

Algorithm Family	Cryptosystems	Security Level (bit)			
		80	128	192	256
Integer factorization	RSA	1024 bit	3072 bit	7680 bit	15360 bit
Discrete logarithm	DH, DSA, Elgamal	1024 bit	3072 bit	7680 bit	15360 bit
Elliptic curves	ECDH, ECDSA	160 bit	256 bit	384 bit	512 bit
Symmetric-key	AES, 3DES	80 bit	128 bit	192 bit	256 bit

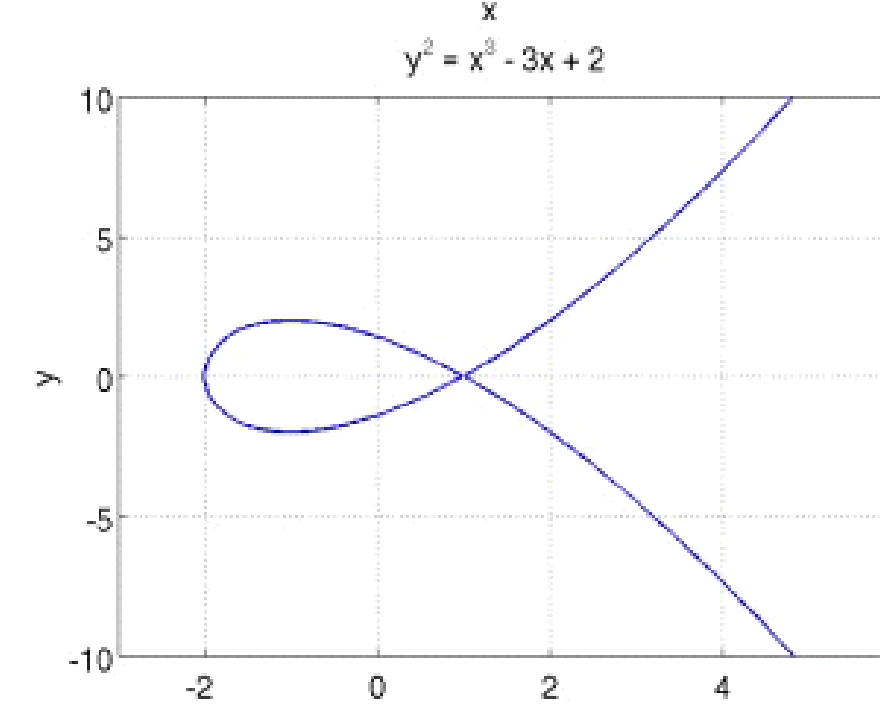
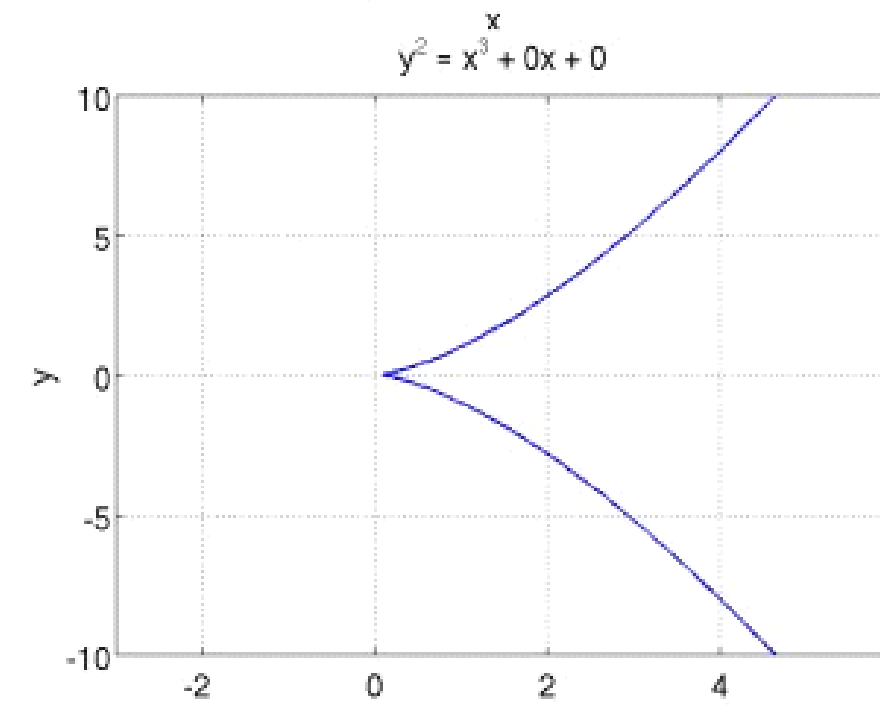
