

CS564

BitWeaving: Fast Scans for Main Memory Data Processing

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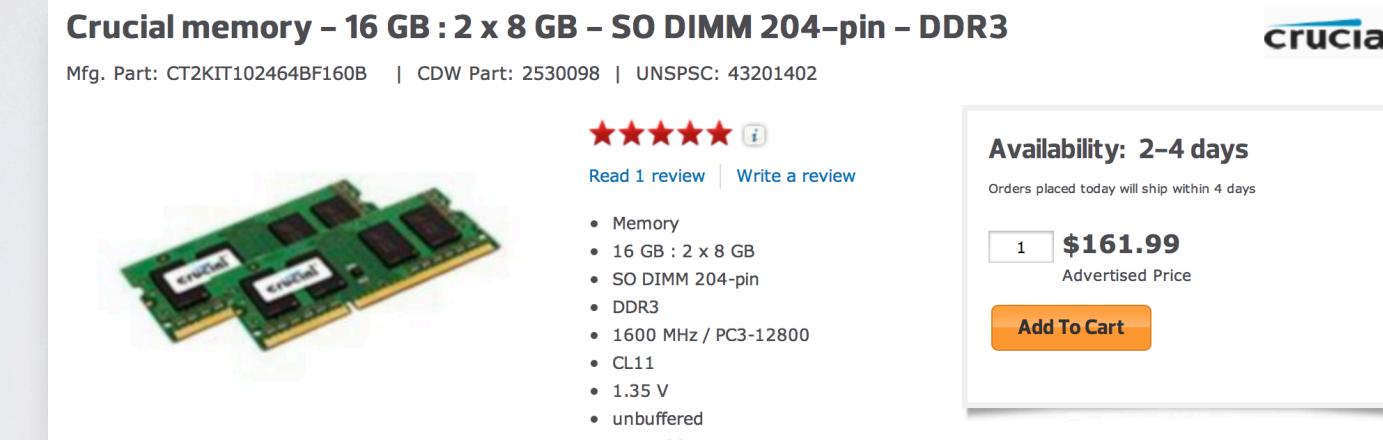


The Problem

- Need interactive analysis (complex SQL queries) run on large volumes of data
- New world in which business decisions are made by analytical engines
 - Speed is king

A common approach

- Disks are slow (yes – even flash is slow), but memory is fast
- Memory densities are increasing and price is dropping



Hence the rise of “in-memory” data processing for data analytics

Research Problem

- We know that for data analytics, using column stores is faster.
- But, can we go even faster than traditional column stores?
- Insights: Need to think about how the CPU sees the “data” and run data operations at the speed of the CPU
 - Recall CPU is the fastest component in the system

Column-store scans: Naïve method

- An example SQL scan query:

```
SELECT COUNT(*)  
FROM Customer  
WHERE age BETWEEN 20 AND 24
```

- An naïve implementation:

```
count = 0;  
FOR EACH value v in column age  
    IF (20 <= v and v <= 24)  
        count++;
```

- Complexity: $O(n)$. Need to run $O(n)$ CPU instructions.
- Better method?

Encoded column values

- Domain size of column is typically small
 - Gender: Male / Female 1 bits
 - Age: 0-122 7 bits
 - States: 50 states 6 bits
- DBMSs converts native column values to **codes**.
- Codes only use as many bits as are needed for fixed-length encoding.

Motivation

```
SELECT count(*)  
FROM Customer  
WHERE age BETWEEN 20 AND 24
```

7 bits

```
count = 0;  
FOR EACH value v in column age  
  IF (20 <= v and v <= 24)  
    count++;
```

CPU register



SIMD word size: 256 bits

~~Word size: 64 bits~~

Code size: 7 bits

Intra-cycle parallelism!

Column-store scans: BitWeaving method

- An example SQL scan query:

```
SELECT COUNT(*)  
FROM Customer  
WHERE age BETWEEN 20 AND 24
```

- An BitWeaving implementation:

```
count = 0;  
FOR EACH group of codes v in column age  
    Evaluate 20 <= v <= 24 in parallel  
    Update count;
```

