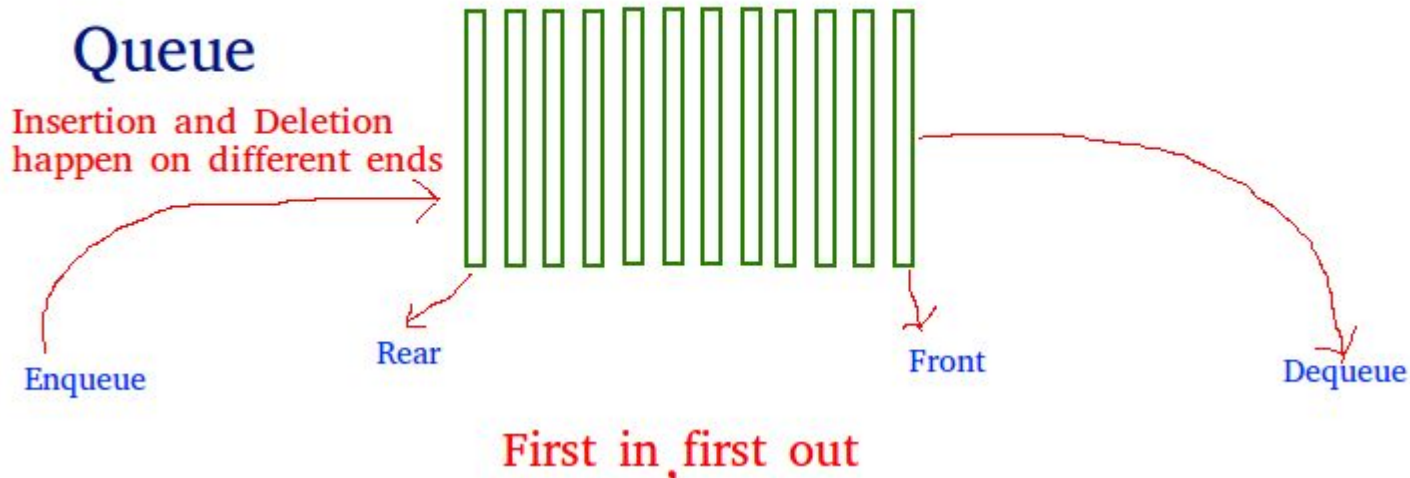


# DSA Queue

## Chapter 3

# Queue

- It is data structure that follows operations on FIFO (First In First Out) order.
- One of the real life example of Queue as it name implies is Queue in Cinema Hall or ticket Counter.



# Basic Operation on Queue

- Enqueue: Insertion of element in a queue

It is similar with the push operation in stack. Enqueue operation is performed from one end i.e. Rear

- Dequeue: Deletion of element from a queue

It is similar with the pop operation in stack. Dequeue operation is performed from another end i.e. front

# Operations on Queue

- Enqueue(n)
- Dequeue()
- front()
- isFull()
- isEmpty()

# Queue as an ADT

```
Class Queue{  
    Int queue[];  
    Int front=-1;  
    Int rear=-1;  
    enqueue();  
    dequeue();  
    display();  
}
```

# Algorithm for enqueue

- **Step 1:** IF  $REAR = MAX - 1$   
    Write Queue is full  
    Go to step  
    [END OF IF]
- **Step 2:** IF  $FRONT = -1$  and  $REAR = -1$   
    SET  $FRONT = REAR = 0$   
    ELSE  
        SET  $REAR = REAR + 1$   
    [END OF IF]
- **Step 3:** Set  $QUEUE[REAR] = NUM$
- **Step 4:** EXIT

# Algorithm for dequeue

- **Step 1:** IF  $\text{FRONT} = -1$  or  $\text{FRONT} > \text{REAR}$   
Write Queue is empty  
ELSE IF :  $\text{FRONT} == \text{REAR}$ ,  $\text{FRONT} = \text{REAR} = -1$   
  
SET  $\text{VAL} = \text{QUEUE}[\text{FRONT}]$   
SET  $\text{FRONT} = \text{FRONT} + 1$   
[END OF IF]
- **Step 2:** EXIT

# Advantages of Queue

- A large amount of data can be managed efficiently with ease.
- Operations such as insertion and deletion can be performed with ease as it follows the first in first out rule.
- Queues are useful when a particular service is used by multiple consumers.
- Queues are fast in speed for data inter-process communication.
- Queues can be used in the implementation of other data structures.



