

# GeoSched: Minimizing Cost of Cloud Workloads in Geo-Distributed Data Centers

## ABSTRACT

Today, many of the leading cloud service providers have geographically distributed data centers to serve millions of users around the world. The need for a geo-distributed data center arises from the need to provide low-latency and high-availability of the different services. The workloads that run on these data centers are also highly diverse. A typical data centers run jobs of different kinds including business-critical workloads—web servers, mail servers, instant messaging services etc, big-data processing, real-time data analytics and HPC jobs. As cloud services grow to serve more customers, there is an increasing need for a workload provisioning mechanism which can minimize the cost of operation of the data center while minimizing SLA violations and also minimizing the user-perceived latency.

We present GeoSched, a job provisioning framework that is workload aware, energy-aware and cooling-aware. GeoSched considers the workload type (batch process, real-time processing, web-service), the energy source of the data center (brown and green energy), electricity pricing of the region and cooling techniques used in the data center (air economizer). We GeoSched on industry workload traces and also on synthetic workloads that can mimic the real-world data center workload.

## 1. RELATED WORK

Cost minimization has been extensively studied both in the context of a single data center and a geo-distributed data center. Most of the work has been targeted towards minimizing the energy consumption of the data center to reduce the operational cost. For example, the work by Gu et al. [1] studies the influence of task assignment, data placement and data movement on the operational expenditure of large-scale geo-distributed data centers for big data applications. The authors jointly optimize these three factors by proposing a 2-D Markov chain to derive the average task completion time and solve the model as a MILP problem. On similar lines Agarwal et al. [2] propose an automated data placement mechanism Volley for geo-distributed cloud services with the consideration of WAN bandwidth cost, data center capacity limits, data inter-dependencies, etc.

Yu et al.[3] propose minimizing energy cost for distributed Internet data centers (IDCs) in smart microgrids by considering system dynamics like uncertainties in electricity price, workload, renewable energy generation, and power outage state. They model the problem as a stochastic program and solve it using Lyapunov optimization technique. The work by Ren and He [4] propose an online algorithm, called COCA to minimize data center operational cost while satisfying carbon neutrality. Unlike some of the existing research, COCA enables distributed server-level resource management: each server autonomously adjusts its processing speed and optimally decides the amount of workloads to process. This work is restricted to a single data center with heterogeneous resources.

Cooling energy is a substantial part of the total energy spent by a data center. According to some reports, this can be as high as 50% [5], [6],[7] of the total energy budget. There has been a lot of work both at the application/middleware level [8], [9] as well as provisioning techniques [10], [11] that try to minimize the cooling energy. Li et al. [12] proposes SmartCool, a power optimization scheme that effectively coordinates different cooling techniques and dynamically manages workload allocation for jointly optimized cooling and server power. Unlike existing work that addresses different cooling techniques in an isolated manner, SmartCool integrates different cooling systems as a constrained optimization problem. Also, since geo-distributed data centers have different ambient temperatures, SmartCool dynamically dispatches the incoming requests among a network of data centers with heterogeneous cooling systems to best leverage the high efficiency of free cooling.

## 2. CLOUD WORKLOAD

With the increase in the adoption of cloud computing by businesses around the world, more and more workloads are moving out of traditional data centers into the cloud. A cloud data center typically handles a range of IT workloads. These workloads include service-critical and latency sensitive interactive applications that run 24x7 such as instant messengers, e-mail applications and other internet services, and batch-style applications such as scientific HPC applications, BigData applications like financial analysis and MapReduce applications. A thoughtful management of these workloads are necessary not only to provide customers with low-latency services but also to reduce the energy consumption of the data centers that run these workloads. A good way to study these workloads are by analyzing real-world or synthetic workload traces. We look at cloud workload traces in detail in section 2.1 and 2.3

### 2.1 Workload Traces

Unfortunately, there is a lack of real-world traces in the cloud computing research community. As a result, there are limited work done to understand the diversity in cloud workloads. The first large-scale workload trace was released by Google in 2011[13], featuring traces from over 12,000 servers over a period of 29 days. Various studies have been carried out to understand the characteristic of the Google trace data. Other real-world traces include Wikipedia request trace [14] and Hadoop logs from Yahoo! [15] There have also been attempts to generate synthetic workload traces using known real-world workload characteristics [16] [17]. For our experiments we have decided to use the workload trace from Google.

