

ANmol KHANNA
 6S140549 - MFD
 Date :
 Page No. .

HOMEWORK - I

Q1

(A) In this case both receptors are store & forward routers.
 time for the packet transmission at each links = $\frac{1KB \cdot 8 \text{ bits/byte}}{1Mbps}$
 $= 8 \text{ ms}$
 $\Rightarrow d_{\text{transmission}} = 8 \text{ ms}$

Propagation delay for each packet on each links = $\frac{100 \text{ km}}{20,000 \text{ km/s}}$
 $= 5 \text{ ms}$

$\Rightarrow d_{\text{propagation}} = 5 \text{ ms}$.

... end to end delay =

$$\begin{aligned}
 & 3(\text{transmission}) + 3(\text{propagation}) + 2(\text{processing}) \\
 & = (3 \times 8 \text{ ms}) + (3 \times 5 \text{ ms}) + (2 \times 5 \text{ ms}) \\
 & = 39.01 \text{ ms.}
 \end{aligned}$$

(b) Here both Routers are not through Routers

$$\begin{aligned}
 \text{End to End delay} &= \text{transmission time} + 3(\text{propagation delay}) \\
 &\quad + 2(\text{processing time}) \\
 &= 8\text{ms} + (3 \times 5\text{ms}) + (2 \times 5\mu\text{s}) \\
 &= 8\text{ms} + 15\text{ms} + 10\mu\text{s} \\
 &= 23.01\text{ms}
 \end{aligned}$$

Calculated transmission time & propagation delay in (A) part

(c) Packet transmission time = $\frac{1\text{kB} \cdot 8\text{bits}/\text{byte}}{1\text{Gbps}}$

$$= 8\mu\text{s}$$

Store and Forward Case:

$$\begin{aligned}
 \text{End to End delay} &= 3 \times (\text{transmission time}) + 3(\text{propagation delay}) \\
 &\quad + 2(\text{processing time}) \\
 &= (3 \times 8\mu\text{s}) + (3 \times 5\text{ms}) + (2 \times 5\mu\text{s}) \\
 &= 24\mu\text{s} + 15\text{ms} + 10\mu\text{s} \\
 &= 15.03\text{ms}
 \end{aligned}$$

Get through case:

$$\text{End-to-end delay} = \frac{\text{transmission time} + 3(\text{propagation delay}) + 2(\text{processing time})}{\text{time}}$$

$$= 3 \times 5 \text{ ms} + 2 \times 5/1000 \text{ ms} + 0.008 \text{ ms}$$

$$= 15 \text{ ms} + 0.01 \text{ ms} + 0.008 \text{ ms}$$

$$= 15.018 \text{ ms.}$$

(d)

For R_1, R_2

$$\text{transmission time } (d_{\text{trans}}) = \frac{1 \text{ KB} \times 8 \text{ bits / byte}}{2 \text{ Mbps}} \\ = 4 \text{ ms}$$

For R_3 ,

$$\text{transmission time } (d_{\text{trans}}) = \frac{1 \text{ KB} \cdot 8 \text{ bits / byte}}{1 \text{ Mbps}} \\ = 8 \text{ ms}$$

Store and forward case:

for packet 1,

$$\text{End to End delay} = 2 \times 4 \text{ ms} + 8 \text{ ms} + 3 \times 5 \text{ ms} \\ = 31 \text{ ms}$$

for packet 2,

$$\text{End to End delay} = 31 \text{ ms} + 4 \text{ ms} \\ = 35 \text{ ms}$$

We have transmission time of 4ms :- when last bit of packet 2 arrives at Router 2, packet 1 is still in process.

for Packet 3,

$$\begin{aligned} \text{End to End delay} &= 31\text{ms} + 8\text{ms} \\ &= 39\text{ms} \end{aligned}$$

Cut-through Case :

Packet 1:

$$\begin{aligned} \text{End to End delay} &= 3 \times 5\text{ms} + 8\text{ms} \\ &= 23\text{ms} \end{aligned}$$

