

2.

(9)

Plain text = 1 1 0 1 0 1 0 1

Key = 0 1 1 1 0 1 0 0 0 1

Key Generation:-(i) After  $P_0 \Rightarrow$  1 0 1 0 1 1 0 0 0 1 $L_0 = 1 0 1 0 1$  $R_0 = 1 0 0 0 1$ 

(ii) After 1-left shift

 $L_1 = 0 1 0 1 1$  $R_1 = 0 0 0 1 1$ 

(iii)

After  $P_1$ :  $P_1 (0 1 0 1 1 0 0 0 1 1) \Rightarrow 0 0 0 1 0 1 1 1$  $\therefore K_1 = 0 0 0 1 0 1 1 1$ 

(iv) After 2-left shift

 $L_2 = 0 1 1 0 1$  $R_2 = 0 1 1 0 0$ 

(v)

After  $P_2$ : 0 1 1 0 1 1 0 0 $K_2 = 0 1 1 0 1 1 0 0$ Encryption:-

(i) Plain text: 1 1 0 1 0 1 0 1

(ii) After IP = 1 1 0 1 1 1 0 0

 $L = 1 1 0 1$  $R = 1 1 0 0$ (iii)  $E/P(R) = 0 1 1 0 1 0 0 1$ 

(iv) 0 1 1 0 1 0 0 1

0 0 0 1 0 1 1 1 ( $K_1$ )

After XOR: 0 1 1 1 1 1 1 0



(v) Output of  $S_0(0111) = 00$   
 output of  $S_1(1110) = 00$

(vi) After  $P_4$  : 0 0 0 0

(vii)

0 0 0 0
1 1 0 1
XOR → 1 1 0 1

(viii) After 1 Round ⇒ 1 1 0 0 1 1 0 1  
 $L = 1 1 0 0$  ,  $R = 1 1 0 1$

(ix)  $E/P(R) = 1 1 1 0 1 0 1 1$

(x)

1 1 1 0 1 0 1 1
0 1 1 0 1 1 0 0
XOR → 1 0 0 0 0 1 1 1

(xi)

$S_0(1000) = 00$	}	0 0 1 1
$S_1(0111) = 11$		

(xii) After  $P_4$  : 0 1 1 0

(xiii)

0 1 1 0
1 1 0 0
XOR → 1 0 1 0

(xiv) Text After Round 2 1 0 1 0 1 1 0 1  
 $IP^{-1}(10101101) = 01110011$

∴ Cipher text = 0 1 1 1 0 0 1 1

(B) Plain text  $\rightarrow$  0 1 0 0 1 1 0 0  
 Key  $\rightarrow$  1 1 1 1 1 1 1 1

Key Generation  $\rightarrow$

After  $P_{10} \rightarrow$  1 1 1 1 1 1 1 1  
 $L \rightarrow$  1 1 1 1 ,  $R \rightarrow$  1 1 1 1

$\rightarrow$  left shift (1)

$L \rightarrow$  1 1 1 1 ,  $R \rightarrow$  1 1 1 1

$\rightarrow$  PS (1 1 1 1 1 1 1 1)  $\rightarrow$  1 1 1 1 1 1 1 1

$K_1 = 1 1 1 1 1 1$

$\rightarrow$  left shift (2)

$L = 1 1 1 1$  ,  $R = 1 1 1 1$

$\rightarrow$  PS (1 1 1 1 1 1 1 1)  $\rightarrow$  1 1 1 1 1 1 1 1

$K_2 = 1 1 1 1 1 1$

Encryption

- Plain text: 0 1 0 0 1 1 0 0  
 - IP(0 1 0 0 1 1 0 0) = 1 1 0 0 0 0 1 0

-  $L = 1 1 0 0$  ,  $R = 0 0 1 0$

-  $E(P(R)) = 0 0 0 1 0 1 0 0$

- 0 0 0 1 0 1 0 0

$\oplus$  1 1 1 1 1 1 1 1 (K<sub>1</sub>)

1 1 1 0 1 0 1 1



$$s_0(1110) = 11 \quad \left. \begin{array}{l} s_1(1011) = 01 \end{array} \right\} 1101$$

$$P_4(1101) = 1101$$

$$\rightarrow 1101 \oplus 1100 = 0001$$

$$\rightarrow \text{After Round 1} \rightarrow 00100001$$

$$L = 0010, \quad R = 0001$$

$$E(P|R) = 10000010$$

$$\begin{array}{r} 10000010 \\ \oplus 11111111 \quad K_2 \\ \hline 01111101 \\ \hline \end{array}$$

$$\left. \begin{array}{l} s_0(0111) = 00 \\ s_1(1101) = 00 \end{array} \right\} 0000$$

$$P_4(0000) = 0000$$

$$\begin{array}{r} 0000 \\ \oplus 0010 \\ \hline 0010 \end{array}$$

$$\text{Round 2} \rightarrow 00100001$$

$$IP^{-1}(00100001) \rightarrow 00100010$$

$$\therefore \text{Cipher text} \rightarrow 00100010$$



(C) Plain text  $\rightarrow$  0 0 0 0 0 0 0 0  
Key  $\rightarrow$  0 0 0 0 0 0 0 0 0 0

Key Generation  $\rightarrow$

- $P_0 = 0000000000$
- $L = 0000$  ,  $R = 0000$
- Left shift (L)  $\rightarrow$  0000  
 $\rightarrow$  (R)  $\rightarrow$  0000

- $P_8 \rightarrow 00000000$

$$K_1 = 00000000$$

- Left shift (2)  $\rightarrow$   
 $L = 0000$   
 $R = 0000$

- $P_8 : 00000000$

$$K_2 = 00000000$$

Encryption:

