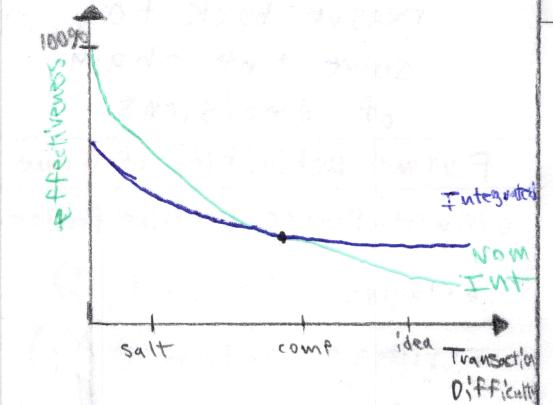


Non-Integrated vs. Integrated Transactions



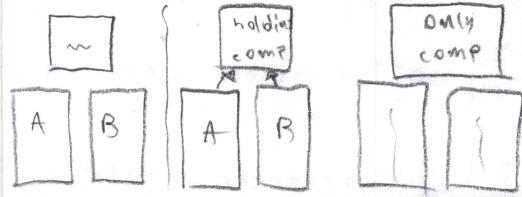
contract
Regime Better
by bosses

Key Points:

- Effectiveness measures what % of value creation you reap, rather than put towards safeguards
- with integrated structure you'll never get 100%

* More complex the good / Transaction the more effective an Integrated structure is

Types of organizations



Fundamentals of Agency Theory

simplest organization { o one Boss
o one agent

ex. shareholder = principal
CEO = "Agent"

What is an Economic Model?

→ simplified description of reality where

- Assumptions are explicit
- robust within environment of interest

Basic Principal-Agent Problem

o Incentive → "cash-payment for a measured outcome"

14.26 Summary Sheet

- o rewards → outcomes that people care abt
- o effort → actions people won't take w/o rewards
- o incentives → link reward & effort

Primary Model Vars.

1. (y) agent's contribution to firm value
2. (a) Action taken by agent to increase firm value
3. (E) events in production process beyond agent's control noise

Contracts

- o (w) is agent's wage a linear function of output

$$wage = S + (b * y)$$

↳ S = salary

↳ b = bonus rate

↳ y = output

Payoffs

- o Principal Payoff (Profit)

$$\Pi = y - w$$

- o c(a) → dollar amount compensating agent for action a

- o Agent's Payoff: $U = w - c(a)$

Event Timing

1. Principal & Agent sign compensation contract $(w = s + b * y)$
2. Agent chooses action (a), but Principal does not observe
3. Events (E) occur beyond control

4. Action & noise determine output y

5. Output observed by the Principal & Agent

6. Agent receives compensation

Risk-Neutral Agent Response

- o Agent chooses action that maximizes

$$\max_a S + (b * a) - c(a)$$

→ higher bonus rate b, stronger incentive to work more

Necessity of performance Measures

- o y not that easily countable

o replace y → p

↳ where p is "alternate performance measure"

→ Dichotomy of p & y makes interesting problems

Multi-Task Agency Model

Technology of production

$$y = f_1 a_1 + f_2 a_2 + \epsilon$$

Technology of Performance Measure

$$p = g_1 a_1 + g_2 a_2 + \phi$$

contract

$$w = s + b p$$

Payoffs { agents:

$$v = w - c(a)$$

Principal: $\Pi = y - w$

where,

$$c(a_1, a_2) = \frac{1}{2} a_1^2 + \frac{1}{2} a_2^2$$

solution

→ agent chooses a_1 & a_2 to maximize

$$\max_{a_1, a_2} \{ S + b(g_1 a_1 + g_2 a_2) - \frac{1}{2} a_1^2 - \frac{1}{2} a_2^2 \}$$

1st order optimization

$$a_i(b) = b * g_i$$

Does paying for θ increase y

- $y = a_1 + \epsilon$ { $\rho = a_1 + \phi$ (agent wants to choose a_1)}
- $y = a_1 + a_2$ { $\rho = a_1 + a_2$ (e.g. agent face time)
- $y = a_1$ { $\rho = a_2 + \epsilon$ (mis-aligned incentives)}

Good performance measure $\text{cov} \left(\frac{\partial y}{\partial a}, \frac{\partial \rho}{\partial a} \right) > 0$

Elemental Cos(θ) Model

- $y = f_1 a_1 + f_2 a_2$ { $\rho = H \leq L$ }
- $\pi = y - w$ (Binary now)
- $U = w - c(a_1, a_2)$ { $w = s + b\rho$ ($\rho = H$) }
-
- Agent has outside opportunity worth V_0

Solution

Step 1: labor maximization (agent)

