

Markov Chains

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Markov chains. Definition and examples

Chapman Kolmogorov equations

Gambler's ruin problem

Queues in communication networks: Transition probabilities

Classes of States

Limiting distributions

Ergodicity

Queues in communication networks: Limit probabilities

- ▶ Consider time index $n = 0, 1, 2, \dots$ & time dependent random state X_n
- ▶ State X_n takes values on a countable number of states
 - ▶ In general denotes states as $i = 0, 1, 2, \dots$
 - ▶ Might change with problem
- ▶ Denote the history of the process $\mathbf{X}_n = [X_n, X_{n-1}, \dots, X_0]^T$
- ▶ Denote stochastic process as $\mathbf{X}_{\mathbb{N}}$
- ▶ The stochastic process $X_{\mathbb{N}}$ is a Markov chain (MC) if

$$P [X_{n+1} = j \mid X_n = i, \mathbf{X}_{n-1}] = P [X_{n+1} = j \mid X_n = i] = P_{ij}$$

- ▶ Future depends only on current state X_n

- ▶ Process's history \mathbf{X}_{n-1} irrelevant for future evolution of the process
- ▶ Probabilities P_{ij} are constant for all times (time invariant)
- ▶ From the definition we have that for arbitrary m

$$P[X_{n+m} | X_n, \mathbf{X}_{n-1}] = P[X_{n+m} | X_n]$$

- ▶ X_{n+m} depends only on X_{n+m-1} , which depends only on X_{n+m-2} ,
... which depends only on X_n
- ▶ Since P_{ij} 's are probabilities they're positive and sum up to 1

$$P_{ij} \geq 0 \quad \sum_{j=1}^{\infty} P_{ij} = 1$$

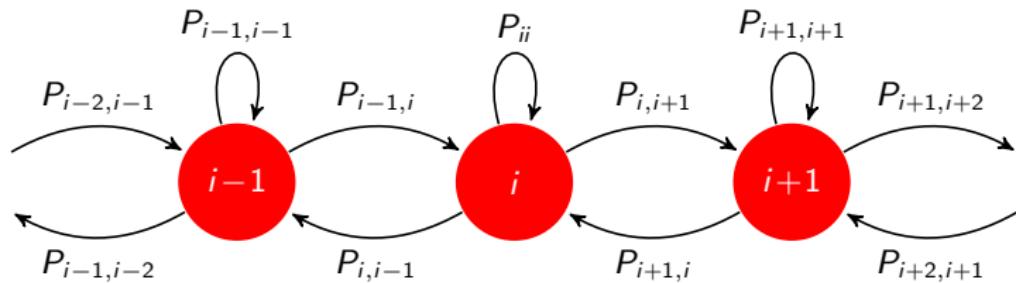
Matrix representation

- ▶ Group transition probabilities P_{ij} in a “matrix” \mathbf{P}

$$\mathbf{P} := \begin{pmatrix} P_{00} & P_{01} & P_{02} & \dots \\ P_{10} & P_{11} & P_{12} & \dots \\ \vdots & \vdots & \vdots & \vdots \\ P_{i0} & P_{i1} & P_{i2} & \dots \\ \vdots & \vdots & \vdots & \ddots \end{pmatrix}$$

- ▶ Not really a matrix if number of states is infinite

- ▶ A graph representation is also used

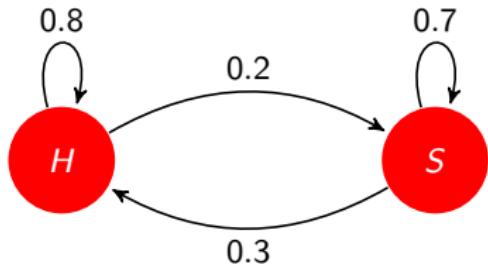


- ▶ Useful when number of states is infinite

Example: Happy - Sad

- ▶ I can be happy ($X_n = 0$) or sad ($X_n = 1$).
- ▶ Happiness tomorrow affected by happiness today only
- ▶ Model as Markov chain with transition probabilities

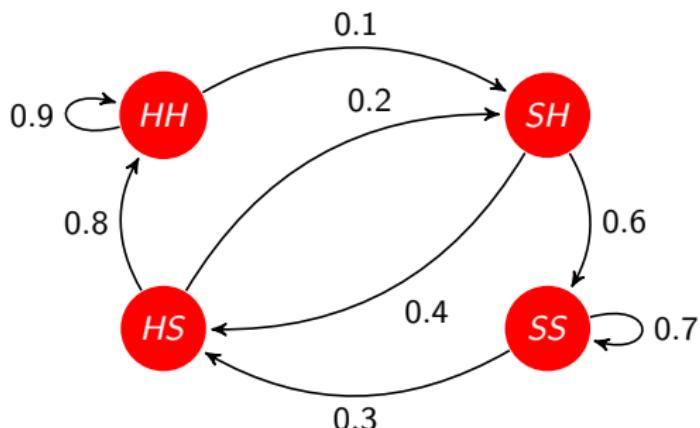
$$\mathbf{P} := \begin{pmatrix} 0.8 & 0.2 \\ 0.3 & 0.7 \end{pmatrix}$$



- ▶ Inertia \Rightarrow happy or sad today, likely to stay happy or sad tomorrow ($P_{00} = 0.8$, $P_{11} = 0.7$)
- ▶ But when sad, a little less likely so ($P_{00} > P_{11}$)

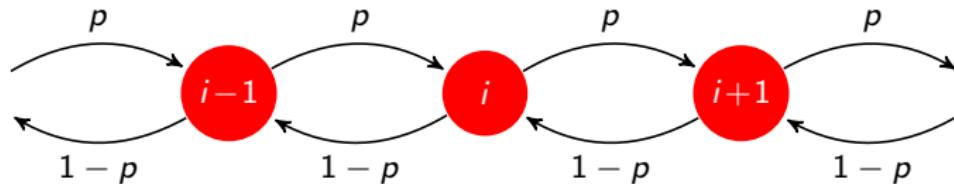
- ▶ Happiness tomorrow affected by today and yesterday
- ▶ Define double states HH (happy-happy), HS (happy-sad), SH, SS
- ▶ Only some transitions are possible
 - ▶ HH and SH can only become HH or HS
 - ▶ HS and SS can only become SH or SS

$$\mathbf{P} := \begin{pmatrix} 0.9 & 0.1 & 0 & 0 \\ 0 & 0 & 0.4 & 0.6 \\ 0.8 & 0.2 & 0 & 0 \\ 0 & 0 & 0.3 & 0.7 \end{pmatrix}$$



- ▶ More time happy or sad increases likelihood of staying happy or sad
- ▶ **State augmentation** \Rightarrow Capture longer time memory

- ▶ Step to the right with probability p , to the left with prob. $(1-p)$



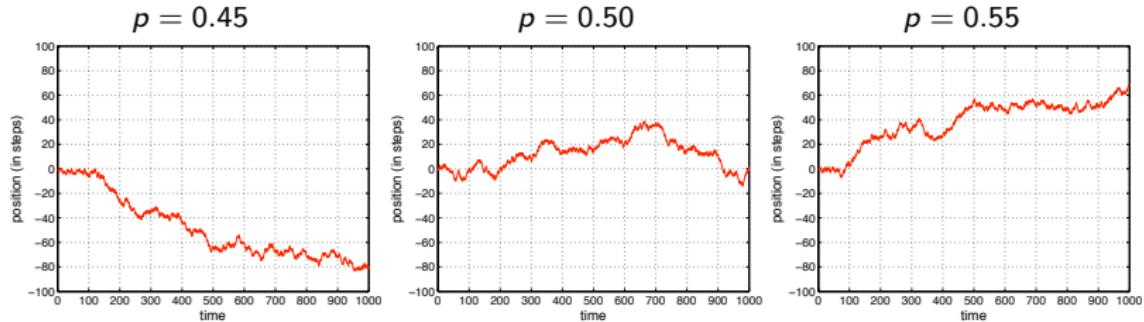
- ▶ States are $0, \pm 1, \pm 2, \dots$, number of states is infinite
- ▶ Transition probabilities are

$$P_{i,i+1} = p, \quad P_{i,i-1} = 1 - p,$$

- ▶ $P_{ij} = 0$ for all other transitions

Random (drunkard's) walk - continued

- ▶ Random walks behave differently if $p < 1/2$, $p = 1/2$ or $p > 1/2$



- ▶ With $p > 1/2$ diverges to the right (grows unbounded almost surely)
- ▶ With $p < 1/2$ diverges to the left
- ▶ With $p = 1/2$ always come back to visit origin (almost surely)
- ▶ Because number of states is infinite we can have all states transient
 - ▶ They are not revisited after some time (more later)

Two dimensional random walk

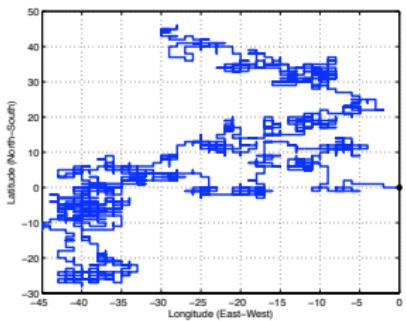
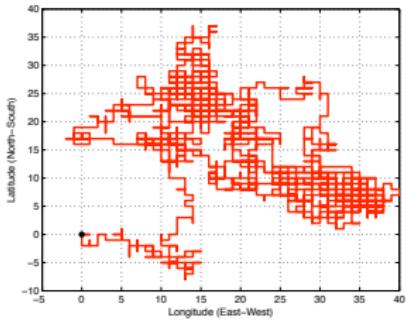
- ▶ Take a step in random direction East, West, South or North
⇒ E, W, S, N chosen with equal probability
- ▶ States are pairs of coordinates (x, y)
 - ▶ $x = 0, \pm 1, \pm 2, \dots$ and $y = 0, \pm 1, \pm 2, \dots$
- ▶ Transition probabilities are not zero only for points adjacent in the grid

$$P[x(t+1) = i+1, y(t+1) = j \mid x(t) = i, y(t) = j] = \frac{1}{4}$$

$$P[x(t+1) = i-1, y(t+1) = j \mid x(t) = i, y(t) = j] = \frac{1}{4}$$

$$P[x(t+1) = i, y(t+1) = j+1 \mid x(t) = i, y(t) = j] = \frac{1}{4}$$

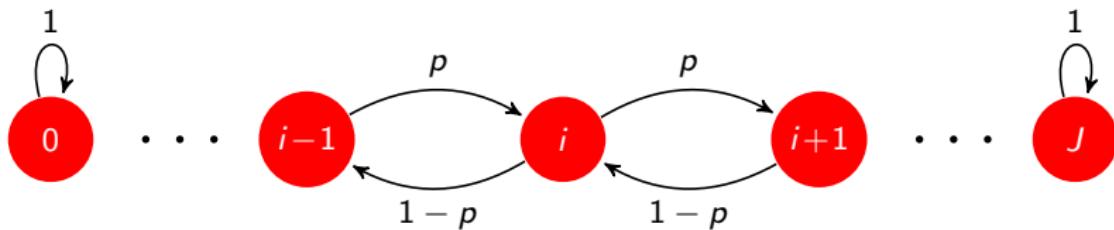
$$P[x(t+1) = i, y(t+1) = j-1 \mid x(t) = i, y(t) = j] = \frac{1}{4}$$



- ▶ Some random facts of life for equiprobable random walks
- ▶ In one and two dimensions probability of returning to origin is 1
- ▶ Will almost surely return home
- ▶ In more than two dimensions, probability of returning to origin is less than 1
- ▶ In three dimensions probability of returning to origin is 0.34
- ▶ Then 0.19, 0.14, 0.10, 0.08, ...