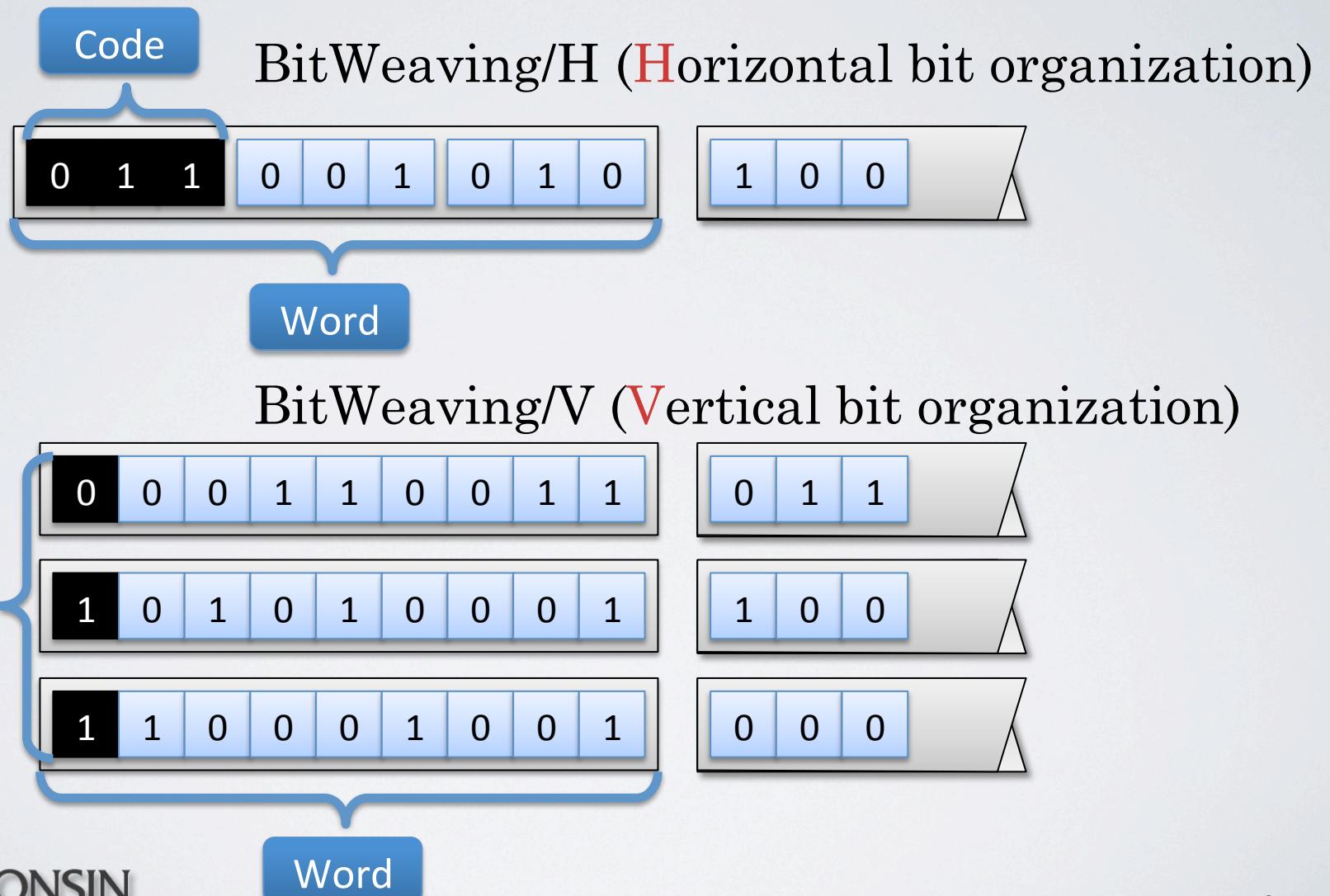
- Complexity:

$O(n/w)$. w: group size.

BitWeaving

- In this lecture, we introduce BitWeaving
 - A fast **scan** method
 - for **column stores**
- Fully exploits intra-cycle parallelism
- How: By “gainfully” using every bit in every processor word.

