

# Types of Cache Misses: *The Three C's*

- 1 Compulsory:** On the first access to a block; the block must be brought into the cache; also called cold start misses, or first reference misses.
- 2 Capacity:** Occur because blocks are being discarded from cache because cache cannot contain all blocks needed for program execution (program working set is much larger than cache capacity).
- 3 Conflict:** In the case of set associative or direct mapped block placement strategies, conflict misses occur when several blocks are mapped to the same set or block frame; also called collision misses or interference misses.

# Miss Rate Reduction Techniques:

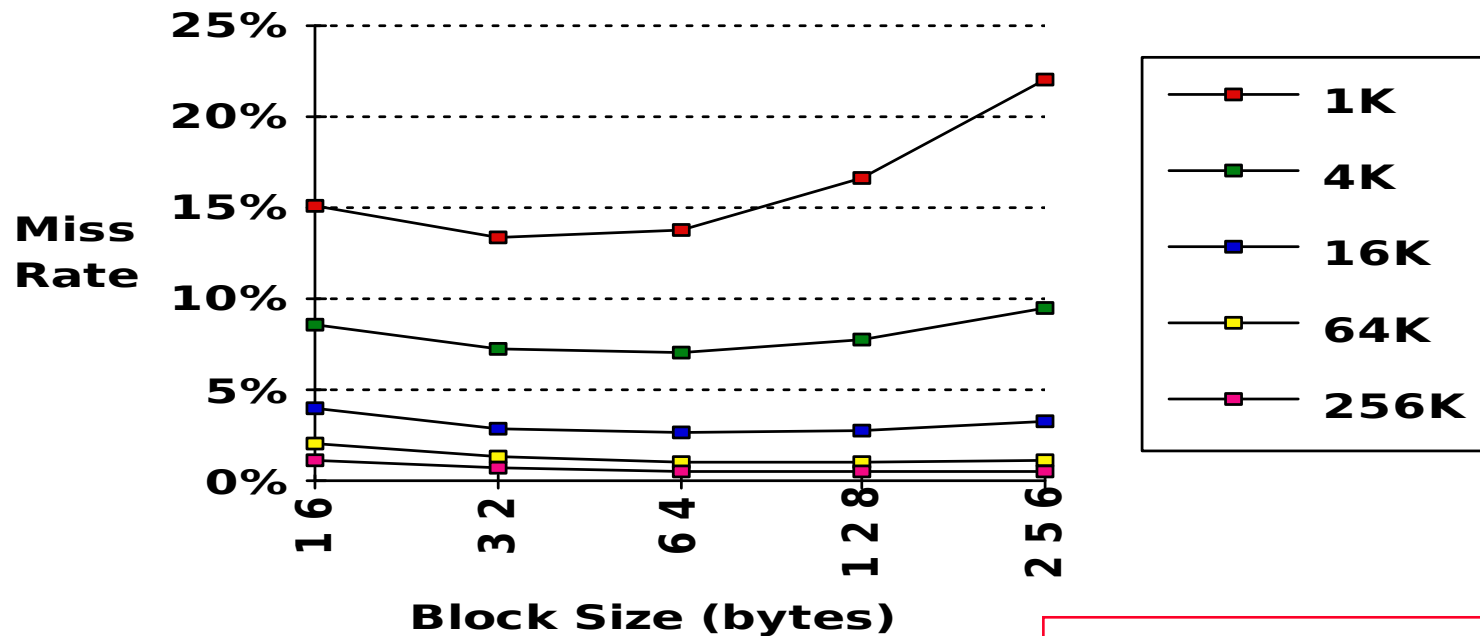
- \* **Increased cache capacity**
- \* **Larger block size**
- \* **Higher associativity**
- \* **Compiler optimizations**

## Miss Rate Reduction Techniques:

# Reduce Misses by Larger Block Size

- A larger block size improves cache performance by taking advantage of spatial locality
- For a fixed cache size, larger block sizes mean fewer cache block frames

Performance keeps improving to a limit when the fewer number of cache block frames increases conflict misses and thus overall cache miss rate



## 2. Reduce Misses by Increasing Cache Size

- **Increasing cache size reduces cache misses**
  - both capacity misses and conflict misses reduced

## Miss Rate Reduction Techniques:

# Higher Cache Associativity

- Reduces conflict misses

Cache Size		Associativity			
(KB)		1-way	2-way	4-way	8-way
1	2.33	2.15	2.07	2.01	
2	1.98	1.86	1.76	1.68	
4	1.72	1.67	1.61	1.53	
8	1.46	<u>1.48</u>	<u>1.47</u>	1.43	
<u>16</u>	<u>1.29</u>	<u>1.32</u>	<u>1.32</u>	<u>1.32</u>	
<u>32</u>	<u>1.20</u>	<u>1.24</u>	<u>1.25</u>	<u>1.27</u>	
<u>64</u>	<u>1.14</u>	<u>1.20</u>	<u>1.21</u>	<u>1.23</u>	
<u>128</u>	<u>1.10</u>	<u>1.17</u>	<u>1.18</u>	<u>1.20</u>	

(Red means A.M.A.T. not improved by more associativity)

# Compiler Optimizations

Compiler cache optimizations improve access locality characteristics of the generated code and include:

- ***Merging Arrays***: Improve spatial locality by single array of compound elements vs. 2 arrays.
- ***Loop Interchange***: Change nesting of loops to access data in the order stored in memory.
- ***Loop Fusion***: Combine 2 or more independent loops that have the same looping and some variables overlap.
- ***Blocking***: Improve temporal locality by accessing “blocks” of data repeatedly vs. going down whole columns or rows.

# Merging Arrays

```
/* Before: 2 sequential arrays */
```

```
int val[SIZE];
```

```
int key[SIZE];
```

```
/* After: 1 array of structures */
```

```
struct merge {
```

```
    int val;
```

```
    int key;
```

```
};
```

```
struct merge merged_array[SIZE];
```

- combines two separate arrays (that might conflict for a single block in the cache) into a single interleaved array.
- This brings together corresponding elements in both arrays, which are likely to be referenced together.
- reduces misses by improving spatial locality.

## Miss Rate Reduction Techniques: Compiler-Based Cache Optimizations

# Loop Interchange Example

```
/* Before */
for (k = 0; k < 100; k = k+1)
    for (j = 0; j < 100; j = j+1)
        for (i = 0; i < 5000; i = i+1)
            x[i][j] = 2 * x[i][j];

/* After */
for (k = 0; k < 100; k = k+1)
    for (i = 0; i < 5000; i = i+1)
        for (j = 0; j < 100; j = j+1)
            x[i][j] = 2 * x[i][j];
```

**Sequential accesses instead of striding through memory every 100 words in this case improves spatial locality.**



# Loop Fusion Example

```
/* Before */
for (i = 0; i < N; i = i+1)
    for (j = 0; j < N; j = j+1)
        a[i][j] = 1/b[i][j] * c[i][j];
for (i = 0; i < N; i = i+1)
    for (j = 0; j < N; j = j+1)
        d[i][j] = a[i][j] + c[i][j];

/* After */
for (i = 0; i < N; i = i+1)
    for (j = 0; j < N; j = j+1)
    {
        a[i][j] = 1/b[i][j] * c[i][j];
        d[i][j] = a[i][j] + c[i][j];
    }
```

- *Many programs have separate loops that operate on the same data.*
- *Combining these loops by grouping operations on the same (cached) data together.*