Packet 2 :

$$\begin{aligned} \text{End to End delay} &= 23\text{ms} + 4\text{ms} \\ &= 27\text{ms} \end{aligned}$$

Packet 3 :

$$\begin{aligned} \text{End to End delay} &= 27\text{ms} + 4\text{ms} \\ &= 31\text{ms} \end{aligned}$$

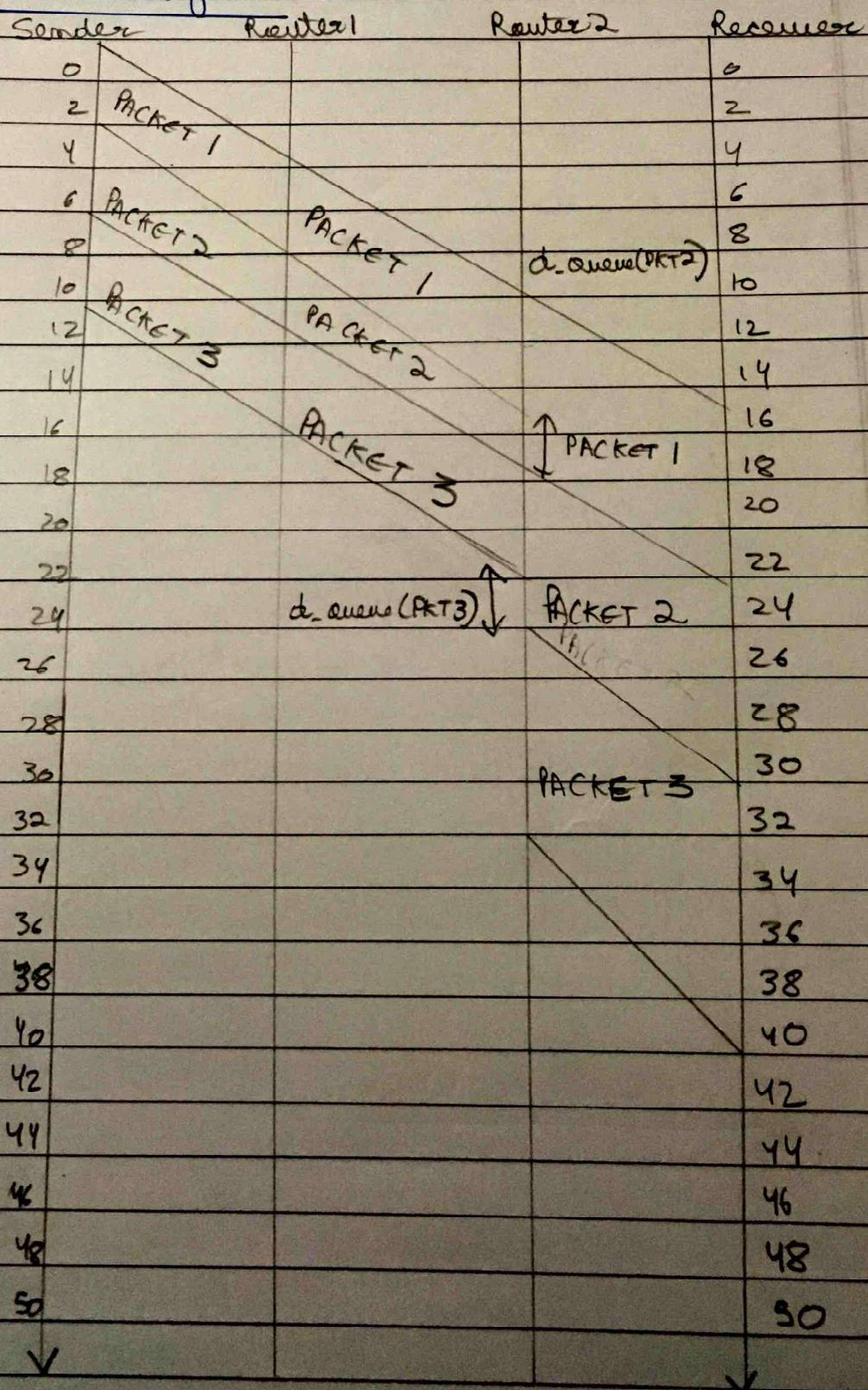
(e)

