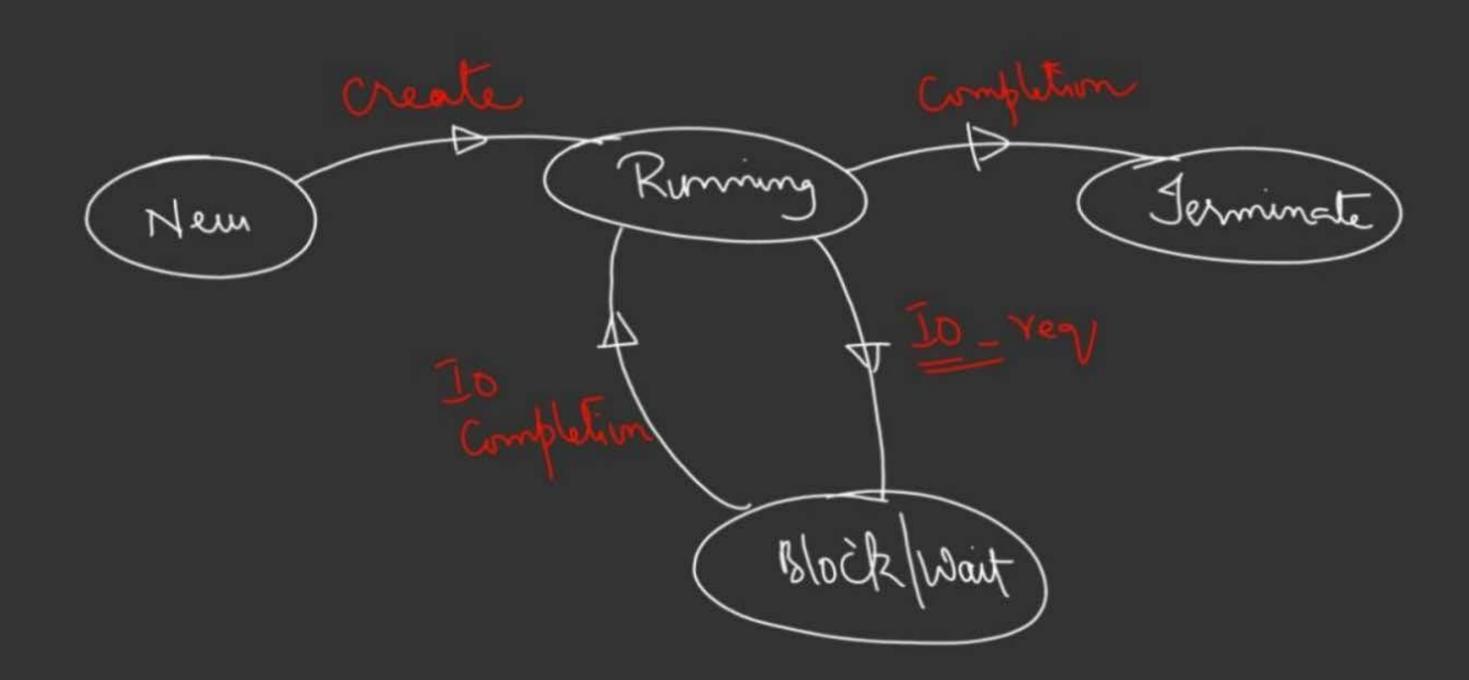
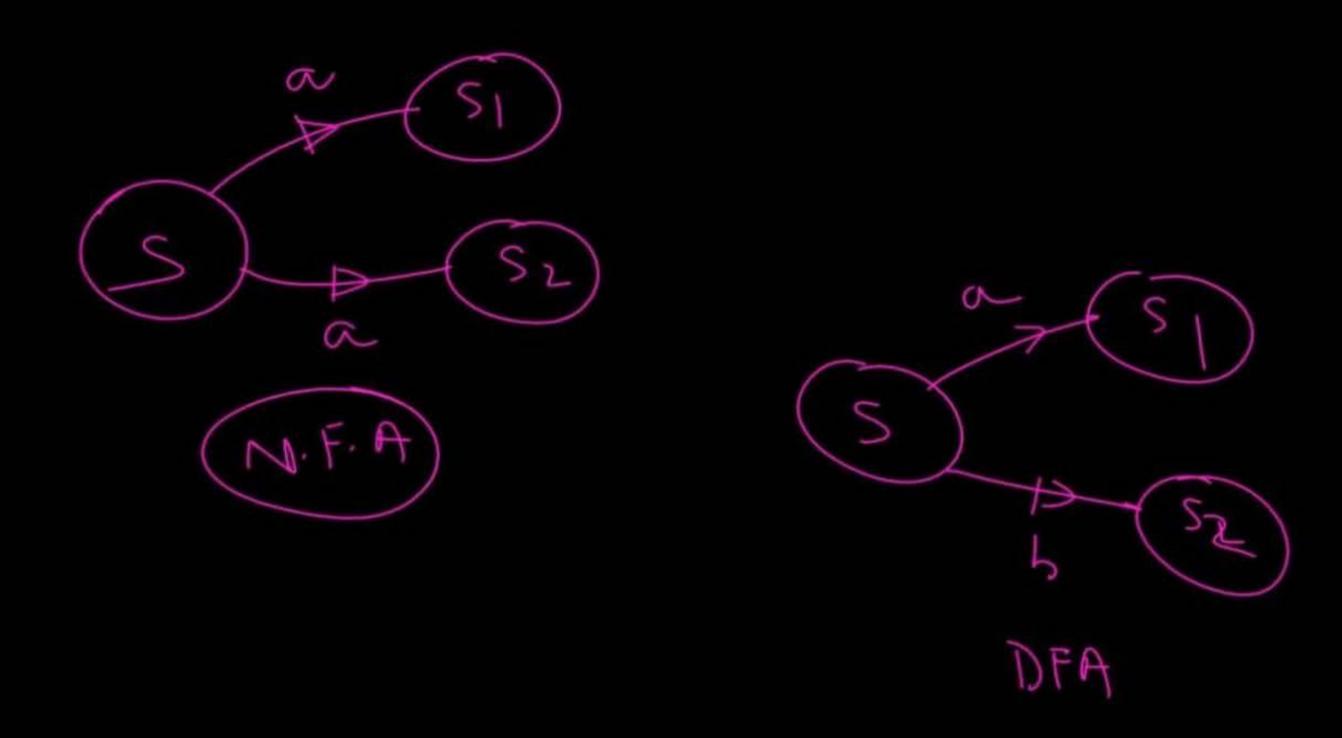
Process States New; Ready; Running; Block
- Suspend Wait + Jalsc 20callo C Suspend Ready Suspend Create Schedule Completion Dispatch Ready New Running Terminate Fys-Gh To Completion Dres-Rre Finite Autombe Sys-Call Completion Block want Surpend Block