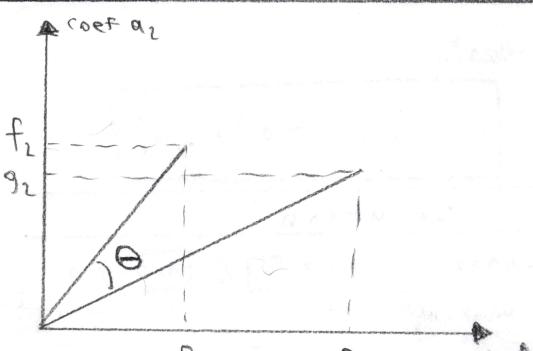
$$\max_{a_1, a_2} s + b(g_1 a_1 + g_2 a_2) - \frac{1}{2}(a_1^2 + a_2^2)$$

$$\Rightarrow \begin{cases} a_1(b) = b g_1 \\ a_2(b) = b g_2 \end{cases}$$

Step 2: Agent bonus Maximization

$$\max_b [f(y) - c(a_1, a_2)] = V_0$$

Graphical interpretation of sol



$$b^* = \frac{\|f\|}{\|g\|} * \cos(\theta)$$

Variations on Problem

Example 1: Manipulation

- $y = f * \rho$
- $\rho = \epsilon + m$
- $w = s + b\rho$
- $c = \left(\frac{1}{2} (\epsilon^2 + \lambda m^2) \right)$

$$\rightarrow \max_{\epsilon, m} E(w) - c(\epsilon, m)$$

$$\rightarrow \max_b E(y) - c(\epsilon, m)$$

Example 2: Principal Principle

- $y = a_1 + a_2$
- $\rho = g_1 a_1 + m$
- Actor chooses a_1, m
- Principal Receives y

$$\rightarrow \text{pays } w_i = s_i + b_i \rho_i + (B_i p_i)$$

$$\rightarrow \max_b E(a_1) - c(a_1, m)$$

\Rightarrow actor w/ lower g_1 should be principal

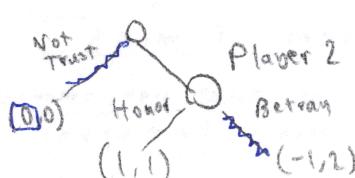
Dynamic Games w/ Complete Info

- Player 1 chooses action a_1 from set A_1
- Player 2 observes Player 1's choice then chooses a_2 from set A_2 .
- Each Player receives payoffs

\rightarrow Money useful metric to capture utility of payouts

Example

Player 1



* To solve, we must reason back to solve the chain of decisions

Formal Definition of Game

- Multi-Player Decision Problem
- players $i = 1, 2, \dots$
- strategies (a_i from set A_i)
- utilities, $u_i(a_1, a_2), u_2(a_1, a_2)$ (aka payoffs)

Physical outcomes

Column n		Not
Sust	Sust, Sust	Sust, Not
Not	Not, Sust	Not, Not

Possible utilities for Row: (supposing sustainable is costlier)

8 4 → Top concern is total
6 2 S

4 2 → Top concern is own cost
8 6

8 6 → Top concern is total S,
4 2 2nd is identity

8 4 → Top preference: solidarity
4 6 2nd is total S

→ utilities reveal players preferences

Best-Response

Given i's belief about what his opponent will play we say as strategy is best response if it gives at least as high utility as does any other strategy

Nash Equilibrium

A profile of strategies is a **Nash Equilibrium** if each player's strategy is a best response to other players' strategy

Criteria

$$u_1(e_i, s) \geq u_1(e_i, t)$$

8

$$u_2(e_j, s) \geq u_2(e_j, r)$$

Specific Games

Game 1.5: Stag Hunt

		column
		L R
Row	T	(2, 2) (0, 1)
	B	(1, 0) (1, 1)

→ No dominated strategy for either player

* 2 Pareto-Ranked Nash Equilibria

Game 1.8: Dating Game

	(2, 1)	(0, 0)
	(0, 0)	(1, 2)

→ 2 Nash equilibria

Game 1.6 : Prisoners' Dilemma

o One Nash Equilibrium

↳ Pareto-Inefficient

Tournaments

Relational contract

- shared understanding of parties' roles & rewards from collaboration

→ Not enforceable

* Equilibrium of repeated game

Repeated Games

- People interact over time, threats & promises of future behavior influence current behavior

Module 1 Review

LARB

(2, 2)	$\begin{pmatrix} 1, 1 \\ 0, 0 \end{pmatrix}$
$\begin{pmatrix} 5, 0 \\ 0, 5 \end{pmatrix}$	(1, 1)

