

Ans→

$$\text{No. of Subnet} = 2^m$$

and No. of host
per network =

$$2^{(8-m)} - 2$$

Ques: If default subnet mask for a network is 255.255.255.0 and if m bits are borrowed from NID then what is its Subnet mask?? (m is less than 8)

Ans→

$$255.255.(2^{8-m}) \cdot 0$$

(8-m)

↔

$$2^6 \quad 2^4 \quad 2^4$$

Ans:→

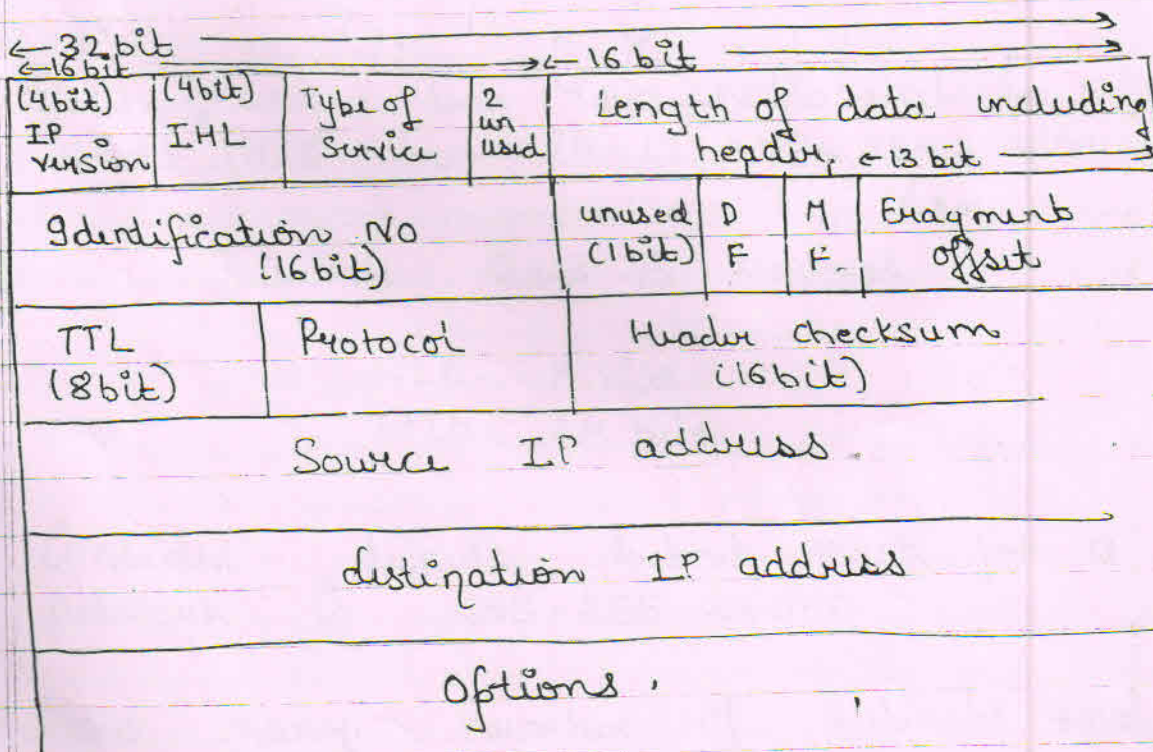
$$255.255.2^{(8-m)} \cdot 0 \quad 255.$$

$$2^{8-m}$$

IP- Header Format :-

→ IP- Layer is connectionless unreliable layer.

- OSI model kahata hai ki network layer should be connection oriented but TCP for network connectionless hote hai.
- TCP model for network-layer ki functionality kaam hote hai as compare to OSI-model.
- DDL layer is the only layer that append both header & trailer.
- Network layer and transport layer append only information in header not in trailer.



i) Version field \rightarrow

- Version field hamne yeh batati hai aap IP ka konsa version use karra hai ho. It is of 4 bit.
- After four bit, another four bit is for IHL (IP-Header Length). This field is necessary because IP-Header is of variable length.
- IHL : Since it is 4 bit long then IHL ki max value 15 and min 0.
- Rya IP-Header sirf 15 byte ka hi ho sakta hai kya??

