

## Location Routing Problem (LRP)

Given, a graph  $G = (D, C, A)$  with set of depot nodes  $D$  with capacity  $q_d$ , operational cost  $\pi_d^o$ , fixed cost  $\pi_d^f$ , and fleet of delivery vehicles  $V_d$  with capacity  $q_v$ , speed  $s_v$ , refueling time  $\tau_v^f$ , service time at the depot node  $\tau_v^d$  (per unit demand), service time at the customer node  $\tau_v^c$ , operational cost  $\pi_v^o$  (per unit distance traveled), fixed cost  $\pi_v^f$ , working hours  $w_v$  for every vehicle  $v \in V_d$  for every depot  $d \in D$ ; set of customer nodes  $C$  with demand  $q_c$ , delivery time-window  $[t_c^e, t_c^l]$  for every customer  $c \in C$ ; set of arcs  $A$  with length  $l_{ij}$  for every arc  $(i, j) \in A$ , the objective is to develop least cost routes from select depot nodes using select vehicles such that every customer node is visited exactly once while also accounting for vehicle capacities.

$$\min z = \sum_{d \in D} \left( \pi_d^f + \pi_d^o \sum_{v \in V_d} \sum_{r \in R_v} z_{cr} q_c \right) y_d + \sum_{v \in V} \left( \pi_v^f + \sum_{(i,j) \in A} \pi_v^o l_{ij} x_{ij}^v \right) y_v$$

Subject to,

Service constraint,

$$\sum_{d \in D} \sum_{v \in V_d} \sum_{r \in R_v} z_{cr} = 1 \quad \forall c \in C$$

Flow conservation,

$$\sum_{i \in T_j} x_{ij}^r = \sum_{k \in H_j} x_{jk}^r \quad \forall j \in \{D \cup C\}; r \in R_v; v \in V_d; d \in D$$

Allocation constraints,

$$z_{cr} = \sum_{v \in V_d} \sum_{j \in H_c} x_{cj}^r \quad \forall c \in C; r \in R_v; v \in V_d; d \in D$$

Resource use constraints,

$$y_v \leq \sum_{r \in R_v} \sum_{j \in H_d} x_{dj}^r \quad \forall v \in V_d; d \in D$$

$$y_d \leq \sum_{v \in V_d} y_v \quad \forall d \in D$$

Route constraints,

$$l_r = \sum_{i \in T_j} x_{ij}^r l_{ij} \quad \forall j \in \{D \cup C\}; r \in R_v; v \in V_d; d \in D$$

$$q_r = \sum_{c \in C} z_{cr} q_c \quad \forall c \in C; r \in R_v; v \in V_d; d \in D$$

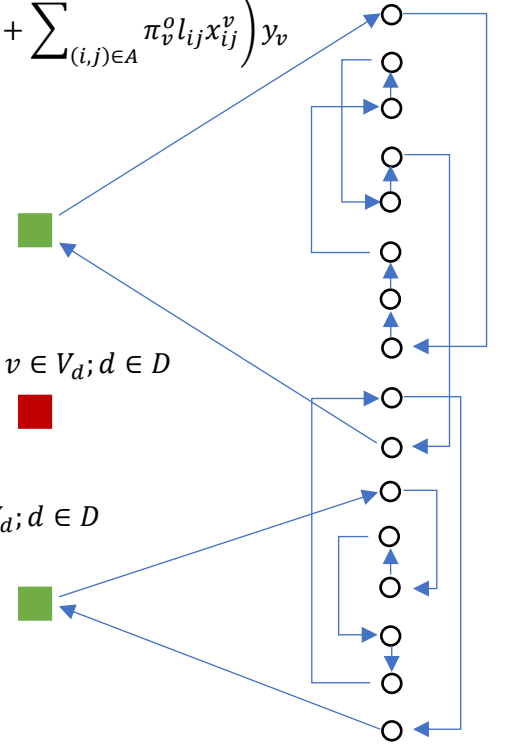


Fig. A typical LRP solution

Capacity constraints,

$$\begin{aligned} \sum_{c \in C} q_c z_{cr} &\leq q_v y_v & \forall v \in V_d; d \in D \\ \sum_{c \in C} q_c \sum_{v \in V_d} \sum_{r \in R_v} z_{cr} &\leq q_d y_d & \forall d \in D \end{aligned}$$

Vehicle range constraint,

$$\sum_{d \in D} \sum_{v \in V_d} \sum_{r \in R_v} l_r \leq l_v \quad \forall v \in V_d; d \in D$$

Route start and end time constraint,

$$\begin{aligned} t_{r_1}^s &= 0 & \forall v \in V_d; d \in D \\ t_{r_k}^s &= t_{r_{k-1}}^e + \tau_v^f l_r / l_v + \tau_v^d q_r & \forall r \in R_v; v \in V_d; d \in D \end{aligned}$$

Driver working hours constraint,

$$t_{r_K}^e \leq w_v \quad \forall v \in V_d; d \in D$$

Customer node arrival and departure time constraint,

$$\begin{aligned} t_c^a + M(1 - x_{ic}^r) &\geq t_i^d + x_{ic}^r l_{ic} / s_v & \forall i \in T_c; c \in C; r \in R_v; v \in V_d; d \in D \\ t_c^d &\geq t_c^a + \max(0, t_c^e - t_c^a) + \tau_c^v & \forall c \in C \end{aligned}$$

Time-window constraint,

$$t_c^a \leq t_c^l \quad \forall c \in C$$

Binary constraints,

Arc use,

$$x_{ij} \in \{0,1\} \quad \forall (i,j) \in A$$

Vehicle use,

$$y_v \in \{0,1\} \quad \forall v \in V_d; d \in D$$

Depot use,

$$y_d \in \{0,1\} \quad \forall d \in D$$

Customer node allocation,

$$z_{cr} \in \{0,1\} \quad \forall c \in C; r \in R_v; v \in V_d; d \in D$$

Where,

$$T_j = \{i; (i,j) \in A\}$$

$$H_j = \{k; (j,k) \in A\}$$

$$R_v = \{r_1, r_2, \dots, r_k, \dots, r_K\} \quad \forall v \in V_d; d \in D$$