

QuickSelect

Warm-up

How can I find the 3rd smallest element in the next array?

data = [10, 3, 6, 9, 2, 4, 15, 23], K = 3rd

Output: 4


The partition method could help.



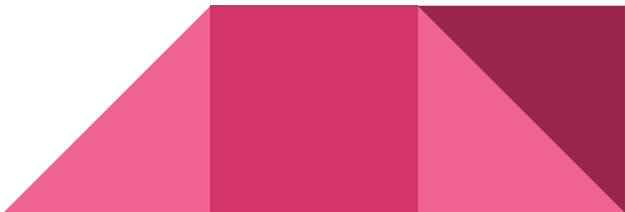
Divide and Conquer

Some algorithms, such as Merge sort and QuickSort and QuickSelect employ a common algorithmic paradigm based on recursion.

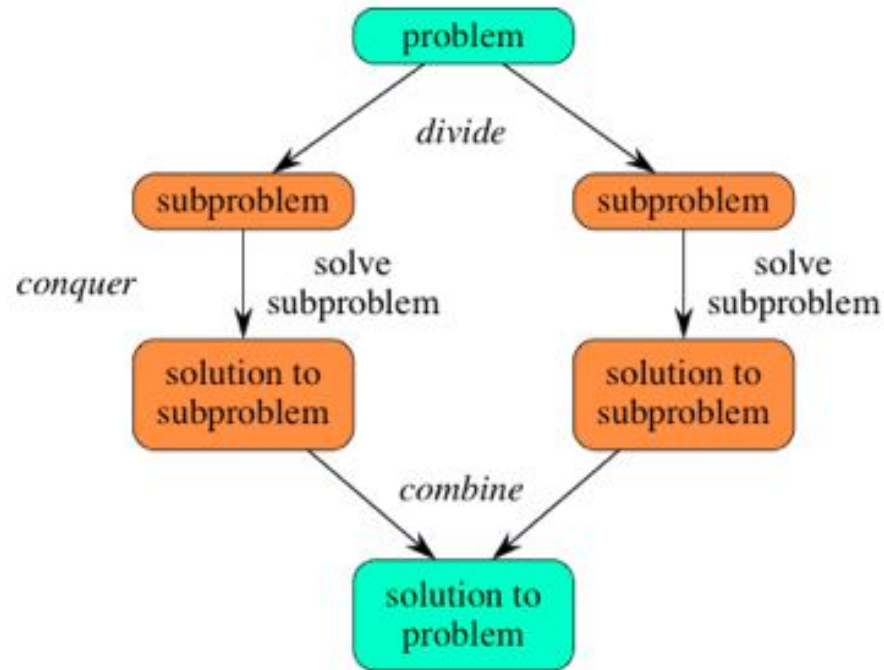
A divide-and-conquer algorithm recursively breaks down a problem into two or more sub-problems of the same or related type, until these become simple enough to be solved directly. (Wikipedia)

