#### System Programming

#### **Abhijit Mane**

abhipucsd.123@gmail.com 8796194950

#### **Contents**

- General Compiler infra-structure
- An Introduction to LEX and YACC
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  - Regular Expressions
  - Lex A lexical analyzer
  - Yacc Yet another compiler compiler
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#### Why Study Compilers?

- To enhance understanding of programming languages
- To have an in-depths knowledge of low-level machine executables
- To write compilers and interpreters for various programming languages and domain-specific languages
  - Examples: Java, JavaScript, C, C++, C#, Modula-3, Scheme, ML, Tcl/Tk,
     Database Query Lang., Mathematica, Matlab, Shell-Command-Languages, Awk, Perl, your .mailrc file, HTML, TeX, PostScript, Kermit scripts, .....
- To learn various system-building tools: Lex, Yacc, ...
- To learn interesting compiler theory and algorithms.
- To learn the beauty of programming in modern programming lang.

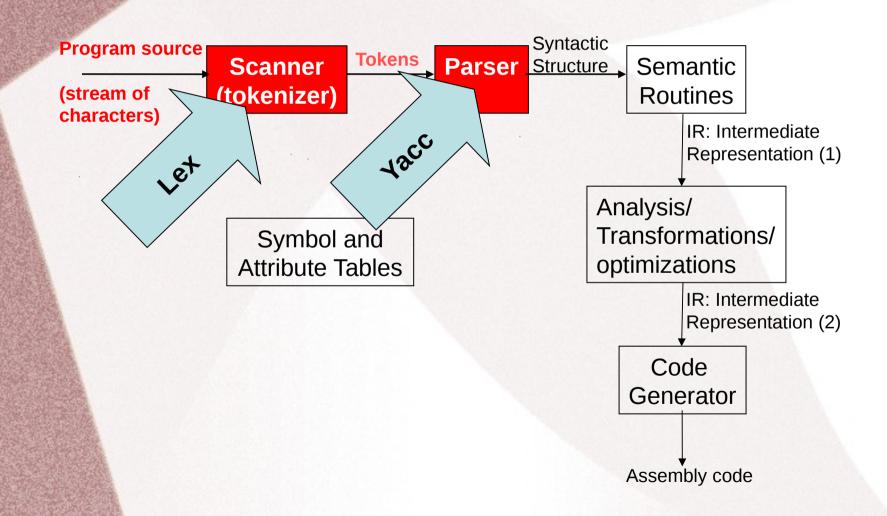
### Compilers are Translators



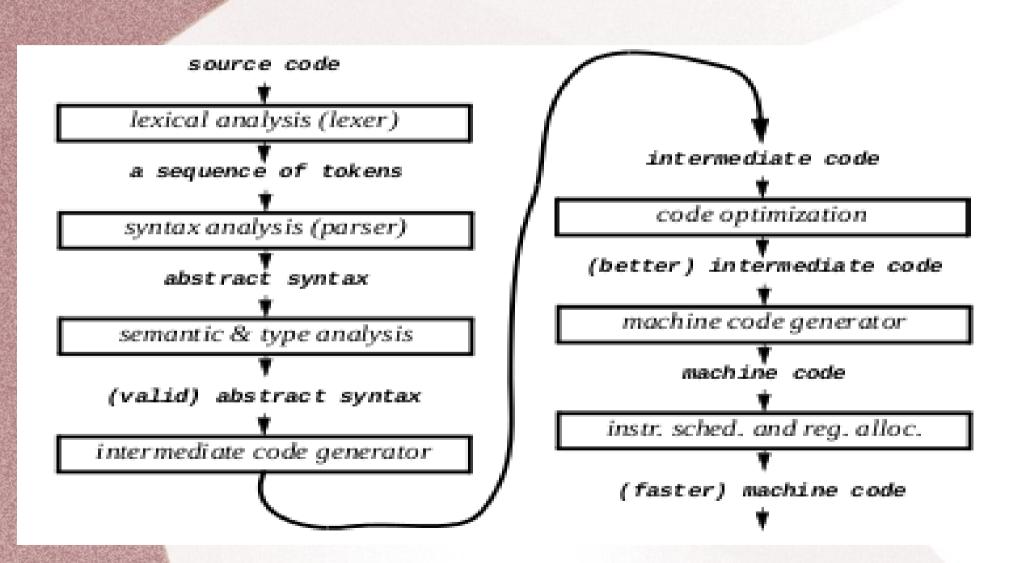
#### **Various forms of translators**

L	L'	translator
C++, ML, Java	assembly/machine code	compiler
assembly lang.	machine code	assembler
"object" code (*.o file)	"executable" code (a.out)	linker/loader
macros/text	text	macro processor (cpp)
troff/Tex/HTML	PostScript	document formatter
any file (e.g., foo)	compressed file (foo.Z)	file compresser

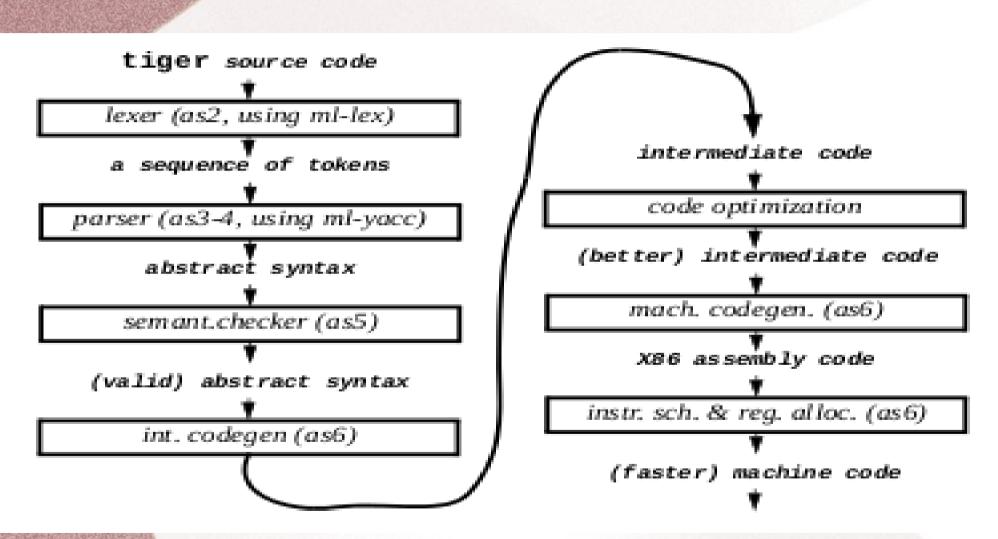
## General Compiler infra-structure



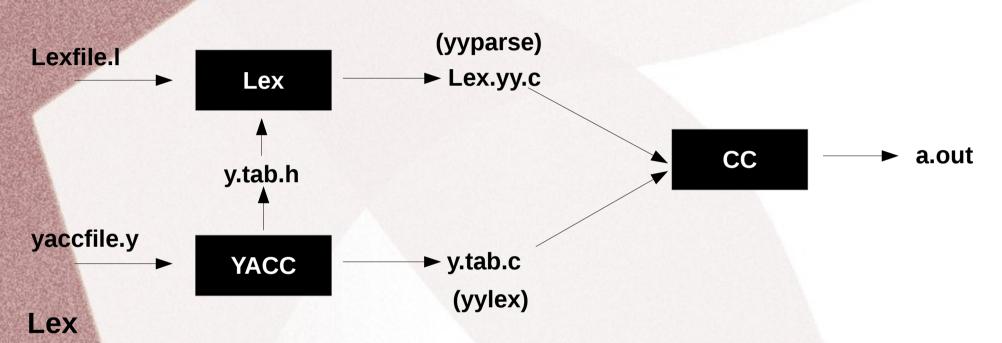
#### **Compilation Phases**



# What happens internally?



#### **General Structure**



generates C code for the lexical analyzer (scanner)