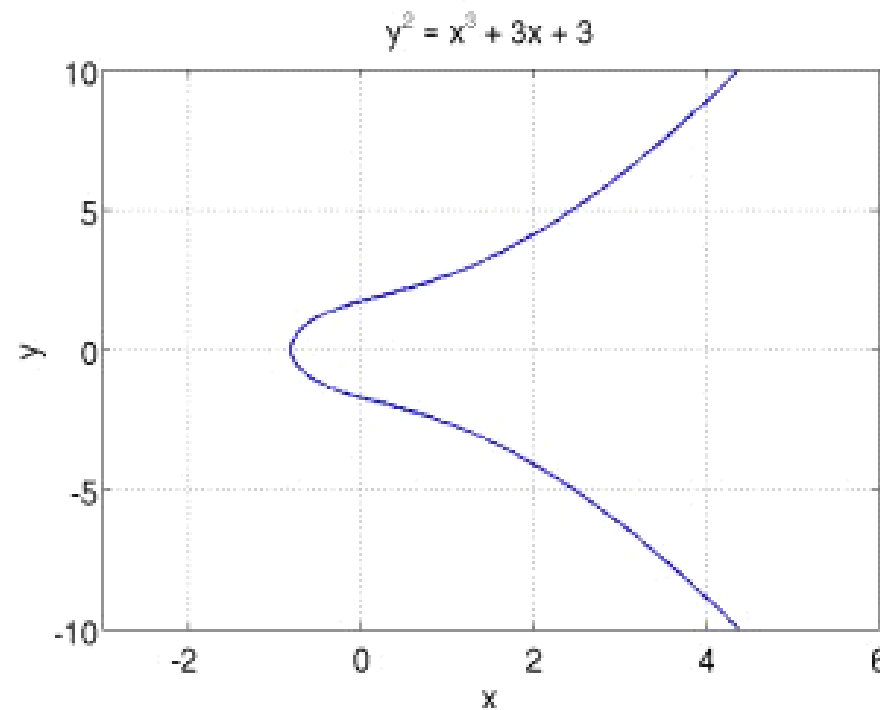
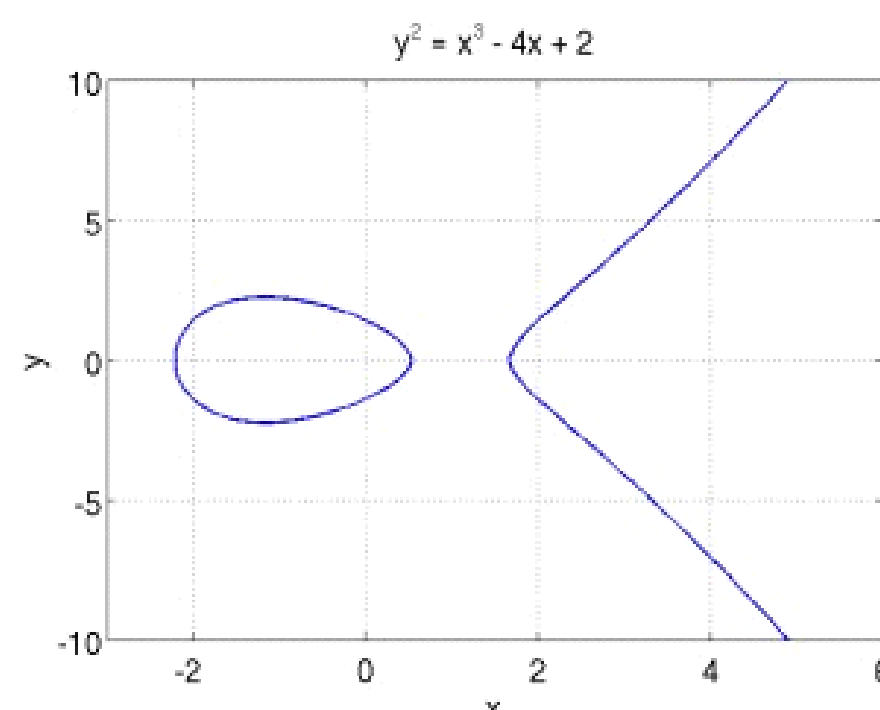
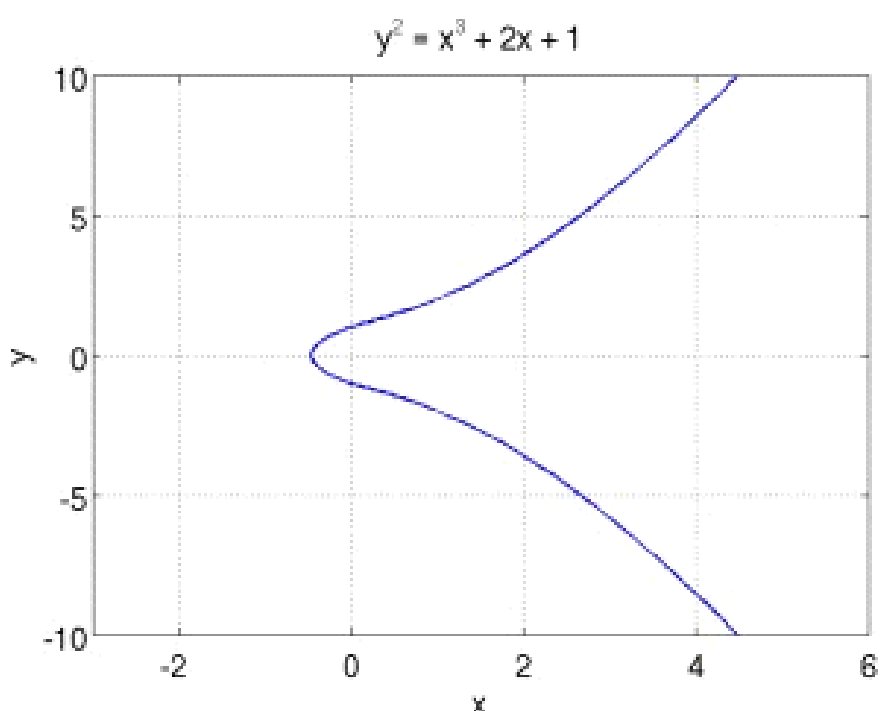
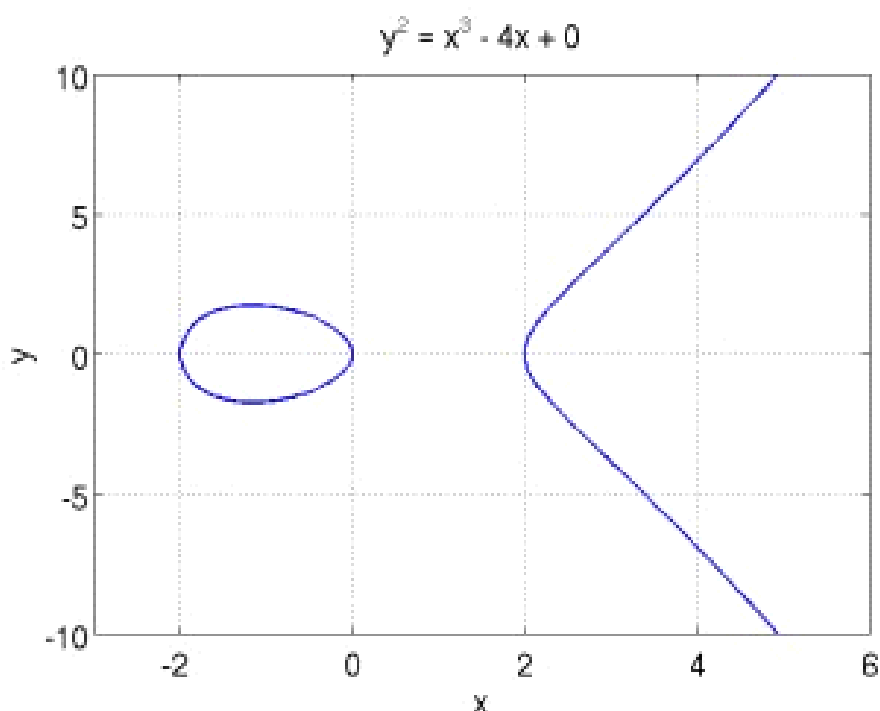
Elliptic Curves