- Plain text = 00000000
- IP = 00000000
- $L = 0000$  ,  $R = 0000$
- $E(P(K)) = 00000000$

$$\begin{array}{r} 00000000 \\ \oplus 00000000 \quad (K) \\ \hline 00000000 \end{array}$$



$$\left. \begin{array}{l} s_0(0000) = 01 \\ s_1(0000) = 00 \end{array} \right\} 0100$$

$$P_4(0100) = 1000$$

$$\rightarrow 1000 \oplus 0000 \Rightarrow 1000$$

$$\rightarrow \text{Round 1: } 0000 \ 1000$$

$$L = 0000, \quad R = 1000$$

$$E/P(R) = 01000001$$

$$\begin{array}{r} \rightarrow \quad 01000001 \\ \oplus \quad 00000000 \quad (K_1) \\ \hline 01000001 \end{array}$$

$$\rightarrow \left. \begin{array}{l} s_0(0100) = 11 \\ s_1(0001) = 10 \end{array} \right\} 1110$$

$$\rightarrow P_4(1110) = 1011$$

$$\rightarrow 1011 \oplus 0000 = 1011$$

$$\rightarrow \text{After Round 2 } 10110000$$

$$\rightarrow IP^{-1}(10110000) = 11110000$$

$$\boxed{\text{Cipher text} = 11110000}$$



(1) Plain text  $\Rightarrow$  1 1 1 1 1 1 1

Key  $\Rightarrow$  1 1 1 1 1 1 1 1

Key word

Pl  $\Rightarrow$  1 1 1 1 1 1 1 1

L  $\Rightarrow$  1 1 1 1 , R = 1 1 1 1

$\Rightarrow$  Shift  $\Rightarrow$

L = 1 1 1 1 , R = 1 1 1 1

Pg = 1 1 1 1 1 1 1

$\therefore$   $R_1 = 1 1 1 1 1 1 1$

$\Rightarrow$  Shift 2  $\Rightarrow$

L = 1 1 1 1 , R = 1 1 1 1

$\Rightarrow$  Pg = 1 1 1 1 1 1 1

$R_2 = 1 1 1 1 1 1 1$

Encryption

- plaintext  $\Rightarrow$  1 1 1 1 1 1 1

- IP  $\Rightarrow$  1 1 1 1 1 1 1

- L  $\Rightarrow$  1 1 1 1 , R = 1 1 1 1

-  $F(P(R)) = 1 1 1 1 1 1 1$

-  $1 1 1 1 1 1 1 \oplus 1 1 1 1 1 1 1 \Rightarrow 0 0 0 0 0 0 0$

-  $S_0(0000) = 01$  }  $0100$   
 $S_1(0000) = 00$  }

- P4 = 1000

-  $1000 \oplus 1111 \Rightarrow 0111$

- After Round 1  $\Rightarrow$  1 1 1 1 0 1 1

- L = 1 1 1 1 , R = 0 1 1 1

-  $F(P(R)) = 10 1 1 1 1 0$

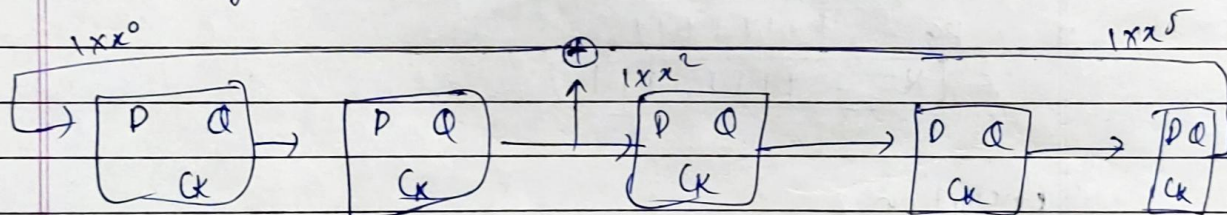
-  $10 1 1 1 1 1 0 \oplus 1 1 1 1 1 1 1 \Rightarrow 01000001$

-  $S_0(0100) = 11$  }  $1110$   
 $S_1(0001) = 10$  }

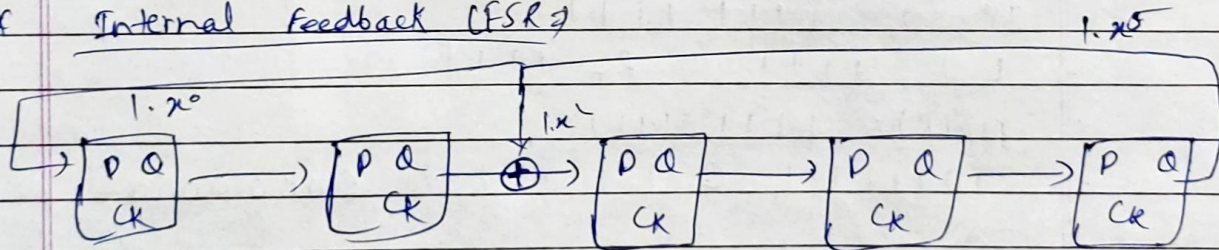
- $P_4 = 1011$
- $1011 \oplus 1111 \neq 0100$
- After Round 2  $\rightarrow 01000111$
- $IP^{-1} = 00001111$
- $\therefore$  Cipher text = 00001111

1. Polynomial  $\Rightarrow x^5 + x^2 + 1$   
degree  $(n) = 5$

\* External feedback LFSR  $\Rightarrow$



\* Internal Feedback LFSR  $\Rightarrow$



$$\text{Periods} \Rightarrow 2^n - 1 = 2^5 - 1$$

$$= 32 - 1$$

$$\boxed{\text{Periods} = 31}$$