### 2.2 Custom Trace Format

To have a common trace format to represent different real-world and synthetic workload traces, we have formulated a trace format similar to standard workload format (swf) [18] defined for grid and supercomputing traces. Table 1 provides information about each field in the trace format.

Field Name	Type	Description
JobID	INTEGER	Unique job identifier for each job
Arrival Time	INTEGER	Arrival time in microseconds
Scheduled Time	INTEGER	Arrival time in microseconds
Finish Time	INTEGER	Arrival time in microseconds
Schedule class	INTEGER	Latency class of the job.
Tasks	INTEGER	Total number of tasks in the job.
Total CPU	INTEGER	Aggregate sum of CPU used by all the tasks in the job
Total Memory	INTEGER	Aggregate sum of memory used by all the tasks in the job
User	INTEGER	Unique user identifier
SLA Param1	FLOAT	SLA parameter placeholder
SLA Param2	FLOAT	SLA parameter placeholder
SLA Param3	FLOAT	SLA parameter placeholder

Table 1: Trace format description

### 2.3 Google Workload Trace

Google trace has job and task information in different files. We consolidated all the required information and generated a single work trace file with the format described in the previous section. For our work, we model five different data centers located at different parts of the globe. A detailed explanation about the data center profile are given in section 3. To simulate workload execution in these five data centers we have extracted five set of traces each worth five days of incoming job requests from the original google trace. This will cover 25 days of traces out of the 29 available days. We label these 5 set of traces from A to E. Additionally, in our experiments we only consider those jobs that end without any failures. Therefore, we clean the original traces by removing all jobs that is either fail or get killed.

