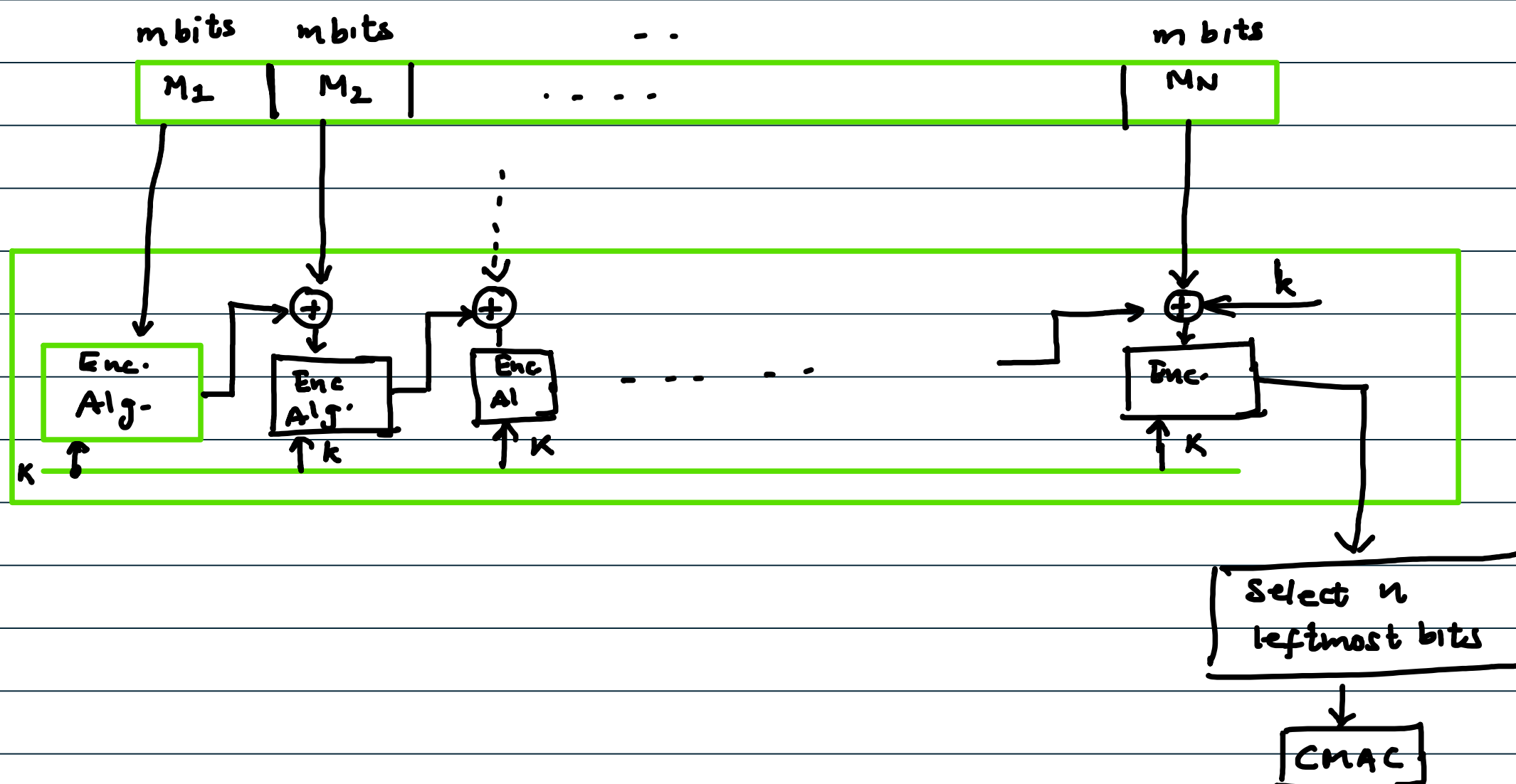


## CMAC (Data Authentication Algorithm)



Step 1: Divide the msg into  $N$  blocks of  $m$  bits

Step 2: Pad the last block with a 1 bit followed by enough 0 bits to  $m$ -bits

Step 3: Find  $C_1 = E_K(M_1)$ ,  $C_2 = E_K(C_1 \oplus M_2)$ ,  $C_3 = E_K(C_2 \oplus M_3)$ ,  $\dots$ ,

$$C_{N-1} = E_K(C_{N-2} \oplus M_{N-1}), \quad C_N = E_K(C_{N-1} \oplus M_N \oplus k)$$

where  $E_K$  is the encryption alg with the key  $K$

Step 4: Select  $n$  leftmost bits of  $C_N$ . It would give us the  $n$ -bit CMAC

### Key Generation for the Last Step

Step 1: Find  $C = E_K(\underbrace{000\dots 0}_{m \text{ bits}})$

Step 2:  $k = \begin{cases} xC & \text{if no padding in } M_N \\ x^2C & \text{if padding is applied in } M_N \end{cases}$

Here, the multiplication is in the field  $GF(2^m)$  with the irreducible poly of degree  $m$  selected by the particular protocol

- Note:
1. MACs provides msg integrity & msg auth.
  2. MACs are much faster than the digital sign.

