

Software Engineering Department

ORT Braude College

Course 61401: Extended Project in Software Engineering

Formal Verification of Specs of Applications

In Partial Fulfillment of the Requirements for

Final Project in Software Engineering (Course 61401)

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## INTRODUCTION

Each program development starts from its specification. Before, one starts implementation, the correctness of the spec must be confirmed. Specs of cellular applications demonstrate very specific character: transfer from one screen to another. We use the specialty of the specs in order to use the machinery of formal verification in order to verify their correctness

**What are we going to do?**

We will build an tool that allows graphical definition of specifications of cellular applications, that's mean Represent the specifications as a graph: nodes are the screens associated with the corresponding values of the parameters.

Our application gets a list of Requirements that user wants to check, then it Uses machinery of formal verification in order to verify the spec, the verification result either Confirmation message or a path were the test Failed.

**Why is it not trivial?**

1. Performance verification at the design stage, that mean before running code we can detect

errors.

**What are the difficulties of the project?**

1. our project has many screens
2. We should storage a lot of nodes and parameters;
3. Building a workspace that allows user to build specifications graph

## 2. THEORY

# 2.1. Background

It is all about money. We are annoyed when our mobile phone malfunctions, or when our video recorder reacts unexpectedly and wrongly to our issued commands. These software and hardware errors do not threaten our lives, but may have substantial financial consequences for the manufacturer.

### 2.1.1. Formal verification

formal verification is the act of proving or disproving the correctness of intended algorithms underlying a system with respect to a certain formal specification or property, using formal methods of mathematics. . This specification prescribes what the system has to do and what not, and thus constitutes the basis for any verification activity.

The verification of these systems is done by providing a formal proof on an abstract mathematical model of the system, the correspondence between the mathematical model and the nature of the system being otherwise known by construction.

One approach and formation is model checking refers to the following problem: Given a model of a system, exhaustively and automatically check whether this model meets a given specification. Typically, one has software systems in mind, whereas the specification contains safety requirements such as the absence of deadlocks and similar critical states that can cause the system to crash. Model checking is a technique for automatically verifying correctness properties of finite-state systems.

In order to solve such a problem algorithmically, both the model of the system and the specification are formulated in some precise mathematical language: To this end, it is formulated as a task in logic, namely to check whether a given structure satisfies a given logical formula. The concept is general and applies to all kinds of logics and suitable structures. A simple model-checking problem is verifying whether a given formula in the propositional logic is satisfied by a given structure.

*//system verification is used to establish that the design or product under consideration possesses certain properties. The properties to be validated can be quite elementary, e.g., a system should never be able to reach a situation in which no progress can be made (a deadlock scenario), and are mostly obtained from the system’s specification*

### 2.1.2. Transition system (TS)

A transition system TS is a tuple (S, Act,→,I, AP, L) where

S is a set of states.

Act is a set of actions,

→ ⊆ S × Act × S is a transition relation,

I ⊆ S is a set of initial states,

AP is a set of atomic propositions, and

L: is a labeling function.

TS is called finite if S, Act, and AP are finite. We can describe behavior of transition system as follows

The transition system starts in some initial state and evolves according to the transition relation →. That is, if s the current state, then a transition originating from s is selected *nondeterministically* and taken, the action α is performed and the transition system evolves from state s into the state q

This selection procedure is repeated in state q and finishes once a state is encountered that has no outgoing transitions. It is important to realize that in case a state has more than one outgoing transition, the “next” transition is chosen in a purely nondeterministic fashion. That is, the outcome of this selection process is not known a priori, and, hence, no statement can be made about Transition Systems 21 the likelihood with which a certain transition is selected. Similarly, when the set of initial states consists of more than one state, the start state is selected nondeterministically.

The labeling function L relates a set L(s) ∈ of atomic propositions to any state s. 1 L(s) intuitively stands for exactly those atomic propositions a ∈ AP which are satisfied by state s.

### 2.1.3. Program graph (PG)

A program graph PG over set Var of typed variables is a tuple (Loc, Act, Effect, →, , ) where

* Loc is a set of locations
* Act is a set of actions,
* Effect:Act×Eval(Var)→Eval(Var) →Eval(Var) is the effect function,
* →⊆Loc×Cond(Var)×Act×Loc↪⊆Loc×Cond(Var)×Act×Loc is the conditional transition relation,
* ⊆Loc is a set of initial locations,
* ∈Cond(Var) is the initial condition.

### 2.1.4. Linear Temporal Logic(LTL)

linear temporal logic (LTL), is a logical formalism that is suited for specifying LT properties. LTL can be used to specify important system properties.

Temporal logic is a formalism par excellence for treating correctness depends on the executions. It extends propositional or predicate logic by modalities that permit to referral to the infinite behavior of a system.

The underlying nature of time in temporal logics is *linear*. i.e. at each moment in time there is a single successor moment, Several model-checking tools use LTL as a property specification language. The model checker SPIN is a prominent example of such an automated verification tool.

### 2.1.5. SPIN

Spin is a popular verification tool of distributed systems, used by thousands of people worldwide. The tool can be used for the formal verification of multi-threaded software applications.

Spin can perform simulations of the system's execution. It was developed at Bell Labs in the Unix group of the Computing Sciences Research Center, starting in 1980. Spin can perform interactive, guided, or random simulations of the system's execution.