Token patterns specified by regular expressions

#### Yacc

generates C code for a LR(1) syntax analyzer (parser) BNF rules for the grammar

#### Lex: A Scanner Generator

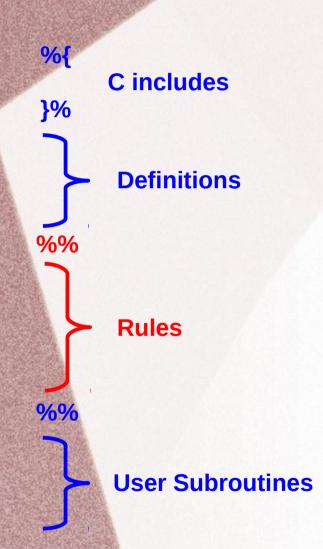
 Helps write programs whose control flow is directed by instances of regular expressions in the input stream.

#### yylex():

- matches the input
   stream against the
   table of regular
   expressions supplied
- carries out the
   associated action
   when a match is
   found.

Table of regular expressions + associated actions LEX yylex() (in file lex.yy.c)

## Structure of Lex Specification File



Red : required Blue : optional

Rules: line oriented: <reg.exp> <whitespace> <action>

<reg . exp > : starts at beginning of line, continues upto
first un-escaped whitespace
<action> : a single C statement (multiple statements:
enclose in braces { }).

unmatched input characters: copied to stdout.

## Lex predefined variables and functions

Name	Function
int yylex(void)	call to invoke lexer, returns token
char *yytext	pointer to matched string
yyleng	length of matched string
yylval	value associated with token
int yywrap(void)	wrapup, return 1 if done, 0 if not done
FILE *yyout	output file
FILE *yyin	input file
INITIAL	initial start condition
BEGIN	condition switch start condition
ECHO	write matched string

### Lex – Pattern Matching Primitives

Metacharacter	Matches
•	any character except newline
\n	newline
*	zero or more copies of the preceding expression
+	one or more copies of the preceding expression
?	zero or one copy of the preceding expression
^	beginning of line
\$	end of line
a b	a or b
(ab)+	one or more copies of ab (grouping)
"a+b"	literal "a+b" (C escapes still work)
[]	character class

## Lex – Pattern Matching Examples

Expression	Matches
abc	abc
abc*	ab abc abcc
abc+	abc abcc
a (bc) +	abc abcbc abcbcbc
a (bc) ?	a abc
[abc]	one of: a, b, c
[a-z]	any letter, a-z
[a\-z]	one of: a, -, z
[-az]	one of: -, a, z
[A-Za-z0-9]+	one or more alphanumeric characters
[ \t\n]+	whitespace
[^ab]	anything except: a, b
[a^b]	one of: a, ^, b
[a b]	one of: a, I, b
a b	one of: a, b

#### Regular Expressions

- What ,Where and How ?
- Write a regular expression that generates each of the following languages over the alphabet  $\Sigma = \{0, 1\}$ . In each case, explain how your answer works.
  - a)  $\{x \in \Sigma * \mid x \text{ begins with a 0 and ends with a 1} \}$
  - b)  $\{x \in \Sigma * | |x| > 3\}$
  - c)  $\{x \in \Sigma * | |x| \text{ is an even integer } \}$
  - d)  $\{x \in \Sigma * | x \text{ contains at least one of the substrings 000 or 111} \}$
  - e)  $\{x \in \Sigma * \mid x \text{ contains both of the substrings 000 and 111} \}$
  - f) Find all patterns that has at least one but no more than 3, 'a's

#### RE Examples

- Write the grep commands for each of the following tasks
  - a) Find all patterns that matches the pattern "ted" or "fred"
  - b) Find all patterns that matches ed, ted or fed
  - c) Find all patterns that does not begin with "g"
  - d) Find all patterns that begins with g or any digit from 0-9
  - e) Find all patterns that begins with "pucsd"
  - f) Find lines in a file where the pattern "sam" occurs at least twice
  - g) Find all lines in a file that contain email addresses
- Write a regex that matches any number between 1000 and 9999
- Write a regex that matches any number between 100 and 9999
- Write a regex that lists all the files in the current directory that was created in Nov and are txt files.