Random walk with borders (gambling)

- ▶ As a random walk, but stop moving when $i = 0$ or $i = J$
 - ▶ Models a gambler that stops playing when ruined, $X_n = 0$
 - ▶ Or when reaches target gains $X_n = J$



- ▶ States are $0, 1, \dots, J$. Finite number of states (J). Transition probs.

$$P_{i,i+1} = p, \quad P_{i,i-1} = 1 - p, \quad P_{00} = 1, \quad P_{JJ} = 1$$

- ▶ $P_{ij} = 0$ for all other transitions
- ▶ States 0 and J are called absorbing. Once there stay there forever
- ▶ The rest are transient states. Visits stop almost surely

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- ▶ What can be said about multiple transitions ?
- ▶ Transition probabilities between two time slots

$$P_{ij}^2 := P [X_{m+2} = j \mid X_m = i]$$

- ▶ Probabilities of X_{m+n} given X_n \Rightarrow n -step transition probabilities

$$P_{ij}^n := P [X_{m+n} = j \mid X_m = i]$$

- ▶ Relation between n -step, m -step and $(m+n)$ -step transition probs.
 - ▶ Write P_{ij}^{m+n} in terms of P_{ij}^m and P_{ij}^n
- ▶ All questions answered by Chapman-Kolmogorov's equations

- ▶ Start considering transition probs. between two time slots

$$P_{ij}^2 = P [X_{n+2} = j \mid X_n = i]$$

- ▶ Using the theorem of total probability

$$P_{ij}^2 = \sum_{k=1}^{\infty} P [X_{n+2} = j \mid X_{n+1} = k, X_n = i] P [X_{n+1} = k \mid X_n = i]$$

- ▶ In the first probability, conditioning on $X_n = i$ is unnecessary. Thus

$$P_{ij}^2 = \sum_{k=1}^{\infty} P [X_{n+2} = j \mid X_{n+1} = k] P [X_{n+1} = k \mid X_n = i]$$

- ▶ Which by definition yields

$$P_{ij}^2 = \sum_{k=1}^{\infty} P_{kj} P_{ik}$$

- ▶ Identical argument can be made (condition on X_0 to simplify notation, possible because of time invariance)

$$P_{ij}^{m+n} = P [X_{n+m} = j \mid X_0 = i]$$

- ▶ Use theorem of total probability, remove unnecessary conditioning and use definitions of n -step and m -step transition probabilities

$$P_{ij}^{m+n} = \sum_{k=1}^{\infty} P [X_{m+n} = j \mid X_m = k, X_0 = i] P [X_m = k \mid X_0 = i]$$

$$P_{ij}^{m+n} = \sum_{k=1}^{\infty} P [X_{m+n} = j \mid X_m = k] P [X_m = k \mid X_0 = i]$$

$$P_{ij}^{m+n} = \sum_{k=1}^{\infty} P_{kj}^n P_{ik}^m$$

- ▶ Chapman Kolmogorov is intuitive. Recall

$$P_{ij}^{m+n} = \sum_{k=1}^{\infty} P_{kj}^n P_{ik}^m$$

- ▶ Between times 0 and $m + n$ time m occurred
- ▶ At time m , the chain is in some state $X_m = k$
 - ⇒ P_{ik}^m is the probability of going from $X_0 = i$ to $X_m = k$
 - ⇒ P_{kj}^n is the probability of going from $X_m = k$ to $X_{m+n} = j$
 - ⇒ Product $P_{ik}^m P_{kj}^n$ is then the probability of going from $X_0 = i$ to $X_{m+n} = j$ passing through $X_m = k$ at time m
- ▶ Since any k might have occurred sum over all k

- ▶ Define matrices $\mathbf{P}^{(m)}$ with elements P_{ij}^m , $\mathbf{P}^{(n)}$ with elements P_{ij}^n and $\mathbf{P}^{(m+n)}$ with elements P_{ij}^{m+n}
- ▶ $\sum_{k=1}^{\infty} P_{kj}^n P_{ik}^m$ is the (i, j) -th element of matrix product $\mathbf{P}^{(m)} \mathbf{P}^{(n)}$
- ▶ Chapman Kolmogorov in matrix form

$$\mathbf{P}^{(m+n)} = \mathbf{P}^{(m)} \mathbf{P}^{(n)}$$

- ▶ Matrix of $(n + m)$ -step transitions is product of n -step and m -step

- ▶ For $m = n = 1$ (2-step transition probabilities) matrix form is

$$\mathbf{P}^{(2)} = \mathbf{P}\mathbf{P} = \mathbf{P}^2$$

- ▶ Proceed recursively backwards from n

$$\mathbf{P}^{(n)} = \mathbf{P}^{(n-1)}\mathbf{P} = \mathbf{P}^{(n-2)}\mathbf{P}\mathbf{P} = \dots = \mathbf{P}^n$$

- ▶ Have proved the following

Theorem

The matrix of n -step transition probabilities $P^{(n)}$ is given by the n -th power of the transition probability matrix \mathbf{P} . i.e.,

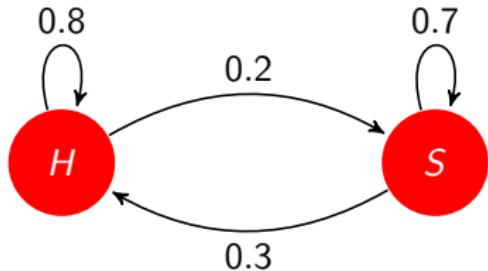
$$\mathbf{P}^{(n)} = \mathbf{P}^n$$

Henceforth we write \mathbf{P}^n

Example: Happy-Sad

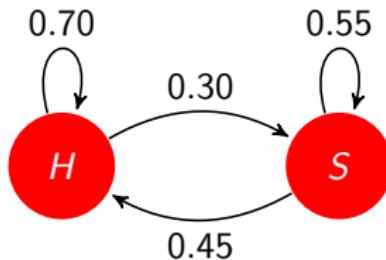
- ▶ Happiness transitions in one day (not the same as earlier example)

$$\mathbf{P} := \begin{pmatrix} 0.8 & 0.2 \\ 0.3 & 0.7 \end{pmatrix}$$



- ▶ Transition probabilities between today and the day after tomorrow?

$$\mathbf{P}^2 := \begin{pmatrix} 0.70 & 0.30 \\ 0.45 & 0.55 \end{pmatrix}$$



- ▶ ... After a week and after a month

$$\mathbf{P}^7 := \begin{pmatrix} 0.6031 & 0.3969 \\ 0.5953 & 0.4047 \end{pmatrix} \quad \mathbf{P}^{30} := \begin{pmatrix} 0.6000 & 0.4000 \\ 0.6000 & 0.4000 \end{pmatrix}$$

- ▶ Matrices \mathbf{P}^7 and \mathbf{P}^{30} almost identical $\Rightarrow \lim_{n \rightarrow \infty} \mathbf{P}^n$ exists
 - ▶ Note that this is a regular limit
- ▶ After a month transition from H to H with prob. 0.6 and from S to H also 0.6
- ▶ State becomes independent of initial condition
- ▶ Rationale: 1-step memory \Rightarrow initial condition eventually forgotten

- ▶ All probabilities so far are conditional, i.e., $P[X_n = j | X_0 = i]$
- ▶ Want unconditional probabilities $p_j(n) := P[X_n = j]$
- ▶ Requires specification of initial conditions $p_i(0) := P[X_0 = i]$
- ▶ Using theorem of total probability and definitions of P_{ij}^n and $p_j(n)$

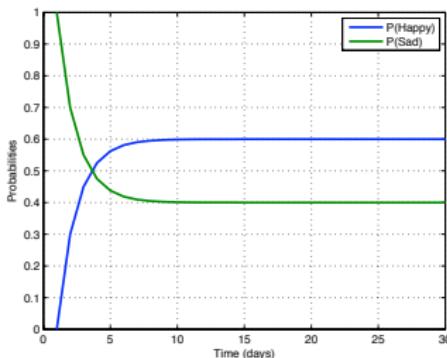
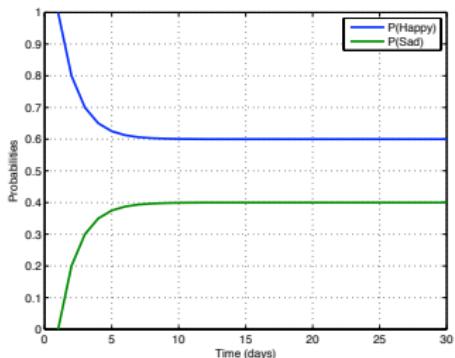
$$\begin{aligned} p_j(n) &:= P[X_n = j] = \sum_{i=1}^{\infty} P[X_n = j | X_0 = i] P[X_0 = i] \\ &= \sum_{i=1}^{\infty} P_{ij}^n p_i(0) \end{aligned}$$

- ▶ Or in matrix form (define vector $\mathbf{p}(n) := [p_1(n), p_2(n), \dots]^T$)

$$\mathbf{p}(n) = \mathbf{P}^{nT} \mathbf{p}(0)$$

Example: Happy-Sad

- ▶ Transition probability matrix $\Rightarrow \mathbf{P} := \begin{pmatrix} 0.8 & 0.2 \\ 0.3 & 0.7 \end{pmatrix}$
- $\mathbf{p}(0) = [1, 0]$
- $\mathbf{p}(0) = [0, 1]$



- ▶ For large n probabilities $\mathbf{p}(t)$ are independent of initial state $\mathbf{p}(0)$

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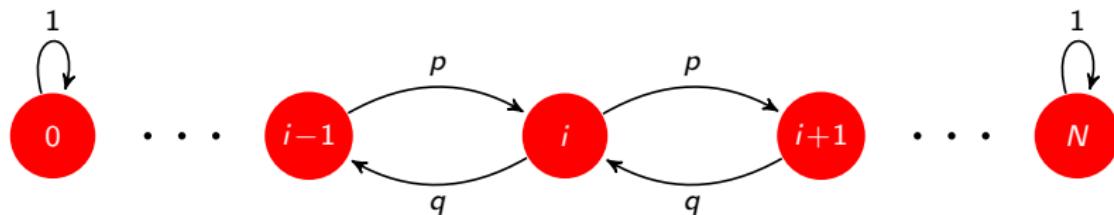
Queues in communication networks: Limit probabilities

- ▶ You place $\$1$ bets,
 - (a) With probability p you gain $\$1$ and
 - (b) With probability $q = (1 - p)$ you loose your $\$1$ bet
- ▶ Start with an initial wealth of $\$i_0$
- ▶ Define bias factor $\alpha := q/p$
 - ▶ If $\alpha > 1$ more likely to loose than win (biased against gambler)
 - ▶ $\alpha < 1$ favors gambler (more likely to win than loose)
 - ▶ $\alpha = 1/2$ game is fair
- ▶ You keep playing until
 - (a) You go broke (loose all your money)
 - (b) You reach a wealth of $\$N$
- ▶ Prob. S_i of reaching $\$N$ before going broke for initial wealth $\$i$?
 - ▶ S stands for success