Two-Stage Game

- Prisoner's Dilemma to be played 2-times

	L1	R1
L2	1, 1	5, 0
R2	0, 5	4, 4

- outcome of 1st play observed

↓ evolves to sum of two payoffs

2, 2	6, 1
1, 6	5, 5

- still a prisoner's dilemma

(no punishment for deviation from global optimal strat)

Module 3

Module 3: 14.26

Cheat sheet

Game w/ Punish

1, 1	5, 0	0, 0
0, 5	4, 4	0, 0
0, 0	0, 0	3, 3



2, 2	6, 1	0, 0
1, 6	7, 7	0, 0
0, 0	0, 0	3, 3

- middle selection reflects belief that if middle option achieved in 1st stage 3, 3 played → otherwise (1, 1)

BUT, would we really ever go for 7, 7 when we have right

Infinitely Repeated Games

Key Point: moves in

1st period become part of 2nd stage equilibrium

Trigger strategies

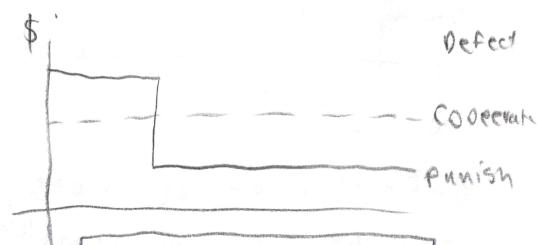
- infinite horizon

⇒ initially will trust, but will deviate forever after if you do

Do you Deviate?

↓
Depends on interest Rate

Decision boils down to 2-time paths:



$$(1 + \frac{1}{r})C > D + \frac{1}{r}P$$

"present value"

Modelling Rational contracts

- Player 1: play Trust until player 2 plays betray; thereafter play not trust

Player 2:

- Message can't be punished
- can't be trusted

Formalization

Nash Equilibrium

- strategy pair (s_1, s_2)
- s_i is i 's best response to s_j

Present Value

→ interest Rate

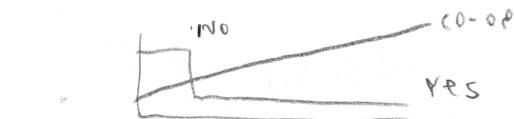
Betray or Trust

Yes if $(1 + \frac{1}{r})C > D + (\frac{1}{r})P$

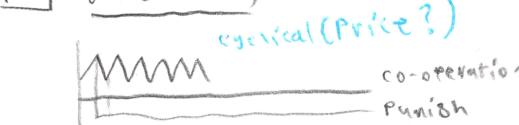
No if other

Alterations to infinite horizon

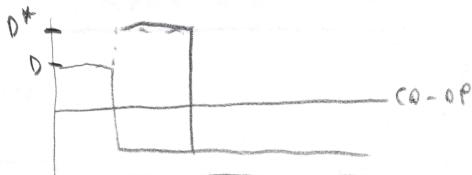
Trends



Seasonality



3 permanent shock



- Defection pay off goes $D \rightarrow D^*$ w/ some probability

Expected Present Value used to determine if trust holds

Adaptation

(A) Timing

- Asset ownership/control rights designed [9]

2.

- state realized $[s]$

- Ex-post decision made

5. Payoffs

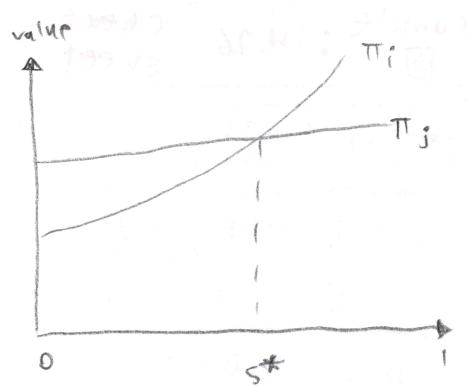
Timing of Relational Adaptation

Once & for all:

- choose Gov. structure
 - control to i: j can pay $i(R)$
 - $s_{+} \in f(s)$ on $[0, 1]$
 - $d \in D$
- "Retainer"
- state
- ex-post decision

4. Payoffs

$$\begin{cases} \Pi_i(d_t/s_r) + R_+ \\ \Pi_j(d_t/s_r) - R_+ \end{cases}$$



First-Best Decision Rule

$$\begin{cases} d_j^*(s) \text{ for } s < s^* \\ d_i^*(s) \text{ for } s > s^* \end{cases}$$

\Rightarrow can't be achieved in one-shot game

Repeated Reln. Adaptation

Q1: If party i has control, for what values of s^* can FB be achieved in repeated game

Q2: Who should have control to achieve first-Best in repeated game at highest r

Suppose A promises d^{FB} :

