PROOF TECHNIQUES

We look at

- Proof by contradiction
- Proof by construction
- Proof by induction

- One common way to prove a theorem is to assume that the theorem is false, and then show that this assumption leads to an obviously false consequence (also called a contradiction)
- This type of reasoning is used frequently in everyday life, as shown in the following example

- Jack sees Jill, who just comes in from outdoor
- Jill looks completely dry
- · Jack knows that it is not raining
- Jack's proof:
 - If it were raining (the assumption that the statement is false), Jill will be wet.
 - The consequence is: "Jill is wet" AND "Jill is dry", which is obviously false
 - Therefore, it must not be raining

By Contradiction [Example 1]

- Let us define a number is rational if it can be expressed as p/q where p and q are integers; if it cannot, then the number is called irrational
- E.g.,
 - 0.5 is rational because 0.5 = 1/2
 - 2.375 is rational because 2.375 = 2375 / 1000

- Theorem: $\sqrt{2}$ (the square-root of 2) is irrational.
- How to prove?
- First thing is ...

Assume that $\sqrt{2}$ is rational

- Proof: Assume that $\sqrt{2}$ is rational. Then, it can be written as p/q for some positive integers p and q.
- In fact, we can further restrict that p and q does not have common factor.
 - If D is a common factor of p and q, we use p' = p/D and q' = q/D so that $p'/q' = p/q = \sqrt{2}$ and there is no common factor between p' and q'
- Then, we have $p^2/q^2 = 2$, or $2q^2 = p^2$.

- Since 2q² is an even number, p² is also an even number
 - This implies that p is an even number (why?)
- So, p = 2r for some integer r
- $2q^2 = p^2 = (2r)^2 = 4r^2$
 - This implies $2r^2 = q^2$
- So, q is an even number
- Something wrong happens... (what is it?)

- We now have: "p and q does not have common factor" AND "p and q have common factor"
 - This is a contradiction
- Thus, the assumption is wrong, so that $\sqrt{2}$ is irrational

By Contradiction [Example 2]

- Theorem (Pigeonhole principle): A total of n+1 balls are put into n boxes. At least one box containing 2 or more balls.
- Proof: Assume "at least one box containing
 2 or more balls" is false
 - That is, each has at most 1 or fewer ball Consequence: total number of balls \leq n Thus, there is a contradiction (what is that?)

Proof By Construction

- Many theorem states that a particular type of object exists
- One way to prove is to find a way to construct one such object
- This technique is called proof by construction

- Theorem: There exists a rational number p which can be expressed as q^r , with q and r both irrational.
- How to prove?
 - Find p, q, r satisfying the above condition
- What is the irrational number we just learnt? Can we make use of it?

By Construction

- What is the following value? $(\sqrt{2}^{\sqrt{2}})^{\sqrt{2}}$
- If $\sqrt{2}$ is rational, then $q = r = \sqrt{2}$ gives the desired answer
- Otherwise, $q = \sqrt{2}^{\sqrt{2}}$ and $r = \sqrt{2}$ gives the desired answer

By Induction

- Normally used to show that all elements in an infinite set have a specified property
- The proof consists of proving two things: The basis, and the inductive step

 Mathematical induction proves that we can climb as high as we like on a ladder, by proving that we can climb onto the bottom rung (the basis) and that from each rung we can climb up to the next one (the inductive step).

We consider only enumerable or countable sets with a least element [well ordered sets]

- 1. The **base case**: prove that the statement holds for the first natural number n. Usually, n = 0 or n = 1;
 - rarely, but sometimes conveniently, the base value of n may be taken as a larger number, or even as a negative number (the statement only holds at and above that threshold).
- 2. The **step case** or **inductive step**: assume the statement holds for some natural number n, and prove that then the statement holds for n + 1.

