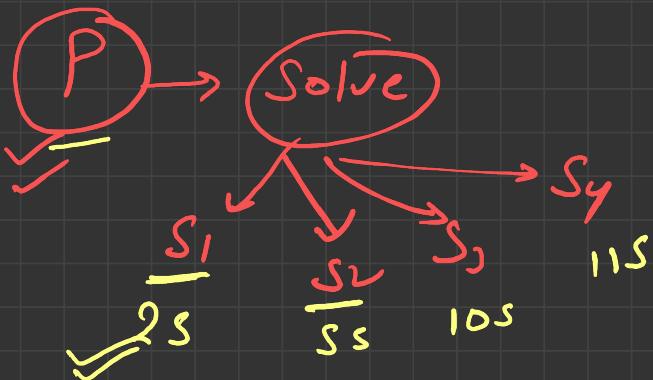
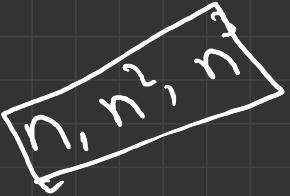


## Time & Space Complexity



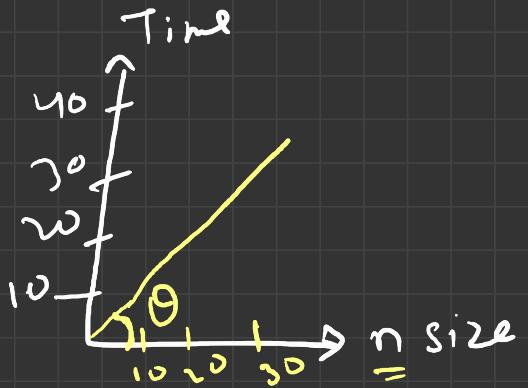
Q. Sum of n number

```
for (i=1 ; i<=n ; i++) {  
    sum += i;  
}
```

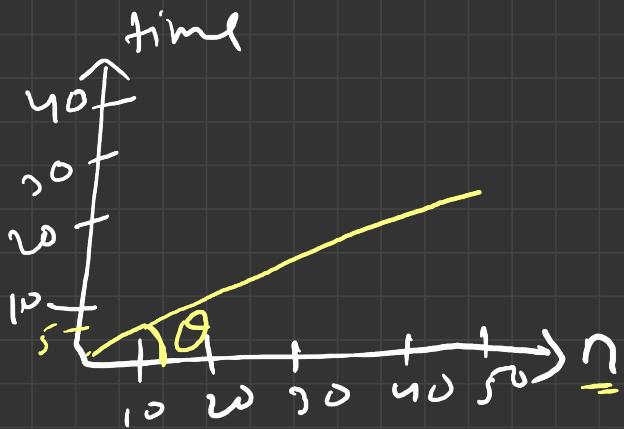
<p>New</p> <p>64 Gb Ram ✓</p> <p><del>2 sec</del> ✓</p> <p><del>10 sec</del></p> <p></p>	<p><math>n = 10 \uparrow</math></p> <p><math>\underline{n = 100000} \uparrow</math></p>	<p>Old</p> <p>4 Gb Ram <math>n = 10</math></p> <p>10 sec ✓</p> <p>100 sec</p>
---	---	---

Time Complexity  $\neq$  Taken taken

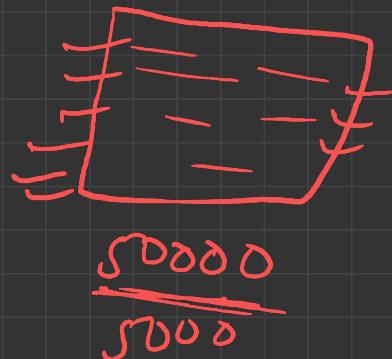
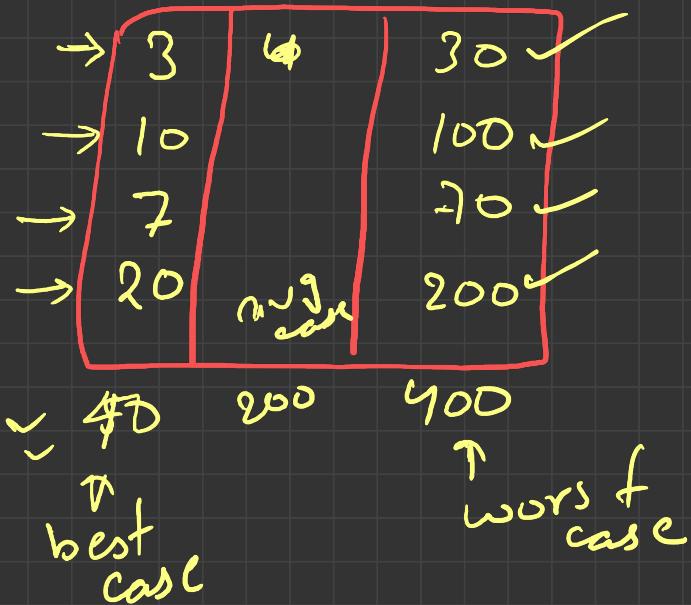
Time Complexity:- It is the total time taken by an algorithm to run as a function of length of the input



