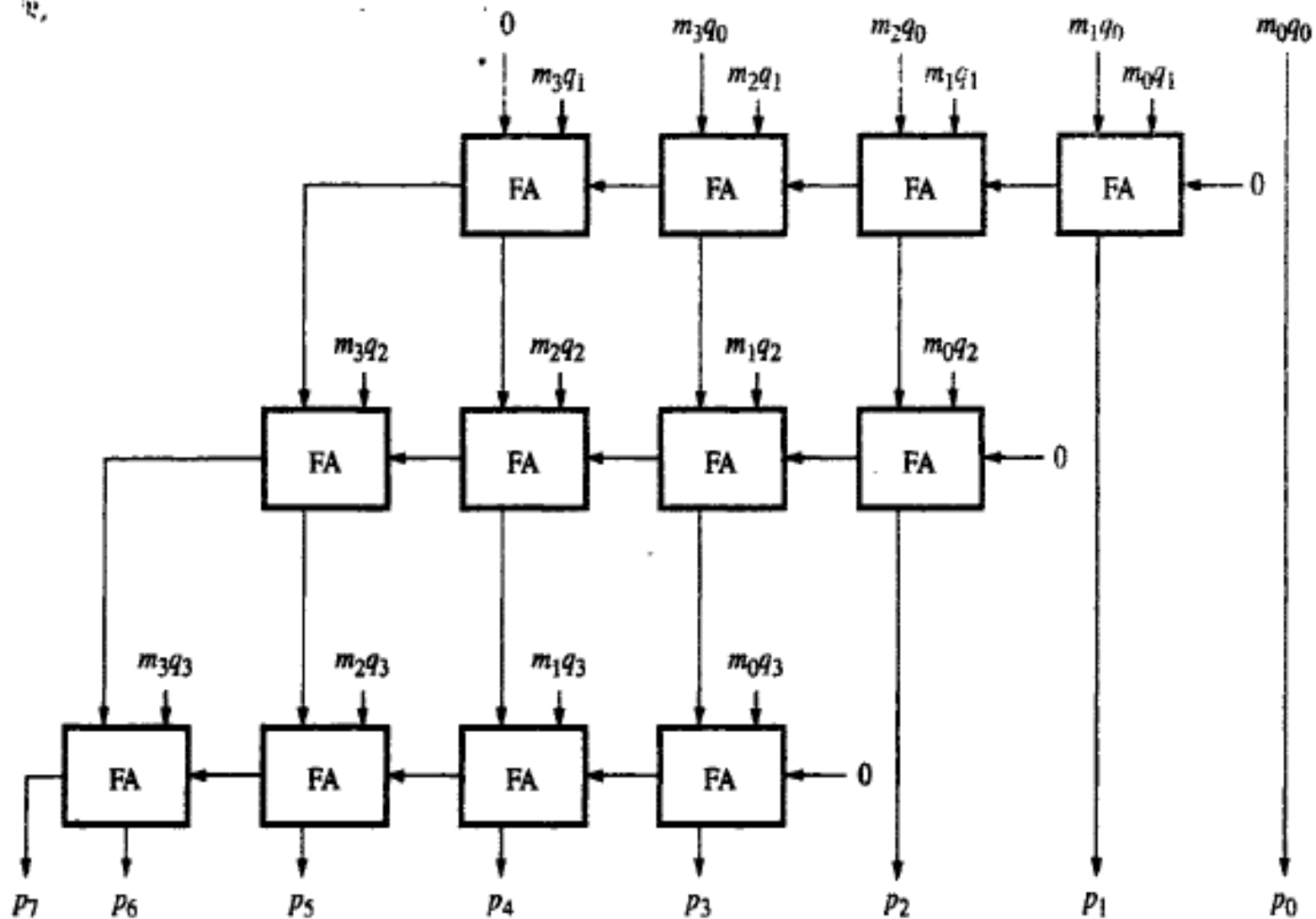
