

P1: Merge K-Sorted Lists

⇒ Use merge function to merge sort two lists sequentially and store result (new merged) list in head node.

① Brute Force :

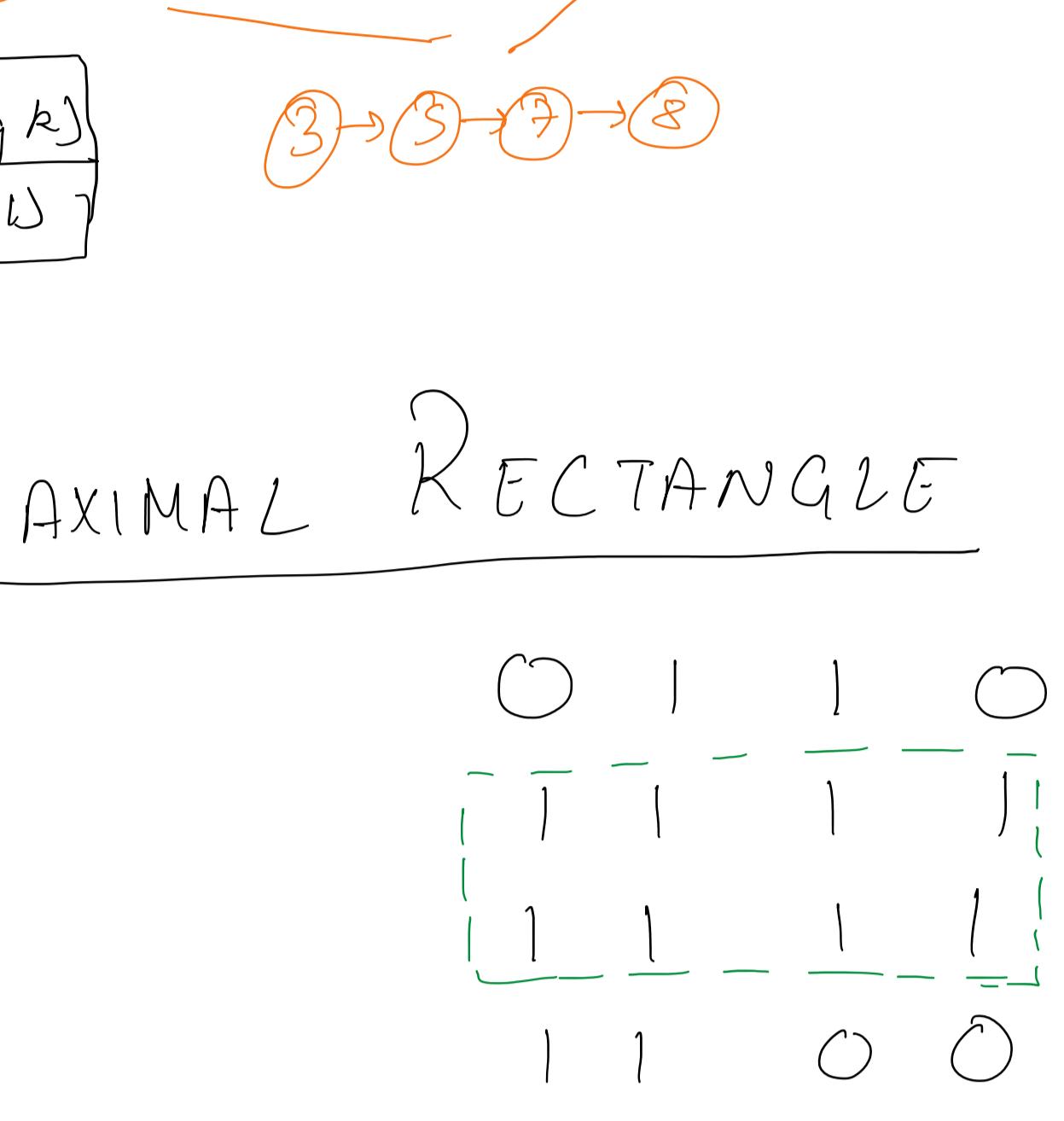
- Traverse all linked list → store nodes value to array
- Sort and iterate over array to get value of nodes
- Create new LL and extend it with new values.