BitWeaving: Two Flavors

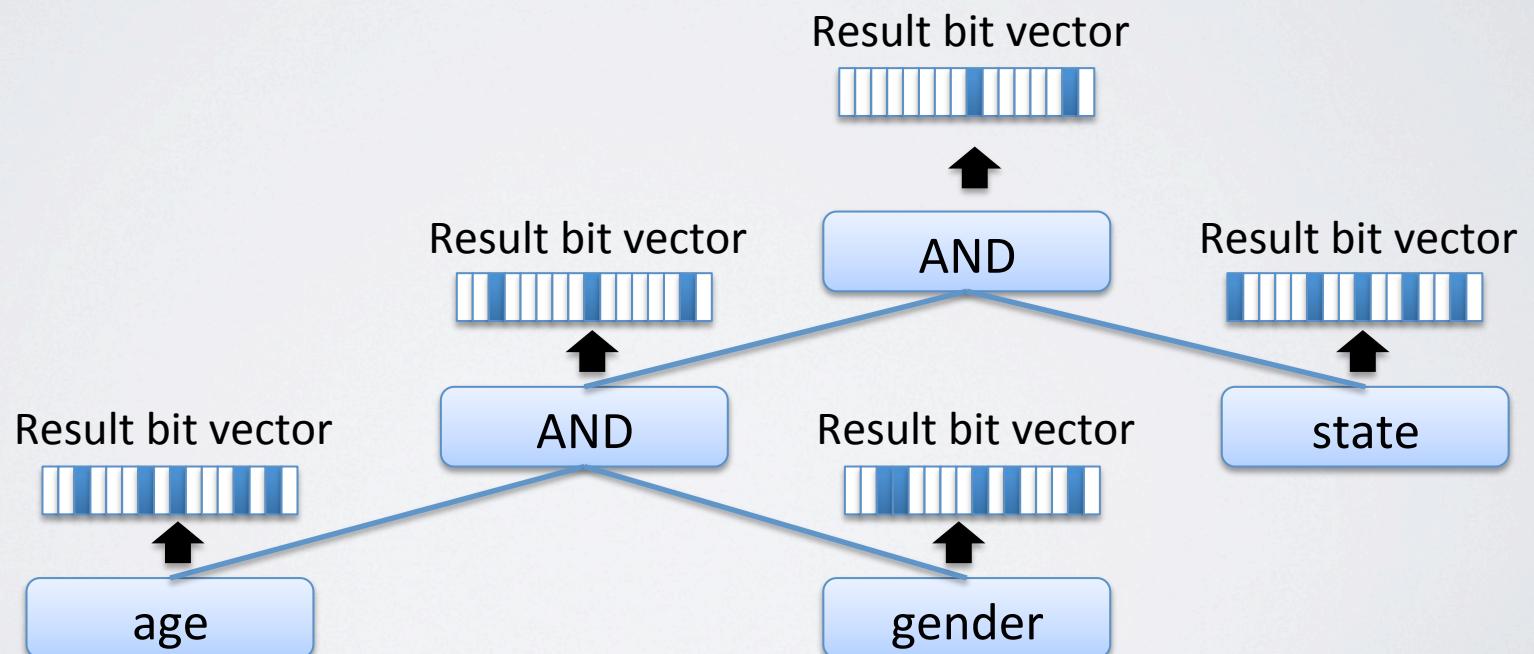


Framework

- Targets single-table scans
- Column-scalar scan: scan on a single column
 - produce a **result bit vector**, with **one** bit for each input tuple to indicate the matching tuples
- Complex predicates in the scan: logical AND and OR operations on these *result bit vectors*

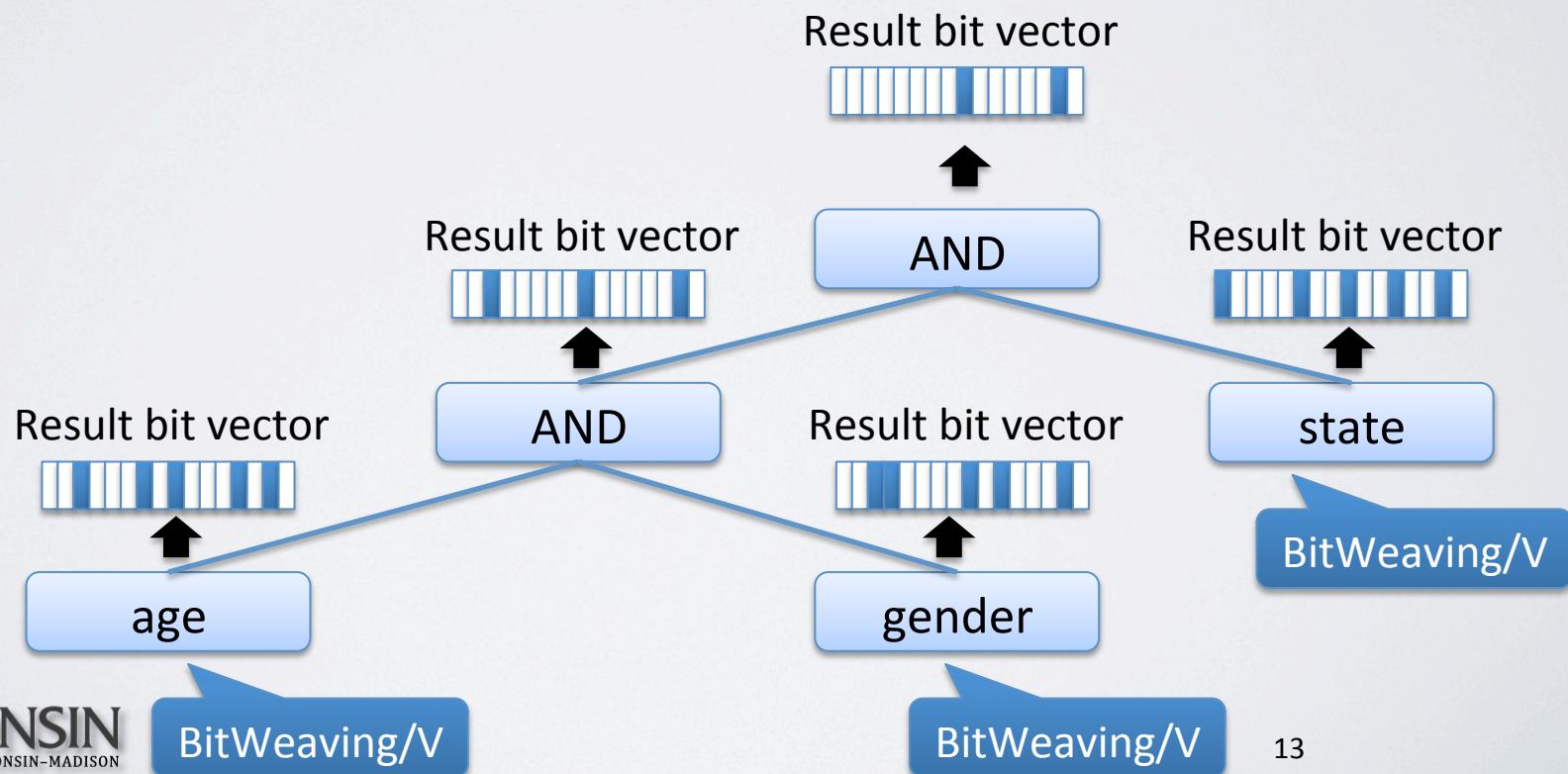
Framework – Example

```
SELECT COUNT(*) FROM Customer  
WHERE age BETWEEN 20 AND 24  
    AND gender = Male  
    AND state = Wisconsin
```