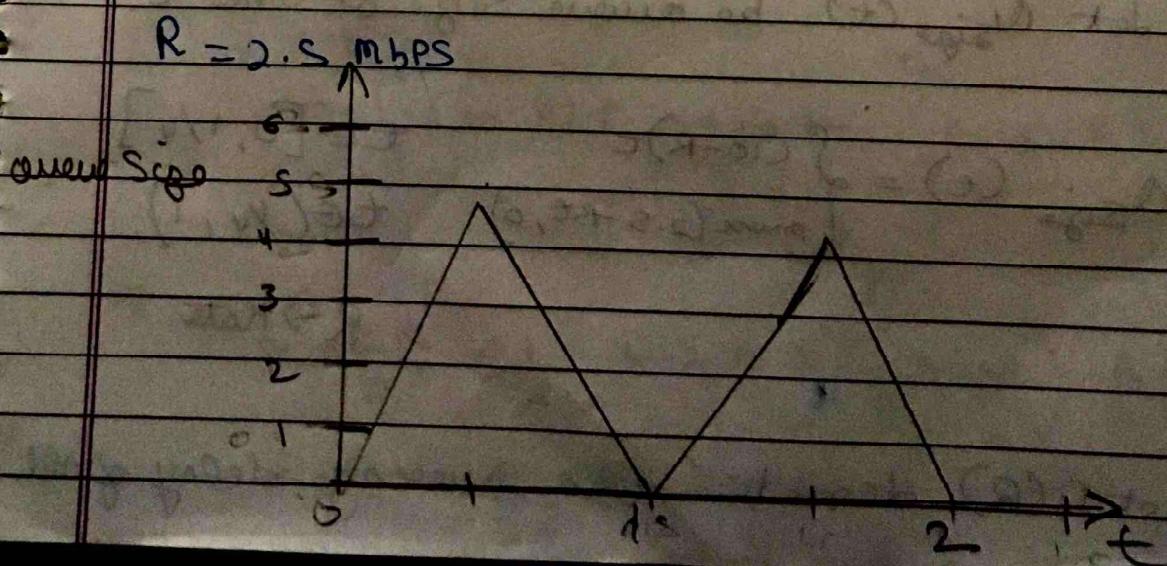
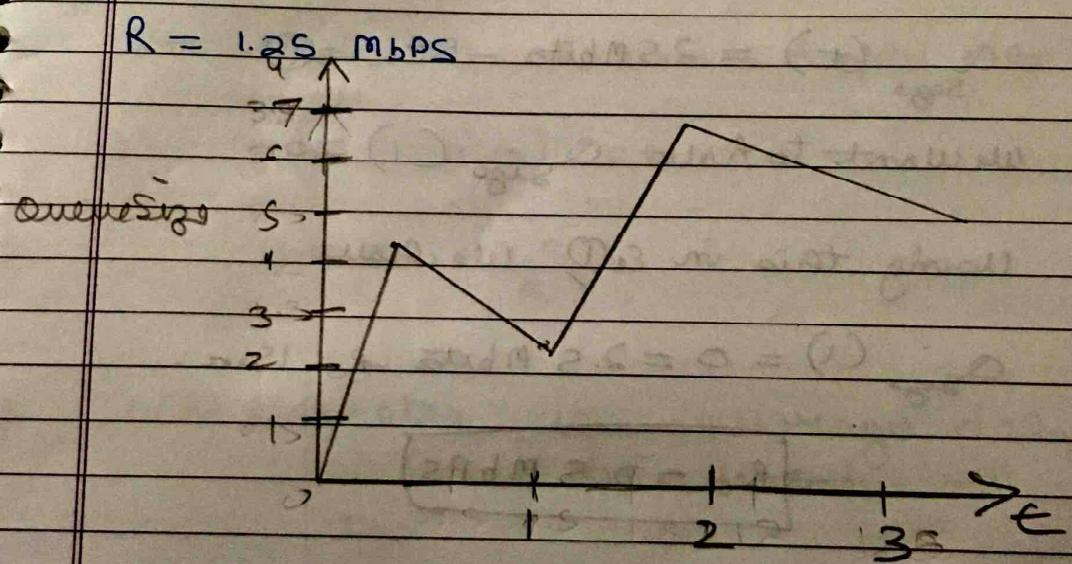
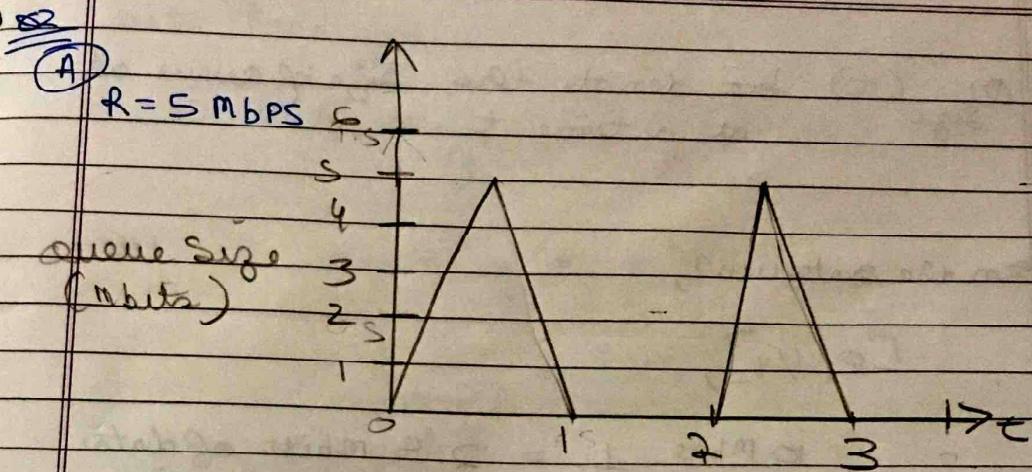
(E)