[Time : $O(N \log N)$, Space : $O(N)$]

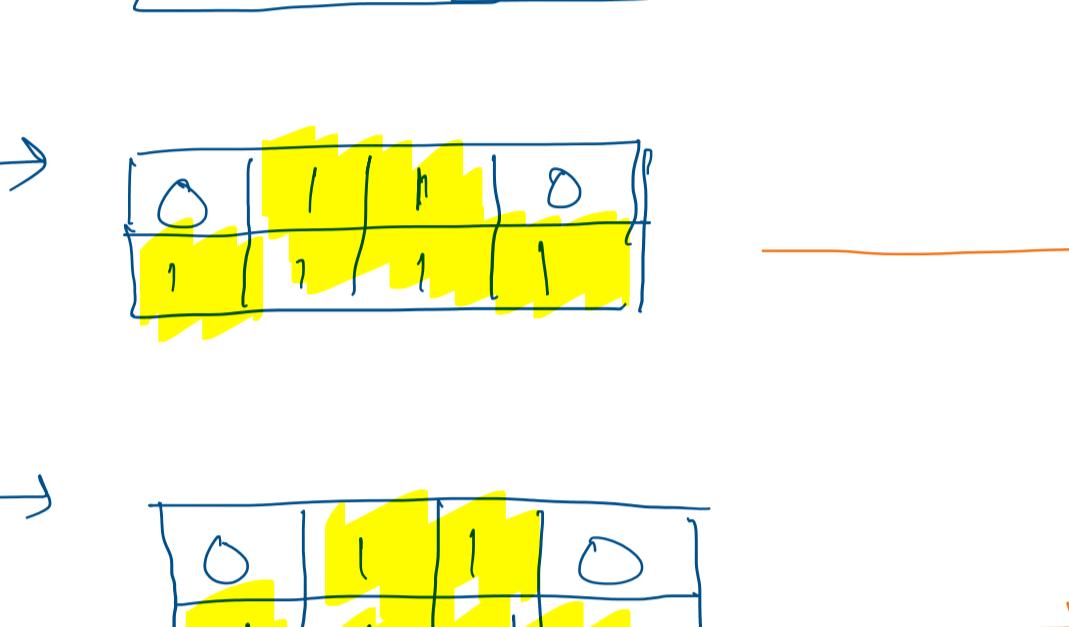
② Compare 1 by 1 :

- Compare all 'K' heads and get the node with smallest value.
- Extend LL with the selected nodes

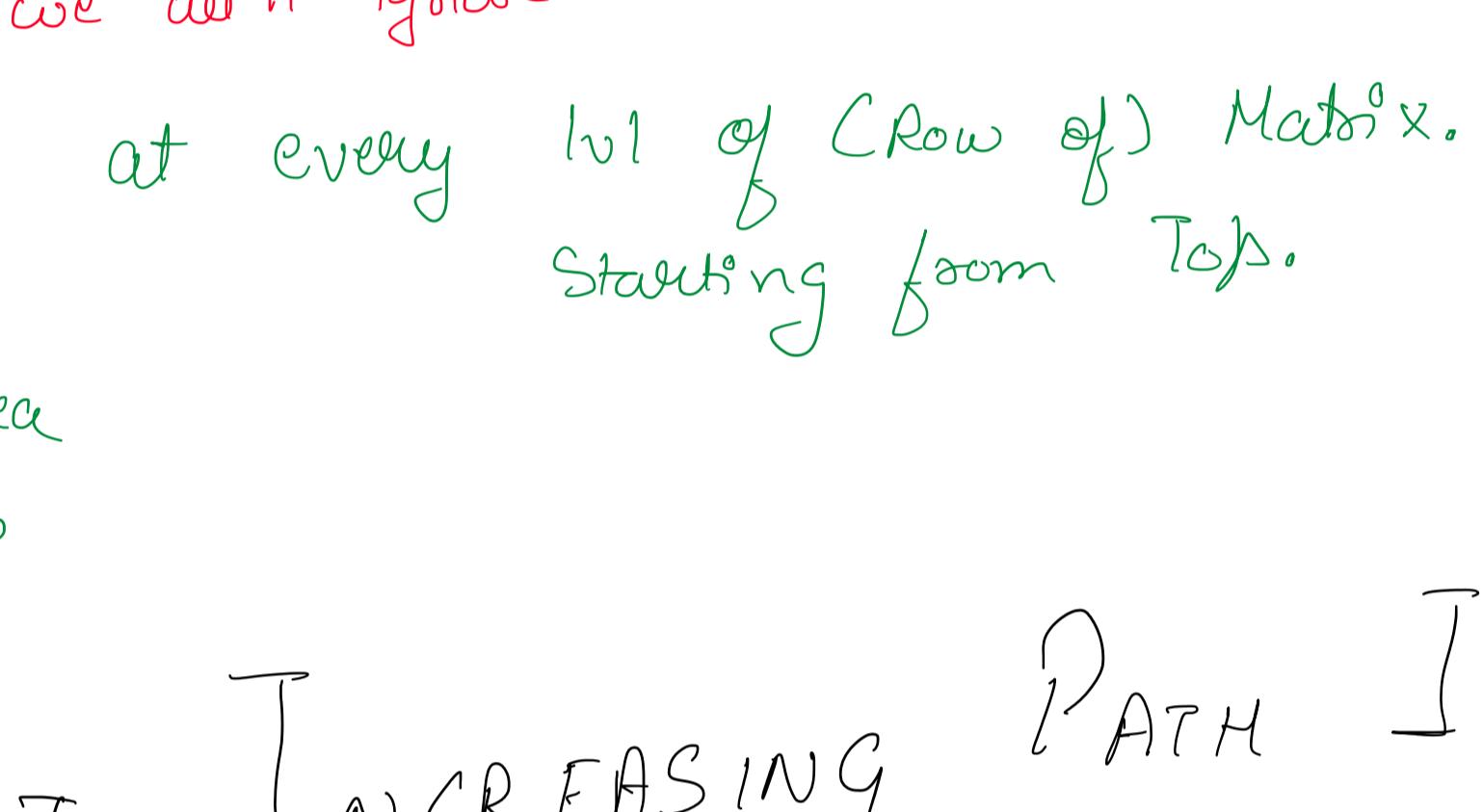
[Time : $O(KN^2)$, Space : $O(N)$]