# Disadvantages of Queue

- The operations such as insertion and deletion of elements from the middle are time consuming.
- Limited Space.
- In a classical queue, a new element can only be inserted when the existing elements are deleted from the queue.
- Searching an element takes  $O(N)$  time.
- Maximum size of a queue must be defined prior.

# Application of Queue

- **Multi programming:** Multi programming means when multiple programs are running in the main memory. It is essential to organize these multiple programs and these multiple programs are organized as queues.
- **Network:** In a network, a queue is used in devices such as a router or a switch. another application of a queue is a mail queue which is a directory that stores data and controls files for mail messages.
- **Job Scheduling:** The computer has a task to execute a particular number of jobs that are scheduled to be executed one after another. These jobs are assigned to the processor one by one which is organized using a queue.
- **Shared resources:** Queues are used as waiting lists for a single shared resource.

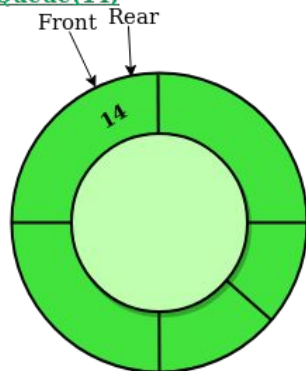
## **Real-time application of Queue:**

- ATM Booth Line
- Ticket Counter Line
- Key press sequence on the keyboard
- CPU task scheduling
- Waiting time of each customer at call centers.

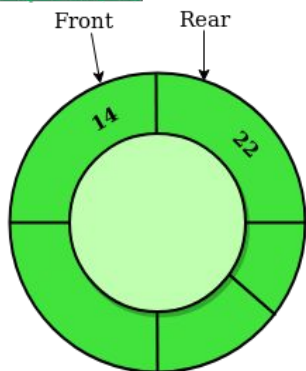
# Circular Queue

Circular Queue is a linear data structure in which the operations are performed based on FIFO (First In First Out) principle and the last position is connected back to the first position to make a circle.

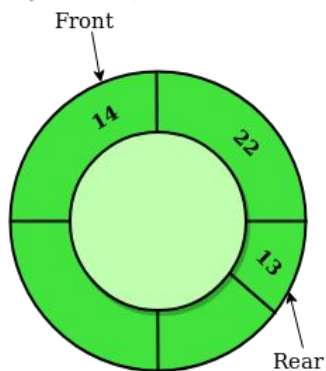
enQueue(14)



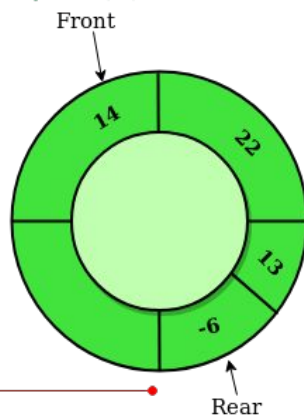
enQueue(22)



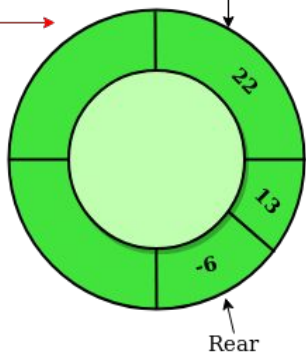
enQueue(13)



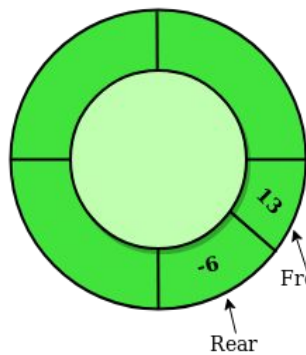
enQueue(-6)



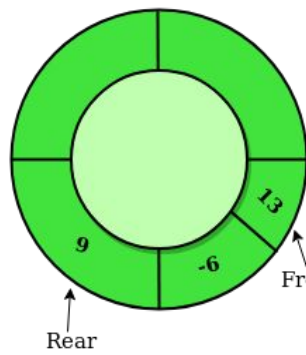
Front



deQueue()

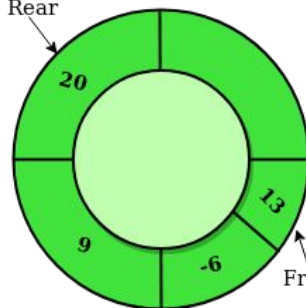


deQueue()

