#### **Chap 4. Inter-Process Communication**

#### Road map

- ◆ 4.1. Intro
- 4.2. API for the Internet Protocols
- 4.3. External data representation and marshalling
- ◆ 4.4. Client-Server Communication
- ◆ 4.5. Group communication (self-read)
- 4.6. Case study (self-read)

#### 4.1. Intro

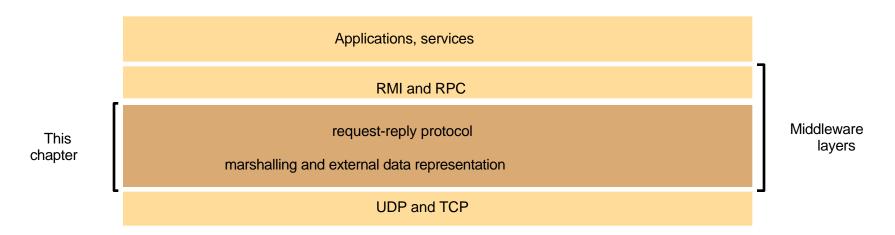
#### Focus:

- Characteristics of protocols for communication between processes to model distributed computing architecture
  - ★ Effective means for communicating objects among processes at language level
- Java API
  - Provides both datagram and stream communication primitives/interfaces – building blocks for communication protocols
- Representation of objects
  - ⋆ providing a common interface for object references
- Protocol construction
  - ★ Two communication patterns for distributed programming: C-S using RMI/RPC and Group communication using 'broadcasting'

#### **4.1.** Intro

- In Chapter 3, we covered Internet transport (TCP/UDP) and network (IP)
   protocols without emphasizing how they are used at programming level
- In Chapter 5, we cover RMI facilities for accessing remote objects' methods

  AND the use of RPC for accessing the procedures in a remote server
- Chapter 4 is on how TCP and UDP are used in a program to effect communication via socket (e.g. Java sockets) – the Middle Layers – for object request/reply invocation and parameter marshalling/representation



- IPC primitives
  - message passing between a pair of processes can be supported by two message communication operations: send and receive
  - send (destination, &msg); receive (source, &buf)
  - Destination and source can be process id or port number (single receiver); Or mailbox (multiple receivers)
    - ⋆ Typically, (IP, port#) pair
- Use of sockets as API for UDP and TCP implementation much more specification can be found at java.net

Message: Header Body

- Synchronous communication
  - Send and receive processes synchronize at every message
  - ◆ Both blocking:
    - whenever a send is issued, blocks until corresponding receive is issued; whenever a receive is issued, blocks until message arrives
- Asynchronous communication
  - Send from client is <u>non-blocking</u>, proceeds as soon as the message has been copied to a local buffer
  - Receive could be non-blocking or <u>blocking</u>, the latter has no disadvantages in system environment supporting multithreading like Java
    - Non-blocking receive is not very useful (you cannot proceed without message)

#### comparison

Blocking

Advantages: Ease of use and low overhead of implementation

Disadvantage: low concurrency

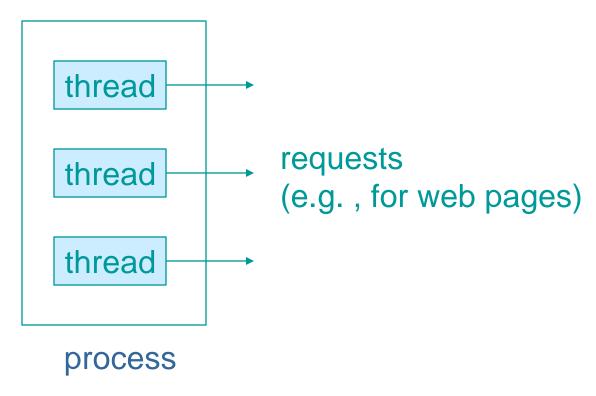
Non-blocking

Advantages: Flexibility, parallel operations

Disadvantages: Programming tricky: Program is timing-dependent (interrupt can occur at arbitrary time, and execution irreproducible)