B promises R

\rightarrow A has control overed

who's player I

- what happens after renege
- what are trigger strategies

* What values of R are trigger strats in equilibrium

Key point two parties want 2 things in different states

Promises Kept

$$E(\Pi_A) = R + \int_{s=s^*}^1 \Pi_A(s) f(s) ds$$

$$E(\Pi_B) = -R + \int_0^{s=s^*} \Pi_B(s) f(s) ds$$

After renege

$$E(\Pi_A) = \int_{s=s^*}^1 \Pi_A(s) f(s) ds$$

$$E(\Pi_B) = 0$$

\rightarrow A's max renege temptation

$$\begin{cases} \Pi_A(s^*) \\ -R \end{cases}$$

\rightarrow B's willingness to participate

$$-R + \int_{s=0}^{s=s^*} \Pi_B(s) f(s) \geq 0$$

\rightarrow A willing to implement first-best decision

$$(R) + \frac{1}{r}(R + \int_{s=s^*}^1 \Pi_A(s) f(s) ds) \geq$$

$$\{ R + \Pi_A(s^*) + \frac{1}{r} \int_{s=s^*}^1 \Pi_A(s) f(s) ds \}$$

\rightarrow greatest moment of temptation occurs after R and stands to gain $\Pi_A(s^*)$

→ Maximum B willing to pay to \boxed{B}

$$= \int_0^{s^*} \pi_B(s) f(s) ds$$

→ maximum \boxed{r} for A to implement FB decision

$$\frac{1}{r} \left[\int_0^{s^*} \pi_B(s) f(s) ds + \int_{s^*}^1 \pi_A(s) f(s) ds \right]$$

$$- \int_0^1 \pi_A(s) f(s) ds \geq \pi_A(s^*)$$

or

$$\frac{V^{FB} - V_A^{SP}}{\pi_A(s^*)} = r_A > r$$

• If critical case holds,

All else holds

Who gets control

• $\pi_B(1) > \pi_A(s^*) \rightarrow \boxed{B}$ gets control

→ b/c magnitude of defection temptation lowest w/ A

• $V_B^{SP} < V_A^{SP} \rightarrow \boxed{B}$ gets control

→ because spot payoff smaller for B

⇒ want to give control to agent who gets us first best when \boxed{B} 's high, agents impatient

Rule: give control to agent w/ highest critical \boxed{s}

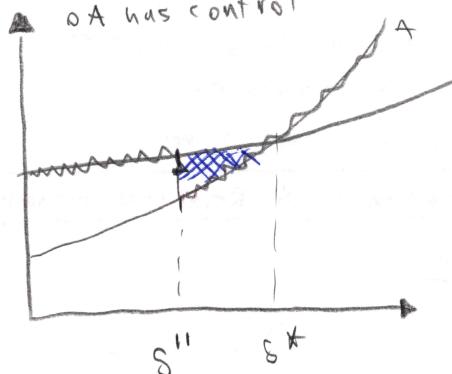
"Second-Best"

what if $r > \max(r_A, r_B)$

control to both would result in betrayal

→ worse than 1st best

⇒ Better than
so far
 ΔA has control



☒ Lost best joint utility picking second-best

→ Basically redetermine critical \boxed{s} to make it acceptable to \boxed{A}

→ control decision still on same criteria

Discretionary Bonuses

→ given relational contract bonus \boxed{b}

Agent problem:

$$\max_a s + (a \cdot b) - c(a)$$

$$\begin{cases} \text{Principal} \\ \text{Profit} \end{cases} \quad \begin{aligned} V(b) &= L + a'(b)[H-L] \\ -c(a'(b)) - w_a \end{aligned}$$

Formalization of Problem

• Production

$$\begin{cases} y = H \text{ or } L \end{cases}$$

$$\text{Prob}(y=H|a) = a$$

• Payoffs

$$P: \pi = y - w \quad \begin{cases} A: U = w - c(a) \end{cases}$$

→ discount at rate r

• First-Best Constraint

$$a^{FB} \text{ solves } \max_a L + a \cdot (H-L) - c(a)$$

• Timing

① P proposes (s, B)

② A accepts

③ A chooses action $a \geq 0$

④ P & A observe y

⑤ P decides whether to pay bonus

• (relational contract)

⇒ P pays w=s if $y=L$

P pays s+B if $y=H$

Q: Is relational contract self-enforcing

Solving Discretionary Bonus Problem

- if relative contract self-enforcing ($\alpha^*(B)$ believes B will be paid)

$$\boxed{\alpha^*(B)} \text{ solves } \max_{\alpha} \text{ s.t. } \alpha^*(B) = c(\alpha)$$