Store and forward:

Sender	Router	Router 2	Receiver
0			0
2	PACKET		2
4			4
6	PACKET 2		6
8			8
10	PACKET 2		10
12	PACKET 3	PACKET	12
14			14
16		done (PKT2)	16
18			18
20			20
22	PACKET 2		22
24	PACKET 3		24
26		PKT1	26
28			28
30		PKT1	30
32		done (PKT3)	32
34		PACKET 2	34
36			36
38			38
40			40
42			42
44		PKT3	44
46			46
48			48

Cut through :





(b)

$Q_{\text{size}}(t)$ denote the size of queue at time t .

In the interval,

$$[0, \frac{1}{4}],$$

10 Mbps. $\frac{1}{4} = 2.5$ mbits of data are received.

$$\Rightarrow Q_{\text{size}}(t) = 2.5 \text{ mbits} - Rt \rightarrow ①$$

We want to have $Q_{\text{size}}(1) = 0$.

Using this in ① we have.

$$Q_{\text{size}}(1) = 0 = 2.5 \text{ mbits} - R_0 \cdot 1 \text{ sec.}$$

$$\Rightarrow R_0 = 2.5 \text{ Mbps}$$

(c) Let $Q_{\text{size}}(t)$ be queue size at time t .

$$Q_{\text{size}}(t) = \begin{cases} (10-R)t & t \in [0, \frac{1}{4}], \\ \max(2.5 - Rt, 0) & t \in (\frac{1}{4}, 1) \end{cases}$$

$R \rightarrow \text{Rate}$

Let $D(R)$ denote the average delay of all

bits received

$$\begin{aligned} \Rightarrow D(R) &= \frac{1}{Y_4} \int_0^{Y_4} (10-R)t dt \\ &= \frac{2(10-R)}{R} t^2 \Big|_0^{Y_4} \\ &= \frac{10-R}{8R} \end{aligned}$$

①

Let at time t_0 , the queue is empty

$$\begin{aligned} \therefore \frac{1}{4} (10 \text{ nbps}) - R t_0 &= 0 \\ \Rightarrow t_0 &= \frac{2S}{R} \end{aligned}$$