#### Lex program

```
%{
... c includes ...
%}
... definitions ...
%%
... rules ...
%%
... subroutines ...
```

```
%{
#include <stdio.h>
#include "y.tab.h"
int c;
extern int yylval;
%}
%%
. .
[a-z] \{ c = yytext[0]; yylval = c - 'a';
return(LETTER); }
[0-9]* { yylval = atoi(yytext);
return(NUMBER); }
[^a-z0-9\b] { c = yytext[0]; return(c); }
```

#### **Examples of Lex Rules**

- int printf("keyword: INTEGER\n");
- [0-9]+ printf("number\n");
- "-"?[0-9]+("."[0-9]+)? printf("number\n");

## Choosing between different possible matches:

- When more than one pattern can match the input, lex chooses as follows:
  - The longest match is preferred.
  - Among rules that match the same number of characters,
     the rule that occurs earliest in the list is preferred.

## Communicating with the user program

yytext: a character array that contains the actual string that matched a pattern.

yyleng: the no. of characters matched.

#### **Example:**

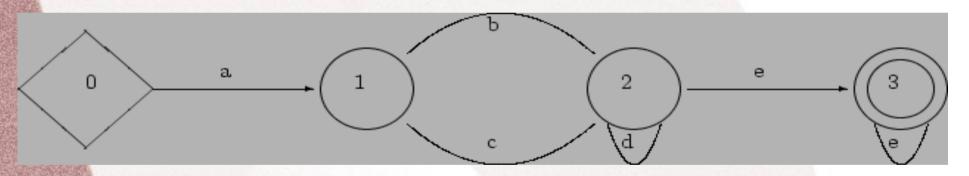
- [a-z][a-z0-9\_]\* printf("ident: %s\n", yytext);
- Counting the number of words in a file and their total size:

```
- [a-zA-Z]+ \{nwords += 1; size += yyleng;\}
```

#### How lexers work?

#### Lexers - Finite State Automata (FSA)

 Lexers are also known as scanners. LEX converts each set of regular expressions into a Deterministic FSA (DFSA) e.g. for a(b|c)d\*e+



which has states 0 to 3, where state 0 is the initial state and state 3 is an accept state that indicates a possible end of the pattern.

#### General Algorithm

```
state= 0; get next input character
     while (not end of input) {
       depending on current state and input character
         match: /* input expected */
          calculate new state; get next input character
         accept: /* current pattern completely matched */
          perform action corresponding to pattern; state= 0
         error: /* input unexpected */
          reset input; report error; state= 0
```

#### Pattern Matching and Action

Match a character in the a-z range



Buffer



[a-z] { c = yytext[0]; yylval = c - 'a'; return(LETTER); } Place the offset c - 'a'

Match a positive integer (sequence of 0-9 digits)



