

Fenwick Tree (BIT)

Given some sequence

+

queries determine range sum $[a, b]$

0	1	4	2	1	3	5	1	2	3
0	1	5	7	8	11	16	17	19	22
/ / / / /									

Normal method $\rightarrow O(N)$

using prefix sum $\rightarrow O(1)$

Dynamic

- $O(1)$ query time

- $O(n)$ recompute time.

\Rightarrow Thus, Fenwick tree.

1	0	0	0	0
	1	1	1	1
	1	1	1	0
	1	1	0	1
	1	1	0	0
	1	0	1	1
	1	0	1	0
	1	0	0	1
	1	0	0	0
	0	1	1	1
	0	1	1	0
	0	1	0	1
	0	1	0	0
	0	0	1	1
	0	0	1	0
	0	0	0	1
				0

Vertical array (respⁿ) range of responsⁿ
(decided by lowest one bit)

* F.T is organized by lowest one bit.
range of responsibility

prefix sum

Ex: Prefix sum till 11th index.

1011 \rightarrow v[11]
 \rightarrow 1010 \rightarrow v[10]
 \rightarrow 1000 \rightarrow v[8]
 \rightarrow 0000 \rightarrow v[0]
 as soln 1 based, it's useless.

$$\Rightarrow \text{pre}[11] = v[11] + v[10] + v[8];$$

Steps

- start at current cell;
- staircase down until no bits are left;
- i.e subtract lowest one bit each time to get to next index in F.T (To move down)

runtime?

- loop runs in at most making zero of all set bits
- no of bits are $O(\log n)$
- complexity = $O(\log n)$

update in fenwick tree:

For Ex:

Δ is change = $v_{\text{new}} - v_{\text{old}}$.
 arr[9] is changed by 2,
 $\rightarrow v[9]$ will also change and
 also all the v[i] in tree
 that took responsibility of '9'

\rightarrow 1001
 1010 [look in table]
 1100
 10000

These are nodes in F-T resp
 for '9'. So we need to add Δ
 to them.

\rightarrow How to find? (To move up)

Soln add '1' to lowest 1 bit.

Find 1st bit.
 $i \& -i$
 Ex:
 $i = 1010$
 $-i \Rightarrow 0101$
 $+1$
 $-i = 0110$
 $(i \& -i)$
 $= (1010) (0110)$
 $= (0010) \checkmark$

- must update all the owners of us.
- move by adding lowest one bit to current index.
- loop runs exactly as number of offbits.
- $O(\log n)$ runtime bounded.

CODE TIME

```
vector<int> t(N, 0);
void bit_update( int i, int delta, int size) {
    while(i <= size) {
        t[i] += delta;
        i += (i & -i);
    }
}
int pre_sum( int i ) {
    int sum = 0;
    while ( i > 0 ) {
        sum += t[i];
        i -= (i & -i);
    }
    return sum ;
}
int range_sum(int l, int r) {
    return pre_sum(r) - pre_sum(l - 1);
}

void solve() {
    int n; cin >> n;
    vector<int> a(n+1, 0);
    for(int i = 1; i <= n; i++ ) {
        cin >> a[i];
        bit_update(i, a[i], n);
    }
```

by how much value changes.

keep in mind to use \leq and not $<$.

if ($l > r$) return 0;

(This is the best part of code.)

Questions Based on BIT $(2-1)=1$

① Counting Inversions. (Also try with Merge Sort)

permutation

7 6 2 3 1 4 5
↑ ↑

$(6, 1)$ is Inversion.

7	6	2	3	1	4	5
6	5	1	1	0	0	0



7
1
6
9

6
1
5

2
1
1

3
1
1

1
0


4
0

↓
5
0

How many no. ahead of me
are smaller than me.

$$\begin{array}{r} 7 \\ 6 \\ 5 \quad 1 \\ \hline 4 \quad 1 \\ \hline 3 \quad 1 \\ 2 \\ \hline 1 \quad 1 \end{array}$$

Implementation



$i < j < k, a[i] < a[j] < a[k]$ find 5.c triplets.
 hint we 2 ~~tree~~ ^{BIT}

hint we 2 ~~tree~~ ¹¹⁷

Q2. Find K^{th} tallest.

$n \rightarrow [1, 10^6]$

update: changes no of people
of a party he.

Quer: find K^{th} tallest $R[1, 10^{18}]$

Ideas:

$R \rightarrow$ 1 2 3 4 5 6
 $f \rightarrow$ 8 2 10 100 1 2

$R \rightarrow$ 8 10 20 (120) 121 123

Ex: (K=6), $k=122$ ✓

So, i can maintain, a variable
prefix sum with BIT, and my
query using binary search.

So complexity $\rightarrow (\log n)(\log n) = (\log n)^2$
 $\uparrow \quad \uparrow$
B.S. BIT

Space complexity $\rightarrow O(N)$

Can we do faster?

Yes we can:

By doing binary search
directly on BIT. find?

For Ex: (sum=20)

11 3
10 1
9 1
→ 8 25
7 2
6 4
5 3
→ 4 16
3 6
2 12
1 8

keep starting
pointer at
max power
of 2. i.e. 8
(i.e. 2³)

→ He can see
that the value is
larger. So we should go

to next smaller power of 2
⇒ 4

* $t[4] < 20$,

so this needs to included
in our found sum.

$sum = 16$

* Now we need to search above it.

* $i = 4 + 2 = 6$
net power of 2

* Check $t[6] \rightarrow 7$, $16 + 7 > 20$
so, No

* $4 + 1 = 5$
 $t[5] \rightarrow 3$, $16 + 3 < 20$ ✓
(pos=5, sum=19)

* Now we are finished with all
powers of 2, so we can stop.

Index with value less than sum is
pos, with sum as sum.

Code:

```
int bS(int n, int sum) {  
    // pos < val;  
    int sm = 0, pos = 0;  
    for (int i = (int)log2(n); i >= 0; i--) {  
        if ((pos + (1 << i)) < n && sm + t[pos + (1 << i)] < sum) {  
            sm += t[pos + (1 << i)];  
            pos += (1 << i);  
        }  
    }  
    return pos+1;  
}
```

Link to this : <https://codeforces.com/blog/entry/61364> and this <https://codeforces.com/blog/entry/11275?locale=en>

<https://codeforces.com/contest/992/problem/E>