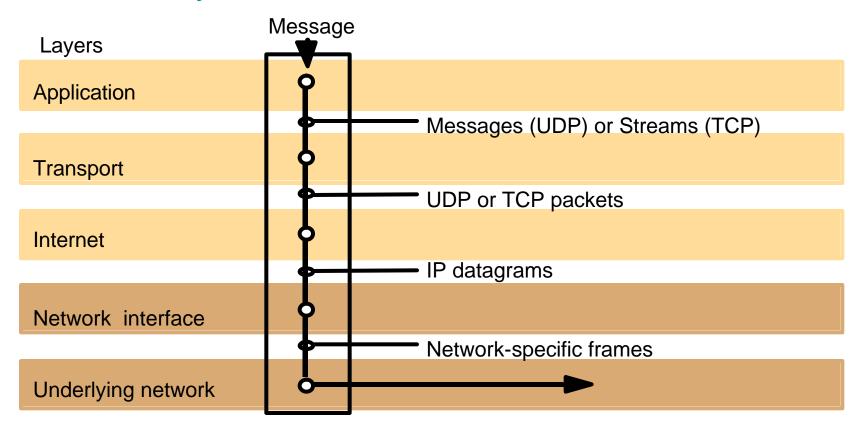


Using blocking operation without penalty

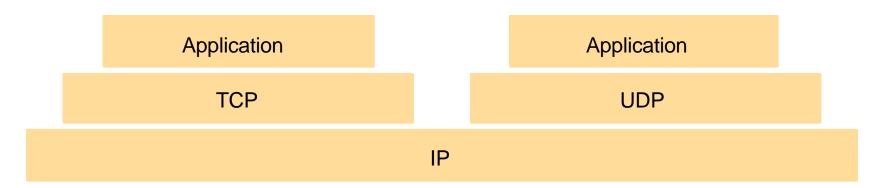


Some threads may be blocked while others continue to be active

#### TCP/IP layers



The programmer's conceptual view of a TCP/IP Internet



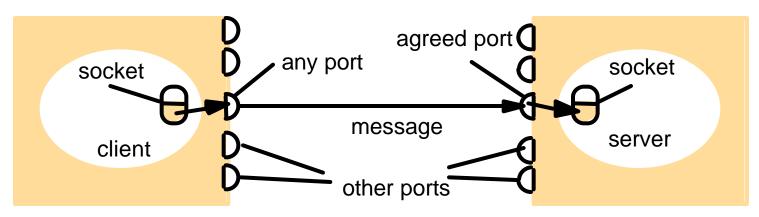
Recall: UDP, TCP, API

- UDP (User Datagram Protocol): offers no guarantee of delivery, a datagram protocol
- TCP (Transmission Control Protocol): reliable connection-oriented protocol, establishment of bidirectional communication channel
- API (Application programming interface)

#### Message destinations

- Messages are sent to a pair (Internet address, local port)
- ◆ Local port: an integer, message destination within a computer
  - ★ Each computer has 2<sup>16</sup> possible ports available to local processes for receiving messages
  - ★ 0-1023: well-known for restricted use of privileged processes
  - ★ 1024-49151: registered ports, holds service descriptions
  - ★ 49152-65535: private purposes
  - ★ In practice, non-restricted ports can be used for private purposes, then, cannot use registered services simultaneously
- a port has exactly one receiver (except for multicast ports)
- receiving process can have many ports for different message types
- server processes usually publish their service-ports for clients

- Sockets : originated from BSD Unix
  - Provide an abstraction of endpoints for both TCP and UDP communication
  - Inter-process communication consists of transmitting a message between a socket in one process and a socket in another process
  - A socket must be bound to a local port and an IP address
  - Processes may use the same socket for sending and receiving messages
  - Sockets are typed/associated with a particular protocol, either TCP or UDP



Internet address = 138.37.94.248

Internet address = 138.37.88.249

- Java API for internet protocols
  - ◆ For either TCP or UDP, Java provides an InetAddress class, which contains a method: getByName(DNS) for obtaining IP addresses, irrespective of the number of address bits (32 bits for IPv4 or 128 bits for IPv6) by simply passing the DNS hostname. For example, a user Java code invokes:

```
InetAddress aComputer =
InetAddress.getByName("www.cs.sfu.ca")
```

 The class encapsulates the details of the representation of the IP address

- UDP Datagram communication
  - Steps:
    - ★ Client finds an available port for UPD connection
    - Client binds the port to local IP (obtained from InetAddress.getByName(DNS))
    - Server finds a designated port, publicizes it to clients, and binds it to local IP
    - Sever process issues a receive method and gets the IP and port # of sender (client) along with the message
  - Issues
    - ★ Message size receiver needs to specify a buffer of certain size to receive a massage. If message too big, truncated on arrival
    - ★ Blocking send is non-blocking, returns when the message gets the UDP and IP layers; receive is blocking until a datagram is received or timeout
    - ★ Timeouts reasonably large time interval can be set on <u>receiver</u> sockets to avoid indefinite blocking if required by program
    - ★ Receive from any no specification of sources (senders)

- UDP Failure Models:
  - Omission failure: due to omission of send or receive (either checksum error or no buffer space at source or destination)
  - Ordering failure: due to out-of-order delivery
  - Applications using UDP are left to provide their own checks to achieve the quality of reliable communication they require (why and how?)
    - ⋆ UDP lacks built-in checks
    - but failure can be modeled by implementing an ACK mechanism