IR constraints

1) Agent Utility constraint

$$U_0 \leq EU(s, B) = s + \alpha^*(B) \cdot B - c(\alpha^*(B))$$

2) Principal Profit constraint

$$\Pi_0 \leq E\Pi(s, B) = (L - s) + \alpha^*(B) \cdot (H - L - B)$$

* constraint governing if principal pays

$$(H - S - B) + \frac{1}{r} \cdot E\Pi(s, B) >$$

$$(H - S - 0) + \frac{1}{r} \cdot \Pi_0$$



$$\boxed{rB < E\Pi(s, B) - \Pi_0}$$

Inequality in terms of value to both parties

$rB < V(B) - V_0$

* sufficient inequality

$$V(B) = E[y] - c[\alpha]$$

$$V_0 = U_0 + \Pi_0$$

FBS w/ Bonuses

• First-Best Bonus

$$B^{FL} = H - L \rightarrow \alpha^*(B) = \alpha^{FB}$$

• Second-Best Bonus

$B^{SB}(r)$ is largest B solving

$$V(B) - V_0 = rB$$

- None { does not exist B s.t. $V(B) - V_0 > rB$

Neighboring Environments

- can control be designed?

→ No or Yes

- When might there be payments
→ at any time!

Formal & Relational contracts

Model:

$$\Pr(y=1) = \alpha_1$$

$$\Pr(\rho=H) = \alpha_1 \cos(\theta) + \alpha_2 \sin(\theta)$$

cost of effort $k(\alpha_1^2 + \alpha_2^2)/2$

agents outside
reputation: \boxed{M}

First-Best

$$\alpha_1 = 1/K$$

$$\alpha_2 = 0$$

$$\text{Total surelty} = 1/(2K)$$

Spot contract case

$$W = S + b \cdot I_{\{\rho=H\}}$$

$$U = S + b [\alpha_1 \cos(\theta) + \alpha_2 \sin(\theta)] - \text{cost}$$

$$\alpha_1 = \left(\frac{b}{K}\right) \cos(\theta), \quad \alpha_2 = \left(\frac{b}{K}\right) \sin(\theta)$$

$$b^* = \cos(\theta)$$

Three Lenses on organizational processes

① Formal

- organizations are Machines
- parts must fit well
- action through Planning

② Political

- organizations are contests
- action through power

③ Cultural

- organizations are institutions

→ action through habit

Team Theory

- 1st economic theory of internal organization

- cost of gathering, communicating, processing info

⇒ organization is rational individual

communication & decision making

- Project value { $y=0$ or 1

- Signal to $i=1, 2$ { $s_i = G$ or B

- Implementation cost \boxed{C}

Assume

$$\Pr(y=1 | G, G) > c > \Pr(y=1 | B, B)$$

→ $\boxed{\Pr(y=1, G, B)}$?

Hierarchy: Unanimity to accept, optimal if $P(y=1|G, B) < c$

Polyarchy: Unanimity to reject, optimal if $P(y=1|G, B) > c$

Imperfect communication

Model:

1. state of world SFS realized from $s \sim f(s)$

2. Player 1 privately observes s

3. Player 1 chooses (c, m)

4. Player 2 observes m w/ probability p , nothing w/ $(1-p)$

5. Player 2 chooses dFR

6. Payoff

$$U_1(s, c, d) = -\phi(c-s)^2 - \beta(d-c)^2$$

→ team-theoretic payoffs (same for both)

→ messages costless

Takeaway: the more accurate player 2 can guess d , the more they both benefit

simulated Run-through

- Player 1 chooses $m = c$
- If m does not arrive
Player 2 chooses $d^* = E(c)$

Player 1 chooses c to maximize

$$-\phi(c-s)^2 - \beta(1-p)(E(c)-c)^2$$

b/c \square % of time can choose $d=c$

$$c^*(s) = \frac{[\phi s + \beta(1-p)E(c)]}{[\phi + \beta(1-p)]}$$

$$E(c) = \frac{[\phi E(s) + \beta(1-p)E(c)]}{[\phi + \beta(1-p)]}$$

Plug back in

• comparative statics on ϕ & p

knowledge hierarchies

• different problems

require different skills to be solved

→ index problems \mathbb{B} skills by \underline{z}

• learning one skill teaches nothing about other

key: problems arrive randomly w/ $P(z)$

• value of solving any problem is \square for any z

• cost of training is c per skill

• workers can be trained on any skill

→ * will be optimal to train one worker in some, and another in others

• Each worker has total time T

→ workers can solve or help

Assumption on Processing times

• When Problem \square arrives, a worker with skill set S_i knows if they can solve → takes P to do so

• "if" similar, worker knows if they can help with it → takes H to do so

• team-theoretic costs

Case 1: workers trained in Disjoint sets

- Some have skill set S_1
- Else, S_2

□ Firm can solve problem z

$$\text{if } z \in S_1 \cup S_2$$

N → # of problems firm will consider

α → fraction of w workers w/ skill set S_1

* Maximize expected number of problems solved

→ Sort each workers expected time

S_1 's Time

$$N \Pr(S_1) P + \alpha W T$$

S_2 's Time

$$(1 - \alpha) W T$$

Firm's expected output:

$$N * \Pr(z \in S_1 \cup S_2)$$

Three Lenses on Org. Processes

Formal → organizations are machines

o An organization is a mechanical system crafted to achieve a defined goal.

