

STORING DATA: DISK AND FILES

CS 564- Fall 2015

ACKs: Dan Suciu, Jignesh Patel, AnHai Doan

MANAGING DISK SPACE

- The disk space is organized into **files**
- Files are made up of **pages**
- Pages contain **records**

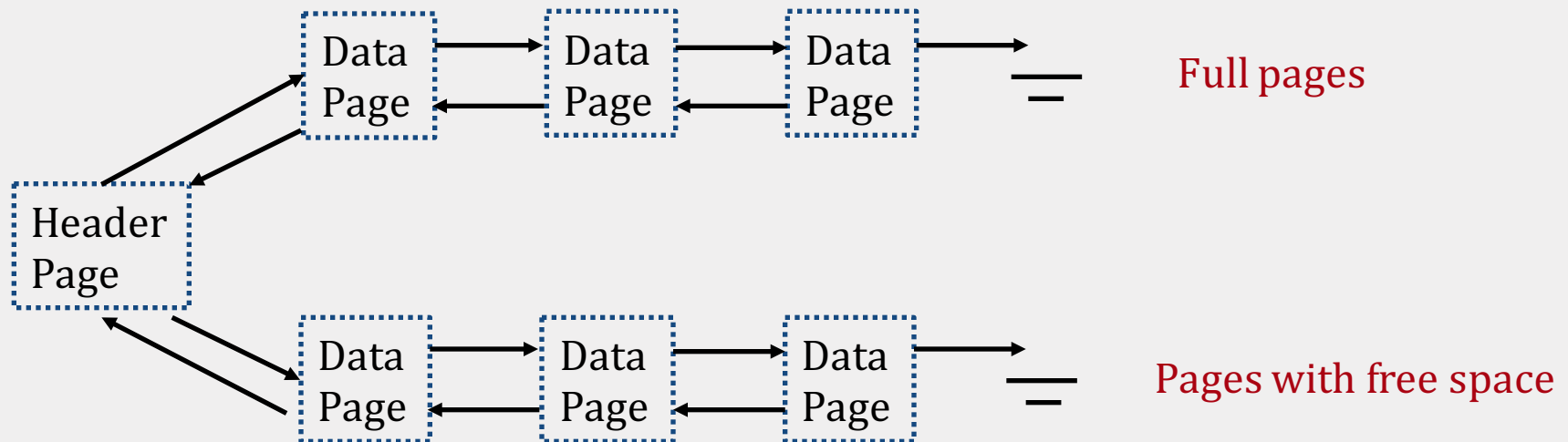
FILE ORGANIZATION

UNORDERED (HEAP) FILES

- Contains records in no particular order
- As file grows/shrinks, disk pages are allocated/deallocated
- To support record level operations, we must keep track of:
 - the pages in a file: page id (**pid**)
 - free space on pages
 - the records on a page: record id (**rid**)