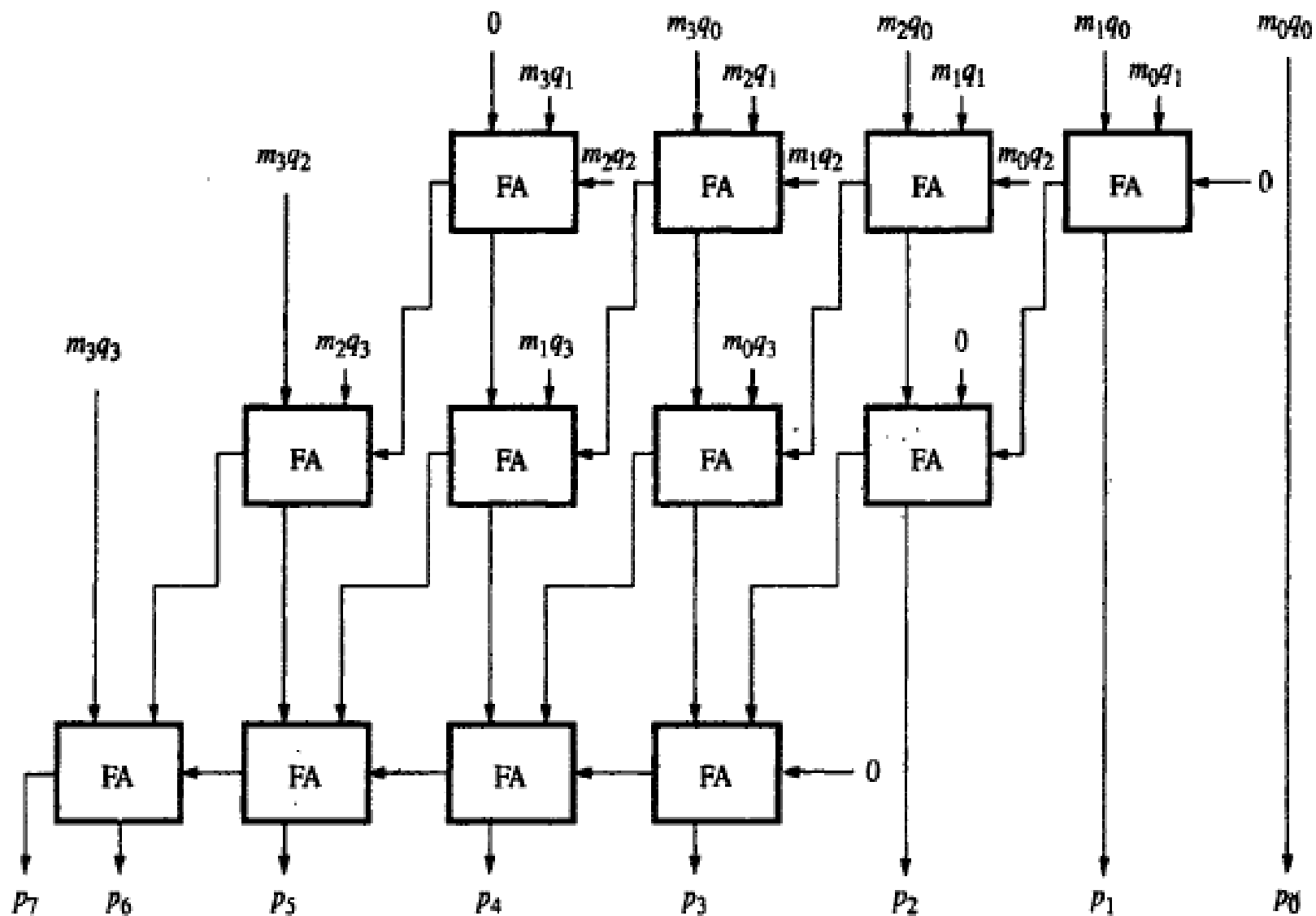