⇒ Action comes through Planning

Political → organizations are contests

o An organization is a social system encompassing diverse, contradictory interests.

⇒ Action comes through

Cultural → organizations are institutions

o A symbolic system of meaning, value, & routines. Informal norms exert strong influence.

⇒ action comes through habit

Team Theory

• 1st economic theory of internal organization

• common interests → no shirking, lying

organization as rational individual

Model 1 (^{communication}_{& Decision Making})

□ Project Value ($V=0$ or 1)

□ Signal to $i=1,2$ ($s_i = G$ or B)

□ Implementation cost c

• Assume $\text{Prob}(y=1|G,G) \geq c$

$$\geq \text{Prob}(y=1|B,B)$$

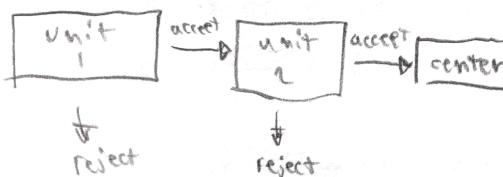
$$Q: \Pr(y=1|G,B) > c \quad ?$$

Module 4: Summary Sheet

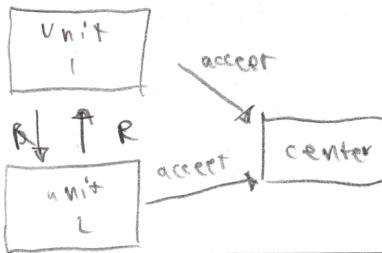
Comm'g & D-M

Building a circuit w/
2 kinds of components

Hierarchy: Unanimity to Accept
optimal if $\underline{\text{P}(y=1|G,B)} < c$



Polyarchy: Unanimity to reject
optimal if $\underline{\text{P}(y=1|G,B)} > c$



Imperfect Communication

1. state of world $[S|S]$

realized: $S \sim f(S)$ on R

2. Player 1 observes $[S]$

3. Player 1 chooses (c, m)

4. Player 2 observes m w/
probability ϕ , but not
w/ $[1-\phi]$

5. Player 2 chooses

$d \in R$

6. Payoffs:

$$U_i(s, c, d) = -\phi(c-s)^2 - \beta(d-c)^2$$

• Adaptation vs.
coordination

• TRR game-theoretic payoffs

• Messages costless but
imperfect

Imperfect comm. (Play-Through)

• Player 1 chooses $[m=c]$

• If m does not arrive
then Player 2 chooses

$$d^* = E(c)$$

• Player 1 chooses $[c]$ to
maximize

$$-\phi(c-s)^2 - \beta(1-\phi)(E[c] - c)^2$$

⇒ maximize utility outcome

$$c^*(s) = [\phi s + \beta(1-\phi) E[c]] / [\phi + \beta(1-\phi)]$$

* weighted avg of adaptation

$$c=5 \quad \text{8 coordination } c=E[c]$$

$$E[c] = [\phi E(s) + \beta(1-\phi) E(c)] / [\phi + \beta(1-\phi)]$$

$$\rightarrow c^*(s) = E(s) + (\phi / [\phi/\phi + \beta(1-\phi)]) \cdot (s - E(s))$$

Comparative Statics

$\phi \rightarrow$ weight on adaptation

$\beta \rightarrow$ weight on coordination

$\epsilon \rightarrow$ probability weight

Knowledge Hierarchies

- different Problems require different skills to be solved

→ index problems by \boxed{z}

(problem \boxed{z} requires skill \boxed{s})

→ nearby values of \boxed{z} , are not necessarily related problems

* Problems arrive randomly, but w/ different probabilities $\boxed{f(z)}$

• Assign \boxed{z} so that more likely problems have smaller \boxed{z}

□ Value of solving any problem is $\boxed{1}$

→ cost of training a skill is \boxed{c}

→ will be optimal to train workers on range of skills

workers [that encompasses whole range of tasks]