- Model as Markov chain $X_{\mathbb{N}}$. Transition probabilities

$$P_{i,i+1} = p, \quad P_{i,i-1} = q, \quad P_{00} = P_{NN} = 1$$

- Realizations $x_{\mathbb{N}}$. Initial state = initial wealth = i_0



- States 0 and N said **absorbing**. Eventually end up in one of them
- Remaining states said **transient** (visits eventually stop)
- Being absorbing states says something about the **limit wealth**

$$\lim_{n \rightarrow \infty} x_n = 0, \text{ or } \lim_{n \rightarrow \infty} x_n = N, \Rightarrow S_i := P \left(\lim_{n \rightarrow \infty} X_n = N \mid X_0 = i \right)$$

- ▶ Prob. S_i of successful betting run depends on current state i only
- ▶ We can relate probabilities of SBR from adjacent states

$$S_i = S_{i+1}P_{i,i+1} + S_{i-1}P_{i,i-1} = S_{i+1}p + S_{i-1}q$$

- ▶ Recall $p + q = 1$. Reorder terms

$$p(S_{i+1} - S_i) = q(S_i - S_{i-1})$$

- ▶ Recall definition of bias $\alpha = q/p$

$$S_{i+1} - S_i = \alpha(S_i - S_{i-1})$$

- ▶ If current state is 0 then $S_i = S_0 = 0$. Can write

$$S_2 - S_1 = \alpha(S_1 - S_0) = \alpha S_1$$

- ▶ Substitute this in the expression for $S_3 - S_2$

$$S_3 - S_2 = \alpha(S_2 - S_1) = \alpha^2 S_1$$

- ▶ Apply recursively backwards from $S_i - S_{i-1}$

$$S_i - S_{i-1} = \alpha(S_{i-1} - S_{i-2}) = \dots = \alpha^{i-1} S_1$$

- ▶ Sum up all of the former to obtain

$$S_i - S_1 = S_1 \left(\alpha + \alpha^2 + \dots + \alpha^{i-1} \right)$$

- ▶ The latter can be written as a geometric series

$$S_i = S_1 \left(1 + \alpha + \alpha^2 + \dots + \alpha^{i-1} \right)$$

- ▶ Geometric series can be summed. Assuming $\alpha \neq 1$

$$S_i = \frac{1 - \alpha^i}{1 - \alpha} S_1$$

- ▶ Write for $i = 1$. When in state N , $S_N = 1$

$$1 = S_N = \frac{1 - \alpha^N}{1 - \alpha} S_1$$

- ▶ Compute S_1 from latter and substitute into expression for S_i

$$S_i = \frac{1 - \alpha^i}{1 - \alpha^N}$$

- ▶ For $\alpha = 1 \Rightarrow S_i = iS_1$, $1 = S_N = NS_1 \Rightarrow S_i = (i/N)$

- ▶ Consider exit bound N arbitrarily large.
- ▶ For $\alpha \geq 1$, $S_i \approx (\alpha^i - 1)/\alpha^N \rightarrow 0$
- ▶ If win prob. does not exceed loose probability will almost surely lose all money
- ▶ For $\alpha < 1$, $P_i \rightarrow 1 - \alpha^i$
- ▶ If win prob. exceeds loose probability might win
- ▶ If initial wealth i sufficiently high, will most likely win
⇒ Which explains what we saw on first lecture and homework

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Queues in communication networks: Limit probabilities

- ▶ Communication systems goal
 - ⇒ Move packets from generating sources to intended destinations
- ▶ Between arrival and departure we hold packets in a memory buffer
- ▶ Want to design buffers appropriately

- ▶ Time slotted in intervals of duration Δt
- ▶ n -th slot between times $n\Delta t$ and $(n + 1)\Delta t$
- ▶ Average arrival rate is $\bar{\lambda}$ packets per unit time
- ▶ During slot of duration Δt probability of **packet arrival** is $\lambda = \bar{\lambda}\Delta t$
- ▶ Packets are transmitted (depart) at a rate of $\bar{\mu}$ packets per unit time
- ▶ During interval Δt **probability of packet departure** is $\mu = \bar{\mu}\Delta t$
- ▶ Assume no simultaneous arrival and departure (no concurrence)
 - ▶ Reasonable for small Δt (μ and λ are likely small)

- ▶ q_n denotes number of packets in queue in n -th time slot
- ▶ \mathbb{A}_n = nr. of packet arrivals, \mathbb{D}_n = nr. of departures (during n -th slot)
- ▶ If there are no packets in queue $q_n = 0$ then there are no departures
- ▶ Queue length at time $n + 1$ can be written as

$$q_{n+1} = q_n + \mathbb{A}_n, \quad \text{if } q_n = 0$$

- ▶ If $q_n > 0$, departures and arrivals may happen

$$q_{n+1} = [q_n + \mathbb{A}_n - \mathbb{D}_n]^+, \quad \text{if } q_n > 0$$

- ▶ $\mathbb{A}_n \in \{0, 1\}$, $\mathbb{D}_n \in \{0, 1\}$ and either $\mathbb{A}_n = 1$ or $\mathbb{D}_n = 1$ but not both
- ▶ Arrival and departure probabilities are

$$P[\mathbb{A}_n = 1] = \lambda, \quad P[\mathbb{D}_n = 1] = \mu$$

- ▶ Future queue lengths depend on current length only
- ▶ Probability of queue length increasing

$$P[q_{n+1} = i+1 \mid q_n = i] = P[\mathbb{A}_n = 1] = \lambda, \quad \text{for all } i$$

- ▶ Queue length might decrease only if $q_n > 0$. Probability is

$$P[q_{n+1} = i-1 \mid q_n = i] = P[\mathbb{D}_n = 1] = \mu, \quad \text{for all } i > 0$$

- ▶ Queue length stays the same if it neither increases nor decreases

$$P[q_{n+1} = i \mid q_n = i] = 1 - \lambda - \mu, \quad \text{for all } i > 0$$

$$P[q_{n+1} = 0 \mid q_n = 0] = 1 - \lambda$$

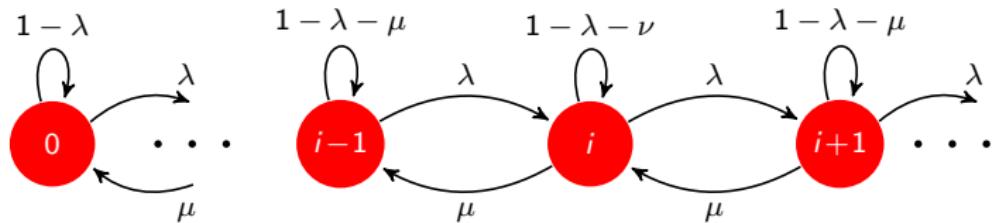
- ▶ No departures when $q_n = 0$ explain second equation

Queue as a Markov chain

- ▶ MC with states $0, 1, 2, \dots$. Identify states with queue lengths
- ▶ Transition probabilities for $i \neq 0$ are

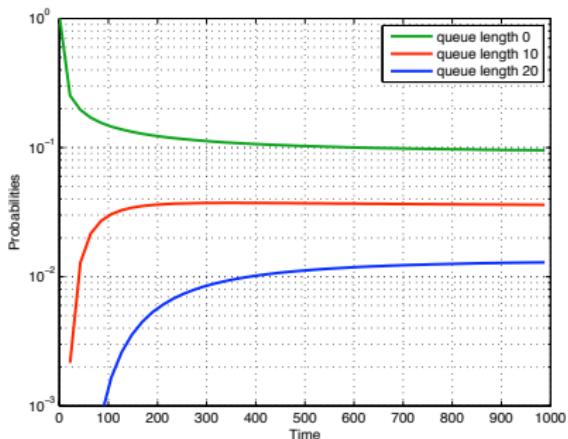
$$P_{i,i-1} = \lambda, \quad P_{i,i} = 1 - \lambda - \mu, \quad P_{i,i+1} = \mu$$

- ▶ For $i = 0$ $P_{0,0} = 1 - \lambda$ and $P_{01} = \lambda$



Numerical example: Probability propagation

- ▶ Build matrix \mathbf{P} truncating at maximum queue length $L = 100$
- ▶ Arrival rate $\lambda = 0.3$. Departure rate $\mu = 0.33$
- ▶ Initial probability distribution $\mathbf{p}(0) = [1, 0, 0, \dots]^T$ (queue empty)



- ▶ Propagate probabilities with product $\mathbf{P}^n \mathbf{p}(0)$
- ▶ Probabilities obtained are

$$P [q_n = i \mid q_0 = 0] = p_i(n)$$

- ▶ A few i 's (0, 10, 20) shown
- ▶ Probability of empty queue ≈ 0.1 .
- ▶ Occupancy decrease with index

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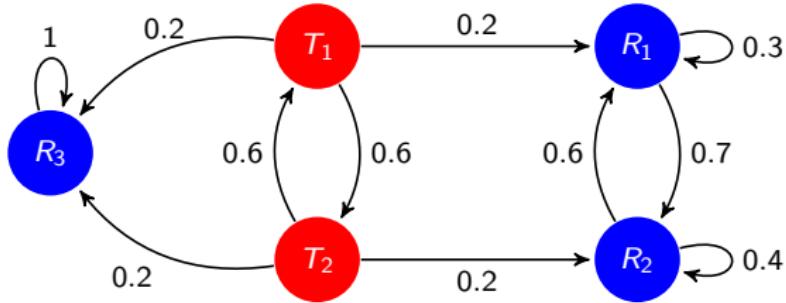
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Queues in communication networks: Limit probabilities

- ▶ States of a MC can be recurrent or transient
- ▶ **Transient states** might be visited at the beginning but eventually visits stop
- ▶ Almost surely, $X_n \neq i$ for n sufficiently large (qualifications needed)
- ▶ Visits to **recurrent states** keep happening forever
- ▶ Fix arbitrary m
- ▶ Almost surely, $X_n = i$ for some $n \geq m$ (qualifications needed)



- ▶ Let f_i be the probability that starting at i , MC ever reenters state i

$$f_i := \mathbb{P} \left[\bigcup_{n=1}^{\infty} X_n = i \mid X_0 = i \right] = \mathbb{P} \left[\bigcup_{n=m+1}^{\infty} X_n = i \mid X_m = i \right]$$

- ▶ State i is recurrent if $f_i = 1$
- ▶ Process reenters i again and again (almost surely). Infinitely often
- ▶ State i is transient if $f_i < 1$
- ▶ Positive probability $(1 - f_i)$ of never coming back to i

Recurrent states example

- ▶ State R_3 is recurrent because $P[X_1 = R_3 \mid X_0 = R_3] = 1$

- ▶ State R_1 is recurrent because

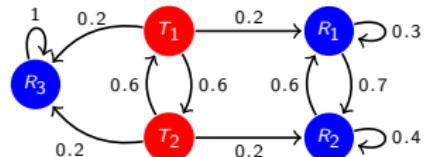
$$P[X_1 = R_1 \mid X_0 = R_1] = 0.3$$

$$P[X_2 = R_1, X_1 \neq R_1 \mid X_0 = R_1] = (0.7)(0.6)$$

$$P[X_3 = R_1, X_2 \neq R_1, X_1 \neq R_1 \mid X_0 = R_1] = (0.7)(0.4)(0.6)$$

⋮

$$P[X_n = R_1, X_{n-1} \neq R_1, \dots, X_1 \neq R_1 \mid X_0 = R_1] = (0.7)(0.4)^{n-1}(0.6)$$

