

# Homework 1 – Linear programming and CVaR

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COMPUTATIONAL ASPECTS OF OPTIMIZATION

## Var & CVaR

Let  $Z$  be a random variables that represent **loss** with p.d.f.  $F_Z(\eta)$ ,  $\alpha \in (0, 1)$  (usually  $\alpha = 0.95$ ).

### Value at Risk (VaR)

$$\text{VaR}_\alpha(Z) \cong F_Z^{-1}(\alpha)$$

### Conditional Value at Risk (CVaR)

$$\text{CVaR}_\alpha(X) \cong \mathbb{E}[Z | Z \geq \text{VaR}_\alpha(Z)]$$

or

$$\text{CVaR}_\alpha(X) \cong \mathbb{E}[Z | Z > \text{VaR}_\alpha(Z)]$$

Exact definitions follow (Rockafellar and Uryasev 2002) ...

## Var & CVaR

**Value at Risk** is defined as a quantile

$$\text{VaR}_\alpha(Z) = \inf\{\eta : P(Z \leq \eta) \geq \alpha\}.$$

For  $Z \in \mathcal{L}_1(\Omega)$ , **Conditional Value at Risk** (CVaR) is defined as the mean of losses in the  $\alpha$ -tail distribution with the distribution function:

$$F_\alpha(\eta) = \begin{cases} \frac{F(\eta) - \alpha}{1 - \alpha}, & \text{if } \eta \geq \text{VaR}_\alpha(Z), \\ 0, & \text{otherwise,} \end{cases}$$

where  $F(\eta) = P(Z \leq \eta)$ . CVaR can be expressed using the following **minimization formula**:

$$\text{CVaR}_\alpha(Z) = \min_{\xi \in \mathbb{R}} \left[ \xi + \frac{1}{1 - \alpha} \mathbb{E} [\max\{Z - \xi, 0\}] \right] \quad (1)$$

with the minimum attained at any  $(1 - \alpha)$ -th quantile.

## Investment problem with CVaR

Solve a simple investment problem

$$\min_{x_i} \text{CVaR}_\alpha \left( - \sum_{i=1}^n x_i R_i \right)$$

$$\text{s.t. } \mathbb{E} \left[ \sum_{i=1}^n x_i R_i \right] \geq r_0,$$

$$\sum_{i=1}^n x_i = 1, \quad x_i \geq 0,$$

where we consider  $n$  assets with random rate of return  $R_i$ . The first constraint ensures minimal expected return  $r_0$ ,  $x_i$  are (nonnegative) portfolio weights which sum to one.

If the distribution of  $R_i$  is discrete with realizations  $r_{is}$  and probabilities  $p_s = 1/S$ , then we can use **linear programming** reformulation

$$\min_{\xi, x_i, u_s} \xi + \frac{1}{(1-\alpha)S} \sum_{s=1}^S u_s,$$

$$\text{s.t. } u_s \geq -\sum_{i=1}^n x_i r_{is} - \xi, \quad s = 1, \dots, S,$$

$$\sum_{i=1}^n x_i \bar{R}_i \geq r_0,$$

$$\sum_{i=1}^n x_i = 1, \quad x_i \geq 0,$$

$$\xi \in \mathbb{R}, \quad u_s \geq 0,$$

where  $\bar{R}_i = 1/S \sum_{s=1}^S r_{is}$ .

## Numerical experiment

- ① Use at least 6 assets and 100 return realizations.
- ② Run the problem for different 11 values  $r_0 \in \{\min_i \bar{R}_i, \dots, \max_i \bar{R}_i\}$ .
- ③ Plot the optimal values  $\text{CVaR}_\alpha$  against the corresponding values of  $r_0$ .