

# Introduction to Separation Logic

Heavily borrows from slides by Cristiano Calcagno, Imperial College  
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# Syntax of Separation Logic

- Given a decidable base-theory  $T$ , the syntax of separation logic  $SL(T)_{Loc, Data}$  is presented
- $Loc$  and  $Data$  represent the type of the address and the values
- E.g Setting  $Loc$  and  $Data$  to be  $Int$ , then our addresses and values are integers

$$\begin{aligned} P, Q ::= & \text{false} \mid P \wedge Q \mid P \vee Q \mid P \rightarrow Q \\ & \mid P * Q \mid P \multimap Q \\ & \mid E = E' \mid E \hookrightarrow E' \mid \text{empty} \end{aligned}$$

We use  $E$  and  $E'$  to denote expressions in the base theory, where pointer indirection is not used.

# Semantics of Separation Logic

The model consists of an interpretation ( $I$ ) and a heap ( $h$ )

$$I : \text{Var} \rightarrow \text{Loc}$$

$$h : \text{Loc} \rightarrow \text{Data}$$

$I, h \models \text{false}$	never satisfied
$I, h \models P \wedge Q$	$I, h \models P$ and $I, h \models Q$
$I, h \models P \vee Q$	$I, h \models P$ or $I, h \models Q$
$I, h \models P \rightarrow Q$	$I, h \models P$ implies $I, h \models Q$
$I, h \models E = E'$	$\llbracket E \rrbracket_I = \llbracket E' \rrbracket_I$

We use  $\llbracket E \rrbracket_I$ , to denote the value of  $E$  under the interpretation  $I$ .

Empty heap

$$\begin{aligned} I, h &\models \text{empty} \\ \text{iff } h &= \phi \end{aligned}$$

Separating conjunction

$$\begin{aligned} I, h &\models P * Q \\ \text{iff } \exists h_1, h_2. &(h_1 \perp h_2) \wedge (h = h_1 \circ h_2) \wedge I, h_1 \models P \wedge I, h_2 \models Q \end{aligned}$$

Where  $h_1 \perp h_2$  denotes that the heaps are disjoint and  $h_1 \circ h_2$  means their union.

## Separating Implication

$$\begin{aligned} I, h &\models P \multimap Q \\ \text{iff } \forall h'. (h \perp h') \wedge (I, h' &\models P) \rightarrow I, h \circ h' \models Q \end{aligned}$$

Interpretation : If we extend the current heap with a disjoint heap satisfying  $P$ , then the new heap satisfies  $Q$ . In some ways, we can imagine that our current heap is only missing the records of  $P$ , to make it satisfy  $Q$ .

Points to

$$\begin{aligned} I, h &\models E \hookrightarrow E' \\ \text{iff } h(\llbracket E \rrbracket_I) &= \llbracket E' \rrbracket_I \end{aligned}$$

# Examples

Points to,

$$F : x \hookrightarrow 10$$

$$I : \{(x, 0)\}$$

$$h : \{(0, 10)\}$$

$$I, h \models F$$

Separating conjunction,

$$F : x \hookrightarrow 10 * y \hookrightarrow 20$$

$$I : \{(x, 0), (y, 1)\}$$

$$h : \{(0, 10), (1, 20)\}$$

$$I, h \models F$$



## Separating Implication

$$\begin{aligned} I, h &\models P \multimap Q \\ \text{iff } \forall h'. (h \perp h') \wedge (I, h' &\models P) \rightarrow I, h \circ h' \models Q \end{aligned}$$

Example,

$$F : (x \hookrightarrow 10) \multimap (x \hookrightarrow 10 * y \hookrightarrow 20)$$

$$I : \{(x, 0), (y, 1)\}$$

$$h' : \{(0, 10)\}$$

$$h : \{(1, 20)\}$$

$$h \circ h' : \{(0, 10), (1, 20)\}$$

$$I, h \models F$$

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# Translating Separation Logic into Pointer Logic

Points to,

$$\begin{aligned} I, h &\models x \hookrightarrow v \\ &\iff \\ L, M &\models *x = v \end{aligned}$$

Separating conjunction,

$$\begin{aligned} I, h &\models x \hookrightarrow v_1 * y \hookrightarrow v_2 \\ &\iff \\ L, M &\models *x = v_1 \wedge *y = v_2 \wedge x \neq y \end{aligned}$$

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