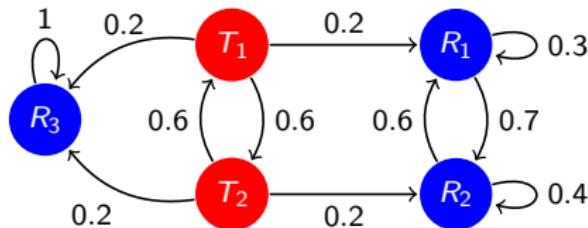


- ▶ Sum up: $f_i = \sum_{n=1}^{\infty} P[X_n = R_1, X_{n-1} \neq R_1, \dots, X_1 \neq R_1 \mid X_0 = R_1]$

$$= 0.3 + 0.7 \left(\sum_{n=1}^{\infty} 0.4^{n-1} \right) 0.6 = 0.3 + 0.7 \left(\frac{1}{1-0.4} \right) 0.6 = 1$$

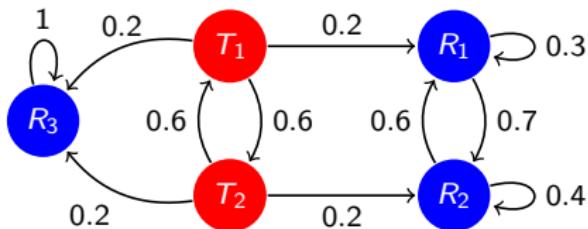
Transient state example

- ▶ States T_1 and T_2 are transient
- ▶ Probability of returning to T_1 is $f_{T_1} = (0.6)^2 = 0.36$
- ▶ Might come back to T_1 only if it goes to T_2 (with prob. 0.6)
- ▶ Will come back only if it moves back from T_2 to T_1 (with prob. 0.6)



- ▶ Likewise, $f_{T_2} = (0.6)^2 = 0.36$

- ▶ State j is accessible from state i if $P_{ij}^n > 0$ for some $n \geq 0$
- ▶ It is possible to enter j if MC initialized at $X_0 = i$
- ▶ Since $P_{ii}^0 = \text{P} [X_0 = 1 \mid X_0 = i] = 1$, state i is accessible from itself



- ▶ All states accessible from T_1 and T_2
- ▶ Only R_1 and R_2 accessible from R_1 or R_2
- ▶ None other than itself accessible from R_3

- ▶ States i and j are said to communicate ($i \leftrightarrow j$) if
 - ⇒ i is accessible from j , $P_{ij}^n > 0$ for some n ; and
 - ⇒ j is accessible from i , $P_{ji}^m > 0$ for some m
- ▶ Communication is an equivalence relation
- ▶ **Reflexivity**: $i \leftrightarrow i$
 - ▶ true because $P_{ii}^0 = 1$
- ▶ **Symmetry**: If $i \leftrightarrow j$ then $j \leftrightarrow i$
 - ▶ If $i \leftrightarrow j$ then $P_{ij}^n > 0$ and $P_{ji}^m > 0$ from where $j \leftrightarrow i$
- ▶ **Transitivity**: If $i \leftrightarrow j$ and $j \leftrightarrow k$, then $i \leftrightarrow k$
 - ▶ Just notice that $P_{ik}^{n+m} \geq P_{ij}^n P_{jk}^m > 0$
- ▶ **Partitions set of states into disjoint classes** (as all equivalences do)
- ▶ What are these classes? (start with a brief detour)

- ▶ Define N_i as the number of visits to state i given that $X_0 = i$

$$N_i := \sum_{n=1}^{\infty} \mathbb{I}\{X_n = i\}$$

- ▶ If $X_n = i$, this is the last visit to i with probability $1 - f_i$
- ▶ Prob. revisiting state i exactly n times is (n visits \times no more visits)

$$\mathbb{P}[N_i = n] = f_i^n (1 - f_i)$$

- ▶ Number of visits N_i has a geometric distribution with parameter f_i
- ▶ Expected number of visits is

$$\mathbb{E}[N_i] = \sum_{n=1}^{\infty} n f_i^n (1 - f_i) = \frac{1}{1 - f_i}$$

- ▶ For recurrent states $N_i = \infty$ almost surely and $\mathbb{E}[N_i] = \infty$ ($f_i = 1$)

- ▶ Another way of writing $\mathbb{E}[N_i]$

$$\mathbb{E}[N_i] = \sum_{n=1}^{\infty} \mathbb{E}[\mathbb{I}\{X_n = i\}] = \sum_{n=1}^{\infty} P_{ii}^n$$

- ▶ Recall that: for transient states $\mathbb{E}[N_i] = 1/(1 - f_i)$
for recurrent states $\mathbb{E}[N_i] = \infty$
- ▶ Therefore proving

Theorem

- ▶ *State i is transient if and only if $\sum_{n=1}^{\infty} P_{ii}^n < \infty$*
- ▶ *State i is recurrent if and only if $\sum_{n=1}^{\infty} P_{ii}^n = \infty$*
- ▶ Number of future visits to transient states is finite
- ▶ If number of states is finite some states have to be recurrent

Theorem

If state i is recurrent and $i \leftrightarrow j$, then j is recurrent

Proof.

- If $i \leftrightarrow j$ then there are l, m such that $P_{ji}^l > 0$ and $P_{ij}^m > 0$
- Then, for any n we have

$$P_{jj}^{l+n+m} \geq P_{ji}^l P_{ii}^n P_{ij}^m$$

- Sum for all n . Note that since i is recurrent $\sum_{n=1}^{\infty} P_{ii}^n = \infty$

$$\sum_{n=1}^{\infty} P_{jj}^{l+n+m} \geq \sum_{n=1}^{\infty} P_{ji}^l P_{ii}^n P_{ij}^m = P_{ji}^l \left(\sum_{n=1}^{\infty} P_{ii}^n \right) P_{ij}^m = \infty$$

- Which implies j is recurrent



Corollary

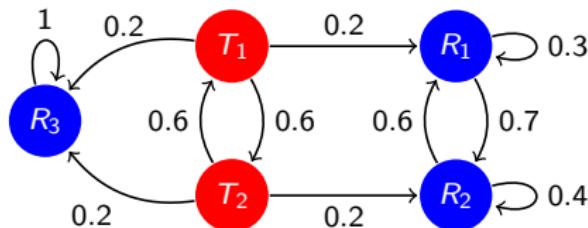
If state i is transient and $i \leftrightarrow j$ then j is transient

Proof.

- ▶ If j were recurrent, then i would be recurrent from previous theorem □
- ▶ Since communication defines classes and recurrence is shared by elements of this class, we say that recurrence is a class property
- ▶ Likewise, transience is also a class property
- ▶ States of a MC are separated in classes of transient and recurrent states

- ▶ A MC is called irreducible if it has only one class
 - ▶ All states communicate with each other
 - ▶ If MC also has finite number of states the single class is recurrent
 - ▶ If MC infinite, class might be transient
- ▶ When it has multiple classes (not irreducible)
 - ▶ Classes of transient states $\mathcal{T}_1, \mathcal{T}_2, \dots$
 - ▶ Classes of recurrent states $\mathcal{R}_1, \mathcal{R}_2, \dots$
 - ▶ If MC initialized in a recurrent class \mathcal{R}_k , stays within the class
 - ▶ If starts in transient class \mathcal{T}_k , might stay on \mathcal{T}_k or end up in a recurrent class \mathcal{R}_l
- ▶ For large time index n , MC restricted to one class
- ▶ Can be separated into irreducible components

Example



- ▶ Three classes
 - ⇒ $\mathcal{T} := \{T_1, T_2\}$, class with transient states
 - ⇒ $\mathcal{R}_1 := \{R_1, R_2\}$, class with recurrent states
 - ⇒ $\mathcal{R}_2 := \{R_3\}$, class with recurrent states
- ▶ Asymptotically in n suffices to study behavior for the irreducible components \mathcal{R}_1 and \mathcal{R}_2

- ▶ States of a MC can be transient or recurrent
- ▶ A MC can be partitioned in classes of states reachable from each other
- ▶ Elements of the class are either all recurrent or all transient
- ▶ A MC with only one class is irreducible
- ▶ If not irreducible can be separated in irreducible components

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Ergodicity

Queues in communication networks: Limit probabilities

- ▶ MCs have one-step memory. Eventually they forget initial state
- ▶ What can we say about probabilities for large n ?

$$\pi_j := \lim_{n \rightarrow \infty} P[X_n = j \mid X_0 = i] = \lim_{n \rightarrow \infty} P_{ij}^n$$

- ▶ Implicitly assumed that limit is independent of initial state $X_0 = i$
- ▶ We've seen that this problem is related to the matrix power \mathbf{P}^n

$$\mathbf{P} := \begin{pmatrix} 0.8 & 0.2 \\ 0.3 & 0.7 \end{pmatrix}$$

$$\mathbf{P}^7 := \begin{pmatrix} 0.6031 & 0.3969 \\ 0.5953 & 0.4047 \end{pmatrix}$$

$$\mathbf{P}^2 := \begin{pmatrix} 0.7 & 0.3 \\ 0.45 & 0.55 \end{pmatrix}$$

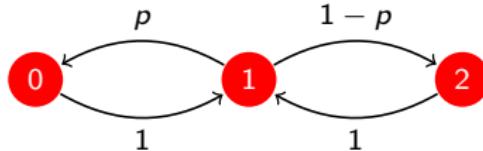
$$\mathbf{P}^{30} := \begin{pmatrix} 0.6000 & 0.4000 \\ 0.6000 & 0.4000 \end{pmatrix}$$

- ▶ Matrix product converges \Rightarrow probs. independent of time (large n)
- ▶ All columns are equal \Rightarrow probs. independent of initial condition

- ▶ The period of a state i is defined as (\dot{d} is set of multiples of d)

$$d = \max \left\{ d : P_{ii}^n = 0 \text{ for all } n \notin \dot{d} \right\}$$

- ▶ State i is periodic with period d if and only if
 - $\Rightarrow P_{ii}^n \neq 0$ only if n is a multiple of d ($n \in \dot{d}$)
 - $\Rightarrow d$ is the largest number with this property
- ▶ Positive probability of returning to i only every d time steps
- ▶ If period $d = 1$ state is aperiodic (most often the case)
- ▶ Periodicity is a class property



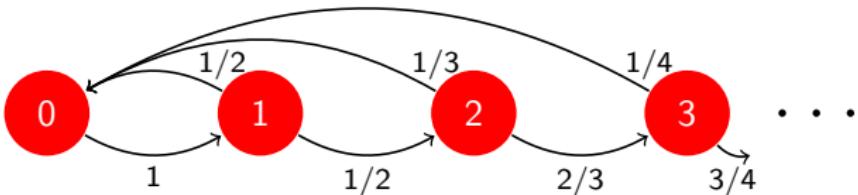
- ▶ State 1 has period 2. So do 0 and 2 (class property)
- ▶ One dimensional random walk also has period 2

- ▶ Recall: state i is recurrent if chain returns to i with probability 1
- ▶ Proved it was equivalent to $\sum_{n=1}^{\infty} P_{ii}^n = \infty$
- ▶ Positive recurrent when expected value of return time is finite

$$\mathbb{E}[\text{return time}] = \sum_{n=1}^{\infty} n P_{ii}^n \prod_{m=0}^{n-1} (1 - P_{ii}^m) < \infty$$

