

Database Management System (DBMS)

Assignment 2 (MC-302)

Delhi Technological University (DTU)

- a) Consider the relation shown in the following table and list all functional dependencies that this relation satisfies.

X	Y	Z
x ₁	y ₁	z ₁
x ₂	y ₁	z ₂
x ₁	y ₂	z ₂
x ₂	y ₂	z ₂

Ans) Now, for a relationship $\alpha \rightarrow \beta$ $\alpha, \beta \subseteq R$ to be satisfied the following needs to be true.

$$\text{If } t_1[\alpha] = t_2[\alpha], \text{ then:} \\ t_1[\beta] = t_2[\beta].$$

Now, the following dependencies can exist.

i) $X \rightarrow Y$

$$X[1] = x_1, \quad X[3] = x_2,$$

$$Y[1] = y_1, \quad Y[3] = y_2, \text{ hence not functionally dependent}$$

i) $X \rightarrow Z$

$$X[1] = x_1, \quad X[3] = x_1$$

 $Z[1] = z_1, \quad Z[3] = z_2$, hence functionally independentii) $Y \rightarrow X$

$$Y[1] = y_1, \quad Y[2] = y_1$$

 $X[1] = x_1, \quad X[2] = x_2$, hence functionally is dependentiii) $Y \rightarrow Z$

$$Y[1] = y_1, \quad Y[2] = y_1$$

 $Z[1] = z_1, \quad Z[2] = z_2$, hence functionally independentiv) $Z \rightarrow X$

$$Z[2] = z_2, \quad Z[3] = z_2$$

 $X[2] = x_2, \quad X[3] = x_1$, hence functionally is dependentv) $Z \rightarrow Y$

$$Z[2] = z_2, \quad Z[3] = z_2$$

 $Y[2] = y_1, \quad Y[3] = y_2$, hence functionally is dependentvi) $XY \rightarrow Z$

This is functionally dependent, and is a dependency

vii) $YZ \rightarrow X$

This is a valid functional dependency.

viii) $ZX \rightarrow Y$

This is a valid functional dependency.

So, the functional dependencies present in the relation are:-

$$F = \{ XY \rightarrow Z, YZ \rightarrow X, ZX \rightarrow Y \}$$

Q2) Given the following set S of functional dependencies

$$F = \{ M \rightarrow O, NP \rightarrow M, O \rightarrow N \}$$

Check if the following dependencies can be deduced from S , if not then give counter example:-

A) $OP \rightarrow M$

$$O \rightarrow N \text{ (F3)}$$

$$OP \rightarrow NP \text{ (Augmentation)}$$

$$NP \rightarrow M \text{ (F2)}$$

$$OP \rightarrow M \text{ (Transitive property)}$$

B) $NO \rightarrow M$

This can't be achieved, as stated in the below example:-

N	O	M	P
n ₁	o ₁	m ₁	p ₁
n ₁	o ₁	m ₂	p ₂

C) $MP \rightarrow N$

$$M \rightarrow O \text{ (F1)}$$

$$O \rightarrow N \text{ (F3)}$$

$$M \rightarrow N \text{ (Transitivity)}$$

$$MP \rightarrow NP \text{ (Augmentation)}$$

$$MP \rightarrow N \text{ (Decomposition)}$$

D) $MO \rightarrow P$

This functional dependency doesn't exist

M	O	P	N
m ₁	o ₁	p ₁	n ₁
m ₁	o ₁	p ₂	n ₁

E) $MN \rightarrow P$

This functional dependency isn't possible

M	N	P	O
m ₁	n ₁	p ₁	o ₁
m ₁	n ₁	p ₂	o ₁

Q4) Consider the relation schema $R(A, B, C, D, E, G)$ with functional dependencies

$F = \{ AB \rightarrow C, AG \rightarrow E, B \rightarrow D, E \rightarrow G \}$. Notice this is the minimum cover of itself

For each of the following decompositions $R(ABCDEG)$,
 determine whether it is lossless, and dependency preserving

i) $\{ABC, CDE, EG\}$

$$F_1 = \{AB \rightarrow C\}$$

$$F_2 = \{\}$$

$$F_3 = \{E \rightarrow G\}$$

Functional dependencies are not preserved, now
 for lossless decomposition:-

	A	B	C	D	E	G
R_1	α	α	α			
R_2						
R_3					α	α

This is clearly lossy.

ii) $\{ABCD, AEG\}$

$$F_1 = \{AB \rightarrow C, B \rightarrow D\}$$

$$F_2 = \{AG \rightarrow E, E \rightarrow G\}$$

This is functional dependency preserving.

	A	B	C	D	E	G	
R_1	α	α	α	α			This is lossy composition
R_2	α				α	α	



iii) $\{ABCE, BD, AEG\}$

$$F_1 = \{AB \rightarrow C\}$$

$$F_2 = \{B \rightarrow D\}$$

$$F_3 = \{AG \rightarrow E, E \rightarrow G\}$$

This satisfies all functional dependencies.

	A	B	C	D	E	G
R_1	x	x	x	x		
R_2		x		x		
R_3	x				x	x

This is lossy decomposition.

iv) $\{AB, ADE, BCG\}$

$$F_1 = \{ \}$$

$$F_2 = \{ \}$$

$$F_3 = \{ \}$$

No functional dependency is covered and it is lossy decomposition.