#### UDP client sends a message to the server & gets a reply

```
import java.net.*; //defines socket-related classes
import java.io.*; //defines stream-related classes
public class UDPClient{
  public static void main(String args[]){ // args give message contents and server hostname
   DatagramSocket aSocket = null;
    try {
     aSocket = new DatagramSocket();
     byte [] m = args[0].getBytes();
     InetAddress aHost = InetAddress.getByName(args[1]);
     int serverPort = 6789:
     DatagramPacket request = new DatagramPacket(m, args[0].length(), aHost, serverPort);
     aSocket.send(request);
     byte[] buffer = new byte[1000];
     DatagramPacket reply = new DatagramPacket(buffer, buffer.length);
     aSocket.receive(reply);
     System.out.println("Reply: " + new String(reply.getData()));
    } catch (SocketException e) {System.out.println("Socket: " + e.getMessage());
    } catch (IOException e) {System.out.println("IO: " + e.getMessage());}
   } finally { if(aSocket != null) aSocket.close();}
```

UDP server repeatedly receives a request and sends it back to the client

```
import java.net.*;
import java.io.*;
public class UDPServer{
   public static void main(String args[]){
   DatagramSocket aSocket = null;
      try{
          aSocket = new DatagramSocket(6789);
          byte[] buffer = new byte[1000];
          while(true){
            DatagramPacket request = new DatagramPacket(buffer, buffer.length);
            aSocket.receive(request):
            DatagramPacket reply = new DatagramPacket(request.getData(),
                    request.getLength(), request.getAddress(), request.getPort());
            aSocket.send(reply);
      }catch (SocketException e){System.out.println("Socket: " + e.getMessage());
     }catch (IOException e) {System.out.println("IO: " + e.getMessage());}
   }finally {if(aSocket != null) aSocket.close();}
```

- TCP Stream Communication issues:
  - Message sizes user application has option to choose how much data it writes to a stream or reads from it
  - Lost messages / Flow control: dealt with by TCP
  - Message duplication and ordering message identifiers are associated with each IP packet so as for recipient to detect and reject duplicates or re-order
  - Message destinations a connection is established first.
     Once established, no IP addresses in packets needed (
     Each connection socket is bidirectional using two streams: output/write and input/read)
    - ★ To establish a connection, client sends connect request to server, then server sends accept request to client
    - ★ Could be a overhead for a single C-S request and reply

#### TCP Stream Communication: other issues

- Matching of data items both client/sender and server/receiver must agree on data types in the stream
- Threads server creates a separate thread in accepting a connection, then it can block waiting for input without delaying other clients

#### Failure Model

- ★ Integrity: uses <u>checksums</u> for detection/rejection of corrupt data and <u>seq #s</u> for detection/rejection of duplicates
- Validity: uses timeout with retransmission techniques to take care of packet losses
- Uses TCP sockets used for such services as: HTTP, FTP, Telnet, SMTP

TCP client makes connection to server, sends request and receives reply

```
import java.net.*; //Defines socket-related classes
import java.io.*; //Defines stream-related classes
public class TCPClient {
   public static void main (String args[]) {
   // arguments supply message and hostname of destination
   Socket s = null:
      try{
          int serverPort = 7896:
          s = new Socket(args[1], serverPort);
          DataInputStream in = new DataInputStream( s.getInputStream());
          DataOutputStream out = new DataOutputStream( s.getOutputStream());
          out.writeUTF(args[0]);
                                   // UTF is a string encoding see Sn 4.3
          String data = in.readUTF();
          System.out.println("Received: "+ data);
             }catch (UnknownHostException e){
                     System.out.println("Sock:"+e.getMessage());
      }catch (EOFException e){System.out.println("EOF:"+e.getMessage());
      }catch (IOException e){System.out.println("IO:"+e.getMessage());}
   }finally {if(s!=null) try {s.close();}catch (IOException
e){System.out.println("close:"+e.getMessage());}}
```

TCP server makes a connection for each client and then echoes the client's request

```
import java.net.*;
import java.io.*;
public class TCPServer {
  public static void main (String args[]) {
   try{
     int serverPort = 7896;
     ServerSocket listenSocket = new ServerSocket(serverPort);
     while(true) {
                Socket clientSocket = listenSocket.accept();
                Connection c = new Connection(clientSocket);
   } catch(IOException e) {System.out.println("Listen :"+e.getMessage());}
// this figure continues on the next slide
```

```
class Connection extends Thread {
  DataInputStream in;
   DataOutputStream out;
   Socket clientSocket;
  public Connection (Socket aClientSocket) {
     try {
    clientSocket = aClientSocket;
    in = new DataInputStream( clientSocket.getInputStream());
    out =new DataOutputStream( clientSocket.getOutputStream());
    this.start();
     } catch(IOException e) {System.out.println("Connection:"+e.getMessage());}
  public void run(){
                                            // an echo server
     try {
    String data = in.readUTF();
    out.writeUTF(data);
     } catch(EOFException e) {System.out.println("EOF:"+e.getMessage());
     } catch(IOException e) {System.out.println("IO:"+e.getMessage());}
     } finally{ try {clientSocket.close();}catch (IOException e){/*close failed*/}}
```