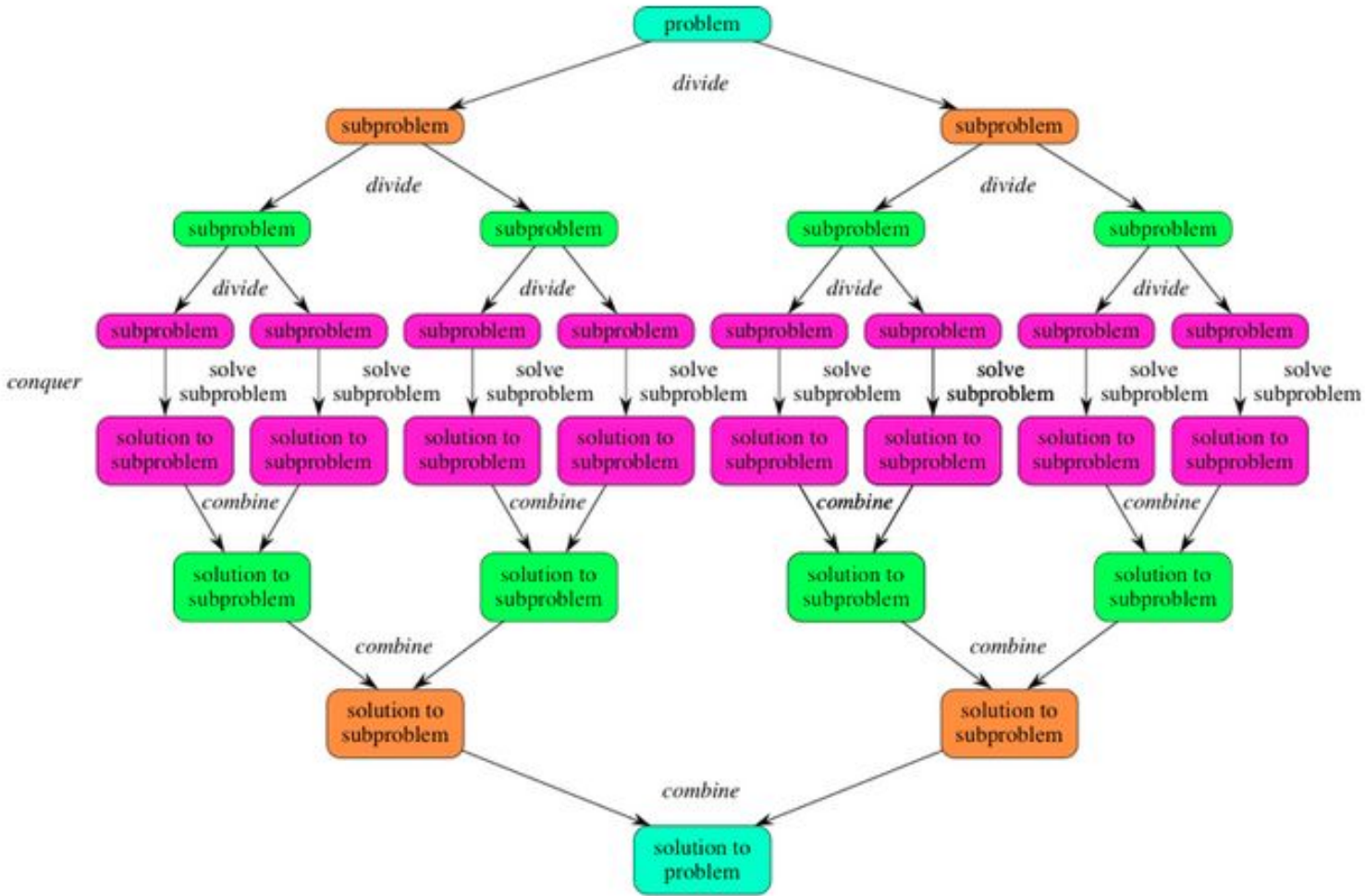


Divide and conquer algorithm parts

1. **Divide** the problem into a number of subproblems that are smaller instances of the same problem.
 2. **Conquer** the subproblems by solving them recursively.
 3. **Combine** the solutions to the subproblems into the solution for the original problem.
- 