4) $\{BDEG, ABC\}$

$$F_1 = \{B \rightarrow D, E \rightarrow G\}$$

$$F_2 = \{AB \rightarrow C\}$$

This doesn't cover all functional dependencies.

	A	B	C	D	E	G
R_1		x		x	x	x
R_2	x	x	x	x		

This is lossy decomposition.

Q3) Consider the relational schema $R(A, B, C, D)$ with functional dependencies

$$F = \{A \rightarrow D, B \rightarrow C, AC \rightarrow D\}$$

i) $A^+ = \{A, D\}$ as $(A \rightarrow D)$

ii) $\{A, B\}^+ = (\{A\}^+ \cup \{B\}^+)^+ = \{A, D, B, C, D\}^+ = \{A, B, C, D\}$

iii) $A_{(A \rightarrow D)}^+ = A \neq A^+ = \{A, D\}$, hence $(A \rightarrow D)$ is not part of the canonical cover.

$$B_{(B \rightarrow C)}^+ = BD \neq B^+ = \{B, C, D\}$$

Hence $(B \rightarrow C)$ is part of the canonical cover.



100)

$$B_{\{B \rightarrow D\}}^+ = \{B, C\} \neq B^+ = \{B, C, D\}$$

Hence $\{B \rightarrow D\}$ is part of the minimal cover.

$$AC_{\{AC \rightarrow D\}}^+ = \{A, C, D\} = \{A, C\}^+ = \{A, C, D\}$$

Hence $\{AC \rightarrow D\}$ is redundant.

$\bar{F} = \{A \rightarrow D, B \rightarrow CD\}$ New canonical cover of the functional dependencies.

d) Candidate keys of R

Non-determined attributes = $\{A, B\}$

These attributes must be part of the candidate key.

$$\{AB\}^+ = \{A, B, C, D\}$$

Any other candidate key isn't possible as that will become a superkey.

Candidate keys = $\{\{A, B\}\}$

Q5) Consider the relational schema $R(A, B, C, D)$ with functional dependencies $\{A \rightarrow B, BC \rightarrow A, B \rightarrow D\}$ which is minimum cover of itself.

i) Find all candidate keys of $R(A, B, C, D)$

Non-determined attributes: $\{C\}$

Hence $\{C\}$ must be part of the candidate key

$$C^+ = \{C\}$$

$$CA^+ = \{A, C, B, D\}$$

$$CB^+ = \{C, B, D, A\}$$

$$CD^+ = \{C, D\}$$

Hence the candidate keys are:-

$$CK = \{\{C, A\}, \{C, B\}\}$$

∴ Is relation R in BCNF? Is it in 3NF?

	$A \rightarrow B$	$BC \rightarrow A$	$B \rightarrow D$
BCNF	X	✓	X
3NF	✓	✓	X
2NF	✓	✓	X
1NF	✓	✓	✓

$\{A \rightarrow B\}$ isn't in BCNF as $\{A\}$ is not a candidate key or a superkey.

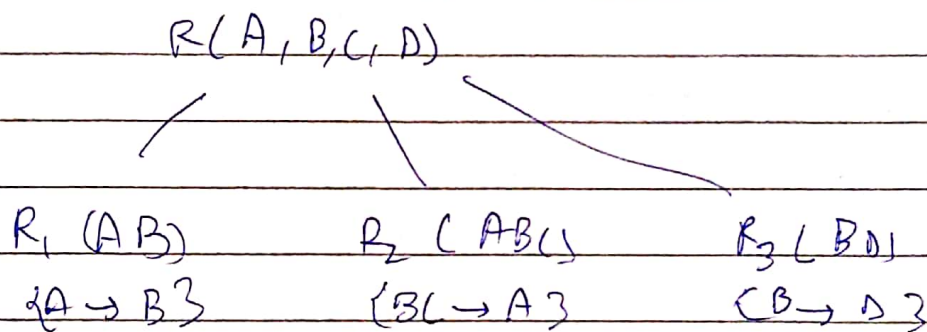
$\{B \rightarrow D\}$ isn't in BCNF as $\{B\}$ is not a candidate key or superkey



$\{B \rightarrow D\}$ isn't in 3NF normal form as $\{B\}$ isn't a candidate key or superkey and $\{D\}$ is a non-prime attribute and there exists a transitive dependency.

$\{A \rightarrow B\}$ and $\{B \rightarrow D\}$ of non-prime attribute is dependent on prime attribute $\{B\}$.

(c) Decompose into dependency-preserving lossless BCNF decomposition.



This preserves functional dependencies.

	A	B	C	D
R_1	α	α		
R_2	α	α	α	α'
R_3		α		α

→ Here this is also a lossless decomposition.

D) Decompose the relation $R(A, B, C, D)$ into collection of 3NF relations, so that decomposition is dependency-preserving lossless.

$R(A, B, C, D)$
 $R_1(A, B, C)$
 $\{A \rightarrow B, BC \rightarrow A\}$
 $R_2(B, D)$
 $\{B \rightarrow D\}$

 Candidate key = $\{B, C\}$

 Candidate key = $\{B\}$

They are both in 3NF form.

$\{A \rightarrow B\}$ doesn't have any transitive dependency.
 $\{BC \rightarrow A\}$ has ~~primary~~ candidate key as determinant
 $\{B \rightarrow D\}$ has candidate key as determinant

 The ~~de~~ decomposition covers all functional dependencies.

	A	B	C	D
R_1	α	α	α	α'_D
R_2		α		α

\rightarrow The decomposition is also lossless.