

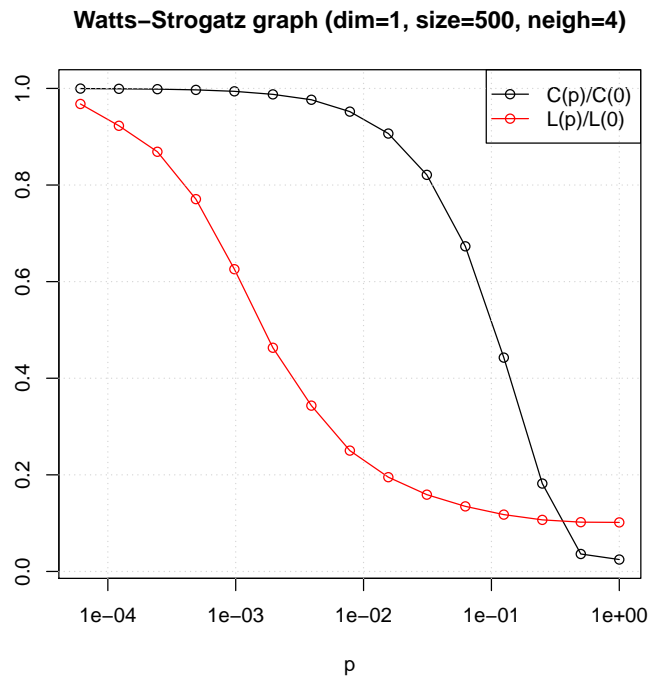
# Lab 1: Introduction to `igraph`

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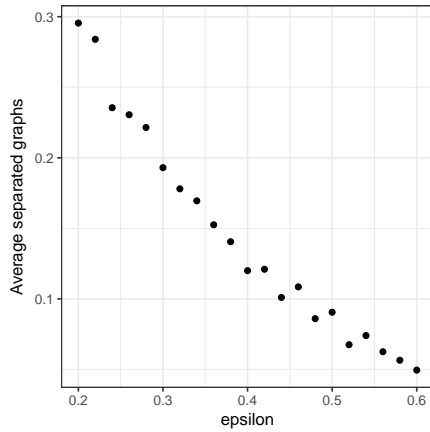
## 1 Watts–Strogatz model

The clustering coefficient  $C(p)$  and average shortest path  $L(p)$  are computed for a Watts–Strogatz graph with dimension 1, size 500, 4 neighbours and probability  $p$ . Each probability is computed as  $p = 2^{-i}$  with  $i \in \{0, 14\}$  (as a logarithmic scale is used). The mean values of 500 random graphs is performed to reduce the variance. Also, both are scaled between 0 and 1 to be compared using a graph with  $p = 0$  as  $C(p)/C(0)$  and  $L(p)/L(0)$ .

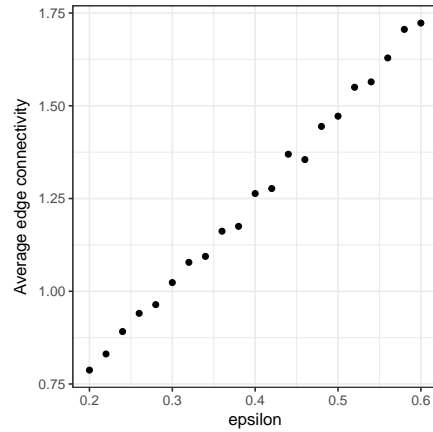


## 2 Erdős–Rényi model

The probability  $p$  of connecting two vertex is set to  $p = (1 + \epsilon) \ln(n)/n$  to keep the graph connected with high probability. In order to find a good  $\epsilon$  value, a set of random graphs are generated, with values of  $\epsilon \in [0.4, 0.6]$ . The edge connectivity  $e$  is measured in a set of  $r = 2000$  runs. The relative number of graphs with  $e = 0$  is plot against  $\epsilon$  in the figure 1a and the mean of  $e$  is plotted in 1b.



(a) The relative number of disconnected graphs against  $\epsilon$



(b) The mean of edge connectivity for graphs with different  $\epsilon$

The chosen value of  $\epsilon = 0.5$  keeps the number of disconnected graphs close to 10%, while the mean edge connectivity is close to 1.5. Finally, the average shortest path is plotted against the number of nodes  $n$  of a Erdős–Rényi graph. The mean values of 10 random graphs are computed to reduce the variance. The number of nodes is set up to 5000, to maintain a reasonable computing time.