• Each worker has total time \boxed{T}

→ work on problems

$h \rightarrow$ help others w/ problem

$$\begin{array}{c} \boxed{t_h} + \boxed{t_p} = \boxed{T} \\ \hline \end{array}$$

Processing Time

If worker can solve it, takes time \boxed{P}
if not, takes \boxed{H}

Two-Skillset Generalization

→ some workers have skills $\boxed{S_1}$ & others disjoint set $\boxed{S_2}$

\boxed{N} = # of problems firm will consider

$\boxed{\alpha}$ = fraction of \boxed{W} workers w/ skillset S_1

→ maximize expected # of problems solved

suppose $\boxed{S_1}$ "answers phone" & $\boxed{S_2}$ helps

$$\begin{aligned} S_1's\ Time &= N \cdot \Pr(S_1) P \\ &+ N[1 - \Pr(S_1)] \Pr(z \in S_2 | z \notin S_1) P \\ &= \boxed{\alpha WT} \end{aligned}$$

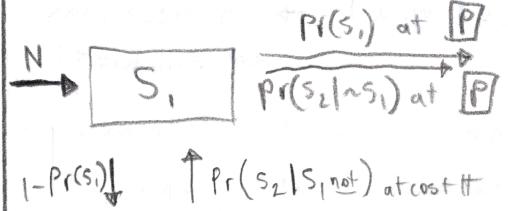
$$\begin{aligned} S_2's\ Time &= N[1 - \Pr(S_1)] \Pr(z \in S_2 | z \notin S_1) P \\ &\quad || \\ &= (1 - \alpha) WT \end{aligned}$$

$$\begin{aligned} \text{Firm's Expected Output} &= N \cdot \Pr(z \in S_1 \cup S_2) \\ &= \boxed{N \cdot \Pr(z \in S_1, US_2)} \end{aligned}$$

$$S_1: \alpha = (NP/WT) * \left\{ \Pr(S_1) + \Pr(S_2) \right\}$$

$$S_2: \quad \quad \quad$$

Pictorially



$\boxed{S_2}$

Solving

$$\text{Recognize } \alpha = \left(\frac{NP}{WT} \right) \cdot \left\{ \Pr(S_1) + \Pr(S_2) \right\}$$

= solve for \boxed{N}

Total Output

$$WT \{ \Pr(S_1) + \Pr(S_2) \} / \left[(\Pr(S_1) H + P \cdot \{ \Pr(S_1) + \Pr(S_2) \}) \right]$$

Suppose Both Answer Phone & Help

- choose \boxed{N} # of problems
- choose $\boxed{\alpha}$ fraction of \boxed{W} workers w/ skillset S_1
- choose $\boxed{\beta}$ fraction of N problems arriving to $\boxed{S_1}$ workers

$$S_1: \boxed{\beta N \Pr(S_1) P + \beta N[1 - \Pr(S_1)]}$$

$$\Pr(z \in S_2 | z \notin S_1) P +$$

$$(1 - \beta) N[1 - \Pr(S_2)] \Pr(z \in S_1 | z \notin S_2)$$

$$H = \alpha WT$$

- ① Problem to S_1, S_1 can solve
- ② Problem + S_1, S_1 solved w/ help
- ③ Problem to S_2, S_1 helps on

- Analogous as before
solve for \boxed{B} \boxed{N}

$$\text{Output} \mid w \in \{P(S_1) + P(S_2)\} / \boxed{D}$$

→ output will be less than before

suppose: All workers receive identical training

- $N \rightarrow \infty$
- optimal output = 0

Takeway: Majority handles majority of problems

→ management handles exceptions

Real Organizations

- Decision Theory predicts relevant info will be gathered & analyzed prior to decision-making

→ organization is not a rational being

Formal vs. Real Authority

- Superior often "rubber-stamps" a subordinate's proposal
- Shareholders → Board of Directors → CEO
- Why and how would actor w/ formal authority cede real authority

A: other actor has better info & similar preferences

The Model

- $\boxed{3}$ potential projects ($K=1, 2, 3$)
- Benefits B_K to \boxed{P} , b_K to \boxed{A}
- one project has $B_K = b_K = -\infty$
- other payoffs are
 $B_K = B$ or 0 ($B > 0$)
 $b_K = b$ or 0 ($b > 0$)

* Alignment probability \boxed{d}

w/ $\Pr(d) \rightarrow (B, b)$ or (0, 0)

w/ $\Pr(1-d) \rightarrow (B, 0)$ or (0, b)