Framework – Example

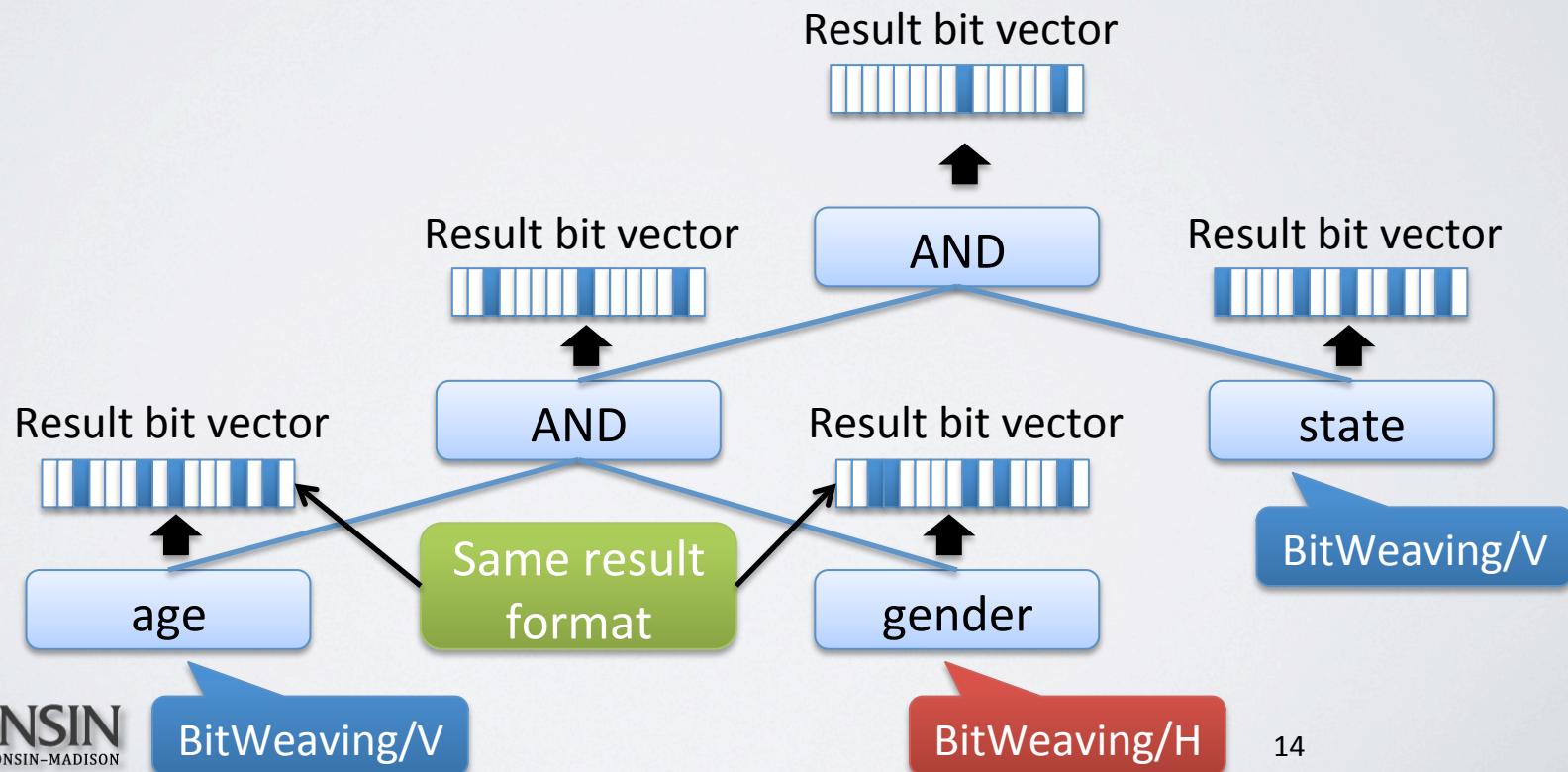
```
SELECT COUNT(*) FROM Customer  
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    AND gender = Male  
    AND state = Wisconsin
```



Framework – Example

```
SELECT COUNT(*) FROM Customer  
WHERE age BETWEEN 20 AND 24  
    AND gender = Male  
    AND state = Wisconsin
```