- ▶ Null recurrent if recurrent but $\mathbb{E}[\text{return time}] = \infty$
- ▶ Positive and null recurrence are class properties
- ▶ Recurrent states in a finite-state MC are positive recurrent
- ▶ Ergodic states are those that are positive recurrent and aperiodic
- ▶ An irreducible MC with ergodic states is said to be an ergodic MC

Example of a null recurrent MC



$$P[\text{return time} = 2] = \frac{1}{2}$$

$$P[\text{return time} = 3] = \frac{1}{2} \times \frac{1}{3}$$

$$P[\text{return time} = 4] = \frac{1}{2} \times \frac{2}{3} \times \frac{1}{4} = \frac{1}{3 \times 4} \quad \dots \quad P[\text{return time} = n] = \frac{1}{(n-1) \times n}$$

- ▶ It is recurrent because probability of returning is 1 (use induction)

$$\sum_{m=2}^n P[\text{return time} = m] = \sum_{m=2}^n \frac{1}{(m-1) \times m} = \frac{n-1}{n} \rightarrow 1$$

- ▶ Null recurrent because expected return time is infinite

$$\sum_{n=2}^{\infty} n P[\text{return time} = n] = \sum_{n=2}^{\infty} \frac{n}{(n-1) \times n} = \sum_{n=2}^{\infty} \frac{1}{(n-1)} = \infty$$

Theorem

For an irreducible ergodic MC, $\lim_{n \rightarrow \infty} P_{ij}^n$ exists and is independent of the initial state i . That is

$$\pi_j = \lim_{n \rightarrow \infty} P_{ij}^n \quad \text{exists}$$

Furthermore, steady state probabilities $\pi_j \geq 0$ are the unique nonnegative solution of the system of linear equations

$$\pi_j = \sum_{i=0}^{\infty} \pi_i P_{ij}, \quad \sum_{j=0}^{\infty} \pi_j = 1$$

- ▶ As observed, limit probs. independent of initial condition exist
- ▶ Simple algebraic equations can be solved to find π_j
- ▶ No periodic states, transient states, multiple classes or null recurrent

- ▶ Difficult part of theorem is to prove that $\pi_j = \lim_{n \rightarrow \infty} P_{jj}^n$ exists
- ▶ To see that algebraic relation is true use theorem of total probability (omit conditioning on X_0 to simplify notation)

$$\begin{aligned} P[X_{n+1} = j] &= \sum_{i=1}^{\infty} P[X_{n+1} = j \mid X_n = i] P[X_n = i] \\ &= \sum_{i=1}^{\infty} P_{ij} P[X_n = i] \end{aligned}$$

- ▶ If limits exists, $P[X_{n+1} = j] \approx P[X_n = j] \approx \pi_j$ (sufficiently large n)

$$\pi_j = \sum_{i=0}^{\infty} \pi_i P_{ij}$$

- ▶ The other equation is true because the π_j are probabilities

- ▶ More compact and illuminating on vector/matrix notation
- ▶ Finite MC with J states
- ▶ First part of theorem says that $\lim_{n \rightarrow \infty} \mathbf{P}^n$ exists and

$$\lim_{n \rightarrow \infty} \mathbf{P}^n = \begin{pmatrix} \pi_1 & \pi_2 & \dots & \pi_J \\ \pi_1 & \pi_2 & \dots & \pi_J \\ \vdots & \vdots & \vdots & \vdots \\ \pi_1 & \pi_2 & \dots & \pi_J \end{pmatrix}$$

- ▶ Same probs. for all rows \Rightarrow independent of initial state
- ▶ Probability distribution for large n .

$$\lim_{n \rightarrow \infty} \mathbf{p}(n) = \lim_{n \rightarrow \infty} \mathbf{P}^T \mathbf{p}(0) = [\pi_0, \pi_1, \dots, \pi_J]^T$$

- ▶ Independent of initial condition

- ▶ Define vector stationary distribution $\pi := [\pi_0, \pi_1, \dots, \pi_J]^T$
- ▶ Limit distribution is unique solution of $(\mathbf{1} = [1, 1, \dots]^T)$

$$\pi = \mathbf{P}^T \pi, \quad \pi^T \mathbf{1} = 1$$

- ▶ π eigenvector associated with eigenvalue 1 of \mathbf{P}^T
 - ▶ Eigenvectors are defined up to a constant
 - ▶ Normalize to sum 1
- ▶ All other eigenvectors of \mathbf{P}^T have modulus smaller than 1
 - ▶ If not, \mathbf{P}^n diverges, but we know \mathbf{P}^n contains n -step transition probs.
 - ▶ π eigenvector associated with largest eigenvalue of \mathbf{P}^T
- ▶ Computing π as eigenvector is computationally efficient and robust in some problems

- ▶ Can also write as (\mathbf{I} is identity matrix, $\mathbf{0} = [0, 0, \dots]^T$)

$$(\mathbf{I} - \mathbf{P}^T) \boldsymbol{\pi} = \mathbf{0} \quad \boldsymbol{\pi}^T \mathbf{1} = 1$$

- ▶ $\boldsymbol{\pi}$ has J elements, but there are $J + 1$ equations \Rightarrow overdetermined
- ▶ If 1 is eigenvalue of \mathbf{P}^T , then 0 is eigenvalue of $\mathbf{I} - \mathbf{P}^T$
 - ▶ $\mathbf{I} - \mathbf{P}^T$ is rank deficient, in fact $\text{rank}(\mathbf{I} - \mathbf{P}^T) = J - 1$
 - ▶ Then, there are in fact only J equations
- ▶ $\boldsymbol{\pi}$ is eigenvector associated with eigenvalue 0 of $\mathbf{I} - \mathbf{P}^T$
 - ▶ $\boldsymbol{\pi}$ spans null space of $\mathbf{I} - \mathbf{P}^T$ (not much significance)

- ▶ MC with transition probability matrix

$$\mathbf{P} := \begin{pmatrix} 0 & 0.3 & 0.7 \\ 0.1 & 0.5 & 0.4 \\ 0.1 & 0.2 & 0.7 \end{pmatrix}$$

- ▶ Does P correspond to an ergodic MC?
 - ▶ All states communicate with state 2 (full row and column $P_{2j} \neq 0$ and $P_{j2} \neq 0$ for all j)
 - ▶ No transient states (irreducible with one recurrent state and finite)
 - ▶ Aperiodic (period of state 2 is 1)
- ▶ Then, there exist π_1 , π_2 and π_3 such that $\pi_j = \lim_{n \rightarrow \infty} P_{jj}^n$
- ▶ Limit is independent of i

- ▶ How do we determine limit probabilities π_j ?
- ▶ Solve system of linear equations $\pi_j = \sum_{i=0}^{\infty} \pi_i P_{ij}$ and $\sum_{j=0}^{\infty} \pi_j = 1$

$$\begin{pmatrix} \pi_1 \\ \pi_2 \\ \pi_3 \\ 1 \end{pmatrix} = \begin{pmatrix} 0 & 0.1 & 0.1 \\ 0.3 & 0.5 & 0.2 \\ 0.7 & 0.4 & 0.7 \\ 1 & 1 & 1 \end{pmatrix} \begin{pmatrix} \pi_1 \\ \pi_2 \\ \pi_3 \end{pmatrix}$$

- ▶ The upper part of matrix above is \mathbf{P}^T
- ▶ There are three variables and four equations
 - ▶ Some equations might be linearly dependent
 - ▶ Indeed, summing first three equations: $\pi_1 + \pi_2 + \pi_3 = \pi_1 + \pi_2 + \pi_3$
 - ▶ Always true, because probabilities in rows of \mathbf{P} sum up to 1
 - ▶ This is because of rank deficiency of $\mathbf{I} - \mathbf{P}^T$
- ▶ Solution yields $\pi_1 = 0.0909$, $\pi_2 = 0.2987$ and $\pi_3 = 0.6104$

- ▶ Limit distributions are sometimes called stationary distributions
- ▶ Select initial distribution such that $P[X_0 = i] = \pi_i$ for all i
- ▶ Probabilities at time $n = 1$ follow from theorem of total probability

$$P[X_1 = i] = \sum_{j=1}^{\infty} P[X_1 = j \mid X_0 = i] P[X_0 = i]$$

- ▶ Definitions of P_{ij} , and $P[X_0 = i] = \pi_i$. Algebraic property of π_j

$$P[X_1 = i] = \sum_{j=1}^{\infty} P_{ij} \pi_j = \pi_i$$

- ▶ Probability distribution is unchanged
- ▶ Proceeding recursively, system initialized with $P[X_0 = i] = \pi_i$,
⇒ Probability distribution constant, $P[X_n = i] = \pi_i$ for all n
- ▶ MC stationary in a probabilistic sense (states change, probs. do not)

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Queues in communication networks: Limit probabilities

- ▶ Define $T_i^{(n)}$ as fraction of time spent in i -th state up to time n

$$T_i^{(n)} := \frac{1}{n} \sum_{m=1}^n \mathbb{I}\{X_m = i\}$$

- ▶ Compute expected value of $T_i^{(n)}$

$$\mathbb{E}\left[T_i^{(n)}\right] = \frac{1}{n} \sum_{m=1}^n \mathbb{E}[\mathbb{I}\{X_m = i\}] = \frac{1}{n} \sum_{m=1}^n \mathbb{P}[X_m = i] \rightarrow \pi_i$$

- ▶ As time $n \rightarrow \infty$, probabilities $\mathbb{P}[X_m = i]$ approach π_i . Then

$$\lim_{t \rightarrow \infty} \mathbb{E}\left[T_i^{(n)}\right] = \lim_{t \rightarrow \infty} \frac{1}{n} \sum_{m=1}^n \mathbb{P}[X_m = i] = \pi_i$$

- ▶ For ergodic MCs same is true without expected value \Rightarrow ergodicity

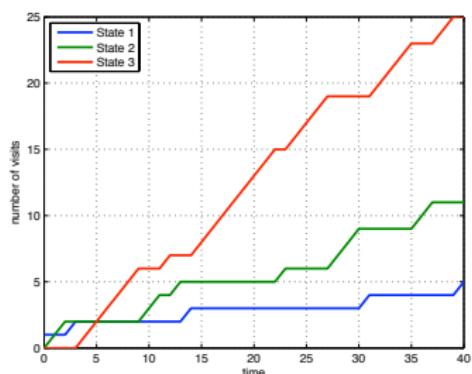
$$\lim_{n \rightarrow \infty} T_i^{(n)} = \lim_{n \rightarrow \infty} \frac{1}{n} \sum_{m=1}^n \mathbb{I}\{X_m = i\} = \pi_i, \quad \text{a.s.}$$

Example: Ergodic Markov chain

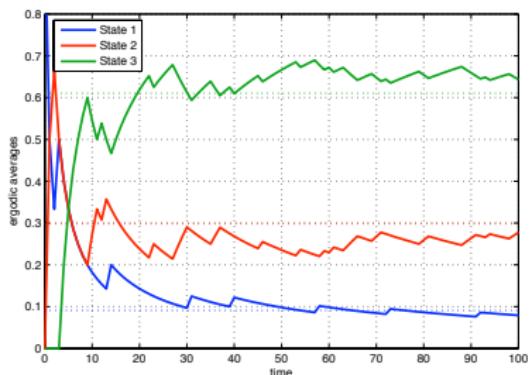
- Recall transition probability matrix

$$\mathbf{P} := \begin{pmatrix} 0 & 0.3 & 0.7 \\ 0.1 & 0.5 & 0.4 \\ 0.1 & 0.2 & 0.7 \end{pmatrix}$$

Visits to states, $nT_i^{(n)}$



Ergodic averages, $T_i^{(n)}$



- Ergodic averages slowly converge to limit probabilities

- ▶ Use of ergodic averages is more general than $T_i^{(n)}$

Theorem

Consider an irreducible Markov chain with states $X_n = 0, 1, 2, \dots$ and stationary probabilities π_j . Let $f(X_n)$ be a bounded function of the state $X(n)$. Then, with probability 1

$$\lim_{n \rightarrow \infty} \frac{1}{n} \sum_{m=1}^n f(X_m) = \sum_{i=1}^{\infty} f(i) \pi_i$$

- ▶ $T_i^{(n)}$ is a particular case with $f(X_m) = \mathbb{I}\{X_m = i\}$
- ▶ Think of $f(X_m)$ as a reward associated with state $X(m)$
- ▶ $(1/n) \sum_{m=1}^n f(X_m)$ is the time average of rewards

Proof.