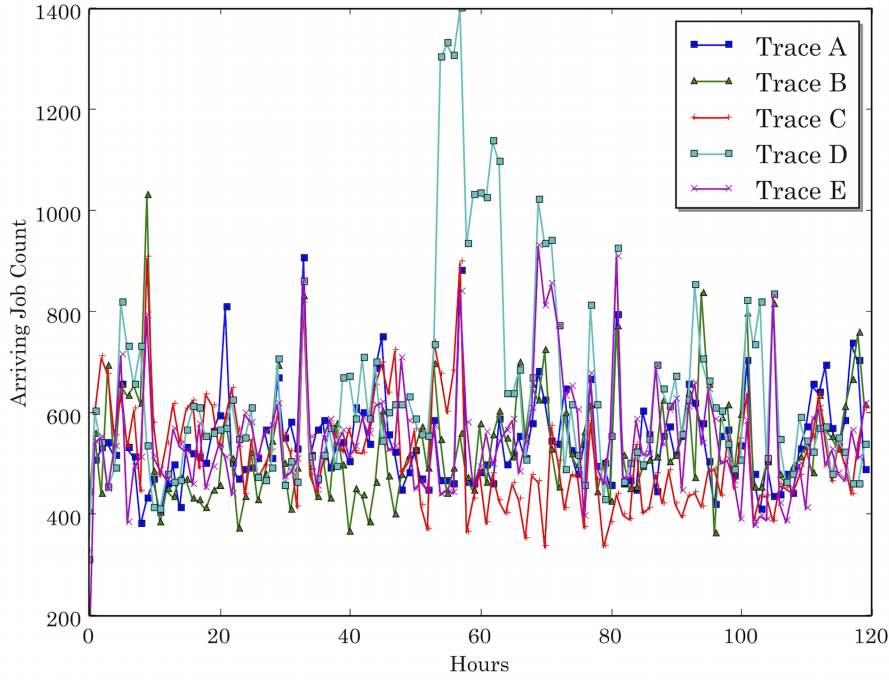


Figure 1: Arrival rate of jobs

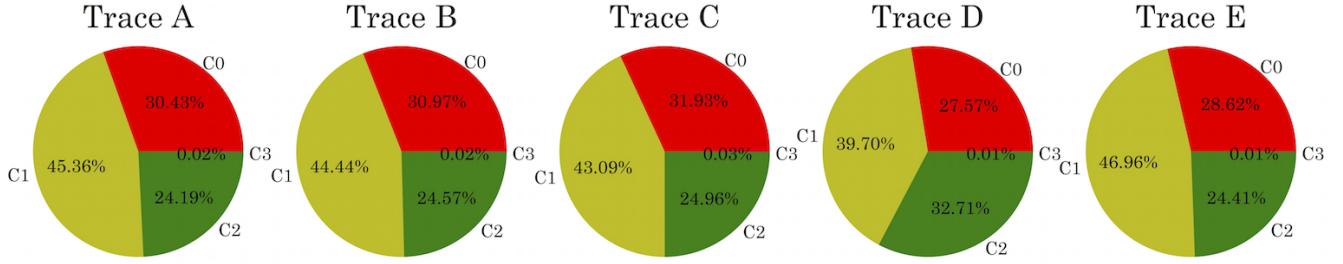


Figure 2: Distribution of job count in various schedule classes

Now we will look at some characteristics of the traces. Table 2 describes the statistics for each of the five traces grouped by the scheduling class of the job. In google traces, CPU and memory usage are normalized relative to the largest capacity of the resource on any machine in the trace (which is 1.0). Most of the class 0 and class 2 jobs are batch jobs with less latency requirements. Generally, these are the kind of jobs that can be provisioned on other data centers as their deadlines are not strict. You will find that as the sensitivity increases the running time of these jobs and the resource usage also increase. High-sensitive jobs are generally service jobs and run of longer periods of time. Figure 2 shows a pie chart of the job counts in different schedule classes for all the 5 traces. You will notice that the job count for class 3 (high latency-sensitive) jobs are low. This is because these jobs generally run for long period of time, typically months. Since we only consider jobs that start and end within the 29 day range such jobs are missed out in our traces. For our work, we will consider class-0 and class-1 jobs as candidates for optimizing the provisioning. We leave, class-2 and class-3 jobs to run on the data center where the job arrived since these are high-sensitive and their SLA demands will be strict.

Figure 1 shows the variation in job arrival rates for the traces. As you can see there is both intra-trace

	Trace A			Trace B			Trace C			Trace D			Trace E		
Sched Class	Avg CPU	Avg Mem	RT (secs)	Avg CPU	Avg Mem	RT (secs)	Avg CPU	Avg Mem	RT (secs)	Avg CPU	Avg Mem	RT (secs)	Avg CPU	Avg Mem	RT (secs)
0	0.075	0.051	415	0.123	0.096	345	0.083	0.065	385	0.093	0.092	373	0.093	0.074	333
1	0.087	0.029	729	0.086	0.032	760	0.082	0.028	711	0.073	0.030	743	0.093	0.043	693
2	0.067	0.043	1028	0.096	0.058	872	0.088	0.057	918	0.073	0.045	707	0.084	0.049	832
3	0.088	0.12	1105	0.099	0.12	82	0.101	0.012	7961	0.081	0.051	17135	0.089	0.061	9021