old



# Best case, Worst Case & Avg case



# 3 terminology?

Big(O)

worst case

$O(n)$

$O(n^2)$

$O(n^3)$

$\Omega(\omega)$

Best  
case

$\Omega(n)$

$\Theta$

Avg case

$\Theta(n)$

$(i=1, i \leq n; i++) \{$

① ③ cout << "

}

$$\begin{array}{r} 1+2+\dots+n \\ 4+3+\dots+3 \\ \hline \end{array}$$

$$\underbrace{1+3+3+\dots+3}_{\text{Sn}+1}$$

constant term  $\rightarrow$  ignore  
in  $n$  term we take highest value

$O(n)$

$$\cancel{8n^3} + \cancel{3n^2} + \cancel{48n} + 64$$

$$\frac{\cancel{n^3} + n^2 + n}{\checkmark \quad \times \quad \times}$$

$\sim^3$

$$\underline{n^3} + \log n$$

```
for (i=0; i<n; i++) {
```

```
    if (arr[i] == x) {
```

```
        cout << " this is result"
```

```
} break;
```

$$\text{avg} = \frac{1+2+3+\dots+n}{n}$$

$$\frac{n \cdot x(n+1)}{\cancel{n}}$$

$$n+1 \\ n$$

$$O(\underline{10}) \\ O(1)$$

5	6	8	7	12
↑		↑		

x=5

best  
case

$$\checkmark \text{count} \ll \frac{n(n+1)}{2}$$

$$O(1)$$

for ( i = 1; i <= n; i++ ) { } }  $\rightarrow n$  }  $n^2$

for ( j = 1; j <= n; j++ ) { } }  $\rightarrow n$

cout << \*

}

}

i = 1

j = 1 to n

i = 2

j = 1 to n

i = 3

j = 1 to n

\* print. n times

$$n + n + n + n \dots + n = n^2$$
$$(n)n = n^2$$

O(n<sup>2</sup>)

for ( $i=1$ ;  $i \leq n$ ;  $i++$ )

  for ( $j=1$ ;  $j \leq i$ ;  $j++$ )

cout << 'x'

$i=2$

$j=1 \text{ to } 2$

2 to print

$i=3$

$j=1 \text{ to } 3$

3 x

$i=1$

$j=1 \text{ to } 1$

x print L time

$i=n$

$j=1 \text{ to } n$

n x

1 + 2 + 3 + ... + n

$$\frac{n(n+1)}{2} = \frac{n^2+n}{2}$$

$$\approx \frac{n^2+n}{2} = O(n^2)$$

Q1 for ( $i = 1$ ;  $i \leq n$ ;  $i = i \times 2$ ) {

    cout << '\*'  
}

Q2 for ( $i = 1$ ;  $i \leq n$ ;  $i++$ ) {

    for ( $j = 1$ ;  $j \leq i^2$ ;  $j++$ ) {

        for ( $k = 1$ ;  $k \leq n/2$ ;  $k++$ ) {

            cout << '\*'  
    }

for (i=1 ; i<=n ; i++) { } ✓

for(j=1 ; j<=m ; j++) { } → m ✓

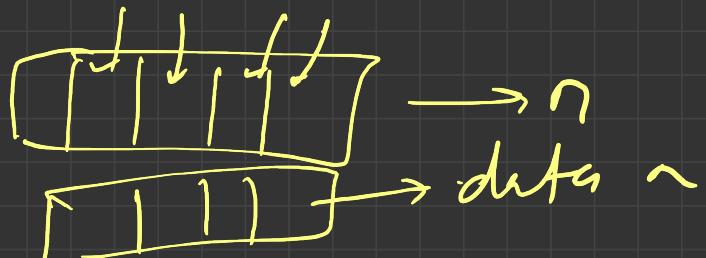
Count C('\*)

O(n+m)

Space Complexity :- It is the amount of space taken by an algorithm as a function ( $f_n$ ) of length of i/p.

→ Auxiliary Space

→ Total Space Complexity



$$\text{Aux} \rightarrow n \rightarrow O(n)$$

$$\text{Total} \rightarrow n + n = 2n \rightarrow O(n)$$

for (i=1; i < n; i += t)      Aux  $\rightarrow i = O(1)$

cout << 'x'

$$\text{Total} = 1 + i = O(1)$$