Mixing of BitWeaving/V BitWeaving/H columns



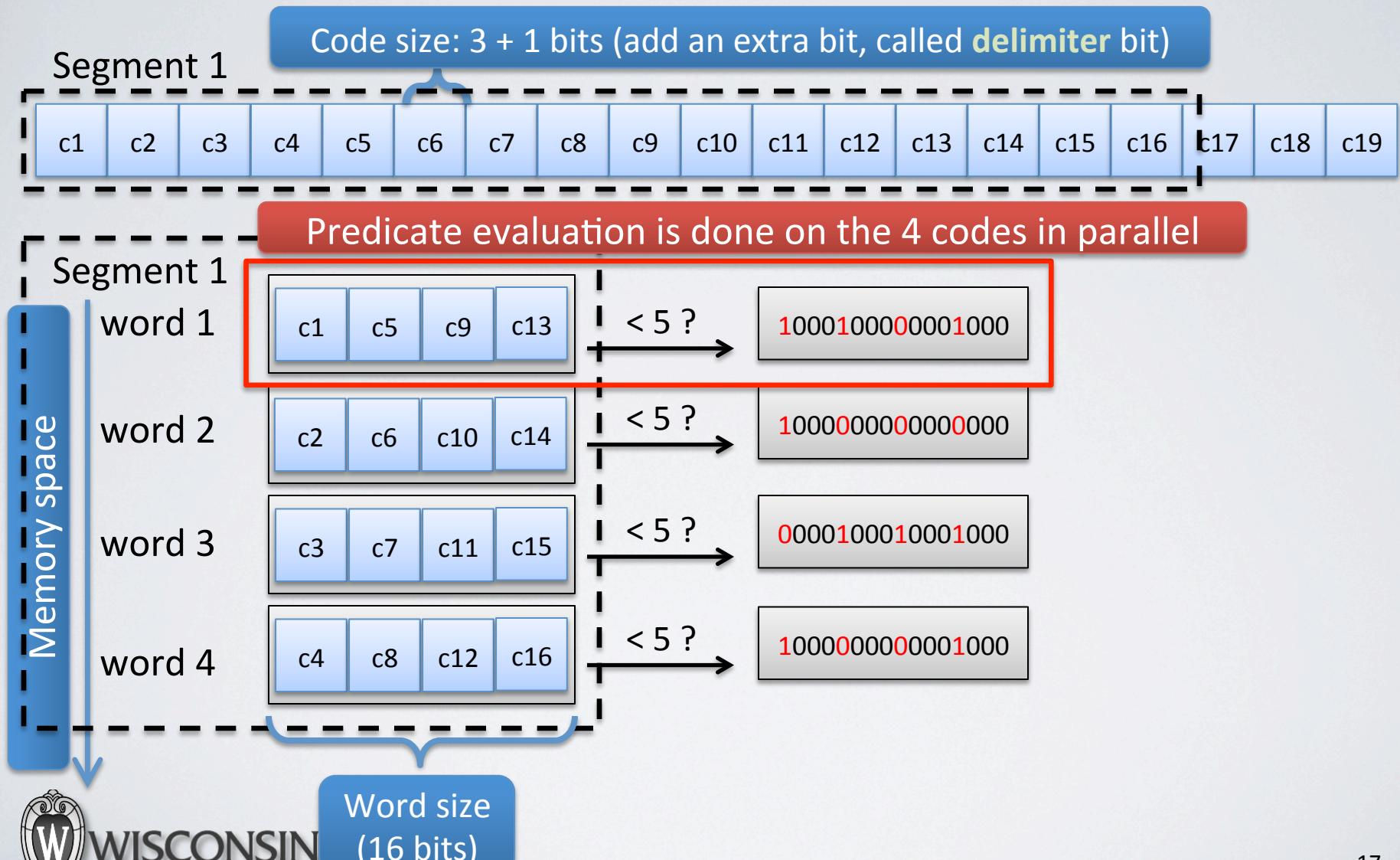
Outline

- Motivation & Overview
- BitWeaving/H
- BitWeaving/V
- Conclusion

BitWeaving/H

- Storage layout
 - Packs codes “horizontally” into processor words
 - Uses an extra bit (**delimiter** bit) in each code
 - Staggers codes across words inside a segment
- Column-scalar scan
 - Parallel predicate evaluation on packed codes

BitWeaving/H - Example

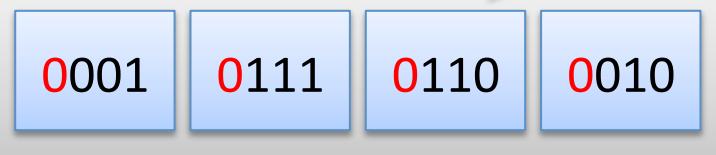


BitWeaving/H Example: Less Than Predicate (< 5)

Word size
(16 bits)

$$X = (c_1 c_5 c_9 c_{13})$$

c₁=1 c₅=7 c₉=6 c₁₃=2



Regular
integer plus

Only use 3 instructions!