$L(R)$: denotes the average delay over a time period.

$$\therefore L(R) = \int_0^{t_0} \text{size}(t) dt$$

$$\begin{aligned} &= \int_0^{Y_4} (10-R)t dt + \int_{Y_4}^{t_0} (2S - R t) dt \\ &= \frac{10-R}{2} \times \frac{1}{16} + 2S \times \left(\frac{2S}{R} - \frac{1}{4} \right) - \frac{R}{2} \left(\frac{(2S)^2}{R^2} - \frac{1}{4} \right) \\ &= \frac{S}{c} + \frac{(2S)^2}{2R} - \frac{2S}{4} \end{aligned}$$

$$D L(R) = \frac{2.5}{R} \left(\frac{10-R}{8} \right)$$

(e) $L(R) = \lambda \cdot D(R)$

$$= 2.5 \times \frac{(10-R)}{8R}$$

Q3

- (A) In this scenario, using a circuit switched network will be better choice because once the application starts, it runs for a longer duration & transmits data at steady rate. Also it would allow to reserve bandwidth for steady transmissions for a longer time.
- (B) Here no congestion control is needed. This is so as the capacities of each link is bigger than sum of application data rates & in the worst case when all the applications transmit data at the same time, each and every link is capable of handling data because of greater capacities.

Q4

(A) The Number of Users that can be Supported is
~~3~~ = $3 \text{ Mbps} / 150 \text{ Kbps} = 3 \times 1000 / 150$
 $= 20$

(B) $P = 0.1$

(C) ${}^{120}C_m \times (P)^m (1-P)^{120-m}$
 ${}^{120}C_m \times (P)^m (1-P)^{120-m}$ Here $P = 0.1$

(D) $1 - \sum_{n=0}^{20} \left({}^{120}C_n \right) \times (P^n) (1-P)^{120-n}$, Here $P = 0.1$

Q9

(A) $\text{total delay} = \frac{\text{queuing delay}}{R} + \frac{\text{transmission delay}}{R}$

$$= \frac{IL}{R(1-I)} + \frac{L}{R}$$

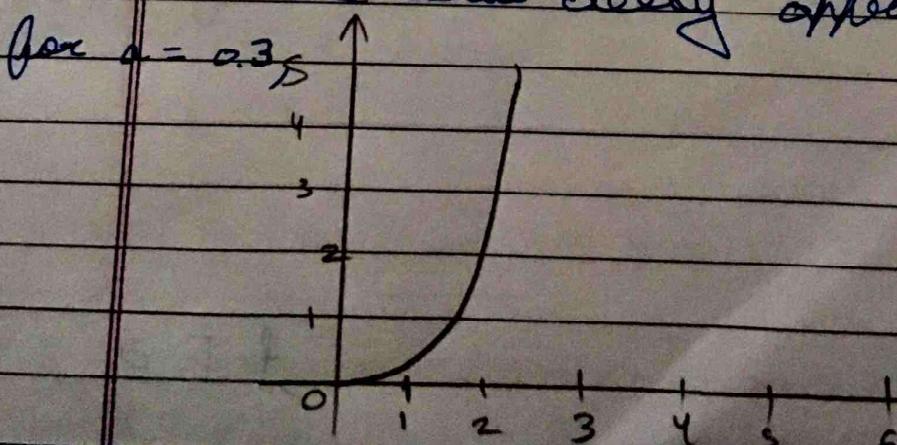
$$= \frac{L/R}{1-I}$$

(B) For plotting total delay as a function of L/R ,

let us assume $x = L/R$.

$$\Rightarrow \text{total delay} = \frac{x}{1-x} \quad \left\{ \because I = La/R \right\}$$

Now we can observe that, for $x=0$ the total delay is 0 & as we increase x , the total delay increases & when x comes closer to value of $1/a$, the total delay approaches infinity.