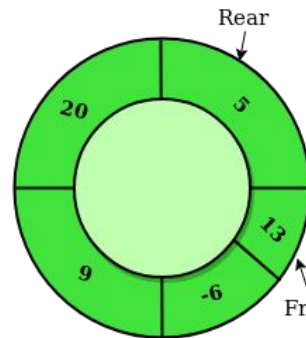


enQueue(9)

Rear



enQueue(20)



enQueue(5)

# Circular Enqueue

1. Check whether queue is Full – Check  $((\text{rear} == \text{SIZE}-1 \ \&\& \ \text{front} == 0) \ || \ (\text{rear} == \text{front}-1))$ .
2. If it is full then display Queue is full. If queue is not full then, check if  $(\text{rear} == \text{SIZE} - 1 \ \&\& \ \text{front} != 0)$  if it is true then set  $\text{rear}=0$  and insert element.

# Circular Enqueue

```
Void Enqueue(int x)
{
    if(front==-1 && rear==-1)
        {
            front=rear=0;
            queue[rear]=X;
        }
    elseif((rear+1)%N==front)
        {
            printf("Circular Queue is full");
        }
    else{
        rear=(rear+1)%N;
        queue[rear]=X;
    }
}
```

# Circular Dequeue

1. Check whether queue is Empty means check ( $\text{front} == -1$ ).
2. If it is empty then display Queue is empty. If queue is not empty then step 3
3. Check if ( $\text{front} == \text{rear}$ ) if it is true then set  $\text{front} = \text{rear} = -1$  else check if ( $\text{front} == \text{size} - 1$ ), if it is true then set  $\text{front} = 0$  and return the element.



# Circular Dequeue

```
Void dequeue(){  
    if(front== -1 && rear == -1)  
        printf("Circular queue is empty");  
    elseif(front==rear)  
        front=rear=-1;  
    else  
        front=(front+1)%N;  
}
```

### **Applications of Circular Queue:**

- In a page replacement algorithm, a circular list of pages is maintained and when a page needs to be replaced, the page in the front of the queue will be chosen.
- Computer systems supply a holding area for maintaining communication between two processes or two programs. This memory area is also known as a ring buffer.
- CPU Scheduling: In the Round-Robin scheduling algorithm, a circular queue is utilized to maintain processes that are in a ready state.

### **Real-time Applications of Circular Queue:**

- Months in a year: Jan – Feb – March – and so on upto Dec- Jan – . . .
- Eating: Breakfast – lunch – snacks – dinner – breakfast – . . .
- Traffic Light is also a real-time application of circular queue.

### **Advantages of Circular Queue:**

- It provides a quick way to store FIFO data with a maximum size.
- Efficient utilization of the memory.
- Doesn't use dynamic memory.
- Simple implementation.
- All operations occur in  $O(1)$  constant time.

### **Disadvantages of Circular Queue:**

- In a circular queue, the number of elements you can store is only as much as the queue length, you have to know the maximum size beforehand.

# Priority Queue

Priority Queue is an extension of the Queue data structure where each element has a particular priority associated with it.

It is based on the priority value, the elements from the queue are deleted.

For our program, we are applying a procedure that the higher the value of the number, higher the priority

# Priority Enqueue

```
void insert_by_priority(int data)
{
    if (rear >= MAX - 1) {
        printf("\nQueue overflow no more elements can be inserted");
        return;
    }
    if ((front == -1) && (rear == -1)) {
        front++;
        rear++;
        pri_que[rear] = data;
        return;
    }
    else
        check(data);
    rear++;
}
```

# Priority Dequeue

```
void delete_by_priority(int data){  
    int i; // Code to check whether queue is empty or not.  
    for (i = 0; i <= rear; i++)    {  
        if (data == pri_que[i])    {  
            for (; i < rear; i++)    {  
                pri_que[i] = pri_que[i + 1];    }  
            pri_que[i] = -99;  
            rear--;  
            if (rear == -1)    front = -1;  
            return;    }  
    }  
    printf("\n%d not found in queue to delete", data); }
```

# Double Ended Queue

The deque stands for Double Ended Queue.

Deque is a linear data structure where the insertion and deletion operations are performed from both ends.

We can say that deque is a generalized version of the queue.