Table 2: Trace set statistics

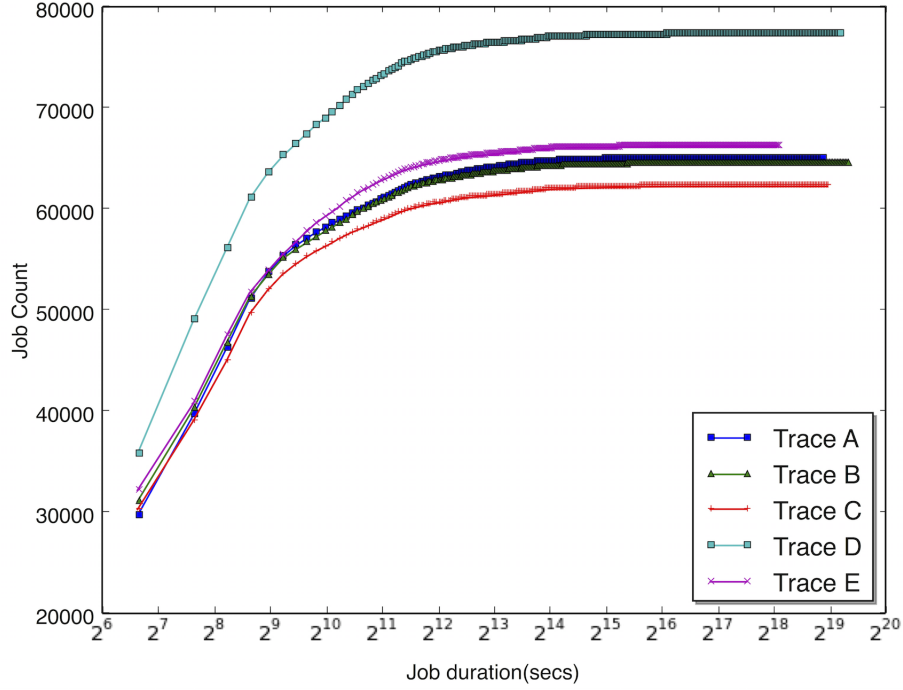


Figure 3: CDF of the job duration

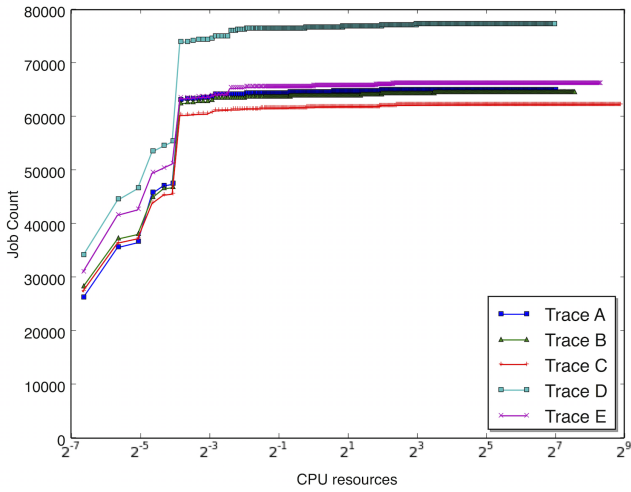
variation as well as inter-trace variation. This is representative of the real-world cloud workload pattern where bursty arrivals are common. There are opportunities to provision jobs to other data centers if the traffic at a particular data center is high. What is also important to notice is that the scheduling decision should be made quickly when the arrival rates of jobs are high, else there are chances for that these jobs miss the SLAs. Hence, even the scheduler should be scalable during high volume traffic.

Figure 3 shows the CDF of the job runtime. As you can see most of the jobs have runtime less than 1024 seconds. This pattern is consistent for all the trace sets. There are a few jobs that have higher run-times. As explained earlier, these jobs are primarily internet-service jobs that are expected to run for longer duration.