- ▶ Because $\mathbb{I}\{X_m = i\} = 1$ if and only if $X_m = i$ we can write

$$\frac{1}{n} \sum_{m=1}^n f(X_m) = \frac{1}{n} \sum_{m=1}^n \left(\sum_{i=1}^{\infty} f(i) \mathbb{I}\{X_m = i\} \right)$$

- ▶ Change order of summations. Use definition of $T_i^{(n)}$

$$\frac{1}{n} \sum_{m=1}^n f(X_m) = \sum_{i=1}^{\infty} f(i) \left(\frac{1}{n} \sum_{m=1}^n \mathbb{I}\{X_m = i\} \right) = \sum_{i=1}^{\infty} f(i) T_i^{(n)}$$

- ▶ Let $n \rightarrow \infty$ in both sides
- ▶ Use ergodic average result for $\lim_{n \rightarrow \infty} T_i^{(n)} = \pi_i$ [cf. page 67]



- ▶ There's more depth to ergodic results than meets the eye
- ▶ Ensemble average: across different realizations of the MC

$$\mathbb{E}[f(X_n)] = \sum_{i=1}^{\infty} f(i)P(X_n = i) \rightarrow \sum_{i=1}^{\infty} f(i)\pi_i$$

- ▶ Ergodic average: across time for a single realization of the MC

$$\bar{f}(n) = \frac{1}{n} \sum_{m=1}^n f(X_m)$$

- ▶ These quantities are fundamentally different but their values coincide asymptotically in n
- ▶ Observing one realization of the MC provides as much information as observing all realizations
- ▶ Practical consequence: Observe/simulate only one path of the MC

- ▶ In some sense, still true if MC is **periodic**
- ▶ For irreducible positive recurrent MC (periodic or aperiodic) define

$$\pi_j = \sum_{i=0}^{\infty} \pi_i P_{ij}, \quad \sum_{j=0}^{\infty} \pi_j = 1$$

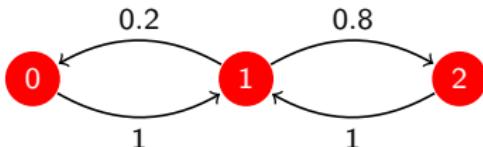
- ▶ A unique solution exists (we say π_j are well defined)
- ▶ The fraction of time spent in state i converges to π_i

$$\lim_{n \rightarrow \infty} T_i^{(n)} = \lim_{n \rightarrow \infty} \frac{1}{n} \sum_{m=1}^n \mathbb{I}\{X_m = i\} \rightarrow \pi_i$$

- ▶ If MC is periodic, probabilities oscillate, but fraction of time spent in state i converges to π_i

- Matrix \mathbf{P} and state diagram of a periodic MC

$$\mathbf{P} := \begin{pmatrix} 0 & 1 & 0 \\ 0.3 & 0 & 0.7 \\ 0 & 1 & 0 \end{pmatrix}$$



- MC has period 2. If initialized with $X_0 = 1$, then

$$\begin{aligned} P_{11}^{2n+1} &= P[X_{2n+1} = 1 \mid X_0 = 1] = 0, \\ P_{11}^{2n} &= P[X_{2n} = 1 \mid X_0 = 1] = 1 \neq 0 \end{aligned}$$

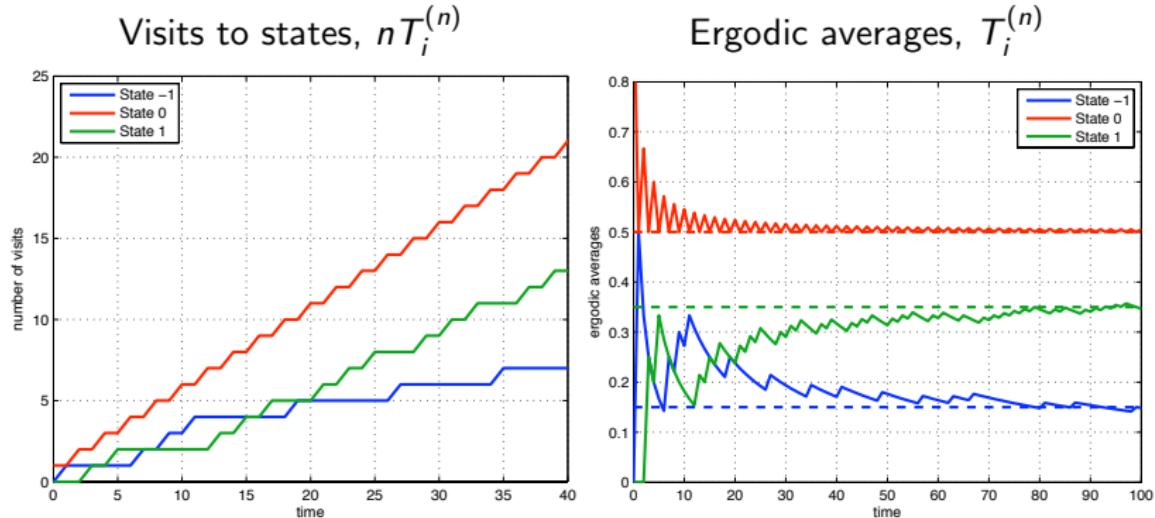
- Define $\boldsymbol{\pi} := [\pi_1, \pi_2, \pi_3]^T$ as solution of

$$\begin{pmatrix} \pi_1 \\ \pi_2 \\ \pi_3 \\ 1 \end{pmatrix} = \begin{pmatrix} 0 & 0.3 & 0 \\ 1 & 0 & 1 \\ 0 & 0.7 & 0 \\ 1 & 1 & 1 \end{pmatrix} \begin{pmatrix} \pi_1 \\ \pi_2 \\ \pi_3 \end{pmatrix}$$

- Normalized eigenvector for eigenvalue 1 ($\boldsymbol{\pi} = \mathbf{P}^T \boldsymbol{\pi}$, $\boldsymbol{\pi}^T \mathbf{1} = 1$)

Example: Periodic irreducible MC (continued)

- ▶ Solution yields $\pi_1 = 0.15$, $\pi_2 = 0.50$ and $\pi_3 = 0.35$



- ▶ Ergodic averages $T_i^{(n)}$ converge to the ergodic limits π_i

- ▶ Powers of the transition probability matrix do **not converge**

$$\mathbf{P}^2 = \begin{pmatrix} 0.3 & 0 & 0.7 \\ 0 & 1 & 0 \\ 0.3 & 0 & 0.7 \end{pmatrix} \quad \mathbf{P}^3 = \begin{pmatrix} 0 & 1 & 0 \\ 0.3 & 0 & 0.7 \\ 0 & 1 & 0 \end{pmatrix} = \mathbf{P}$$

- ▶ In general we have $\mathbf{P}^{2n} = \mathbf{P}^2$ and $\mathbf{P}^{2n+1} = \mathbf{P}$
- ▶ At least one other eigenvalue of the transition probability matrix has modulus 1

$$|\text{eig}_2(\mathbf{P}^T)| = 1$$

- ▶ In this example, eigenvalues of \mathbf{P}^T are 1, -1 and 0

- ▶ If MC is not irreducible it can be decomposed in transient (\mathcal{T}_k), ergodic (\mathcal{E}_k), periodic (\mathcal{P}_k) and null recurrent (\mathcal{N}_k) components
 - ▶ All of these are class properties
- ▶ Limit probabilities for transient states are null $P[X_n = i] \rightarrow 0$, for all $X_n \in \mathcal{T}_k$
- ▶ For arbitrary ergodic component \mathcal{E}_k , define conditional limits

$$\pi_i = \lim_{n \rightarrow \infty} P[X_n = i \mid X_0 \in \mathcal{E}_k], \quad \text{for all } i \in \mathcal{E}_k$$

- ▶ Result in page 58 is true with this (re)defined π_i

- ▶ Likewise, for arbitrary periodic component \mathcal{P}_k (re)define π_j as

$$\pi_j = \sum_{i \in \mathcal{P}_k} \pi_i P_{ij}, \quad \sum_{j \in \mathcal{P}_k} \pi_j = 1, \quad \text{for all } j \in \mathcal{P}_k$$

- ▶ A conditional version of the result in page 72 is true

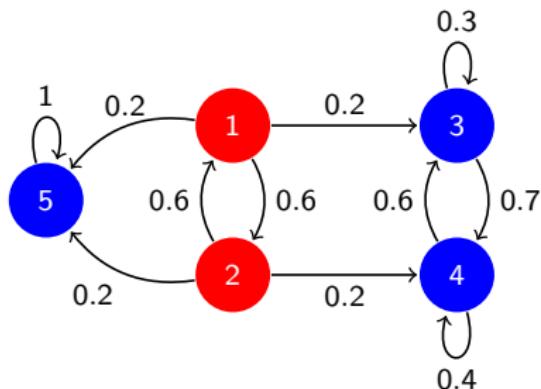
$$\lim_{n \rightarrow \infty} T_i^{(n)} := \lim_{n \rightarrow \infty} \frac{1}{n} \sum_{m=1}^n \mathbb{I}\{X_m = i \mid X_0 \in \mathcal{P}_k\} \rightarrow \pi_i$$

- ▶ For null recurrent components limit probabilities are null
 $P[X_n = i] \rightarrow 0$, for all $X_n \in \mathcal{N}_k$

Example: Reducible Markov chain

- ▶ Transition matrix and state diagram of a reducible MC

$$\mathbf{P} := \begin{pmatrix} 0 & 0.6 & 0.2 & 0 & 0.2 \\ 0.6 & 0 & 0 & 0.2 & 0.2 \\ 0 & 0 & 0.3 & 0.7 & 0 \\ 0 & 0 & 0.6 & 0.4 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$



- ▶ States 1 and 2 are transient $\mathcal{T} = \{1, 2\}$
- ▶ States 3 and 4 form an ergodic class $\mathcal{E}_1 = \{3, 4\}$
- ▶ State 5 is a separate ergodic class $\mathcal{E}_2 = \{5\}$