~~QF~~

$$N = \text{total Number of packets in system}$$
$$= 10 + 1 = 11$$

{ \therefore 2t applies
8 packets in buffer
1 packet
trans being transmitted }

also,

$$N = a \times (\text{queuing delay} + \text{transmission delay})$$

$$\Rightarrow 11 = a \times (0.01 + 1/100)$$

$$\Rightarrow 11 = a \times (0.01 + 0.01)$$

$$\Rightarrow 11 = a \times (0.02)$$

$$\Rightarrow a = \frac{11 \times 100}{2} = \frac{1100}{2} = 550 \text{ Packets/Sec.}$$

Q8

(A) Bandwidth Delay product = Rx Propagation delay

$$= 2 \text{ Mbps} * (20000 \text{ km}) / 2.5 \times 10^8 \text{ m/s}$$

$$= 2 \times 1000000 \times 20000 \times 10^3 / 2.5 \times 10^8 \text{ bits}$$

$$= 4000000 / 25$$

$$= 160,000 \text{ bits}$$

(B) The maximum Number of bits that will be present in a link at a given time : 160,000 bits

(C) The bandwidth delay product is defined as the maximum number of bits that can be in the link.

(D) Width of a bit = $\frac{\text{length of a link}}{\text{bandwidth delay product}}$

Width of a bit is longer than a football field because 1 bit is 125 meters long

(E) Width of a bit = S/R

7a.

Sample 1

```
storm:6% traceroute www.amazon.com
traceroute to www.amazon.com (23.202.75.20), 30 hops max, 60 byte packets
 1  128.227.205.193 (128.227.205.193)  0.560 ms  0.501 ms  0.488 ms
 2  csevl-core-msfc-1-v104-1.ns.ufl.edu (128.227.254.73)  0.431 ms  0.378 ms  0.391 ms
 3  ctx36-nexus-msfc-1-v41-1.ns.ufl.edu (128.227.236.13)  0.344 ms  0.320 ms  0.301 ms
 4  ctx36-pel-asr9001-1-v17-1.ns.ufl.edu (128.227.236.207)  1.091 ms  1.228 ms  1.308 ms
 5  ctx36-ewan-msfc-1-v18-1.ns.ufl.edu (128.227.236.209)  0.472 ms  0.436 ms  0.441 ms
 6  jax-flrcore-asr9010-1-te0104-1804.net.flrnet.org (108.59.29.242)  6.554 ms  6.300 ms  6.341 ms
 7  orl-flrcore-asr9010-1-hu0701-1.net.flrnet.org (108.59.31.151)  8.671 ms  8.654 ms  8.616 ms
 8  xe-11-3-0.bar2.Orlando1.Level3.net (67.30.142.5)  8.558 ms  8.495 ms  8.453 ms
 9  * * *
10  nap-brdr-01.inet.qwest.net (63.146.26.137)  13.031 ms  13.008 ms  12.869 ms
11  nap-edge-04.inet.qwest.net (205.171.19.53)  13.208 ms  13.002 ms  12.925 ms
12  a23-202-75-20.deploy.static.akamaitechnologies.com (23.202.75.20)  13.897 ms  13.915 ms  13.894 ms
```

Sample 2

```
storm:15% traceroute www.amazon.com
traceroute to www.amazon.com (23.202.75.20), 30 hops max, 60 byte packets
 1  128.227.205.193 (128.227.205.193)  0.500 ms  0.460 ms  0.524 ms
 2  csevl-core-msfc-1-v104-1.ns.ufl.edu (128.227.254.73)  0.421 ms  0.384 ms  0.535 ms
 3  ctx36-nexus-msfc-1-v41-1.ns.ufl.edu (128.227.236.13)  0.489 ms  0.424 ms  0.503 ms
 4  ctx36-pel-asr9001-1-v17-1.ns.ufl.edu (128.227.236.207)  0.949 ms  1.171 ms  3.410 ms
 5  ctx36-ewan-msfc-1-v18-1.ns.ufl.edu (128.227.236.209)  0.485 ms  0.471 ms  0.470 ms
 6  jax-flrcore-asr9010-1-te0104-1804.net.flrnet.org (108.59.29.242)  6.545 ms  6.570 ms  6.529 ms
 7  orl-flrcore-asr9010-1-hu0701-1.net.flrnet.org (108.59.31.151)  8.785 ms  8.737 ms  8.691 ms
 8  xe-11-3-0.bar2.Orlando1.Level3.net (67.30.142.5)  8.585 ms  8.573 ms  8.532 ms
 9  * * *
10  nap-brdr-01.inet.qwest.net (63.146.26.137)  12.924 ms  12.906 ms  12.860 ms
11  nap-edge-04.inet.qwest.net (205.171.19.53)  13.081 ms  13.038 ms  12.980 ms
12  a23-202-75-20.deploy.static.akamaitechnologies.com (23.202.75.20)  13.905 ms  13.946 ms  13.909 ms
```