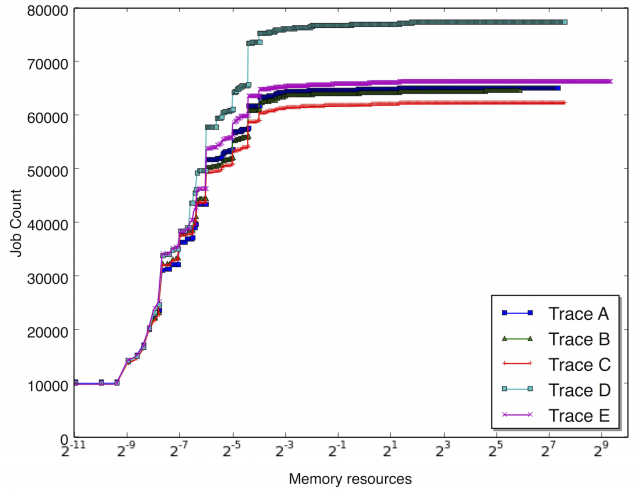
In google workload trace, it is tough to study the resource usage since these value are scaled down and therefore and accurate picture of the resource utilization is tough to generate. In figure 4 we can see that most jobs request a small portion of the cluster resources. Most of the resources in the google cluster are used by the few service jobs that are present. The large number of batch jobs only use a small percentage of the resources [19] [20]. From the above statistics, we believe that google workload trace provides an accurate representation of the real-world cloud datacenter workload.

### 3. DATA CENTER PROFILING

In order to run our experiments, we need to collect information about data centers that will help us make better scheduling decisions. We call this process data center profiling. The major energy costs in data



(a) CPU resource



(b) RAM resource

Figure 4: CDF of resource requested by jobs

center are from running the IT resources and from cooling the data centers. We create a profile for each of the data center with the information that effects the energy costs.

Location	Label	Associated Trace
Berkeley County, South Carolina	DC1	A
Lenoir, North Carolina	DC2	B
Singapore	DC3	C
Quilicura, Chile	DC4	D
Dublin, Ireland	DC5	E

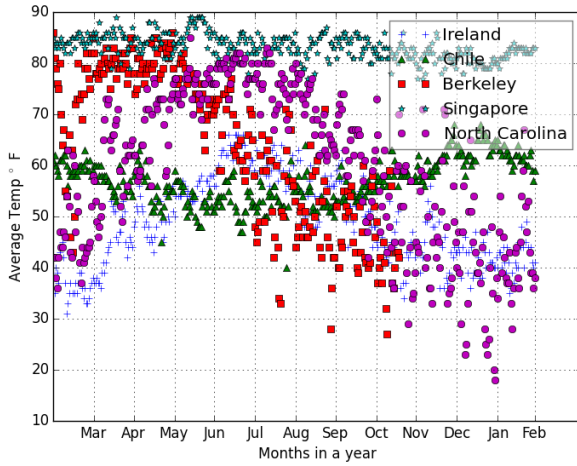
Table 3: Data center location information

### 3.1 Geo-Distributed Data Center

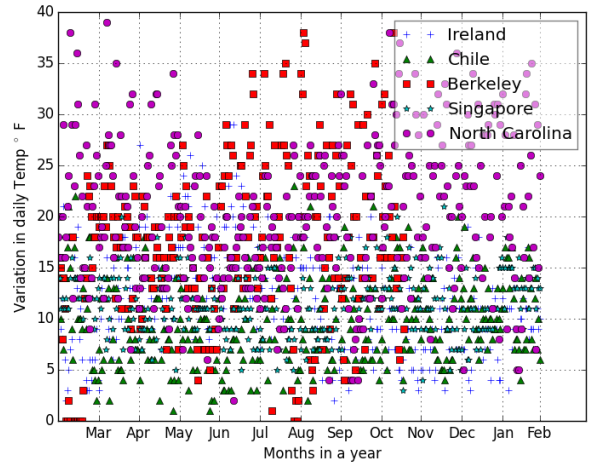
Most cloud service providers including Google, Amazon, Facebook and Apple have data centers distributed across the globe. For our experiments, we consider five locations where Google has its data centers[13]. Table 3 provides the location and labels for each of the data center. We also replicate the computing resource used in the original google trace across all the data centers. The configuration of the resources along with the normalized CPU and memory are given in table 4

