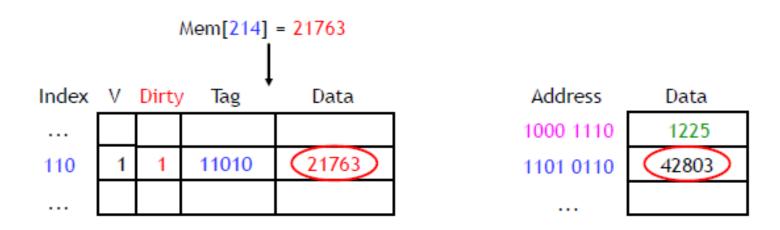
- Review:
 - Write-through?
 - Write-back?

Write-back caches

- In a write-back cache, the memory is not updated until the cache block needs to be replaced (e.g., when loading data into a full cache set).
- For example, we might write some data to the cache at first, leaving it inconsistent with the main memory as shown before.
 - The cache block is marked "dirty" to indicate this inconsistency



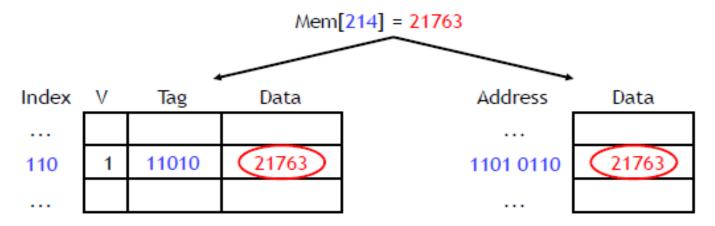
- Subsequent reads to the same memory address?
- Multiple writes to same block?

Write-back cache discussion

- The advantage of write-back caches is that not all write operations need to access main memory, as with write-through caches.
 - If a single address is frequently written to, then it doesn't pay to keep writing that data through to main memory.
 - If several bytes within the same cache block are modified, they will
 only force one memory write operation at write-back time.

Write-through caches

 A write-through cache forces all writes to update both the cache and the main memory.



- This is simple to implement and keeps the cache and memory consistent.
- The bad thing is that forcing every write to go to main memory, we use up bandwidth between the cache and the memory.

Write misses

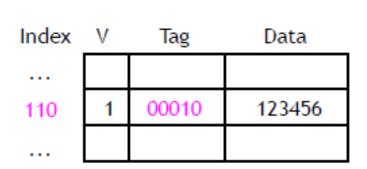
- A second scenario is if we try to write to an address that is not already contained in the cache; this is called a write miss.
- Let's say we want to store 21763 into Mem[1101 0110] but we find that address is not currently in the cache.

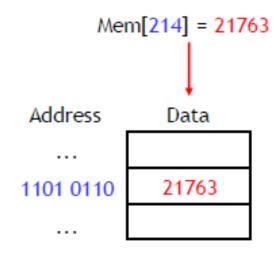
Index	٧	Tag	Data	Address	Data
110	1	00010	123456	1101 0110	6378

When we update Mem[1101 0110], should we also load it into the cache?

Write around caches (a.k.a. write-no-allocate)

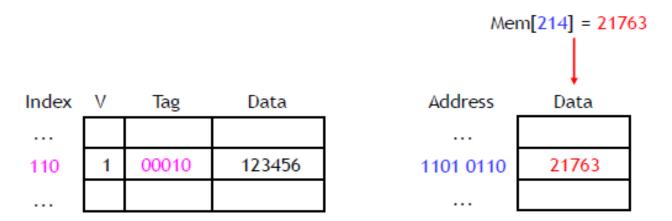
 With a write around policy, the write operation goes directly to main memory without affecting the cache.





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 With a write around policy, the write operation goes directly to main memory without affecting the cache.

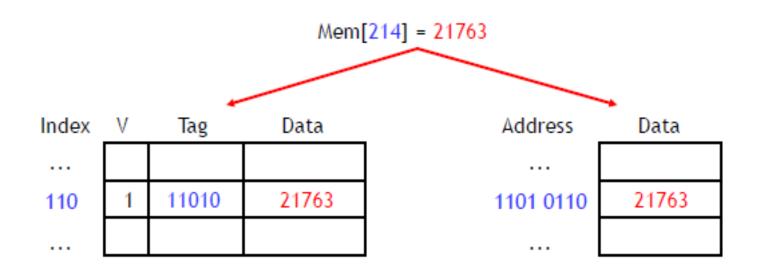


 This is good when data is written but not immediately used again, in which case there's no point to load it into the cache yet.

```
for (int i = 0; i < SIZE; i++)
a[i] = i;
```

Allocate on write

 An allocate on write strategy would instead load the newly written data into the cache.



If that data is needed again soon, it will be available in the cache.