Ans. In this format IHL min joh value store hoti hai woh actual header length nahi hoti.

$$\text{Actual Header Length} = 4 \times \text{IHL (in bytes)}$$

- Minimum value of IHL is 5, less than 5 represent error.
- Compulsory IP-Header 20 bytes ka hota hai aur maximum 60 byte ka.
- Next 6 byte is for type of service & next 2 bit is unused but actual min type of service field bhi unused hoti hai.

- Modern Routers inko use nahi karte.
- Type of Service indicates kis - type ki service hai, reliable hai, unreliable hai, connection oriented hai, connectionless hai.
- For IP is field ko dek K apne working mein kahi change nahi karta hai.
- Now next 16 bit for Total length of IP packet including header.

$$\text{Packet ki max size} = 2^{16}$$

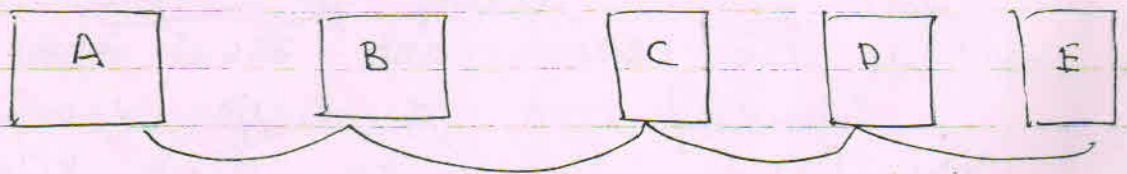
$$\text{aur data ki} = (2^{16} - 20) \text{ bytes.}$$

- But theoretical we have max datagram size 64K bytes (2^{16}) but in practically they are usually around 1500 bytes.
- Now next 16 bits for identification number.

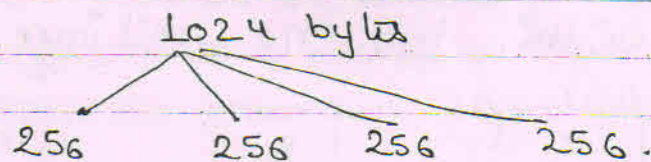
Identification matlab, when ultimate host divide segments into packets then it assign some identification number to each packet.

- Har packet ka identification number hota hai.

- Suppose A wants to send data of 10kb to host E, sabse pahle A ka data, segments mein divide hua then segment packet/data gram mein divide hua.
- Now source system (host) ki network layer har packet ko ek unique identification number assign karke hai.
- Packet ki size underline wire ki capacity ko dek kar decide ki jate hai.
- Now suppose B k pass gaye packet gaya, now suppose B se C se jane wali link ki underline capacity slow hai as compare to A and B. Toh packets k aur packets hogye for identification number jis packet ka job hai woh same rahega.
- Bhooliye aap packets ko kitana bhi divide karlo.
- But packets ki merging ultimate destination par hi hogey, bich mein packet ki assembly nahi hogi, bahi aapko achi bandwidth kyo na mil jaye.



- Now we know that ki network layer connectionless hai, toh identification number 5 wale packet k agar 5 parts hue hai toh konsa bhi part pahale aasakta hai, since connectionless layer hai to ultimate destination ko munge karke samay kaise pata Chalga ki yeh subpacket original ka pahale subpacket hai.
- To solve this problem, fragment offset came into existence, woh field 13 bit ka hota hai.
- Fragment offset ki help se hum yeh bata sakte hai ki y subpacket konsi packet ka konsa wala hai, pahala hai, dusra hai, thirsa hai like this.
- Fragment offset hame kya batata hai 1st byte of this packet is which byte of original packet.
- Suppose 1kb ka packet agar four parts mein divide hua hai



Toh pahale packet ka offset kitana hoga

① - - - 255

Answer

- Next wala ka kitana hoga 256 and so on.
- Sabse pahale murging karna hai toh kaha se karegey packet with fragment offset 0.
- Hum actual fragment offset ^{store} + nahi karte hai hum usko 8 se divide karke store karke hai

$$\text{Actual} = 8 * \text{Stored offset offset}$$

- Fragment offset hamesa 8 se divisible hona chahiye, agar nahi hai toh woh apka fragment offset banana yog nahi hai, bus last wala fragment ko chor k sab 8 se divisible hone chahiye.
- Nahi hai toh error aayega.

