

Recall:

Definition [Newton-Local (NL) condition]: f satisfies the NL condition for $\delta_0 < 1$

if $\forall 0 \leq \delta \leq \delta_0$, $\forall (x, y)$ s.t. $\|y - x\|_2 \leq \delta$, we have that

$$(1 - 3\delta) H_f(x) \preceq H_f(y) \preceq (1 + 3\delta) H_f(x)$$

PSD ordering

Theorem [analysis of Newton's iteration]: Let $f: \mathbb{R}^n \rightarrow \mathbb{R}$ be a strictly convex function that satisfies the NL condition for $\delta_0 = \frac{1}{6}$.

Then if $x_1 := x_0 + n(x_0)$

$$\left(\text{recall } n(x_0) := -H_f(x_0)^{-1} \nabla f(x_0) \right)$$

$$\text{s.t. } \|n(x_0)\|_{x_0} \leq \frac{1}{6}$$

$$\text{then } \|n(x_1)\|_{x_1} \leq 3 \|n(x_0)\|_{x_0}^2$$

Already proved in last lecture:

Lemma: Suppose x, y are such that $\|y - x\|_2 \leq \frac{1}{6}$ (Also f satisfies the NL condition with $\delta_0 = \frac{1}{6}$). Then $\forall u$

$$\textcircled{1} \quad \frac{1}{2} \|u\|_x \leq \|u\|_y \leq 2 \|u\|_x$$

$$\textcircled{2} \quad \frac{1}{2} \|u\|_{H_f(x)^{-1}} \leq \|u\|_{H_f(y)^{-1}} \leq 2 \|u\|_{H_f(x)^{-1}}$$

Proof of Theorem [analysis of Newton's iteration]:

$$\begin{aligned}
 \|n(x_1)\|_{x_1} &= \sqrt{\langle n(x_1), H_f(x_1) n(x_1) \rangle} \\
 &= \sqrt{\langle H_f(x_1)^{-1} \nabla f(x_1), H_f(x_1) H_f(x_1)^{-1} \nabla f(x_1) \rangle} \\
 &= \|\nabla f(x_1)\|_{H_f(x_1)^{-1}}
 \end{aligned}$$

$$\|n(x_0)\|_{x_0} = \|\nabla f(x_0)\|_{H_f(x_0)^{-1}}$$

Because of the lemma, it suffices to prove that

$$\|\nabla f(x_1)\|_{H_f(x_0)^{-1}} \leq \frac{3}{2} \|\nabla f(x_0)\|_{H_f(x_0)^{-1}}$$

Fundamental theorem of calculus: $g(1) = g(0) + \int_0^1 g'(t) dt$

$$h(t) = \nabla f(x_0 + t(x_1 - x_0))$$

$$\nabla f(x_1) = \nabla f(x_0) + \int_0^1 H_f(x_0 + t(x_1 - x_0))(x_1 - x_0) dt$$

$$= \nabla f(x_0) - \int_0^1 H_f(x_0 + t(x_1 - x_0)) H_f(x_0)^{-1} \nabla f(x_0) dt$$

$$= \left(H_f(x_0) - \int_0^1 H_f(x_0 + t(x_1 - x_0)) dt \right) H_f(x_0)^{-1} \nabla f(x_0)$$

$\underbrace{_{M(x_0)}$

$$\|\nabla f(x_1)\|_{H_f(x_0)^{-1}} = \sqrt{\langle M(x_0) H_f(x_0)^{-1} \nabla f(x_0), H_f(x_0)^{-1} M(x_0) H_f(x_0)^{-1} \nabla f(x_0) \rangle}$$

$$= \sqrt{\underbrace{\langle H_f(x_0)^{-1/2} M(x_0) H_f(x_0)^{-1/2}, H_f(x_0)^{-1/2} \nabla f(x_0) \rangle}_{N(x_0)} + \underbrace{\langle H_f(x_0)^{-1/2} M(x_0) H_f(x_0)^{-1/2}, H_f(x_0)^{-1/2} \nabla f(x_0) \rangle}_{H_f(x_0)^{-1/2} \nabla f(x_0)}}$$

$$\begin{aligned} &= \| N(x_0) H_f(x_0)^{-1/2} \nabla f(x_0) \|_2 \\ &\leq \| N(x_0) \| \| H_f(x_0)^{-1/2} \nabla f(x_0) \|_2 \\ &= \| N(x_0) \| \| \nabla f(x_0) \|_{H_f(x_0)^{-1}} \end{aligned}$$