0101

$$(Y + (X \oplus M1)) \wedge M2$$

$$M1 = 0111 \ 0111 \ 0111 \ 0111$$

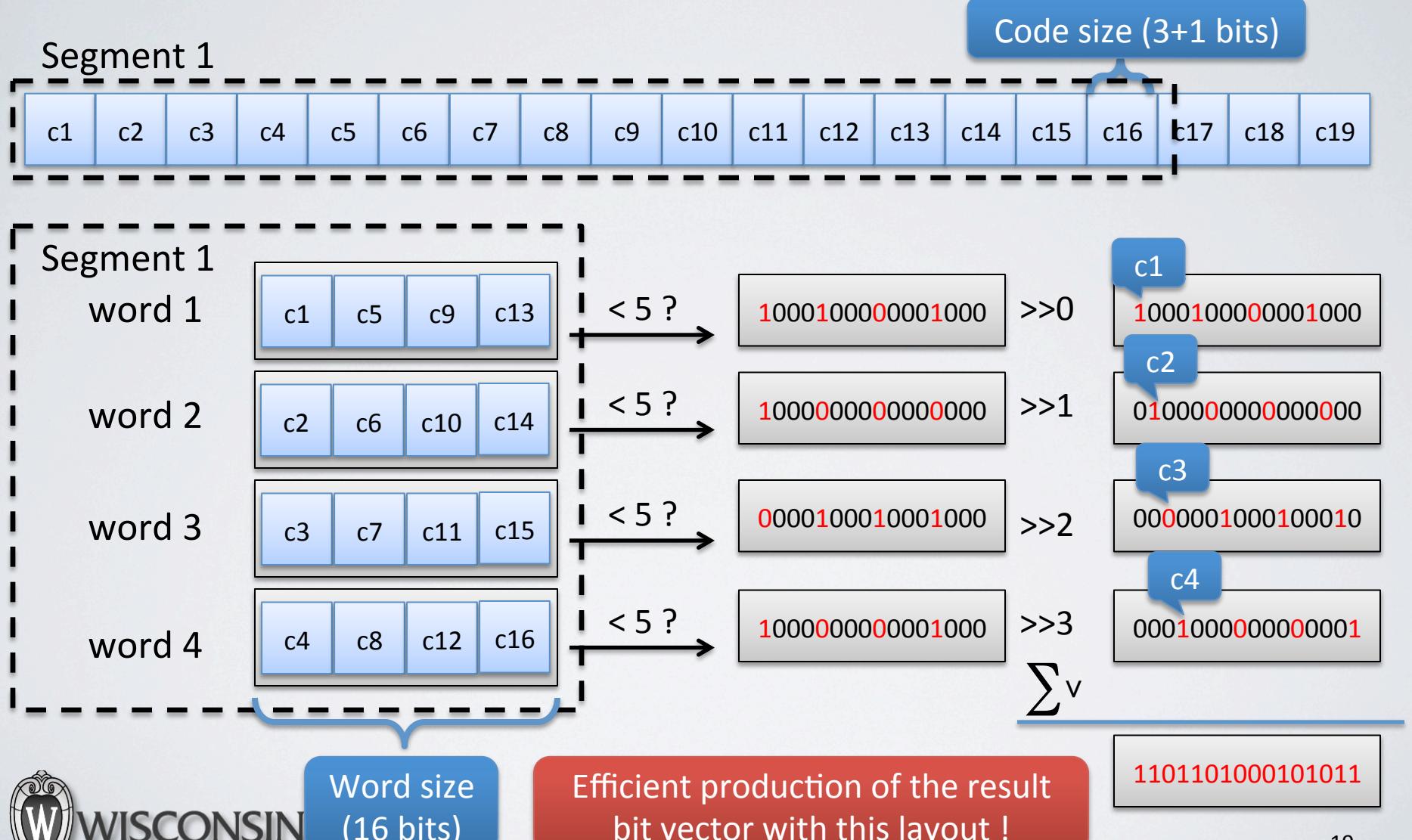
$$M2 = 1000 \ 1000 \ 1000 \ 1000$$



Works for arbitrary code sizes & word sizes!

Curious about why? See our paper!

BitWeaving/H – Example...