## Hash Function

This is a function that creates a fixed sized digest out of a variable size msg.

### Desirable Properties

- 1 Hash function should be computationally efficient.
- 2 Hash function should be of fixed length and it should be independent of the input length.

(In practice, output length of the hash value is b/w 128 bits to 512 bits)

- 3 The computed hash value should be highly sensitive to all input bits

### Security Requirements:

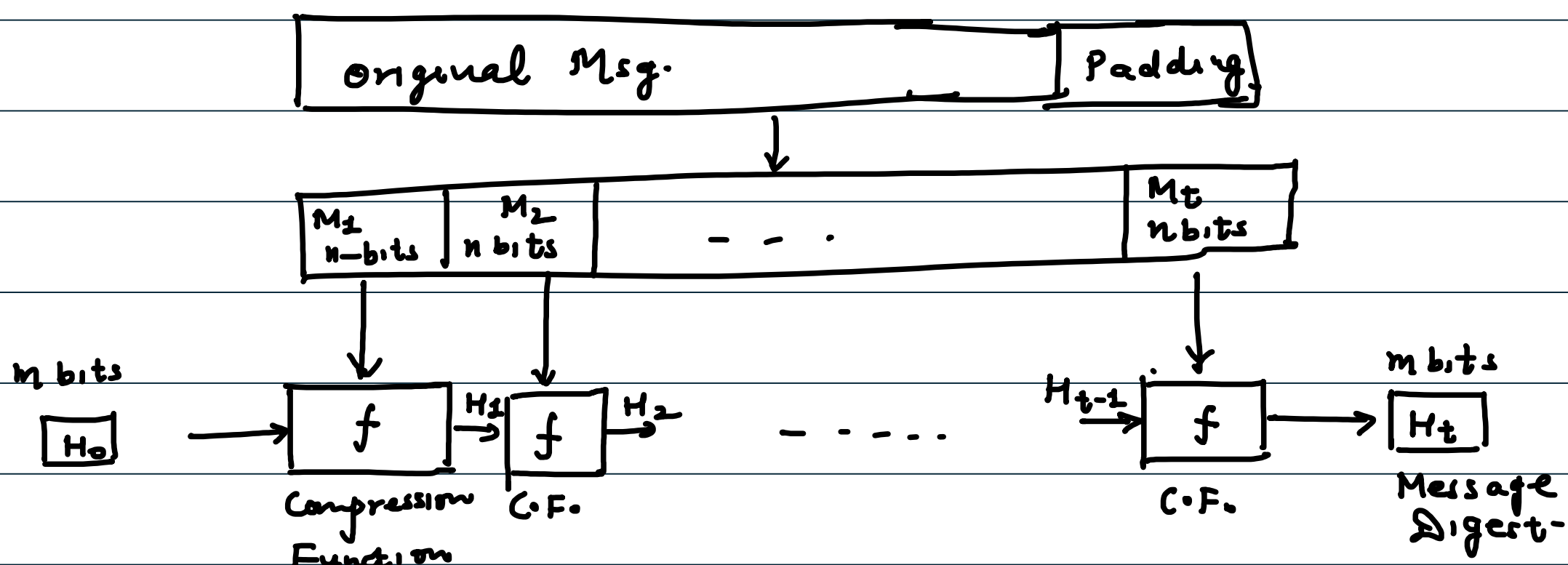
- 1 Preimage Resistance
- 2 Second Preimage Resistance
- 3 Collision Resistance

1. Given  $h(x)$ , it should be computationally infeasible to find  $x$

2. Given  $h(x)$ , it should " " " " " "  $x'$  s.t.  
 $h(x) = h(x')$

3. " " " " " "  $x$  &  $x'$  s.t.  
 $h(x) = h(x')$

### Merkle-Damgard Scheme



Step 1: Divide the msg into  $t$  blocks of  $n$  bits each <sup>after</sup> apply the padding if required

Step 2: Compute  $H_i = f(H_{i-1}, M_i)$ ,  $i = 1, 2, \dots, t$

Here,  $H_0$  is a fixed value which is called the initial value or initial vector

$H_t$  is the cryptographic hash function of the original msg.

$$H_t = h(M)$$

Note: If the compression function  $f$  is collision res. then hash function is collision resistant

Message Digest (MD) (Family of Hash alg.)

Ex. MD2, MD4, MD5

SHA (Secure Hash Algorithm)