Info \otimes Control

→ A pays $\boxed{c_A(e)}$ for soft info

→ $\Pr(e)$: A learns all own payoffs

→ $\Pr(1-e)$: A learns nothing

• P pays cost $c_P(e)$ for soft info

→ $\Pr(E)$: P learns all own payoffs

→ $\Pr(1-E)$: P learns nothing

• \boxed{P} formal authority ✓

• \boxed{A} formal authority ✓

P-Formal Authority

- P informed → chooses B
- P uninformed, but A informed → chooses b

Expected Payoffs:

$$U_P = E B + (1-E) \epsilon d B - c_P$$

$$U_A = E b + (1-E) \epsilon b - c_A$$

* Note: if neither informed stick to status quo (0, 0)

→ w/ probability $(1-E)\epsilon$
A has real authority

Best responses for \boxed{P} \otimes \boxed{A}

$$\max_E E B + (1-E) \epsilon d B - c_P$$

$$\max_E E b + (1-E) \epsilon b - c_A$$

• Suppose functional cost forms

$$\left\{ \begin{array}{l} c_P = (K/2) E \\ c_A = (m/2) \epsilon^2 \end{array} \right.$$

$$\Rightarrow E = B(1-\epsilon)/K$$

$$\epsilon = (1-E)b/m$$

$$E^*$$

$$E^*$$

use to solve for

Empowerment

Timing of one-shot game:

1. Principal offers (s, y^*) to agent

$$s = \text{salary}$$

y^* → promise to implement
 $y \geq y^*$

2. agent accepts (paid) or rejects (U_0, π_0)

3. If Accept, A chooses

$$a \in [0, 1] \text{ at cost } c(a)$$

• P does not observe a

• Project delivered w/ probability α

$$\begin{aligned} c(0) &= 0, c'(0) = 0 \\ \Rightarrow c'(a) &\rightarrow \infty \text{ as } a \rightarrow 1 \end{aligned}$$

* certainty of completion carries high price

4. P & A observe a project if it exists & payoffs (x, π)

$$x > 0$$

$$y_L < y_M < 0 < y_H \text{ w/ prob } s$$

$$\alpha_L, \alpha_M, \alpha_H$$

5. If Project exists, principal chooses whether or not to implement

6.

$$\begin{cases} \text{if implemented} & \begin{cases} \pi = y - s & U = x + s - c(a) \\ \pi = -s & U = s - c(a) \end{cases} \\ \text{If not} & \end{cases}$$

P will ONLY implement

$$y^* \rightarrow A \text{ chooses } a^*$$

$$\text{Solving } \max_a \begin{cases} s + \alpha_H ax - c(a) \\ \dots \end{cases}$$

Result of one-shot game

→ unwilling to give high effort if low-chance of implement

Repeated Game Scenario

can $y^* = y_M$ in repeated-game

• In one-shot, P will not do $y < 0$

⇒ Agent's effort higher when y^* is lower

• Let a_{HM} denote solution of

$$\max_a s + (\alpha_{HM} + \alpha_{HT}) ax - c(a)$$

• Agent will accept provided

$$s + \alpha_{HM}(\alpha_H + \alpha_H) - c(a_{HM}) \geq U_0$$

$$E[\pi_{HM}] = \alpha_{HM}(\alpha_H + \alpha_H) - c(a_{HM}) - U_0$$

Expected pay off to principal

$$E[\pi_{HM}] = \alpha_{HM}[\alpha_{HM}(x+y_M) + \alpha_{HT}(x+y_H) - c(a_{HM}) - U_0]$$

Relational contract aspect

• will P implement y_M ?

$$C_H + \frac{1}{r} E[C_T] \geq D_H + \left(\frac{1}{r}\right) E(P_T)$$

$$(y_M - s) + \frac{1}{r} E[\pi_{HM}] \geq -s + \left(\frac{1}{r}\right) \cdot \pi_0$$

$$\begin{cases} \text{condition to not renege} \\ \left(\frac{1}{r}\right)[V_{HM} - V_0] \geq -y_M \end{cases}$$

positive #

→ principal always tempted to renege

Influence Activities

• Care what decisions are taken

→ care what decision makers believe

→ P wants $d_P^*(s)$

→ A wants $d_A^*(s)$, can influence

$$\sigma = s + \alpha_\sigma + \epsilon$$

→ A distracted from my goal

Influence (Model)

1. Nature determines state S

$$\Rightarrow S \sim N(m, h)$$

2. Agent chooses "lobbying"
A at cost $K(\lambda)$

3. Principal observes signal

$$O = S + \lambda + \varepsilon$$

4. Principal chooses decision d

5. Parties receive payoffs

$$U_i(S, d)$$

Key points

• S not observed by

P or A

• λ not observed by P

Key Questions

□ How will Principal update belief $f(S|O)$

□ What decision d will Principal choose given belief

□ What lobbying λ will agent choose

1761 7/3 moved to E

1761 7/3 moved to E