General Eqⁿ: $y^2 + b_1xy + b_2y = x^3 + a_1x^2 + a_2x + a_3$

Elliptic curves on real no.

$y^2 = x^3 + ax + b$ — (1)
 a & b are real constants

(1) is also called Weierstrass equation of characteristic zero



Discriminant of a Polynomial of degree n : Product of squares of difference of polynomial roots-

$$D_n = \prod_{i < j}^n (r_i - r_j)^2, \quad r_i \text{'s are roots of the polynomial.}$$

In case of a polynomial of degree 3

$$D_3 = (r_1 - r_2)^2 (r_1 - r_3)^2 (r_2 - r_3)^2, \quad r_1, r_2, r_3 \text{ are roots of the poly}$$

If $D_n = 0$ then the poly. (P_n) is called singular and
" $D_n \neq 0$ " " " " " non-singular

\Rightarrow A polynomial P is non-singular if all of its roots are distinct.

In Cryptography we use ECs whose RHS is a non-singular polynomial

Note: Let $P = x^3 + ax + b$ then

$$D = -16(4a^3 + 27b^2)$$

$\Rightarrow P$ is non-singular if $4a^3 + 27b^2 \neq 0$

* EC are sym. about x -axis.

Defining a Group:

Set :

Let $E(a, b)$ denotes the set of all points on the curve

$$y^2 = x^3 + ax + b \quad \text{i.e.}$$

$$E(a, b) = \{(x, y) \in \mathbb{R} \times \mathbb{R} \mid y^2 = x^3 + ax + b\}$$

Operation We define an operation '+' on $E(a, b)$ as follows

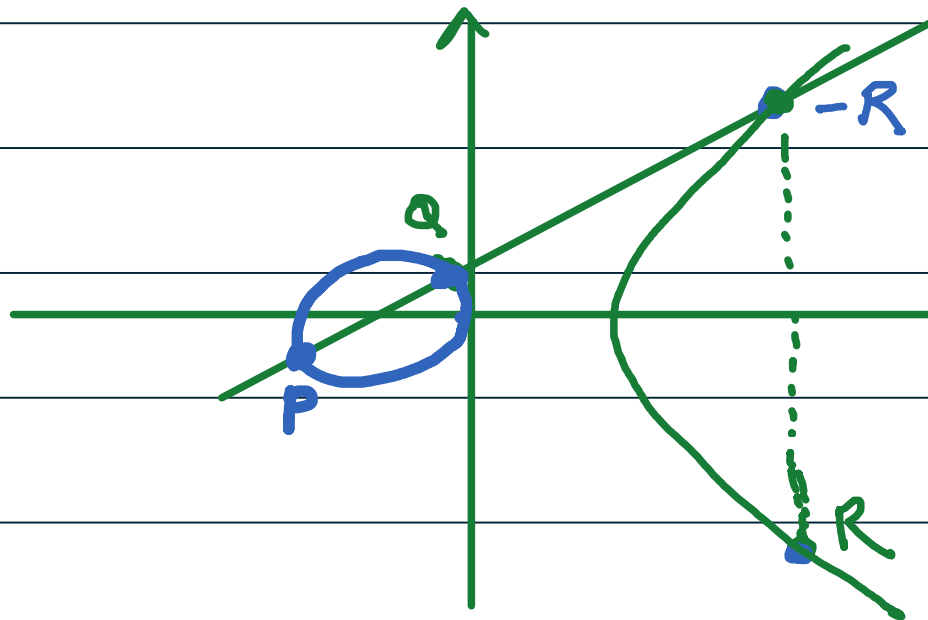
Let $P, Q, R \in E(a, b)$, $P = (x_1, y_1)$, $Q = (x_2, y_2)$,
 $R = (x_3, y_3)$