To complete the proof, need $\| N(x_0) \| \leq \frac{3}{2} \| \nabla f(x_0) \|_{H_f(x_0)^{-1}}$

$$N(x_0) = H_f(x_0)^{-1/2} M(x_0) H_f(x_0)^{-1/2}$$

[Easy to prove fact: $\| A^{-1/2} B A^{-1/2} \| \leq \alpha$
 $\text{iff } -\alpha A \leq B \leq \alpha A$]

$$\delta := \| \nabla f(x_0) \|_{H_f(x_0)^{-1}} = \| x_1 - x_0 \|_{x_0}$$

suffices to prove that $-\frac{3}{2}\delta H(x_0) \leq M(x_0) \leq \frac{3}{2}\delta H(x_0)$

$$M(x_0) = H(x_0) - \int_0^1 H(x_0 + t(x_1 - x_0)) dt$$

$$z_t := x_0 + t(x_1 - x_0)$$

$$\| z_t - x_0 \|_{x_0} = t \| x_1 - x_0 \|_{x_0} = t \delta$$

$$\Rightarrow (1 - 3t\delta) H(x_0) \leq H(z_t) \leq (1 + 3t\delta) H(x_0)$$

(NL condition)

$$\Rightarrow (1-3t\delta) H(x_0) \leq H(z_t) \leq (1+3t\delta) H(x_0) \quad (\text{NL condition})$$

$$\Rightarrow \frac{-3t\delta}{H(x_0)} H(x_0) \leq H(x_0) - H(z_t) \leq 3t\delta H(x_0)$$

Integrating on t from 0 to 1 finishes the proof.

Interior point methods for Linear Programming

$$P = \left\{ x \in \mathbb{R}^n : \underbrace{\langle a_i, x \rangle \leq b_i}_{i \in [m]} \right\}$$

Linear program: $\min \langle c, x \rangle \rightarrow Ax \leq b$
 s.t. $x \in P$

Theorem: There is an algorithm that given the description of the LP (A, b, c) with n variables, m constraints & total bit complexity L , $\varepsilon > 0$, & assuming that the polytope P is full-dimensional & non-empty, outputs $\hat{x} \in P$ s.t.

$$\langle c, \hat{x} \rangle \leq \langle c, x^* \rangle + 2\varepsilon \quad \begin{array}{l} \text{(running time)} \\ \text{poly}(L, \log(1/\varepsilon)) \end{array}$$

$\langle c, x^* \rangle$ optimal solution

Remarks:

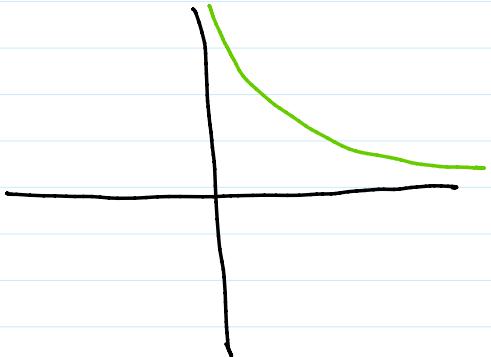
- ① Possible to remove the full-dimensionality assumption.
- ② Possible to solve the LP exactly. Set $\varepsilon = 2^{-\text{poly}(L)}$ & then do some rounding.

Barrier functions: $K \subseteq \mathbb{R}^n$ convex set.

$F: \text{int}(K) \rightarrow \mathbb{R}$ is said to be a barrier function for K if F approaches $+\infty$ as we approach the boundary of K . Also require F to be strictly convex.