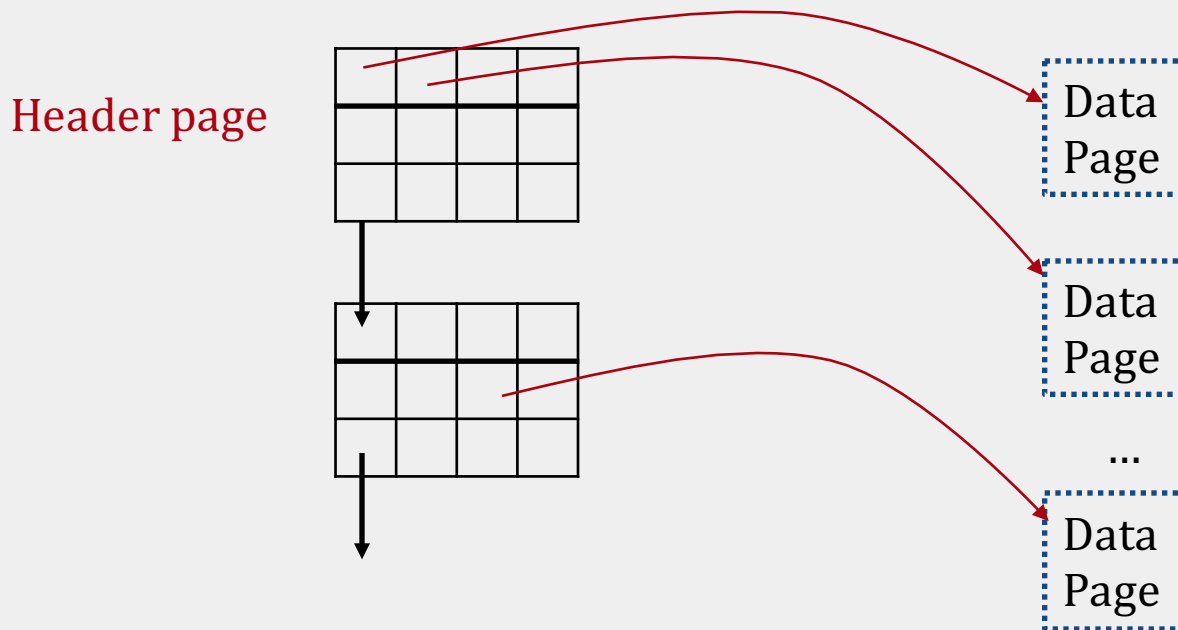
HEAP FILE AS LINKED LIST

- (heap file name, header page id) stored somewhere
- Each page has 2 pointers + data
- Pages in the free space list have “some” free space



HEAP FILE AS PAGE DIRECTORY

- Entry for a page:
 - Free/full
 - Number of free bytes
- Can locate pages for new tuples faster!



PAGE ORGANIZATION

FILES OF RECORDS

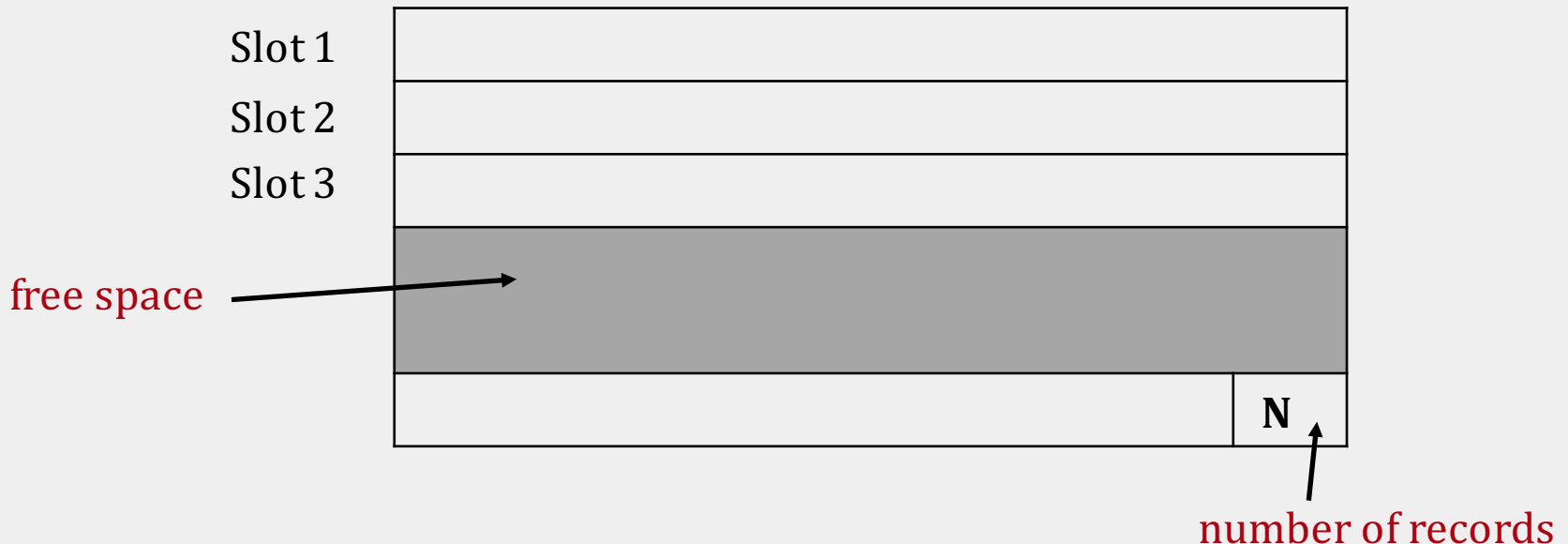
- Page or block is OK for I/O, but higher levels operate on records, and files of records
- File operations:
 - insert/delete/modify record
 - read a record (specified using the record id)
 - scan all records (possibly with some conditions on the records to be retrieved)