| | | | | | | | |
|-----|------|------|------|-----------|-----------|-----------|-------------|
| | | | | m3 | m2 | m1 | m0 |
| | | | | q3 | q2 | q1 | q0 |
| | | | | m3q0 | m2q0 | m1q0 | m0q0 |
| | | | m3q1 | m2q1 | m1q1 | m0q1 | |
| | | m3q2 | m2q2 | m1q2 | m0q2 | | |
| | m3q3 | m2q3 | m1q3 | m0q3 | | | |
| p7= | p6= | p5= | p4= | p3= | p2= | p1= | P0= m0q0 |

2.



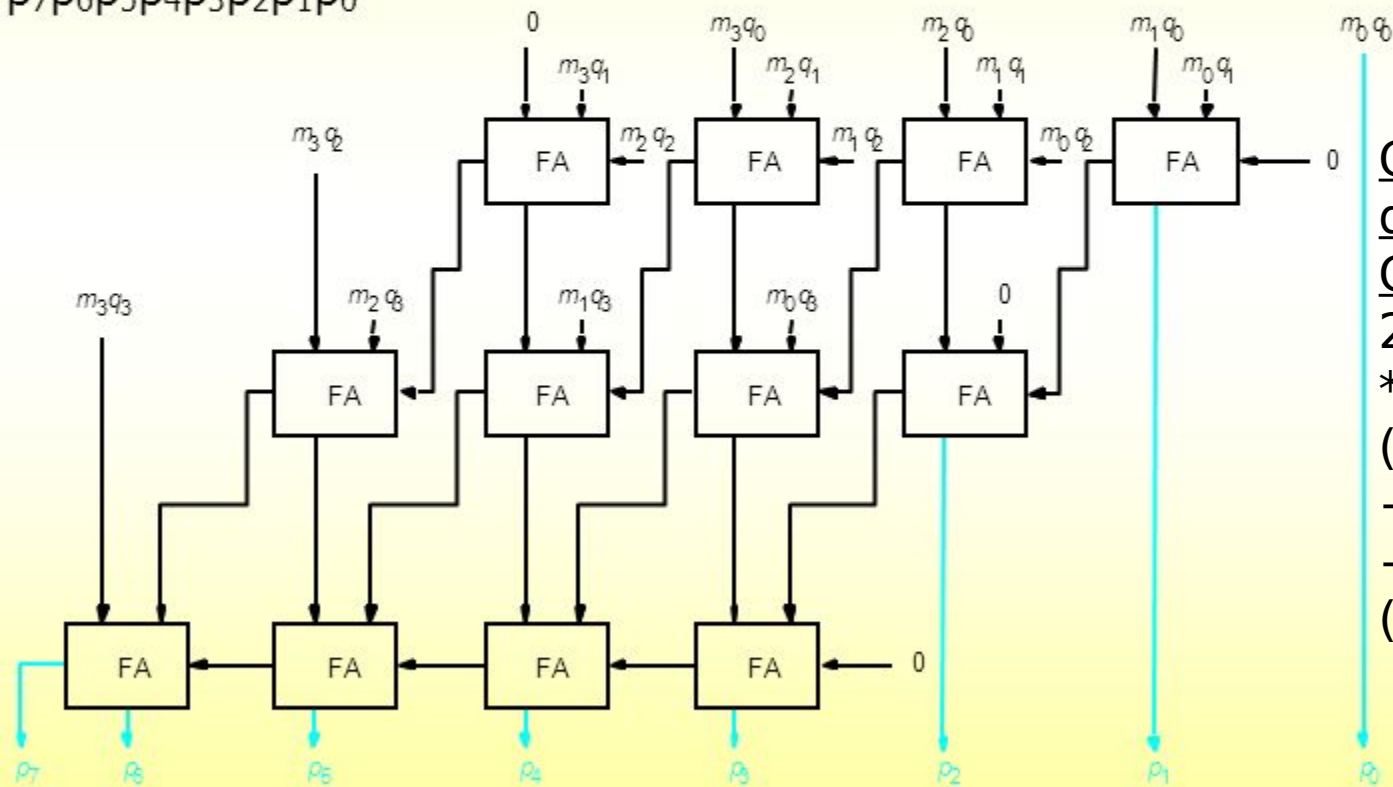
(a) Ripple-carry array (Figure 6.6 structure)



(b) Carry-save array

Carry Save Array

- For the multiplication operation $M \times Q = P$ for 4-bit operands
- M: $m_3m_2m_1m_0$
- Q: $q_3q_2q_1q_0$
- P: $p_7p_6p_5p_4p_3p_2p_1p_0$