Node count	Architecture	CPU	Memory
6732	B	0.5	0.50
3863	B	0.5	0.25
1001	B	0.5	0.75
795	C	1	1
126	A	0.25	0.25
52	B	0.5	0.12
5	B	0.5	0.3
5	B	0.5	0.97
3	C	1	0.50
1	B	0.5	0.06

Table 4: Data center IT resource configuration



(a) Daily Average Temperature



(b) Daily Maximum Variation

Figure 5: Temperature variation for the year 2013

### 3.2 Profile Parameters

We consider two major parameters that has the most effect on the energy cost. They are a) Temperature b) Electricity. We collect temperature and electricity prices for these locations from historical archives. If time permits, we wish to also consider green energy available at the data center locations. We look at each of these parameters and how they effect energy costs in sections 3.2.1, 3.2.2 and 3.2.3.

#### 3.2.1 Temperature

Data centers located in cold regions have the advantage of using outside air to cool the data center. An air-side economizer brings outside air into a building and distributes it to the servers. Instead of being re-circulated and cooled, the exhaust air from the servers is simply directed outside. This can result in million dollars of saving annually. A difference in temperature across data centers mean different cooling costs at these data centers. We wish to utilize this difference by scheduling jobs in low temperature data centers. In this case we make our assumption concrete by doing an empirical analysis of historical climate data. We extracted these historical climate data from various data repositories of the National Climate Data Center [21] for all five locations covering the entire one-year period of 2013.

Figure 5 a plots the daily average temperature for five different locations in United States, Europe, South-America, and Asia. As expected we see high diversity in temperature at different locations. For example, North Carolina, Ireland seems to be equally favorable for cooling during the month of November - March. Chile and Singapore show almost no variation in temperature across the year. However, Chile is favorable for cooling conditions compared to other data centers in the month of May - July.

Figure 5 b plots the daily maximum variation in temperature across the day i.e. difference of maximum temperature and minimum temperature in a day. It is worth noting that almost every location faces a variation of up to 40°F during different periods. These highs and lows do not occur at the same time across different regions due to time differences. It is evident from the results that there is no single location that is always cooling efficient. So we conclude that we need a dynamic scheduler, which caters to the locations that have lower temperature and henceforth save overall energy costs.

#### 3.2.2 Electricity Pricing

Creation of power trading zones have resulted in wholesale power purchase under dynamic pricing schemes. Such markets are popular in USA [22], Germany [23] and other countries. Based on supply and demand bids there can be high fluctuations in prices [24]. There are many ways to predict these prices [25] [26].

If we can predict the electricity prices for the job duration, then we will be able to schedule the job such that the running costs go down.

### 3.2.3 Energy Source

Solar and wind energy are available in most of the modern data centers. Similar to electricity price predictions green energy prediction can also be done [27] [28]. Since, this part is in exploratory stage, we will provide more information during the final report.

### 3.2.4 Case Study: Google Data Centers

Google data-centers show 50% less energy consumption than a typical data center. Google is able to achieve a trailing twelve-month (TTM) PUE of 1.12 across all data centers, in all seasons, including all sources of overhead. There are five techniques that Google typically uses to minimize its energy consumption and electricity costs. Frequent measuring of Power Usage Effectiveness is the key for determining the power usage for non-computing functions like cooling and power distribution. Secondly they manage airflow at data-centers, which allows cooling the hot spots available in data-centers. Moreover by setting thermostat to 80°F it saves significant amount of electricity consumption and costs.

## 4. SCHEDULER SIMULATOR

We plan to implement our simulator using CloudSim [29]. CloudSim is an extensible simulation toolkit that enables modeling and simulation of Cloud computing systems and application provisioning environments. The CloudSim toolkit supports both system and behavior modeling of Cloud system components such as data centers, virtual machines (VMs) and resource provisioning policies.

