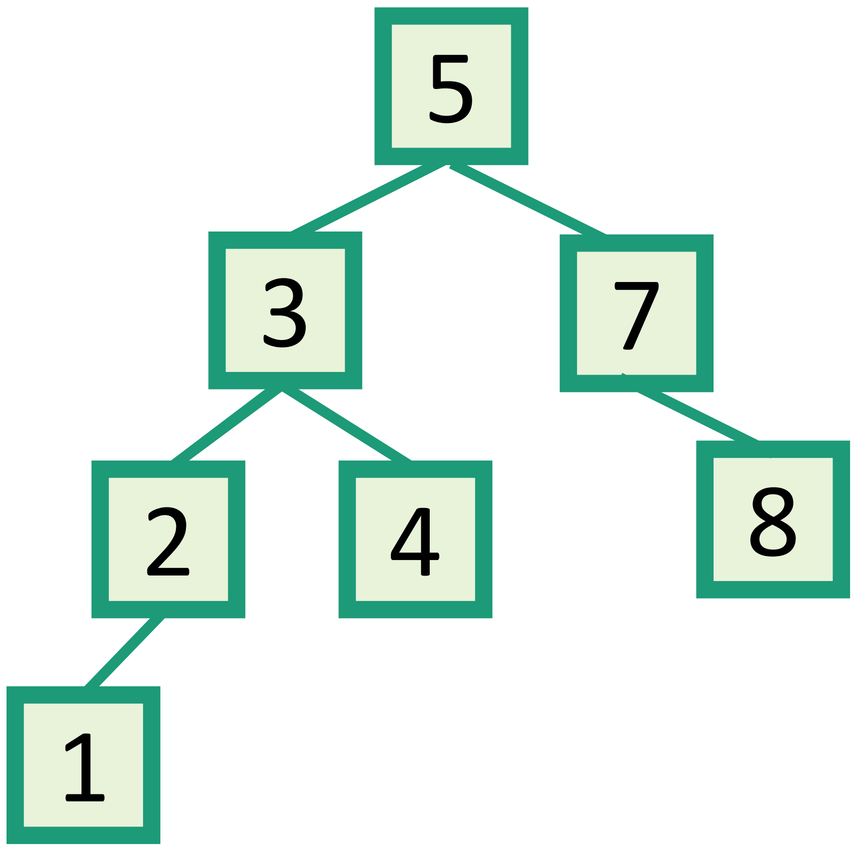


Do them in this order:



Another use of DFS that we've already seen

- In-order enumeration of binary search trees



Do DFS and print a node's label when you are done with the left child and before you begin the right child.

Acknowledgement

- Stanford University