Gate delay
CSA:-
2(perFA)
*2
(levels)
+ 1(and)
+ 4
(CLA)=9

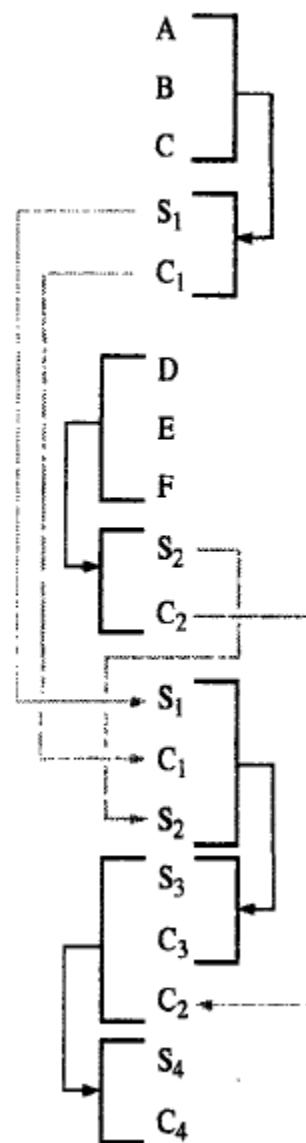


| | | | | | | | | | | | | | | |
|-------|---|---|---|---|---|---|---|---|---|---|---|---|---------|---------|
| | | | | | | | 1 | 0 | 1 | 1 | 0 | 1 | (45) | M |
| | | | | | | x | 1 | 1 | 1 | 1 | 1 | 1 | (63) | Q |
| <hr/> | | | | | | | | | | | | | | |
| | | | | | | | 1 | 0 | 1 | 1 | 0 | 1 | A | |
| | | | | | | | 1 | 0 | 1 | 1 | 0 | 1 | B | |
| | | | | | | | 1 | 0 | 1 | 1 | 0 | 1 | C | |
| | | | | | | | 1 | 0 | 1 | 1 | 0 | 1 | D | |
| | | | | | | | 1 | 0 | 1 | 1 | 0 | 1 | E | |
| | | | | | | | 1 | 0 | 1 | 1 | 0 | 1 | F | |
| <hr/> | | | | | | | | | | | | | | |
| 1 | 0 | 1 | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | (2,835) | Product |

Figure 6.17 A multiplication example used to illustrate carry-save addition as shown in Figure 6.18.

$$\begin{array}{r}
 1\ 0\ 1\ 1\ 0\ 1\ M \\
 \times 1\ 1\ 1\ 1\ 1\ 1\ Q \\
 \hline
 \end{array}$$

$$\begin{array}{r}
 1\ 0\ 1\ 1\ 0\ 1 \\
 1\ 0\ 1\ 1\ 0\ 1 \\
 1\ 0\ 1\ 1\ 0\ 1 \\
 \hline
 1\ 1\ 0\ 0\ 0\ 0\ 1\ 1 \\
 0\ 0\ 1\ 1\ 1\ 1\ 0\ 0 \\
 \hline
 1\ 0\ 1\ 1\ 0\ 1 \\
 1\ 0\ 1\ 1\ 0\ 1 \\
 1\ 0\ 1\ 1\ 0\ 1 \\
 \hline
 1\ 1\ 0\ 0\ 0\ 0\ 1\ 1 \\
 0\ 0\ 1\ 1\ 1\ 1\ 0\ 0 \\
 \hline
 1\ 1\ 0\ 0\ 0\ 0\ 1\ 1 \\
 0\ 0\ 1\ 1\ 1\ 1\ 0\ 0 \\
 1\ 1\ 0\ 0\ 0\ 0\ 1\ 1 \\
 \hline
 1\ 1\ 0\ 1\ 0\ 1\ 0\ 0\ 0\ 1\ 1 \\
 0\ 0\ 0\ 0\ 1\ 0\ 1\ 1\ 0\ 0\ 0 \\
 0\ 0\ 1\ 1\ 1\ 1\ 0\ 0 \\
 \hline
 0\ 1\ 0\ 1\ 1\ 1\ 0\ 1\ 0\ 0\ 1\ 1 \\
 + 0\ 1\ 0\ 1\ 0\ 1\ 0\ 0\ 0\ 0\ 0 \\
 \hline
 1\ 0\ 1\ 1\ 0\ 0\ 0\ 1\ 0\ 0\ 1\ 1\ \text{Product}
 \end{array}$$