$$P + Q = R$$

(i) $x_1 \neq x_2$ & $x_2 \neq x_3$

$$R = (\lambda^2 - x_1 - x_2, \lambda(x_1 - x_3) - y_1)$$

where $\lambda = \frac{y_2 - y_1}{x_2 - x_1}$

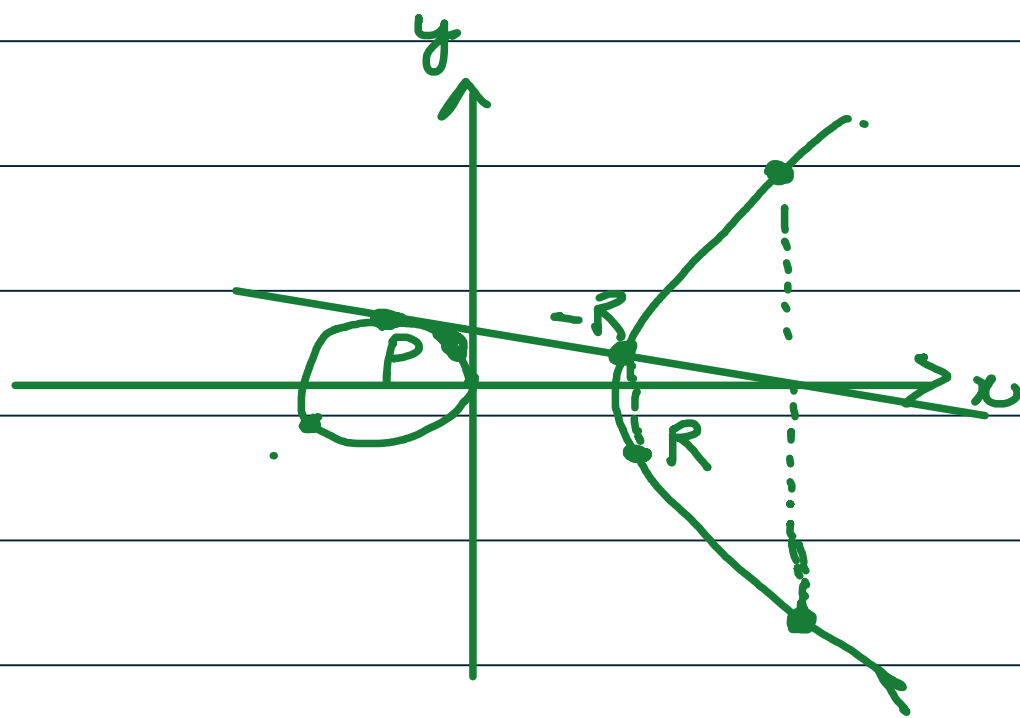


(ii) $P = Q$ i.e. $x_1 = x_2$, $y_1 = y_2$

$$P + P = R$$

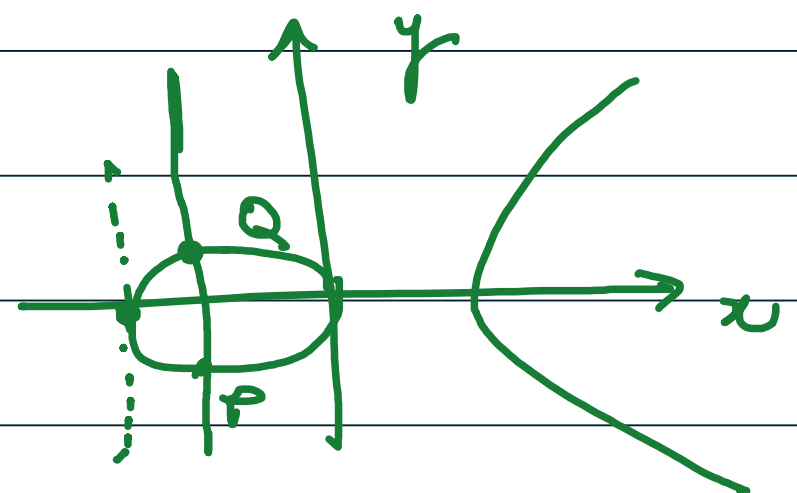
$$R = (\lambda^2 - x_1 - x_2, \lambda(x_1 - x_3) - y_1) \checkmark$$

$$\lambda = \frac{3x_1^2 + a}{2y_1} \checkmark$$



(iii) If $P = (x_1, y_1)$, $Q = (x_1, -y_1)$

Then, we define $R = O$
 \downarrow
 An imaginary point
 at inf.