PAGE FORMATS

- A page is collection of records
- Slotted page format
 - Page a collection of slots
 - Each slot contains a record
- **rid** = <page id, slot number>
- Many slotted page organizations
- Support for:
 - search, insert, delete records on a page

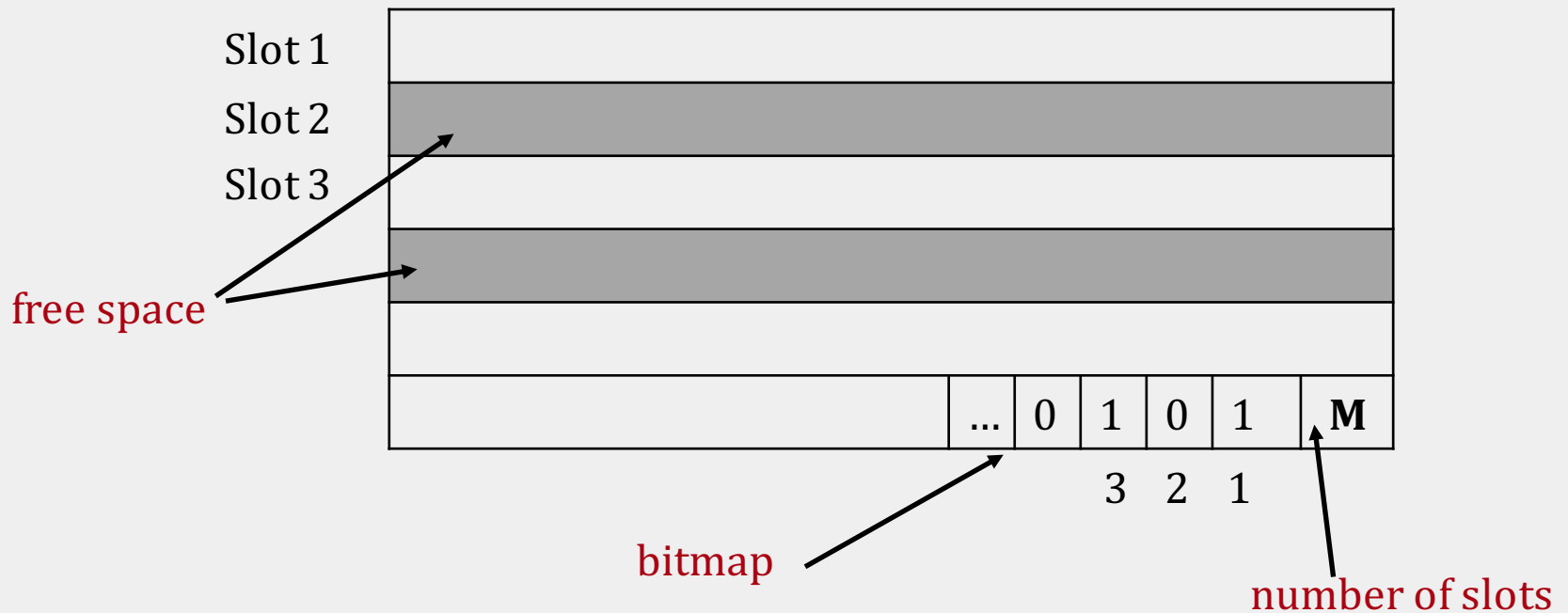
FIXED LENGTH RECORDS (1)

- **packed** organization: records are always stored in the first N slots
- problem when there are references to records!