[0-9]\* { yylval = atoi(yytext); return(NUMBER); }



Place the integer value In the stack

In the stack

### Steps to Execute Lex Program

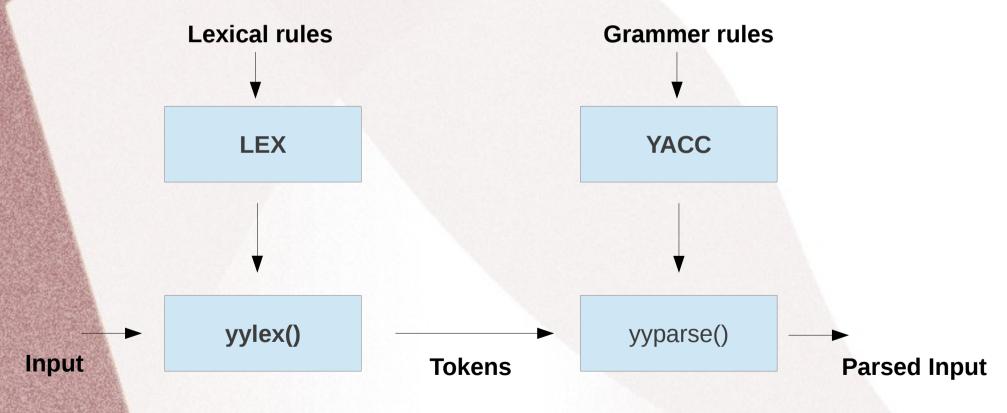
- lex <pgm name>
- cc lex.yy.c -II
- > ./a.out

## Examples: Lexical Analysis

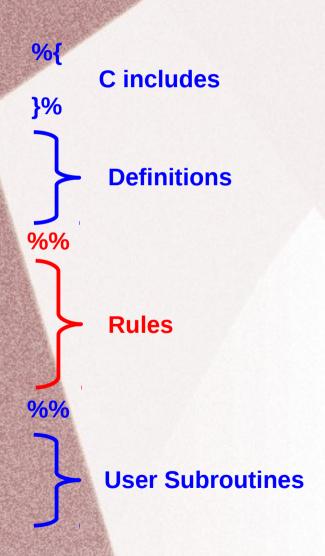
- 1.Remove white space (very simple example)
- 2.Character, line, word counting Lex example from the slides.
- 3.Lex example using states: removing comments from code
- 4. Transforming input

#### Yacc: A Parser Generator

 Takes a specification for a CFG, produces an LALR parser.



#### Structure of Yacc Specification File



Red: required Blue: optional

#### Yacc: Grammar Rules

- Terminals (tokens): Names must be declared:
  - %token name 1 name 2 ...

Any name not declared as a token in the declarations section is assumed to be a nonterminal.

- Start symbol :
  - may be declared, via: %start name
  - if not declared explicitly, defaults to the nonterminal on the LHS of the first grammar rule listed.

#### **Yacc Grammar Rules**

Productions : A grammar production A → B1
 B2 · · · Bn is written as

```
- A: B 1 B2 ⋅ ⋅ ⋅ B n;
```

Note: Left-recursion is preferred to right-recursion for efficiency reasons.

Example:

```
- stmt : KEYWD_IF '(' expr ')' stmt ;
```

#### Actions

 the user may associate actions to be performed each time the rule is recognized in the input process, eg:

- \$ is special!
  - \$n → psuedo-variables which refer to the values returned by the components of the right hand side of the rules.
  - \$\$ → The value returned by the left-hand side of a rule.

```
Expr: '(' expr')' { $$ = $2;}
```

Default return type is integer.

#### **Declarations**

- %start: means the whole input should match line
- %union: lists all possible types for values associated with parts of the grammar and gives each a field-name
- the type is generated and must be included into the lex source so that types can be associated with tokens.

```
typedef union {
   body of union ...
} YYSTYPE;
```

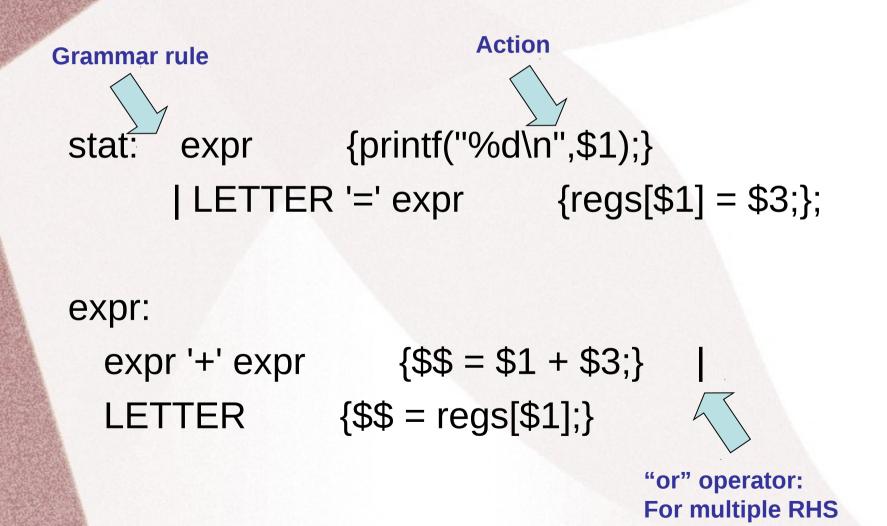
 %type: gives an individual type for the values associated with each part of the grammar, using the field-names from the %union declaration