- ▶ 10-step and 20 step transition probabilities

$$\mathbf{P}^5 = \begin{pmatrix} 0 & 0.08 & 0.24 & 0.22 & 0.46 \\ 0.08 & 0 & 0.19 & 0.27 & 0.46 \\ 0 & 0 & 0.46 & 0.54 & 0 \\ 0 & 0 & 0.46 & 0.54 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{pmatrix} \quad \mathbf{P}^{10} = \begin{pmatrix} 0.00 & 0 & 0.23 & 0.27 & 0.50 \\ 0 & 0.00 & 0.23 & 0.27 & 0.50 \\ 0 & 0 & 0.46 & 0.54 & 0 \\ 0 & 0 & 0.46 & 0.54 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

- ▶ Transition into transient states is vanishing (columns 1 and 2)
- ▶ Transition from 3 and 4 into 3 and 4 only
 - ▶ If initialized in ergodic class $\mathcal{E}_1 = \{3, 4\}$ stays in \mathcal{E}_1
- ▶ Transition from 5 only into 5
- ▶ From transient states $\mathcal{T} = \{1, 2\}$ can go into either ergodic component $\mathcal{E}_1 = \{3, 4\}$ or $\mathcal{E}_2 = \{5\}$

- Matrix \mathbf{P} can be separated in blocks

$$\mathbf{P} = \begin{pmatrix} 0 & 0.6 & 0.2 & 0 & 0.2 \\ 0.6 & 0 & 0 & 0.2 & 0.2 \\ 0 & 0 & 0.3 & 0.7 & 0 \\ 0 & 0 & 0.6 & 0.4 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{pmatrix} = \begin{pmatrix} \mathbf{P}_{\mathcal{T}} & \mathbf{P}_{\mathcal{T}\mathcal{E}_1} & \mathbf{P}_{\mathcal{T}\mathcal{E}_2} \\ 0 & \mathbf{P}_{\mathcal{E}_1} & 0 \\ 0 & 0 & \mathbf{P}_{\mathcal{E}_2} \end{pmatrix}$$

- Block $\mathbf{P}_{\mathcal{T}}$ describes transition between transient states
- Blocks $\mathbf{P}_{\mathcal{E}_1}$ and $\mathbf{P}_{\mathcal{E}_2}$ describe transitions in ergodic components
- Blocks $\mathbf{P}_{\mathcal{T}\mathcal{E}_1}$ and $\mathbf{P}_{\mathcal{T}\mathcal{E}_2}$ describe transitions from \mathcal{T} to \mathcal{E}_1 and \mathcal{E}_2
- Powers of n can be written as

$$\mathbf{P}^n = \begin{pmatrix} \mathbf{P}_{\mathcal{T}}^n & \mathbf{Q}_{\mathcal{T}\mathcal{E}_1} & \mathbf{Q}_{\mathcal{T}\mathcal{E}_2} \\ 0 & \mathbf{P}_{\mathcal{E}_1}^n & 0 \\ 0 & 0 & \mathbf{P}_{\mathcal{E}_2}^n \end{pmatrix}$$

- The transient transition block converges to 0, $\lim_{n \rightarrow \infty} \mathbf{P}_{\mathcal{T}}^n = \mathbf{0}$

- ▶ As n grows the MC hits an ergodic state with probability 1
- ▶ Henceforth, MC stays within ergodic component

$$P[X_{n+m} \in \mathcal{E}_i \mid X_n \in \mathcal{E}_i] = 1, \quad \text{for all } m$$

- ▶ For large n suffices to study ergodic components
 - ⇒ MC behaves like a MC with transition probabilities $\mathbf{P}_{\mathcal{E}_1}$
 - ⇒ Or like one with transition probabilities $\mathbf{P}_{\mathcal{E}_2}$
- ▶ We can think that all MCs are ergodic
- ▶ Ergodic behavior cannot be inferred a priori (before observing)
- ▶ Becomes known a posteriori (after observing sufficiently large time)

Culture micro: Something is known a priori if its knowledge is independent of experience (MCs exhibit ergodic behavior). A posteriori knowledge depends on experience (values of the ergodic limits). They are inherently different forms of knowledge (search for Immanuel Kant)

Markov chains. Definition and examples

Chapman Kolmogorov equations

Gambler's ruin problem

Queues in communication networks: Transition probabilities

Classes of States

Limiting distributions

Ergodicity

Queues in communication networks: Limit probabilities

- ▶ Communication system: Move packets from source to destination
- ▶ Between arrival and transmission hold packets in a memory buffer
- ▶ Example problem, buffer design: Packets arrive at a rate of 0.45 packets per unit of time and depart at a rate of 0.55. How many packets the buffer needs to hold to have a drop rate smaller than 10^{-6} (one packet dropped for every million packets handled)
- ▶ Time slotted in intervals of duration Δt
- ▶ During each time slot n
 - ⇒ A packet arrives with prob. λ , arrival rate is $\lambda/\Delta t$
 - ⇒ A packet is transmitted with prob. μ , departure rate is $\mu/\Delta t$
- ▶ No concurrence: No simultaneous arrival and departure (small Δt)

- ▶ Future queue lengths depend on current length only
- ▶ Probability of queue length increasing

$$P[q_{n+1} = i+1 \mid q_n = i] = \lambda, \quad \text{for all } i$$

- ▶ Queue length might decrease only if $q_n > 0$. Probability is

$$P[q_{n+1} = i-1 \mid q_n = i] = \mu, \quad \text{for all } i > 0$$

- ▶ Queue length stays the same if it neither increases nor decreases

$$P[q_{n+1} = i \mid q_n = i] = 1 - \lambda - \mu, \quad \text{for all } i > 0$$

$$P[q_{n+1} = 0 \mid q_n = 0] = 1 - \lambda$$

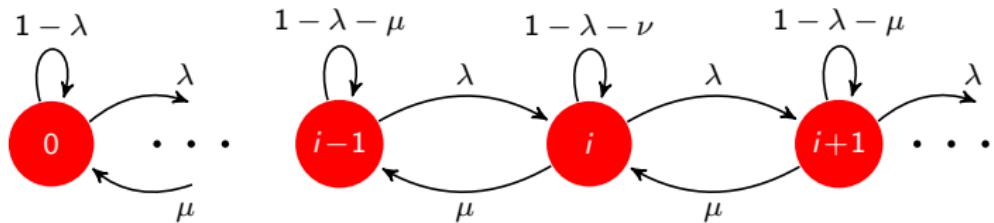
- ▶ No departures when $q_n = 0$ explain second equation

Queue as a Markov chain (reminder)

- ▶ MC with states $0, 1, 2, \dots$. Identify states with queue lengths
- ▶ Transition probabilities for $i \neq 0$ are

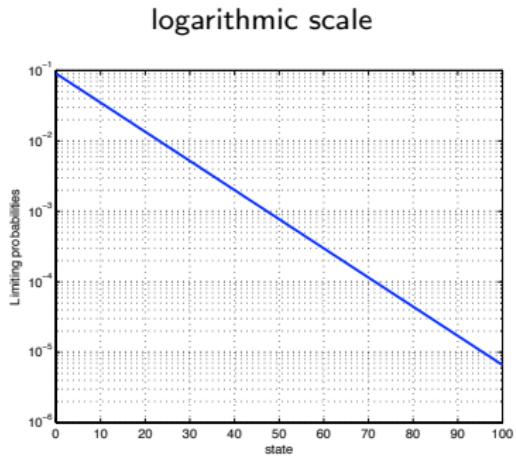
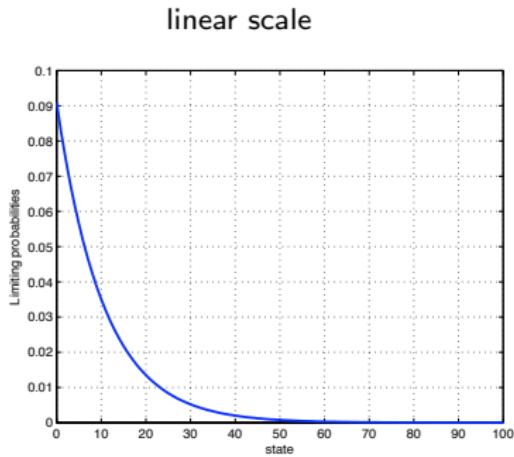
$$P_{i,i-1} = \lambda, \quad P_{i,i} = 1 - \lambda - \mu, \quad P_{i,i+1} = \mu$$

- ▶ For $i = 0$ $P_{0,0} = 1 - \lambda$ and $P_{01} = \lambda$

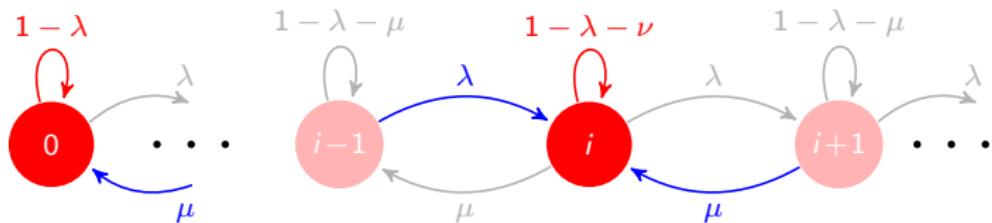


Numerical example: Limit probabilities

- ▶ Build matrix \mathbf{P} truncating at maximum queue length $L = 100$
- ▶ Arrival rate $\lambda = 0.3$. Departure rate $\mu = 0.33$
- ▶ Find eigenvector of \mathbf{P}^T associated with largest eigenvalue (i.e., 1)
- ▶ Yields limit probabilities $\pi = \lim_{n \rightarrow \infty} \mathbf{p}(n)$



- ▶ Limit probabilities appear linear in logarithmic scale
⇒ Seemingly implying an exponential expression $\pi_i \propto \alpha^i$



- ▶ Limit distribution equations for state 0 (empty queue)

$$\pi_0 = (1 - \lambda)\pi_0 + \mu\pi_1$$

- ▶ For the remaining states

$$\pi_i = \lambda\pi_{i-1} + (1 - \lambda - \mu)\pi_i + \mu\pi_{i+1}$$

- ▶ Propose candidate solution $\pi_i = c\alpha^i$

- ▶ Substitute candidate solution $\pi_i = c\alpha^i$ in equation for π_0

$$c\alpha^0 = (1 - \lambda)c\alpha^0 + \mu c\alpha^1 \quad \Rightarrow \quad 1 = (1 - \lambda) + \mu\alpha$$

- ▶ The above equation is true if we make $\alpha = \lambda/\mu$
- ▶ Does $\alpha = \lambda/\mu$ verify the remaining equations ?
- ▶ From the equation for generic π_i (divide by $c\alpha^{i-1}$)

$$\begin{aligned} c\alpha^i &= \lambda c\alpha^{i-1} + (1 - \lambda - \mu)c\alpha^i + \mu c\alpha^{i+1} \\ \mu\alpha^2 - (\lambda + \mu)\alpha + \lambda &= 0 \end{aligned}$$

- ▶ The above quadratic equation is satisfied by $\alpha = \lambda/\mu$
 - ▶ And $\alpha = 1$, which is irrelevant

Compute normalization constant

- ▶ Determine c so that probabilities sum 1 ($\sum_{i=0}^{\infty} \pi_i = 1$)

$$\sum_{i=0}^{\infty} \pi_i = \sum_{i=0}^J c(\lambda/\mu)^i = \frac{c}{1 - \lambda/\mu} = 1$$