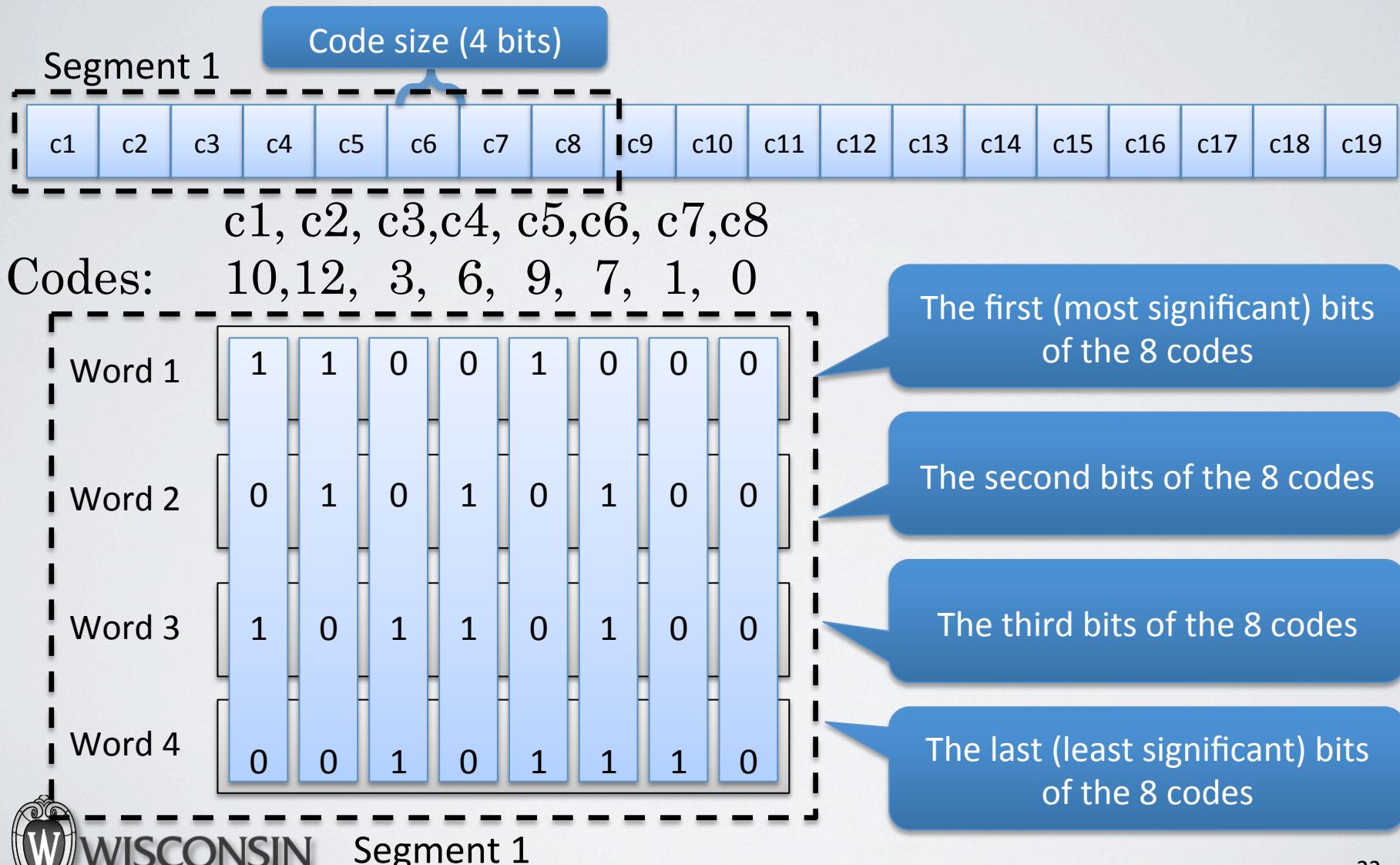
Outline

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BitWeaving/V

- Storage layout
 - Bit-level columnar data organization, i.e. its like a **bit-level columnar store**.
- Column-scalar scan
 - Predicate evaluation is converted to logical computation on these “words of bits”

BitWeaving/V – Storage Layout



BitWeaving/V – Column-scalar Scan

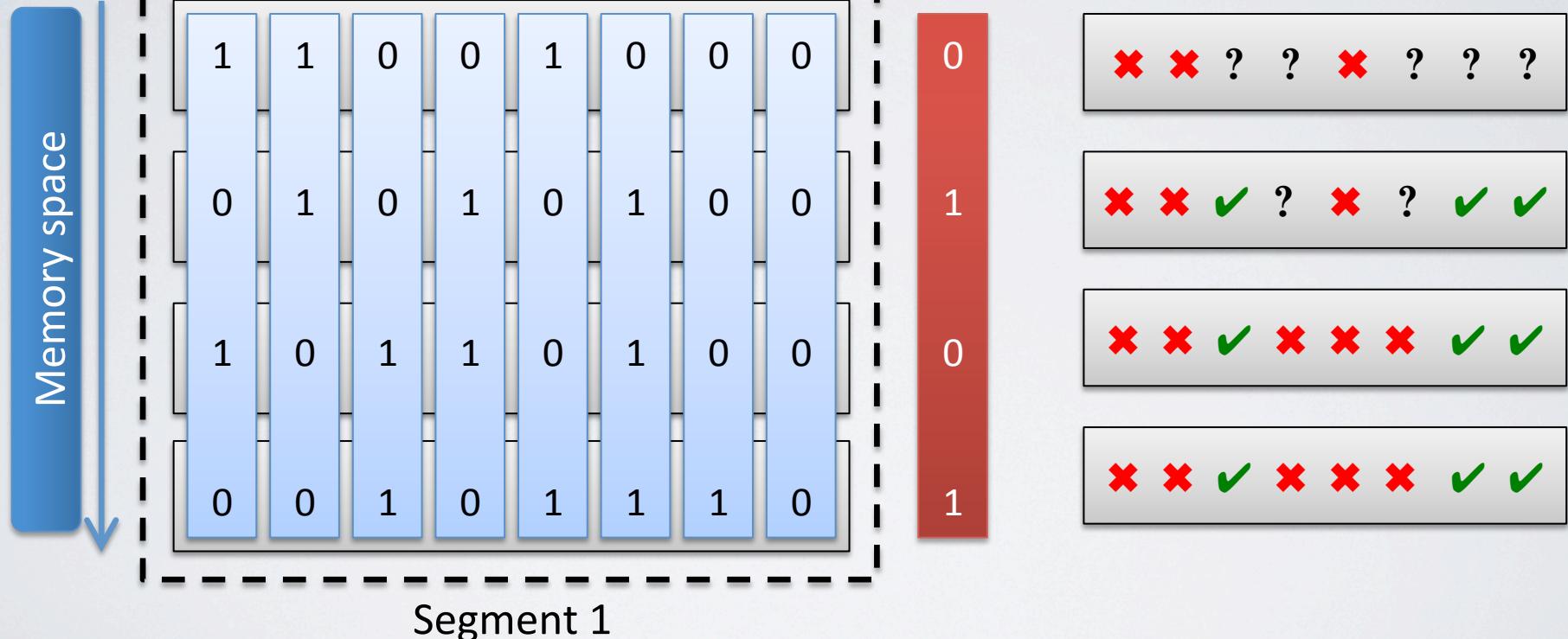
Column codes:

10, 12, 3, 6, 9, 7, 1, 0

Constant

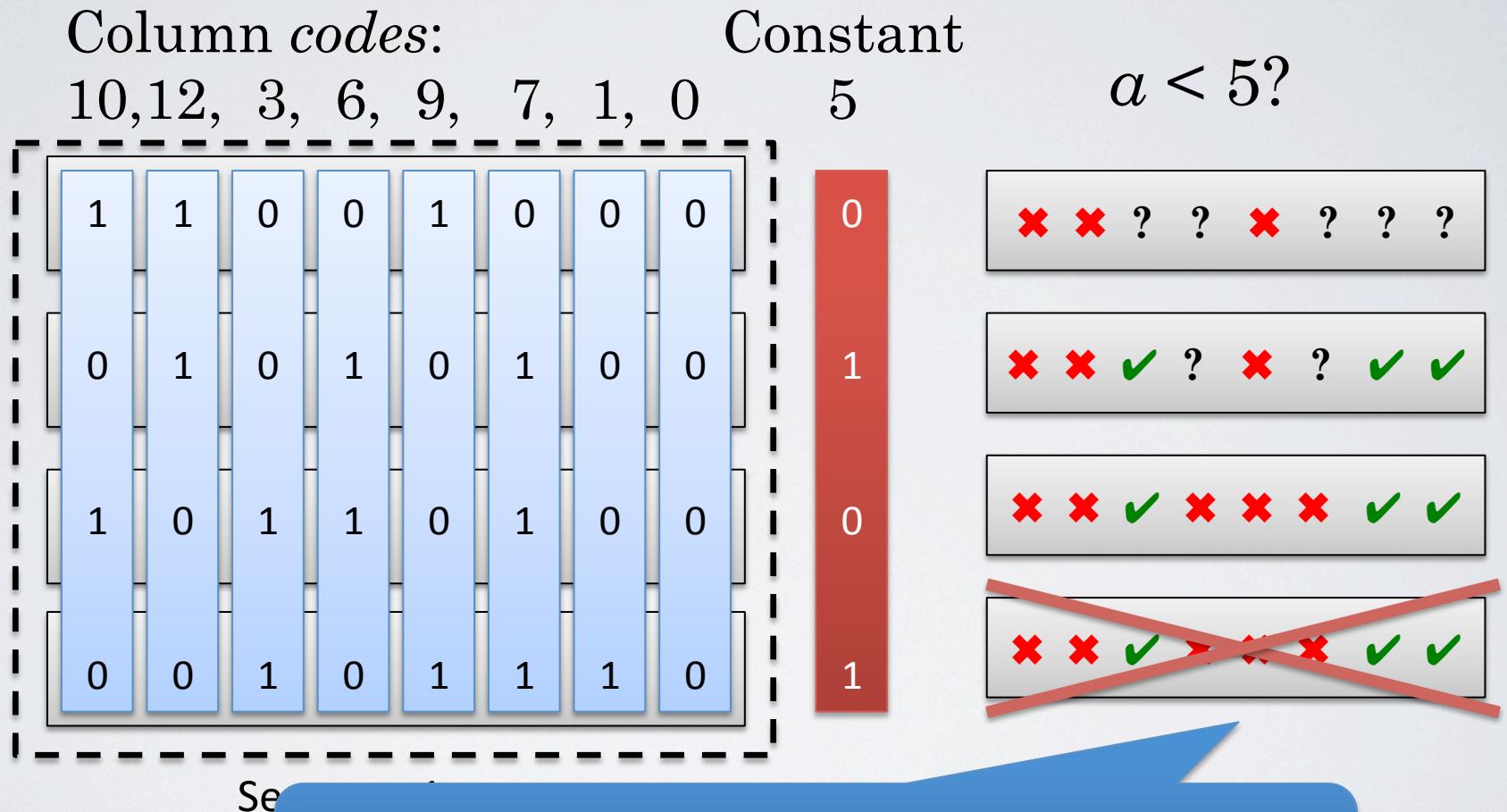
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$a < 5?$



The layout of the segment exactly matches the access pattern of column-scalar scans

BitWeaving/V – Early pruning



Outline

- Motivation & Introduction
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Conclusions

BitWeaving: A new method to use all the bits in a processor word gainfully.

Two flavors: BitWeaving/H and BitWeaving/V.

BitWeaving are faster than state-of-the-art scan methods, in some cases by an order of magnitude.

Resource

- Blog article:
 - <http://bigfastdata.blogspot.com/>
- Paper:

BitWeaving: Fast Scans for Main Memory Data Processing. In SIGMOD 2013.