③ Merge List one by one :

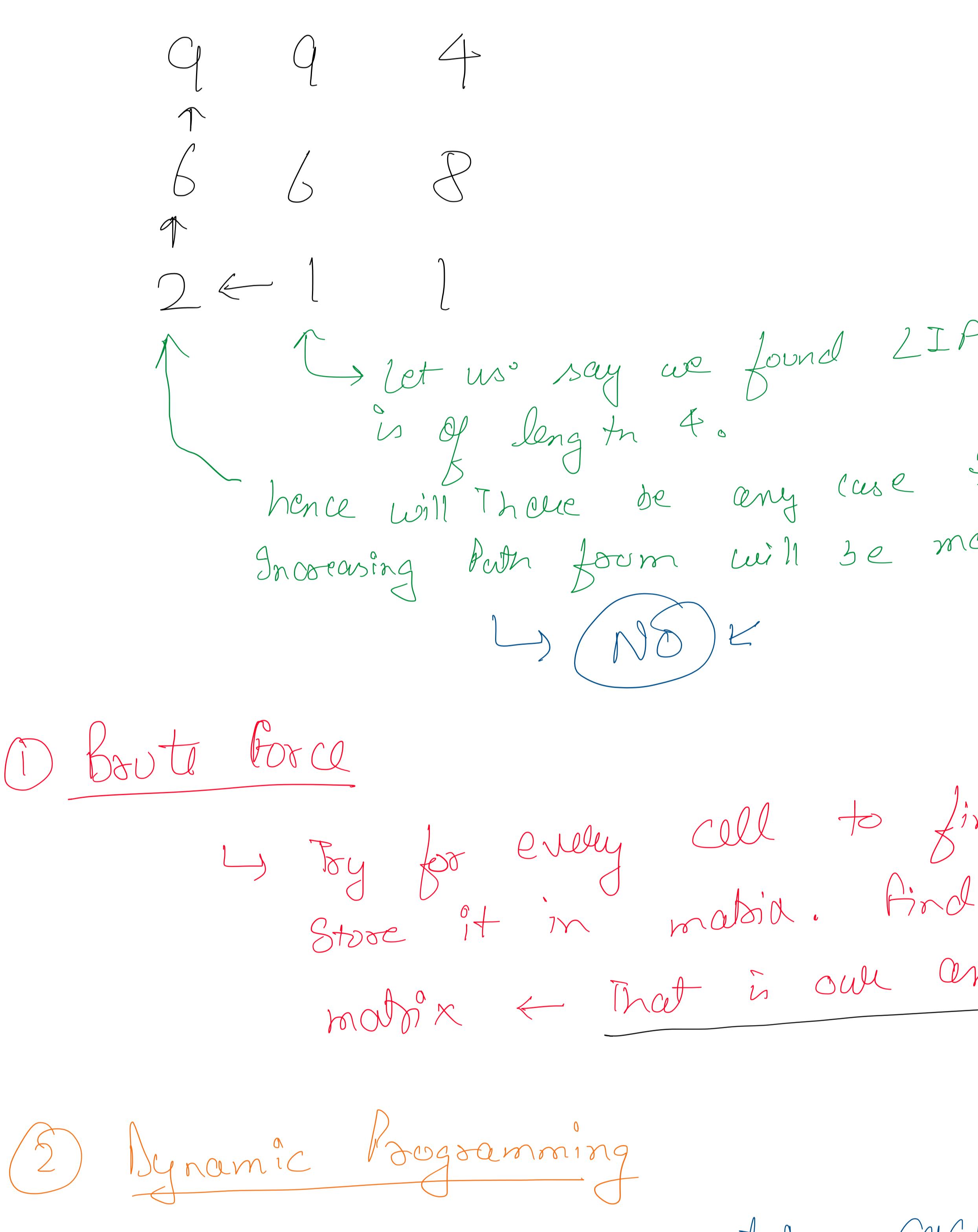
[Time : $O(NK)$, Space : $O(1)$]

