CS 381 HW 5

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Q1

After running DFS on the graph of figure 22.6, the start and finish times of the vertices are :

q(1, 16)
r(17, 20)
s(3, 6)
t(8, 15)
u(18, 19)
v(4, 5)
w(2, 7)
x(11, 14)
y(9, 10)
z(12, 13)

After reversing G and running DFS on the vertices in order of decreasing finishing time:

dfs(r) discovers r
dfs(u) discovers u
dfs(q) discovers q, y, t
dfs(x) discovers x, z
dfs(w) discovers w, v, s

Thus, the SCCs produced are {r}, {u}, {q, y, t}, {x, z}, {w, v, s}.

Q2

Let N_{ij}^k denote the number of different shortest paths from i to j for which all intermediate vertices are in the set {1,2,...k}. Using the same convention for w_{ij} as the book, we get

$N_{ij}^0 =$

$$\begin{cases} 1 & w_{ij} < \infty \\ 0 & w_{ij} = \infty \end{cases}$$

since the number of shortest paths between i and j using no intermediate vertices is 1 if there exists an edge between i and j and 0 otherwise.

For k > 0, we can split N_{ij}^k into three cases:

$N_{ij}^k =$

$$\begin{cases} N_{ij}^{k-1} & d_{ij}^{k-1} < d_{ik}^{k-1} + d_{kj}^{k-1} \\ N_{ij}^{k-1} + \delta(k) & d_{ij}^{k-1} = d_{ik}^{k-1} + d_{kj}^{k-1} \\ \delta_{ij}^k & d_{ij}^{k-1} > d_{ik}^{k-1} + d_{kj}^{k-1} \end{cases}$$

where $\delta_{ij}^k = N_{ik}^{k-1}$ x N_{kj}^{k-1}

The reasoning is as follows:

1. If the shortest path between i and j using intermediate vertices {1,2,...k} doesn't use vertex k (i.e $d_{ij}^{k-1} < d_{ij}^{k-1} + d_{kj}^{k-1}$), the the number of shortest paths using intermediate vertices {1,2,...k} is just the number of shortest paths using intermediate vertices {1,2,...k-1}.

2. If $d_{ij}^{k-1} = d_{ik}^{k-1} + d_{kj}^{k-1}$, then the number of shortest paths using intermediate vertices {1,2,...k} is the number of shortest paths using intermediate vertices {1,2,...k-1} plus the number of shortest paths going through intermediate vertex k. The number of shortest paths going through vertex k is δ_{ij}^k $= N_{ik}^{k-1}$ x N_{kj}^{k-1} (number of shortest paths from i to k multiplied by number of shortest paths from k to j) Obviously, k cannot be an intermediate vertex on a path from i to k or k to j (else we'd have a negative cycle contrary to our hypothesis) so we can use the set of intermediate vertices {1,2,...k-1} to calculate δ_{ij}^k.

3. If $d_{ij}^{k-1} > d_{ik}^{k-1} + d_{kj}^{k-1}$, then the number of shortest paths from i to j is just the number of shortest paths that go through k which is δ_{ij}^k.

Note that we can drop the superscripts. Why? If having dropped the superscripts, we were to compute and store N_{ik} or N_{jk} before using these values to compute δ_{ij}, then we might have one of the following situations:

$\delta_{ij}^k =$

$$\begin{cases} N_{ik}^k \text{ x } N_{kj}^{k-1} \\ N_{ik}^{k-1} \text{ x } N_{kj}^k \\ N_{ik}^k \text{ x } N_{kj}^k \end{cases}$$

However, k cannot be an intermediate vertex on any shortest path from i to k (else negative cycle contrary to hypothesis), so the number of shortest paths from i to k using intermediate vertices {1,2,...k} is just the number of shortest paths from i to k using vertices {1,2,...k-1}. Thus, $N_{ik}^k = N_{ik}^{k-1}$. Similarly, $N_{kj}^k = N_{kj}^{k-1}$. Thus, we can drop the superscripts.

Algorithm 1 Floyd-Warshall with total number of shortest paths

function FLOYD-WARSHALL(W)
    n = W.rows
    D = W
    Initialize N ($n_{ij} = N_{ij}^0$ as described above)
    for k = 1 to n **do**
        for i = 1 to n **do**
            if $d_{ij} < d_{ik} + d_{kj}$ **then**
                $n_{ij} = n_{ij}$
            else if $d_{ij} = d_{ik} + d_{kj}$ **then**
                $n_{ij} = n_{ij} + n_{ik}$ x n_{kj}
            else
                $n_{ij} = n_{ik}$ x n_{kj}
            end if
            $d_{ij} = \min(d_{ij}, d_{ik} + d_{kj})$
        end for
    end for
    return D, N
end function