By Induction [Example 1]

- Let F(k) be a sequence defined as follows:
- F(1) = 1
- F(2) = 1
- for all $k \ge 3$, F(k) = F(k-1) + F(k-2)
- Theorem: For all $n \ge 1$, F(1)+F(2) + ... + F(n) = F(n+2) - 1

By Induction

- Let P(k) means "the theorem is true when n = k"
- Basis: To show P(1) is true.
 - F(1) = 1, F(3) = F(1) + F(2) = 2
 - Thus, F(1) = F(3) 1
 - Thus, P(1) is true
- Inductive Step: To show for $k \ge 1$, $P(k) \rightarrow P(k+1)$
 - P(k) is true means: F(1) + F(2) + ... + F(k) = F(k+2) 1
 - Then, we have

$$F(1) + F(2) + ... + F(k+1)$$

= $(F(k+2) - 1) + F(k+1)$
= $F(k+3) - 1$

- Thus, P(k+1) is true if P(k) is true

Variants

- There can be many other types of basis and inductive step, as long as by proving both of them, they can cover all the cases
- For example, to show P is true for all k > 1, we can show
 - Basis: P(1) is true, P(2) is true
 - Inductive step: $P(k) \rightarrow P(k+2)$

Variants

 Complete (strong) induction: (in contrast to which the basic form of induction is sometimes known as weak induction)

makes the inductive step easier to prove by using a stronger hypothesis: one proves the statement P(m + 1) under the assumption that P(n) holds for all $n, n \leq m$.

Example: forming dollar amounts by coins

- Assume an infinite supply of 4 and 5 dollar coins.
- Prove that any whole amount of dollars greater than 12 can be formed by a combination of such coins.
- In more precise terms, we wish to show that for any amount $n \ge 12$ there exist natural numbers a and b such that n = 4a + 5b, where 0 is included as a natural number.
- The statement to be shown true is thus:

$$S(n): n \geq 12 \Rightarrow \exists a,b \in \mathbb{N}. \ n = 4a + 5b$$

Base case: Show that S(k) holds for k = 12, 13, 14, 15.

$$4 \cdot 3 + 5 \cdot 0 = 12$$

$$4 \cdot 2 + 5 \cdot 1 = 13$$

$$4 \cdot 1 + 5 \cdot 2 = 14$$

$$4 \cdot 0 + 5 \cdot 3 = 15$$

The base case holds.

Induction step:

For $j=12,13,\ldots,15,\ldots,k$ we assume that the theorem is true.

For j = k + 1, we show that the theorem is true.

Since for j = k, k - 1, k - 2, k - 3 the theorem is true (why?).

So, k - 3 = 4a + 5b, for some nonnegative integers a and b.

Since k + 1 = (k - 3) + 4,

we have, k + 1 = 4a + 5b + 4 = 4(a + 1) + 5b. Q.E.D.

•	The follow	wing is not	a valid pr	oof by inc	duction!

By Induction?

- CLAIM: In any set of h horses, all horses are of the same color.
- PROOF: By induction. Let P(k) means
 "the claim is true when h = k"
- Basis: P(1) is true, because in any set of 1 horse, all horses clearly are the same color.

By Induction?

Inductive step:

- Assume P(k) is true.
- Then we take any set of k+1 horses.
- Remove one of them. Then, the remaining horses are of the same color (because P(k) is true).
- Put back the removed horse into the set, and remove another horse
- In this new set, all horses are of same color (because P(k) is true).
- Therefore, all horses are of the same color!
- What's wrong?

More on Pigeonhole Principle

- Theorem: For any graph with more than two vertices, there exists two vertices whose degree are the same.
- · How to prove?

For connected graphs

First, suppose that G is a connected finite simple graph with n vertices. Then every vertex in G has degree between 1 and n-1 (the degree of a given vertex cannot be zero since G is connected, and is at most n-1 since G is simple). Since there are n vertices in G with degree between 1 and n-1, the pigeon hole principle lets us conclude that there is some integer k between 1 and n-1 such that two or more vertices have degree k.

For arbitrary graphs

Now, suppose G is an arbitrary finite simple graph (not necessarily connected). If G has any connected component consisting of two or more vertices, the above argument shows that that component contains two vertices with the same degree, and therefore G does as well. On the other hand, if G has no connected components with more than one vertex, then every vertex in G has degree zero, and so there are multiple vertices in G with the same degree.