(iv) If the tangent at a point P has the slope $\infty/2$ and
 $Q = P$ then

$$P + Q = \underline{P + P} = 2P = \underline{P}$$

Let $E(a, b)$ also includes 'O'.

Then $(E(a, b), +)$ is a group.

(i) closed

(ii) Assoc

(iii) 'O' is the identity

(iv) Inverse of $P = (x_1, y_1)$ $-P = (x_1, -y_1)$.

Elliptic curve over \mathbb{Z}_p ($p > 3$)

ECs over \mathbb{Z}_p ($p > 3$) is the set of all points $x, y \in \mathbb{Z}_p$
s.t. $y^2 = (x^3 + ax + b) \bmod p$

together with an imaginary point O at inf
where $a, b \in \mathbb{Z}_p$ & $4a^3 + 27b^2 \not\equiv 0 \bmod p$.

$$E_p(a, b) = \{ (x, y) \in \mathbb{Z}_p \times \mathbb{Z}_p \mid y^2 = x^3 + ax + b, a, b \in \mathbb{Z}_p \} \cup \{ \underline{O} \}$$

$(E_p(a, b), +)$ is a group where $+$ is define as follow.

$P, Q, R \in E_p(a, b)$ where $P = (x_1, y_1)$, $Q = (x_2, y_2)$, $R = (x_3, y_3)$
Then $P + Q = R = \left(\underbrace{(\lambda^2 - x_1 - x_2) \bmod p}_{x_3}, \underbrace{(\lambda(x_1 - x_3) - y_1) \bmod p}_{y_3} \right)$

$$\text{where } \lambda = \begin{cases} \frac{y_2 - y_1}{x_2 - x_1} \bmod p, & P \neq Q \\ \frac{3x_1^2 + a}{2y_1} \bmod p, & P = Q \end{cases}$$

and if $Q = (x_1, -y_1)$ then

$$P + Q = O$$

and if the slope of the tangent at P is ∞ then

$$P + P = \underline{P}.$$

Ex. $E_{23}(1, 1)$

$$4 \times 1^3 + 27 \times 1^2 = 31 \bmod 23 \neq 0$$

Let $P = (3, 10)$, $Q = (9, 7)$ then

$$\underline{P + Q ?} \quad \lambda = \frac{7 - 10}{9 - 3} \bmod 23 = \frac{-3}{6} \bmod 23 = \frac{-1}{2} \bmod 23$$

$$= (-1) \cdot 2^{-1} \bmod 23$$

$$= 11$$

$$x_3 = (\lambda^2 - x_1 - x_2) \bmod 23 = (11^2 - 3 - 9) \bmod 23 = 17$$
$$y_3 = 20 \quad R = (17, \underline{20})$$

$$y^2 = x^3 + 2x + 2 \pmod{17}$$

$$P = (5, 1)$$

$$P + P = \underline{(6, 3)}$$

$$2P = (5, 1) + (5, 1) = (6, 3)$$

$$3P = 2P + P = (10, 6)$$

$$4P = (3, 1)$$

$$5P = (9, 16)$$

$$6P = (16, 13)$$

$$7P = (0, 6)$$

$$8P = (13, 7)$$

$$9P = (7, 6)$$

$$10P = (7, 11)$$

$$11P = (13, 10)$$

$$12P = (0, 11)$$

$$13P = (16, 4)$$

$$14P = (9, 1)$$

$$15P = (3, 16)$$

$$16P = (10, 11)$$

$$17P = (6, 14)$$

$$18P = (5, 16)$$

$$19P = \mathcal{O}$$

No. of points N in $E_p(a, b)$ is bounded by

$$p+1 - 2\sqrt{p} \leq N \leq p+1 + 2\sqrt{p}$$

$$\text{Thus, } |E_p(a, b)| \approx |\mathbb{Z}_p|$$

Thus, $E_p(a, b)$ has app p elements.