Example: $K \subseteq \mathbb{R} = \{x : x > 0\}$

$$F(x) = 1/x$$



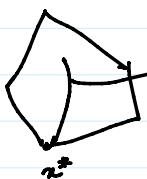
An approach for $\min_{\substack{\text{s.t. } x \in K}} \langle c, x \rangle$

Consider a family of objective functions

$$f_\eta(x) := \eta \langle c, x \rangle + F(x)$$

as η gets big enough, the optimal solution for f_η gets close to the optimal sol'n x^* for the above program

$$x_\eta^* := \underset{x \in \text{int}(K)}{\operatorname{argmin}} \eta \langle c, x \rangle + F(x)$$



curve of x_n^*
(central path)

Path following Interior Point Methods
(IPM)

$$P = \left\{ x \in \mathbb{R}^n : \langle a_i, x \rangle \leq b_i \forall i \in [m] \right\}$$

Definition [Logarithmic barrier]: $F(x) := - \sum_{i=1}^m \log(b_i - \langle a_i, x \rangle)$
 (defined only in the interior of P)

The idea is to follow the central path.

- ① Initialization: How to find some initial $\eta_0 > 0$ & a point close to x_η^* ?
- ② Following the central path: How to follow the central path using discrete steps?
- ③ Termination: When to stop? How to guarantee that we have reached close to x^* ?

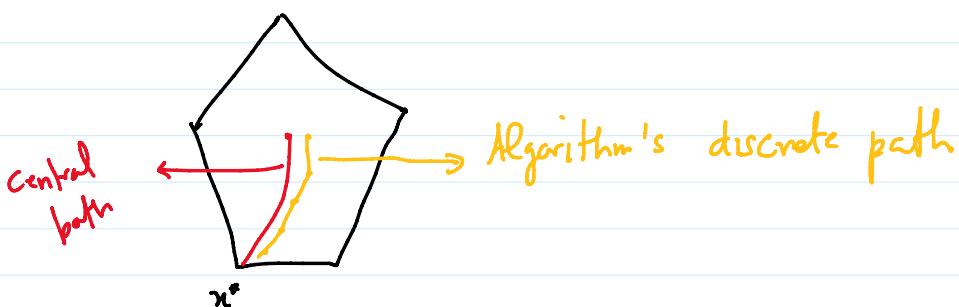
Suppose we are given x_0 close to x_{η_0} for some $\eta_0 > 0$.

We will use Newton's method to update the current point.

$$x_1 := x_0 + N_{\eta_0}(x_0), \quad N_\eta(x) := -H_{f_\eta}(x)^{-1} \nabla f_\eta(x) \\ \text{first step} \qquad \qquad \qquad = -H_F(x)^{-1} \nabla f_\eta(x)$$

$$\text{more generally } x_{t+1} := x_t + N_{\eta_t}(x_t)$$

$$\text{& then also update } \eta_{t+1} := \eta_t(1+\gamma)$$



What is the notion of closeness to maintain?
 whatever came out of the analysis of Newton's method.

Invariant that we'll maintain : $\| N_{\eta_t}(x_t) \|_{x_t} := \| N_{\eta_t}(x_t) \|_{H_F(x_t)} \leq \frac{1}{6}$

IPM Algorithm :

1. Initialization : Find some $\eta_0 > 0$ & x_0 s.t. $\| N_{\eta_0}(x_0) \|_{x_0} \leq \frac{1}{6}$.
2. T s.t. $\eta_T := \left(1 + \frac{1}{2\sqrt{m}}\right)^T \eta_0 > m/\varepsilon$

for $t = 0, 1, \dots, T$

3. Newton step : $x_{t+1} := x_t + N_{\eta_t}(x_t)$
4. Update η : $\eta_{t+1} := \eta_t \left(1 + \frac{1}{2\sqrt{m}}\right)$

Termination : $\hat{x} :=$ Two Newton steps from x_T w.r.t. f_{η_T} .
 return \hat{x}

Theorem : Above algorithm after $T = O\left(\sqrt{m} \log\left(\frac{m}{\varepsilon \eta_0}\right)\right)$ iterations outputs a point \hat{x} s.t.

$$\langle c, \hat{x} \rangle \leq \langle c, x^* \rangle + 2\varepsilon$$

Each iteration requires solving a linear system of the kind $H_F(x) y = z$ (y is the unknown).