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Assignment 4  
CIS 415

1. OSC 10.11 (Currently at 2150; Previously 1805; Ranges from 0-4999;  
FIFO: 2069, 1212, 2296, 2800, 544, 1618, 356, 1523, 4965, 3681)
  - a. FCFS  
Order: 2150, 2069, 1212, 2296, 2800, 544, 1618, 356, 1523, 4965, 3681  
Distance:  $(2150-2069) + (2069-1212) + (2296-1212) + (2800-2296) + (2800-544) + (1618-544) + (1618-356) + (1523-356) + (4965-1523) + (4965-3681) = \mathbf{13011 \text{ cylinders}}$
  - b. SSTF  
Order: 2150, 2069, 2296, 2800, 3681, 4965, 1618, 1523, 1212, 544, 356  
Distance:  $(2150-2069) + (2296-2069) + (2800-2296) + (3681-2800) + (4965-3681) + (4965-1618) + (1618-1523) + (1523-1212) + (1212-544) + (544-356) = \mathbf{7586 \text{ cylinders}}$
  - c. SCAN  
Order: 2150, 2296, 2800, 3681, 4965, 4999, 2069, 1618, 1523, 1212, 544, 356, 0  
Distance:  $(2296-2150) + (2800-2296) + (3681-2800) + (4965-3681) + (4999-4965) + (4999-2069) + (2069-1618) + (1618-1523) + (1523-1212) + (1212-544) + (544-356) = \mathbf{7492 \text{ cylinders}}$
  - d. LOOK  
Order: 2150, 2296, 2800, 3681, 4965, 2069, 1618, 1523, 1212, 544, 356  
Distance:  $(2296-2150) + (2800-2296) + (3681-2800) + (4965-3681) + (4965-2069) + (2069-1618) + (1618-1523) + (1523-1212) + (1212-544) + (544-356) = \mathbf{7424 \text{ cylinders}}$
  - e. C-SCAN  
Order: 2150, 2296, 2800, 3681, 4965, 4999, 0, 356, 544, 1212, 1523, 1618, 2069  
Distance:  $(2296-2150) + (2800-2296) + (3681-2800) + (4965-3681) + (4999-4965) + (356-0) + (544-356) + (1212-544) + (1523-1212) + (1618-1523) + (2069-1618) = \mathbf{4918 \text{ cylinders}}$
  - f. C-LOOK  
Order: 2150, 2296, 2800, 3681, 4965, 356, 544, 1212, 1523, 1618, 2069  
Distance:  $(2296-2150) + (2800-2296) + (3681-2800) + (4965-3681) + (544-356) + (1212-544) + (1523-1212) + (1618-1523) + (2069-1618) = \mathbf{4528 \text{ cylinders}}$

2. Consider a file system that uses inodes to represent files.
  - a. Max size of disk (in bytes) for which one can use this file system?  
The inode layout does not impact the max size of the disk; it only impacts the max size of the file.
  - b. What is the max size of file (in bytes) that can be stored in this file system?  
 8 direct block pointers  $\rightarrow 8 \times 1024 = 8\text{kB}$   
 1 single indirect  $\rightarrow (1024/4) \times 1024 = 256\text{kB}$   
 1 double indirect  $\rightarrow (1024/4) \times (1024/4) \times 1024 = 2^6 \text{kB} = 2^6 \text{MB} = 64\text{MB}$   
 1 triple indirect  $\rightarrow (1024/4) \times (1024/4) \times (1024/4) \times 1024 = 2^{24} \text{kB} = 16\text{GB}$   
 total:  $8 \times 1024 + (1024/4) \times 1024 + (1024/4) \times (1024/4) \times 1024 + (1024/4) \times (1024/4) \times (1024/4) \times 1024 = \mathbf{17247248384 \text{ bytes}}$
3. The processor for which you are designing your application as L1i and L1d virtual caches
  - a. Type of data does cache hold?  
 L1i  $\rightarrow$  holds instruction cache  
 L1d  $\rightarrow$  holds data cache
  - b. Describe in detail the activities of the cache + memory system when executing the instruction  
 To get the contents of "virtual address", CPU will check the cache first. If not in cache, the OS will need to the corresponding page table to find out the physical frame needed. If the page table entry is not in cache, CPU will access memory to read the page table entry to find the physical frame number. If the physical frame is not in memory, OS will page-fault to read the memory frame, and then fetch the contents of "virtual address" – this data will also be cached for future references.
  - c. Assume that the above instruction is executed many times in a loop, and that the instruction itself is in the cache. Also assume that memory access costs  $\tau$   $\mu\text{s}$ , and cache access costs  $\tau/15$   $\mu\text{s}$ . What cache hit rate  $\rho$  for "virtual address" is required for the memory system to run 5 times faster than with no caching at all? Show your work.  

$$(\rho * \tau/15) + ((1 - \rho) * \tau) = \tau/5$$

$$5(\rho/15) + 5(1 - \rho) = 1$$

$$\rho/3 + 5 - 5\rho = 1$$

$$\rho + 15 - 15\rho = 3$$

$$12 = 14\rho$$

$$\rho = .857 = \mathbf{85.7\%}$$
  - d. Suppose we have a memory system that has a main memory, a single-level cache, and demand paging virtual memory. The three levels of the memory system have the following access times
    - i. effective memory access time:  $(0.95 \times 2) + (0.05 \times 100) = \mathbf{6.9\text{ns}}$

- ii. new effective memory access time:  $(0.95 \times 2) + (.05 \times .00001 \times (100 + 10 \times 10^6)) + (.05 \times 0.9999 \times 100) = \mathbf{11.9 \text{ ns}}$