# %{ ... c includes ... %} ... definitions ... %% ... rules ... %% ... subroutines ...

#### Yacc program

```
%{
#include <stdio.h>
int regs[26];
int base;
%}
%token NUMBER LETTER
%left '+' '-'
%left '*' '/'
%%
list: | list stat '\n' |list error '\n' {yyerrok;};
stat: expr {printf("\%d\n",$1);}
      | LETTER '=' expr \{regs[$1] = $3;\};
expr:
'(' expr ')' {$$ = $2;}
expr'+'expr {$$ = $1 + $3;}
LETTER
             $$ = regs[$1];
%%
main(){return(yyparse());}
yyerror(CHAR *s){fprintf(stderr, "%s\n",s);}
yywrap(){ return(1);}
```

## Rule Reduction and Action



#### Communication between Scanner and Parser

- The user must supply an integer-valued function yylex() that implements the lexical analyzer (scanner).
- If there is a value associated with the token, it should be assigned to the external variable yylval.
- The token error is reserved for error handling.
- Token numbers: These may be chosen by the user if desired. The default is:
  - chosen by yacc [in a file y.tab.h]
  - the token no. for a literal is its ASCII value
  - other tokens are assigned numbers starting at 257
  - the endmarker must have a number zero or negative.
- Generate y.tab.h using 'yacc -d'

#### **Using Yacc**

- Suppose the grammar spec is in a file foo.y. Then:
  - The command 'yacc foo.y' yields a file y.tab.c containing the parser constructed by yacc.
  - The command 'yacc -d foo.y' constructs a file y.tab.h that can be #include'd into the scanner generated bylex.
  - The command 'yacc -v foo.y' additionally constructs a file y.output containing a description of the parser (useful for debugging).
- The user needs to supply a function main() to driver, and a function yyerror() that will be called by the parser if there is an error in the input.

## Conflicts and Ambiguities

- Conflicts may be either shift/reduce or reduce/reduce:
  - In a shift/reduce conflict, the default is to shift.
  - In a reduce/reduce conflict, the default is to reduce using the first applicable grammar rule.
- Arithmetic Operators : associativity and precedence can be specified:
  - Associativity: use %left, %right, %nonassoc
- Precedence (Binary Operators):
  - Specify associativity using %left etc.
  - Operators within a group have same precedence. Between groups, precedence increases going down.

# Conflicts and Ambiguities cont'd

- Precedence (Unary Operators): use %prec keyword. This changes the precedence of a rule to be that of the following token.
- Example:

```
%left '+' '-'
%left '*' '/'
.....

expr : expr '+' expr
| expr '*' exp
| '-' expr %prec '*'
| ID
```

# Steps to execute YACC program:

- yacc -d <yacc\_pgm name>
- lex <lex\_pgm\_name>
- cc y.tab.c lex.yy.c -ly -ll
- > ./a.out

#### **Examples**

 Program to recognize strings 'aaab', 'abbb', 'ab' and 'a' using grammar(a^nb^n, n>=0)

#### Conclusions

- Yacc and Lex are very helpful for building the compiler front-end
- A lot of time is saved when compared to handimplementation of parser and scanner
- They both work as a mixture of "rules" and "C code"
- C code is generated and is merged with the rest of the compiler code