Sample 3

```
storm:17% traceroute www.amazon.com
traceroute to www.amazon.com (23.202.75.20), 30 hops max, 60 byte packets
 1  128.227.205.193 (128.227.205.193)  0.477 ms  0.465 ms  0.457 ms
 2  csevl-core-msfc-1-v104-1.ns.ufl.edu (128.227.254.73)  0.520 ms  0.431 ms  0.376 ms
 3  ctx36-nexus-msfc-1-v41-1.ns.ufl.edu (128.227.236.13)  0.427 ms  0.383 ms  0.351 ms
 4  ctx36-pel-asr9001-1-v17-1.ns.ufl.edu (128.227.236.207)  1.135 ms  3.557 ms  3.491 ms
 5  ctx36-ewan-msfc-1-v18-1.ns.ufl.edu (128.227.236.209)  0.529 ms  0.494 ms  0.463 ms
 6  jax-flrcore-asr9010-1-te0104-1804.net.flrnet.org (108.59.29.242)  6.301 ms  6.392 ms  6.354 ms
 7  orl-flrcore-asr9010-1-hu0701-1.net.flrnet.org (108.59.31.151)  8.687 ms  8.646 ms  8.613 ms
 8  xe-11-3-0.bar2.Orlando1.Level3.net (67.30.142.5)  8.523 ms  8.505 ms  8.447 ms
 9  * * *
10  nap-brdr-01.inet.qwest.net (63.146.26.137)  12.994 ms  12.981 ms  13.053 ms
11  nap-edge-04.inet.qwest.net (205.171.19.53)  13.039 ms  13.149 ms  12.913 ms
12  a23-202-75-20.deploy.static.akamaitechnologies.com (23.202.75.20)  13.866 ms  14.043 ms  13.935 ms
```

Mean, Standard Deviation for each sample for Intra continent

1. 13.902, 0.011

2. 13.91, 0.03

3. 13.948, 0.089

b. There are 12 routers in the path . No, the path didn't change.

c. There are 4 ISP . Yes, the largest delays occurred at the peering interfaces between adjacent ISPs.

d. Intercontinent

Sample 1

```
storm:19% traceroute www.amazon.in
traceroute to www.amazon.in (104.92.16.18), 30 hops max, 60 byte packets
 1 128.227.205.193 (128.227.205.193) 0.450 ms 0.452 ms 0.476 ms
 2 csev1-core-msfc-1-v104-1.ns.ufl.edu (128.227.254.73) 10.319 ms 10.278 ms 10.228 ms
 3 ctx36-nexus-msfc-1-v41-1.ns.ufl.edu (128.227.236.13) 0.675 ms 0.630 ms 0.583 ms
 4 ctx36-pel-asr9001-1-v17-1.ns.ufl.edu (128.227.236.207) 1.027 ms 3.523 ms 3.488 ms
 5 ctx36-ewan-msfc-1-v18-1.ns.ufl.edu (128.227.236.209) 0.556 ms 0.492 ms 0.456 ms
 6 ssrb230a-ewan-msfc-1-v723-1.ns.ufl.edu (128.227.0.97) 0.392 ms 0.535 ms 0.475 ms
 7 wsip-184-188-101-186.at.at.cox.net (184.188.101.186) 1.320 ms 1.129 ms 1.324 ms
 8 ip68-105-160-3.ga.at.cox.net (68.105.160.3) 0.890 ms 0.903 ms 0.859 ms
 9 maribprj02.ae0.0.rd.at.cox.net (68.1.0.251) 14.292 ms 14.362 ms 14.297 ms
10 68.105.31.182 (68.105.31.182) 2498.640 ms 2498.625 ms 2498.584 ms
11 a104-92-16-18.deploy.static.akamaitechnologies.com (104.92.16.18) 14.844 ms 14.804 ms 14.956 ms
```