Ques:- How to find fragment offset of packet if current fragment offset is given -

$$\text{Fragment offset of next packet} = \text{Current} + \frac{(\text{Total Length} - (\text{IHL} * 4))}{8}$$

Ques:- If fragment offset of current packet is 8. Total length of packet is 268 and value of IHL is 5 then what will be the fragment offset of next packet??

Ans → $8 + (268 - 20) / 8 = 39 \leftarrow \text{Ans}$

Ans:→ If current packet is second packet with fragment offset = 8 what is the size of first packet??

Ans
↳



Total 64 byte data + 20 IP header
= 84 byte.

- Agar offset 8 se divisible nahi hai, toh something went wrong in the question data.
- Last fragment k fragment offset ko chack kar
same fragment offset 8 se divisible hona chahiye.
- Now question yeh kaise bata karuqy ki
yeh last fragment hai.
- MF bit :- More fragment :->
 - Agar MF = 1 then is fragment k baad is identification k aur bhi fragment exist karke hai.
 - Is packet ki MF bit 0 woh packet ka last fragment hoga woh original packet ka last fragment hoga.
 - DF → Don't fragment :-
 - Agar DF field is 1 then us packet k aur fragment nahi ho sakta.

- DF \rightarrow Source Router, Set karta hai.
- IP-Layer jaise jab bhi koi error aati hai woh ICMP handle karta hai.
- ICMP stands for Internet Control Message Protocol.
- DF = 1 set hai aur hamne us fragment ko aur divide karna h tabhi hum send kar sakte ho toh this situation is handle by ICMP, aur ICMP, ek error message "Unreachable destination" Source Router ko send karega.
- Now next 8 bit for TTL (Time to Live)
- \rightarrow TTL :- Every Router decrement TTL by one & Router on which TTL become 0, will discard the packet. (sabar pahale decrement then check)
- TTL measured in seconds (destination bhi decrement karta hai)
- To avoid the indefinite traversing TTL field is used
- As TTL is of 1 byte its value is 255, at most ek packet at most 255 Router ko cross kar sakte hai.
- Now next 8 bit for Protocol is field min highest layer i.e transport layer

Ka protocol konsa used karna hai woh
mahata hai.

- Then 16 bit for Header checksum.

↳ Har router checksum compute karta
hai aur check bhi karta hai. because TTL
change hoga to checksum bhi change hoga
har router k liye.

- Header checksum check bhi karta hai aur
laagay wale ka checksum bhi inkal k
sakta hai.

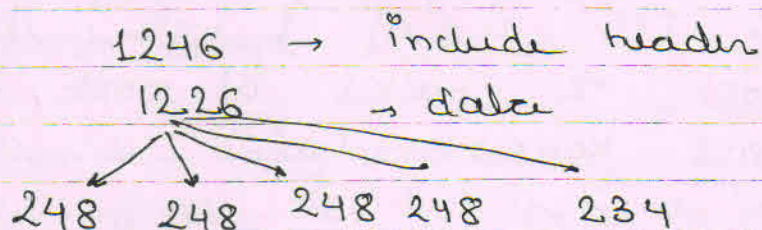
- If Header of the same layer is already
added then system first remove the header
and then divide the data and then again
add header to each divided packet.

Ques: Consider a 1246 byte packet is received
by router that has a PPP link to next
hop for the packet. The MTU (excluding
IP) for the PPP link is 255 byte.

Ans *

② How many fragments will be there?

Ans



5 fragments

{ 255 divisible nahi h 8 se
isliye 248 hoga }