④ Divide and Conquer :

[Time : $O(N \log K)$, Space : $O(1)$]

P2: MAXIMAL RECTANGLE

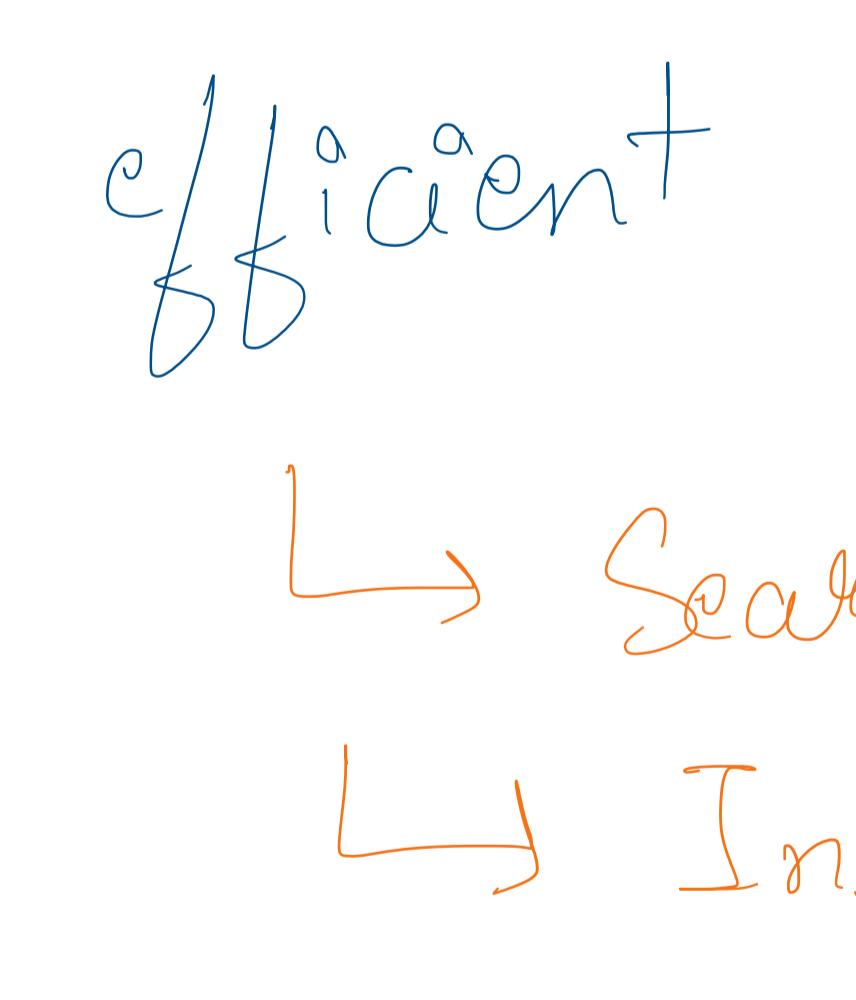
⇒ We will need to make histogram of given matrix :-



⇒ Apply [MAX] at every lvl of (Row of) Matrix.
↓
Maximum Area of Histograms

P3: LONGEST INCREASING PATH IN MATRIX

Observation: We cannot reuse a valuable



Let us say we found LIP from here is of length 4.
hence will there be any case that longest increasing path from will be more than 3?
↳ (No) ↳

① Brute Force

↳ Try for every cell to find LIP and store it in matrix. Find max of that matrix ← That is our answer.

② Dynamic Programming

↳ Use that matrix caching to reduce no. of DFS calls.

[Time : $O(n.m)$, Space : $O(n.m)$]

TRIE

efficient for following operations :-

↳ Search

↳ Prefix Search

↳ Insert

↳ Lexicographically ordering of words

↳ Delete

→ Implementing dictionaries

→ Search : $O(\text{word_len})$

Insert : $O(\text{word_len})$

Delete : "

Prefix Search : $O(\text{prefix_len} + \text{output_len})$

Lexographical Ordering : $O(\text{output_len})$

→ Abstract Trie Node {

TrieNode * Child[26];

bool isEnd;

TrieNode() {

isEnd = false;

for (i=0, i<26, i++)

Child[i] = NULL;

}

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