Divide and Conquer





Quick Select

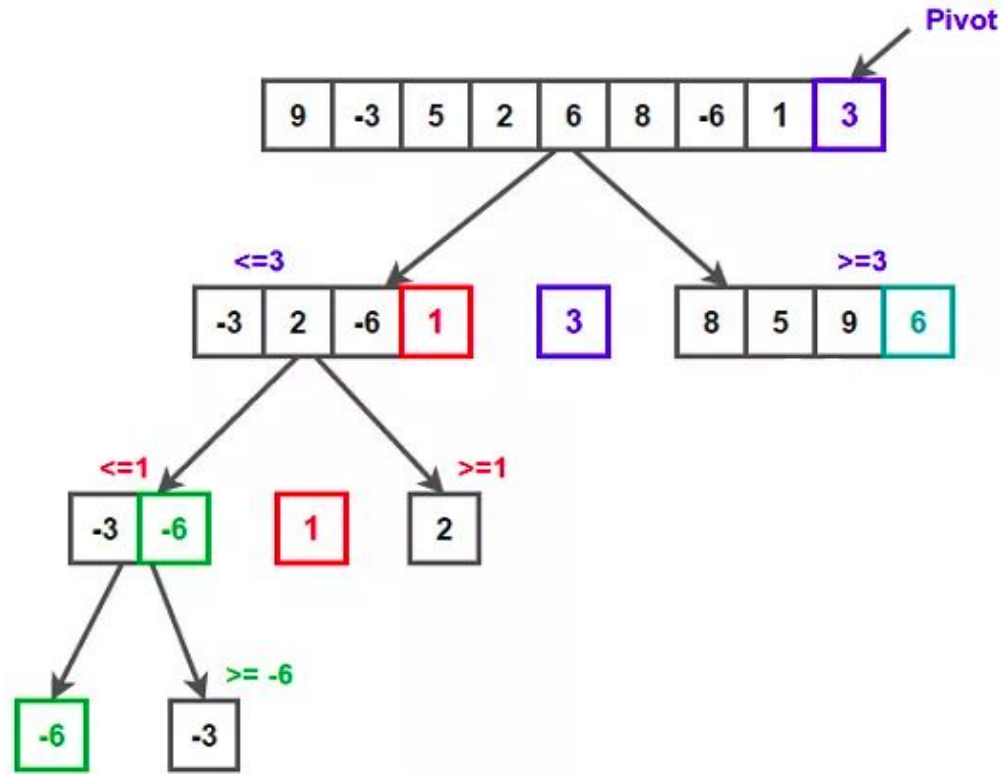
It is a selection algorithm. It looks for the k -th smallest element in an unsorted array.

It uses a partition strategy to find if index of the partitioned element is more than k , then we recur for the left part. If index is the same as k , we have found the k -th smallest element and we return. If index is less than k , then we recur for the right part.

Using a random pivot when defining partitions we can avoid the worst case (largest/smallest element picked as pivot).



Quick Select



Quick Select - Algorithm

QuickSelect method to find the Kth smallest element

```
    pivot = partition(data, start, end)
```

```
    if pivot > k - 1
```

```
        Recursive call to your QuickSelect - parameters: start=?, end=?
```

```
    Else if pivot < k - 1
```

```
        Recursive call to your QuickSelect - parameters: start=?, end=?
```

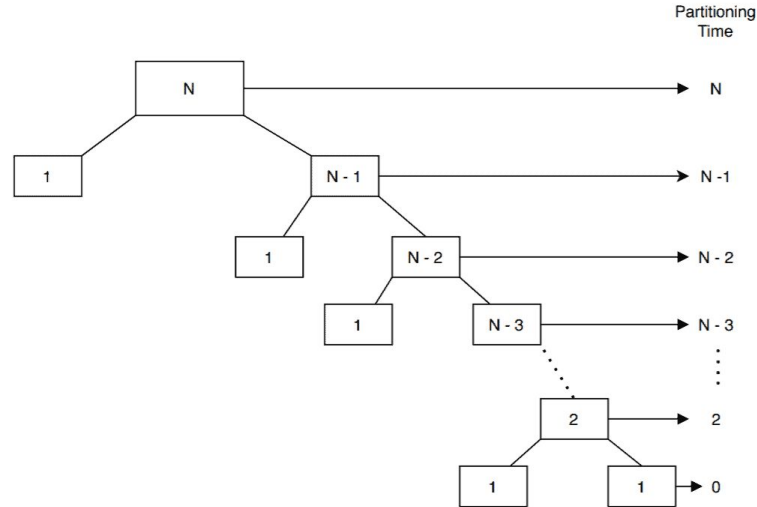
```
    Else
```

```
        Kth element found return the value [k-1]
```



Quick Select - Time Complexity

Worst case: larger/smallest element picked as pivot



$$O(N^2)$$

The height of the tree will be n and in top node we will be doing N operations then $n-1$ and so on until 1

Quick Select - Time Complexity

Best Case: when we partition the list into two halves and continue with only the half we are interested in.

$$N + N/2 + N/4 + N/8 + N/16$$

The sum of that series will never reach $2*N$.

$O(n)$



Coding Time!

Write your code here: classwork/xx_quick/Quick.java

You must implement the following method:

```
public static int quickSelect(int[] data, int k, int start, int end){  
    }  
}
```

Copy the partition method to your Quick.java

Call the method like this:

```
// find the third smallest element (test for different Kth elements)  
quickSelect(yourArray, 3, 0, yourArray.length-1));
```