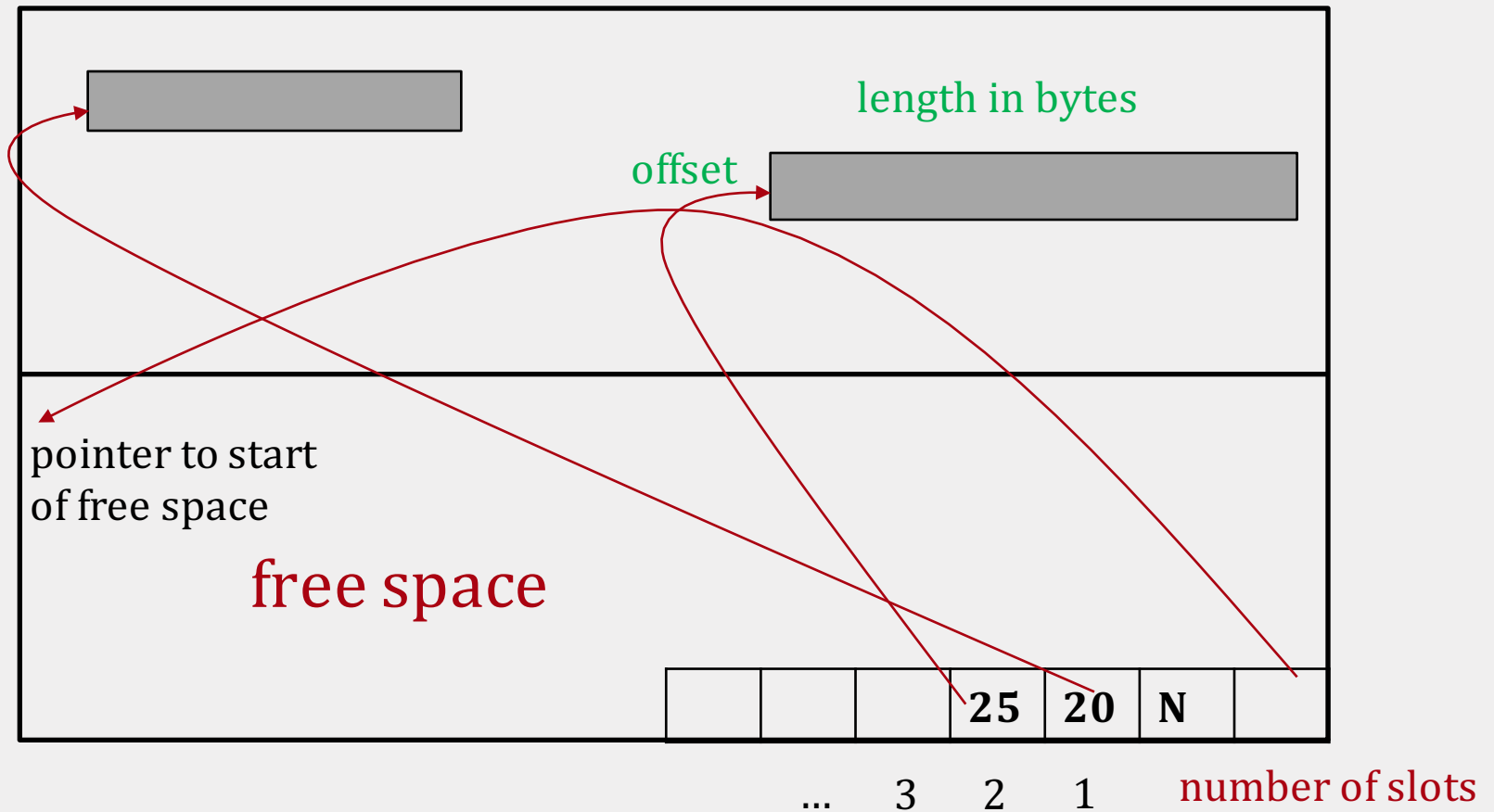


FIXED LENGTH RECORDS (2)