### 4.1 Components

- **Datacenter:** This component will simulate a data center. This composed of 1) scheduler that implements a specific provisioning algorithm, 2) data center profile database, 3) a work trace reader and 4) proxy elements of other data center for inter-data center communication
- **DataCenter Proxy:** All interactions with the data center will be through a proxy. Each datacenter will hold the proxies to other data centers.
- **Scheduler:** This component will implement the optimization technique and schedule jobs onto data centers.
- **Datacenter Profile DB:** This components contains the temperature and electricity data. Scheduler will access this to make provisioning decisions.

#### 4.1.1 Component Interaction

Figure 6 gives the interaction model between the components. The datacenter object will start a thread which will run the scheduling algorithm, scheduling jobs from the trace file. Each data center had proxies to other data centers if it needs to communicate resource availability or job provisioning to other data centers.

## 5. CONCLUSION

We have completed the work trace analysis and data center profiling parts of the project. Currently the simulator is being implemented, following which we will work on the energy minimization techniques and the final evaluation of these techniques. If time permits, we would also like to consider green energy in our provisioning mechanism. Section 6 shows the current state sand future directions of the project.

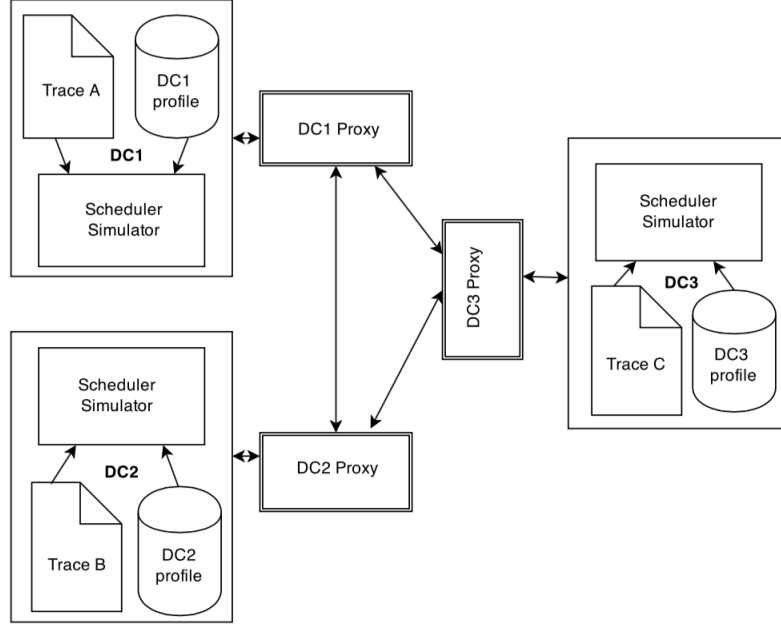


Figure 6: Simulation components and their interaction

## 6. MILESTONES

Item	Deadline	Description	Status
Study real-world and synthetic workloads	10 <sup>th</sup> March	We have started studying workloads traces for google cluster-data which provides traces from a 12k-machine cell for one month duration in May 2011.	Completed
Data Center profile	20 <sup>th</sup> March	In this process we will note the data center specific parameters such as electricity cost, available energy, economizer, cooling energy that is useful for simulating geo-distributed data centers.	Completed
Simulator design and implementation	12 <sup>th</sup> April	Implement the simulator infrastructure for multi-datacenter environment.	In progress
Feature Extraction and Optimization	25 <sup>th</sup> April	Here we identifying all the necessary features that contribute to the energy cost and build an optimization technique which is adaptive to workload characteristics and data center's resource capacity and electricity costs	TBD
Evaluation	7 <sup>th</sup> May	This phase we will be simulating data centers with previously gathered data center profile and using the optimization technique mentioned above	TBD

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