Sample 2

```
storm:20% traceroute www.amazon.in
traceroute to www.amazon.in (104.92.16.18), 30 hops max, 60 byte packets
 1 128.227.205.193 (128.227.205.193) 0.445 ms 0.429 ms 0.569 ms
 2 csev1-core-msfc-1-v104-1.ns.ufl.edu (128.227.254.73) 0.439 ms 0.388 ms 0.352 ms
 3 ctx36-nexus-msfc-1-v41-1.ns.ufl.edu (128.227.236.13) 5.799 ms 5.783 ms 5.738 ms
 4 ctx36-pel-asr9001-1-v17-1.ns.ufl.edu (128.227.236.207) 1.017 ms 1.271 ms 1.260 ms
 5 ctx36-ewan-msfc-1-v18-1.ns.ufl.edu (128.227.236.209) 0.524 ms 0.506 ms 0.506 ms
 6 ssrb230a-ewan-msfc-1-v723-1.ns.ufl.edu (128.227.0.97) 0.398 ms 0.412 ms 0.446 ms
 7 wsip-184-188-101-186.at.at.cox.net (184.188.101.186) 1.564 ms 1.544 ms 1.511 ms
 8 ip68-105-160-3.ga.at.cox.net (68.105.160.3) 0.804 ms 0.873 ms 0.839 ms
 9 maribprj02.ae0.0.rd.at.cox.net (68.1.0.251) 24.310 ms 24.275 ms 24.157 ms
10 * 68.105.31.182 (68.105.31.182) 408.861 ms *
11 a104-92-16-18.deploy.static.akamaitechnologies.com (104.92.16.18) 14.901 ms 14.865 ms 14.785 ms
```

Sample 3

```
storm:21% traceroute www.amazon.in
traceroute to www.amazon.in (104.92.16.18), 30 hops max, 60 byte packets
 1 128.227.205.193 (128.227.205.193) 0.510 ms 0.491 ms 0.533 ms
 2 csev1-core-msfc-1-v104-1.ns.ufl.edu (128.227.254.73) 0.466 ms 0.416 ms 0.367 ms
 3 ctx36-nexus-msfc-1-v41-1.ns.ufl.edu (128.227.236.13) 0.401 ms 0.386 ms 0.360 ms
 4 ctx36-pel-asr9001-1-v17-1.ns.ufl.edu (128.227.236.207) 1.033 ms 1.263 ms 1.248 ms
 5 ctx36-ewan-msfc-1-v18-1.ns.ufl.edu (128.227.236.209) 0.491 ms 0.501 ms 0.471 ms
 6 ssrb230a-ewan-msfc-1-v723-1.ns.ufl.edu (128.227.0.97) 0.412 ms 0.512 ms 0.451 ms
 7 wsip-184-188-101-186.at.at.cox.net (184.188.101.186) 1.231 ms 1.188 ms 1.160 ms
 8 ip68-105-160-3.ga.at.cox.net (68.105.160.3) 0.992 ms 0.957 ms 0.917 ms
 9 maribprj02.ae0.0.rd.at.cox.net (68.1.0.251) 14.276 ms 14.278 ms 14.245 ms
10 68.105.31.182 (68.105.31.182) 354.808 ms 354.793 ms 354.750 ms
11 a104-92-16-18.deploy.static.akamaitechnologies.com (104.92.16.18) 14.870 ms 14.845 ms 14.804 ms
```

Mean, Standard Deviation for each Sample

1. 14.868, 0.07879
2. 14.85, 0.059
3. 14.8396, 0.0333