Though the insertion and deletion in a deque can be performed on both ends, it does not follow the FIFO rule. The representation of a deque is given as follows -



**Representation of deque**

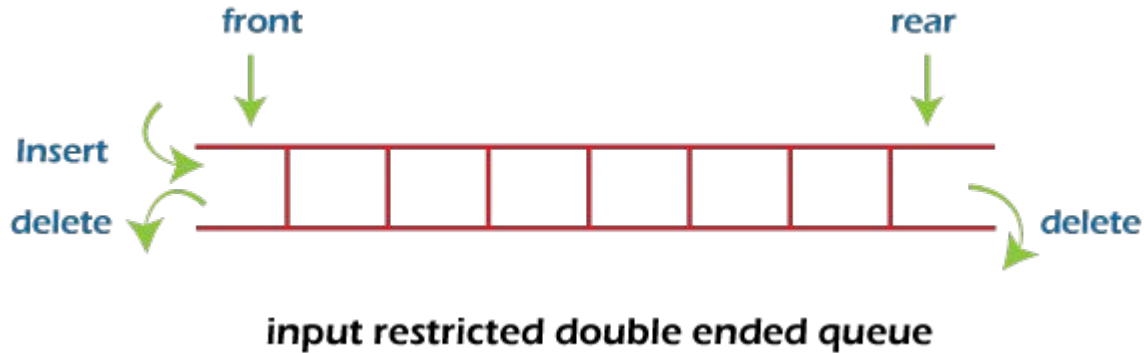
# Types of Double ended Queue

1. Input Restricted
2. Output Restricted



# Input Restricted

In input restricted queue, insertion operation can be performed at only one end, while deletion can be performed from both ends.



# Output Restricted Deque

In output restricted queue, deletion operation can be performed at only one end, while insertion can be performed from both ends.



# Operations on Deque

1. Insert at front
2. Delete at front
3. Insert at rear
4. Delete at rear

# Pseudo Code for insert at front

```
void insert_front(int x) {  
    if((f==0 && r==size-1) || (f==r+1)) {  
        printf("Overflow");  
    }  
    else if((f==-1) && (r==-1)) {  
        f=r=0;  
        deque[f]=x;  
    }  
}
```

```
else if(f==0) {  
    f=size-1;  
    deque[f]=x;  
}  
else {  
    f=f-1;  
    deque[f]=x;  
} }
```

# Pseudo Code for insert at rear

```
void insert_rear(int x)
```

```
{
    if((f==0 && r==size-1) || (f==r+1))
    {
        printf("Overflow");
    }
    else if((f==size-1) && (r==0))
    {
        r=0;
        deque[r]=x;
    }
    else if(r==size-1)
    {
        r=0;
        deque[r]=x;
    }
    else
    {
        r++;
        deque[r]=x;
    }
}
```

# Pseudo Code for insert at rear

```
void delete_front() {  
    if((f==-1) && (r==-1)) {  
        Display "Deque is empty" ;  
    }  
    else if(f==r)  
    {  
        Display "\nThe deleted element is  
deque[f];  
        f=-1;  
        r=-1;  
    }  
    else if(f==(size-1))  
    {  
        Display "The deleted element is  
deque[f];  
        f=0;  
    }  
    else  
    {  
        Display "The deleted element is  
deque[f];  
        f=f+1;  
    }  
}
```

# Pseudo Code for delete at front

```
void delete_front() {  
    if((f==-1) && (r==-1)) {  
        Display "Deque is empty" ;  
    }  
    else if(f==r)  
    {  
        Display "\nThe deleted element is  
deque[f];  
        f=-1;  
        r=-1;  
    }  
    else if(f==(size-1))  
    {  
        Display "The deleted element is  
deque[f];  
        f=0;  
    }  
    else  
    {  
        Display "The deleted element is  
deque[f];  
        f=f+1;  
    }  
}
```

# Pseudo Code for delete at rear

```
void delete_rear() {
    if((f==-1) && (r==-1)) {
        printf("Deque is empty");
    }
    else if(f==r) {
        printf("\nThe deleted element is %d", deque[r]);
        f=-1;
        r=-1;
    }
    else if(r==0)
    {
        printf("\nThe deleted element is %d", deque[r]);
        r=size-1;
    }
    else
    {
        printf("\nThe deleted element is %d", deque[r]);
        r=r-1;
    }
}
```