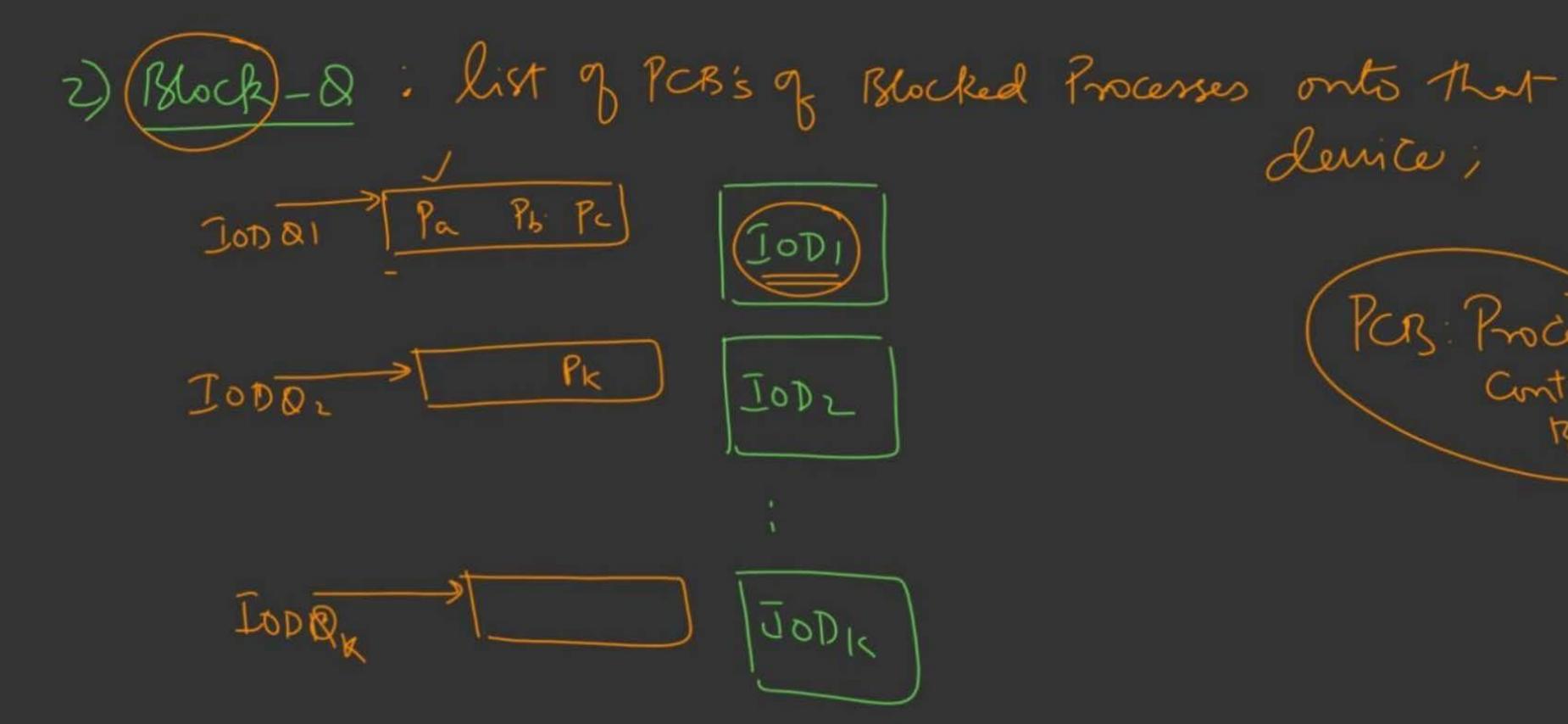
-> When the Process is in Ready, Running, Block states
-> When the Process is in Ready, Running, Block states then it is in Main Memory
- There can be many Ready Processes,
many Block Robers and one Running
Process (for one cpu) processes may get souspended from Memory onto
-> processes may get surpended from Memory onto
sisk for performance reason;
If a process has to complete IO/Back in Surpendrytock, him
If a process has to complete IO Block in Surpendistock, then Jempornily, it is brought in Memory, Satisfy the request of
again surfended I a ready Process is prempted of its resource
again surfended _ If a ready Process is prempted of its resource then it gets Blocked;

Uni-programmed os State Diagram





Scheduling Queues & State-Quewing diagram 3) Imput - Q Job - Q on-018/2/ In- Memory 4) Suspend - Q 2) Block-Q P5 PK--1 > Contains list of PCB's of Ready Processes Pusi



Mode

derrice;

m-Disk-Q's Input -Q Job-Q: Programs that are ready to be Loaded in Memory; P1 12 83 sisk

whend Disk

Processes that get Surfended from Memory are Stored in Surfend-Q