Carry Save Adder Tree for 6 Operands

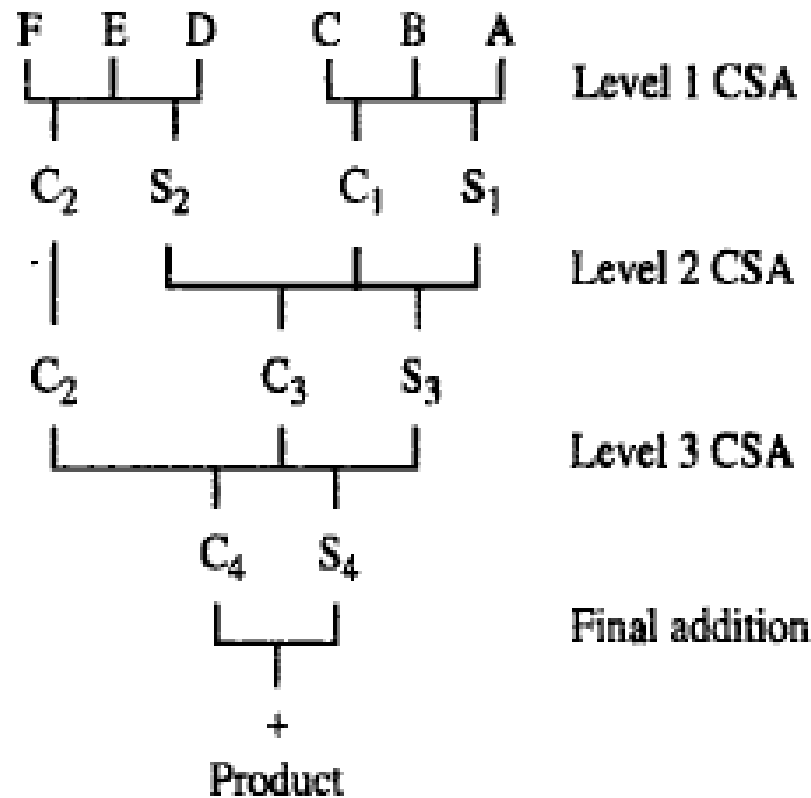
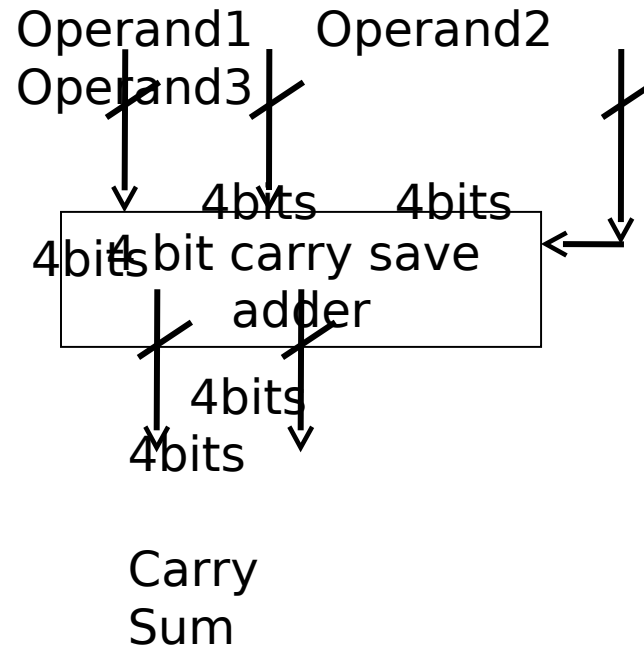


Figure 6.19 Schematic representation of the carry-save addition operations in Figure 6.18.

Exercise

- How many levels are needed to reduce
 - 8 summands to 2 ?
 - 5 summands to 2?

Carry Save Adder (CSA)



Levels in the CSA Tree

Approx $1.7\log_2 k - 1.7$ to $1.7\log_2 k - 3.4$ levels of CSA steps needed to reduce k summands to 2 vectors

TABLE 5.1 THE NUMBER OF LEVELS IN A CSA TREE FOR k OPERANDS.

| Number of operands | Number of levels |
|---------------------|------------------|
| 3 | 1 |
| 4 | 2 |
| $5 \leq k \leq 6$ | 3 |
| $7 \leq k \leq 9$ | 4 |
| $10 \leq k \leq 13$ | 5 |
| $14 \leq k \leq 19$ | 6 |
| $20 \leq k \leq 28$ | 7 |
| $29 \leq k \leq 42$ | 8 |
| $43 \leq k \leq 63$ | 9 |