- **unpacked** organization: use a **bitmap** to locate records in the page



VARIABLE LENGTH RECORDS



VARIABLE LENGTH RECORDS

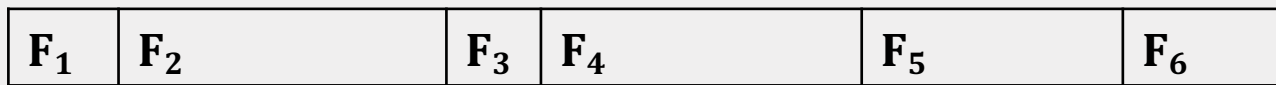
- **Deletion:**
 - offset is set to -1
- **Insertion:**
 - use any available slot
 - if no space is available, reorganize
- rid remains unchanged when we move the record (since it is defined by the slot number)

RECORD FORMAT

- How do we organize the field **within** a record?
 - fixed length
 - variable length
- Information common to all records of a given type is kept in the **system catalog**:
 - number of fields
 - field type

RECORD FORMAT: FIXED LENGTH

- All records have the same length and same number of fields
- The address of any field can be computed from info in the system catalog!



L_2 = length of field F_2

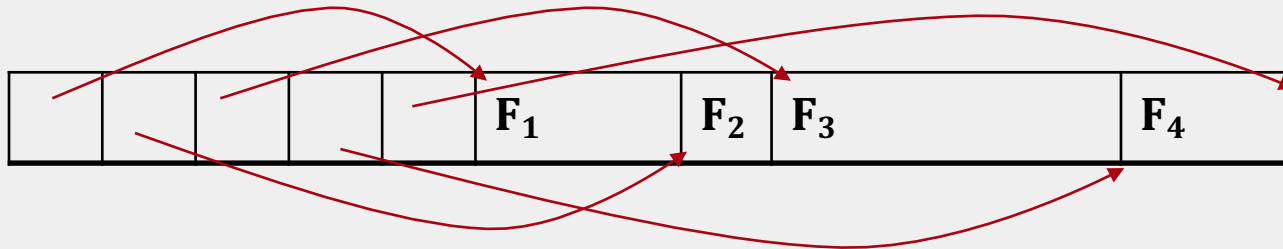
RECORD FORMAT: VARIABLE LENGTH (1)

- store fields consecutively
- use **delimiters** to denote the end of a field
- need a scan of the whole record to locate a field

F₁	\$	F₂	\$	F₃	\$	F₄	\$	F₅
----------------------	-----------	----------------------	-----------	----------------------	-----------	----------------------	-----------	----------------------

RECORD FORMAT: VARIABLE LENGTH (2)

- store fields consecutively
- use an array of integer offsets in the beginning



BONUS: COLUMN STORES

- Consider a table:
 - **Foo** (a INTEGER, b INTEGER, c VARCHAR(255))
- and the query:
 - **SELECT a FROM Foo WHERE a > 10**
- What could be the problem when we read using the previous record formats?

BONUS: COLUMN STORES

- We can instead store data **vertically** !
- Each column of a relation is stored in a **different file** (and can be compressed as well)

1234	45	Here goes a very long sentence 1
4657	2	Here goes a very long sentence 2
3578	45	Here goes a very long sentence 3

row-store

column-store

1234	45
4657	2
3578	45

Here goes a very long sentence 1
Here goes a very long sentence 2
Here goes a very long sentence 3