- ▶ Used geometric sum
- ▶ Solving for c and substituting in $\pi_i = c\alpha^i$ yields

$$\pi_i = (1 - \lambda/\mu) \left(\frac{\lambda}{\mu}\right)^i$$

- ▶ The ratio μ/λ is the queues' stability margin
- ▶ Larger μ/λ \Rightarrow larger probability of having few queued packets

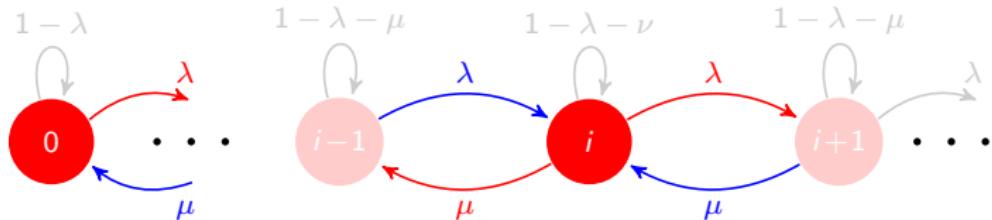
- ▶ Rearrange terms in equation for limit probabilities [cf. page 87]

$$\lambda\pi_0 = \mu\pi_1$$

$$(\lambda + \mu)\pi_i = \lambda\pi_{i-1} + \mu\pi_{i+1}$$

- ▶ $\lambda\pi_0$ is average rate at which the queue **leaves** state 0
- ▶ Likewise $(\lambda + \mu)\pi_i$ is the rate at which queue **leaves** state i
- ▶ $\mu\pi_0$ is average rate at which the queue **enters** state 0
- ▶ $\lambda\pi_{i-1} + \mu\pi_{i+1}$ is rate at which queue **enters** state i
- ▶ Limit equations prove validity of queue balance equations

Rate at which leaves = Rate at which enters



- ▶ Packets may arrive and depart in same time slot (concurrence)
- ▶ Queue evolution equations remain the same, [cf. 35]
- ▶ But queue probabilities change [cf. 84]
- ▶ Probability of queue length increasing (for all i)

$$P[q_{n+1} = i + 1 \mid q_n = i] = P[\mathbb{A}_n = 1] P[\mathbb{D}_n = 0] = \lambda(1 - \mu)$$

- ▶ Queue length might decrease only if $q_n > 0$ (for all $i > 0$)

$$P[q_{n+1} = i - 1 \mid q_n = i] = P[\mathbb{D}_n = 1] P[\mathbb{D}_n = 0] = \mu(1 - \lambda)$$

- ▶ Queue length stays the same if it neither increases nor decreases

$$P[q_{n+1} = i \mid q_n = i] = \lambda\mu + (1 - \lambda)(1 - \mu) = \nu, \quad \text{for all } i > 0$$

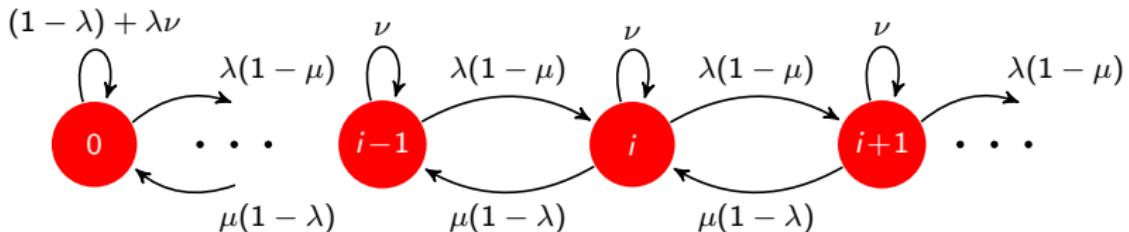
$$P[q_{n+1} = 0 \mid q_n = 0] = (1 - \lambda) + \lambda\mu$$

- ▶ Write limit distribution equations \Rightarrow queue balance equations
- ▶ Rate at which leaves = rate at which enters

$$\lambda(1 - \mu)\pi_0 = \mu(1 - \lambda)\pi_1$$

$$(\lambda(1 - \mu) + \mu(1 - \lambda))\pi_i = \lambda(1 - \mu)\pi_{i-1} + \mu(1 - \lambda)\pi_{i+1}$$

- ▶ Propose exponential solution $\pi = c\alpha^i$



- ▶ Substitute candidate solution in equation for π_0

$$\lambda(1 - \mu)c = \mu(1 - \lambda)c\alpha \quad \Rightarrow \quad \alpha = \frac{\lambda(1 - \mu)}{\mu(1 - \lambda)}$$

- ▶ Same substitution in equation for generic π_i

$$\mu(1 - \lambda)c\alpha^2 + (\lambda(1 - \mu) + \mu(1 - \lambda))c\alpha + \lambda(1 - \mu)c = 0$$

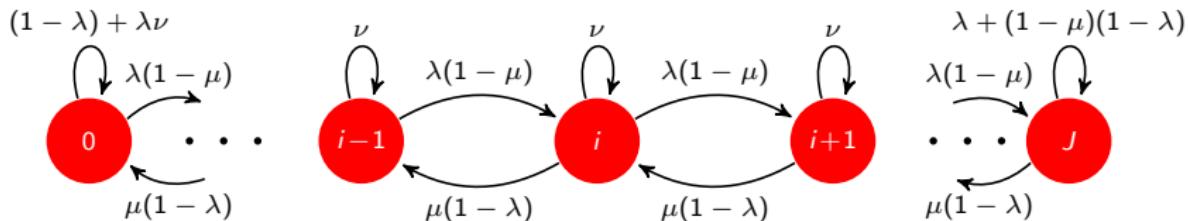
- ▶ which as before is solved for $\alpha = \lambda(1 - \mu)/\mu(1 - \lambda)$
- ▶ Find constant c to ensure $\sum_{i=0}^{\infty} c\alpha^i = 1$ (geometric series). Yields

$$\pi_i = (1 - \alpha)\alpha^i = \left(1 - \frac{\lambda(1 - \mu)}{\mu(1 - \lambda)}\right) \left(\frac{\lambda(1 - \mu)}{\mu(1 - \lambda)}\right)^i$$

- ▶ Packets dropped if there are too many packets in queue
- ▶ Too many packets in queue, then delays too large, packets useless when they arrive. Also preserve memory
- ▶ Equation for state J requires modification (rate leaves = rate enters)

$$\mu(1 - \lambda)\pi_J = \lambda(1 - \mu)\pi_{J-1}$$

- ▶ $\pi_i = c\alpha^i$ with $\alpha = \lambda(1 - \mu)/\mu(1 - \lambda)$ also solve this equation (Yes!)



- ▶ Limit probabilities are not the same because constant c is different
- ▶ To compute c , sum a finite geometric series

$$1 = \sum_{i=0}^J c\alpha^i = c \frac{1 - \alpha^{J+1}}{1 - \alpha} \quad \Rightarrow \quad c = \frac{1 - \alpha}{1 - \alpha^{J+1}}$$

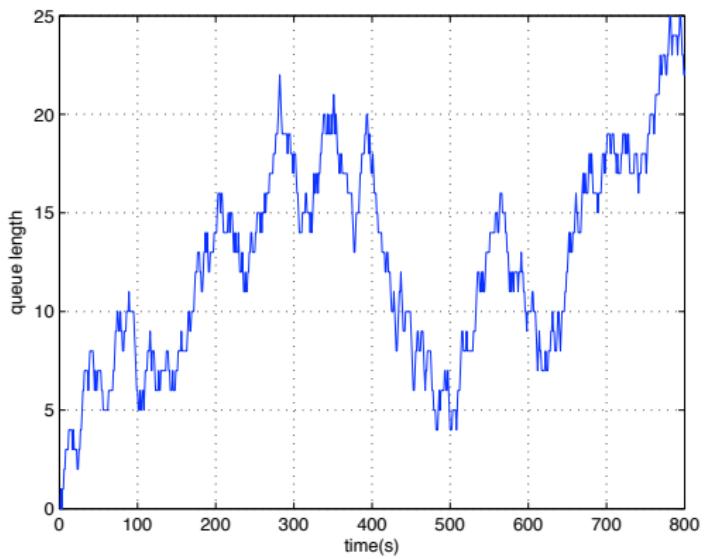
- ▶ Limit distributions for the finite queue are then

$$\pi_i = \frac{1 - \alpha}{1 - \alpha^{J+1}} \alpha^i \approx (1 - \alpha) \alpha^i$$

- ▶ with $\alpha = \lambda(1 - \mu)/\mu(1 - \lambda)$, and approximation valid for large J
- ▶ Approximation for large J yields same result as infinite length queue
 - ▶ As it should

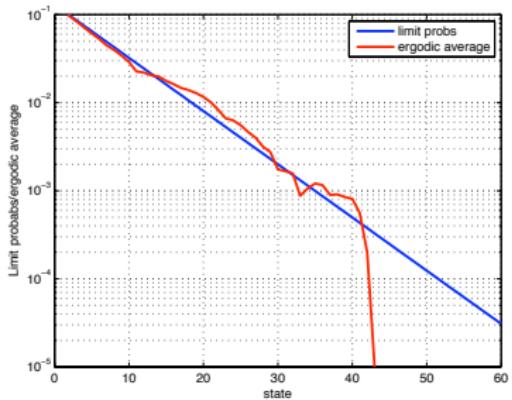
- ▶ Arrival rate $\lambda = 0.3$. Departure rate $\mu = 0.33$. Resulting $\alpha \approx 0.87$
- ▶ Maximum queue length $J = 100$. Initial state $q_0 = 0$ (queue empty)
 - ▶ Not the same as initial probability distribution

Queue lenght as function of time

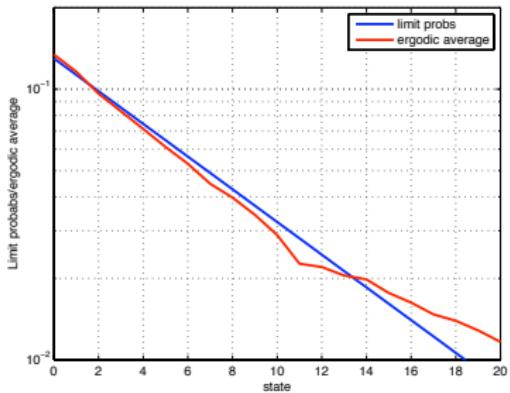


- ▶ Average time spent at each queue state is predicted by limit distribution
- ▶ For $i = 60$ occupancy probability is $\pi \approx 10^{-5}$.
 - ▶ Explains inaccurate prediction for large i

60 states



20 states



- ▶ If $\lambda = 0.45$ and $\mu = 0.55$ how many packets the buffer needs to hold to have a drop rate smaller than 10^{-6} (one packet dropped for every million packets handled)
- ▶ What is the probability of buffer overflow?
- ▶ Packet discarded if queue is in state J and a new packet arrives

$$P[\text{overflow}] = \lambda \pi_J = \frac{1 - \alpha}{1 - \alpha^{J+1}} \lambda \alpha^J \approx (1 - \alpha) \lambda \alpha^J$$

- ▶ With $\lambda = 0.45$ and $\mu = 0.55$, $\alpha \approx 0.82 \Rightarrow J \approx 57$
- ▶ A final caveat
 - ▶ Still assuming only 1 packet arrives per time slot
 - ▶ Lifting this assumption requires introduction of continuous time Markov chains