

#1.  $G = (D, C, A, X) \rightarrow$  Properties of the Graph -  $G$   
 $\hookrightarrow$  Customer Nodes  
 $\hookrightarrow$  Depot Nodes

Decision variables:  $x_{ij}^v$ ,  $y_d$ ,  $z_c^v$

$x_{ij}^v$   $\rightarrow$  Whether vehicle  $v$  traverses  $(i, j)$  arc

$y_d$   $\rightarrow$  Whether depot node  $d$  is open (operational)

$z_c^v$   $\rightarrow$  Load of vehicle  $v$  on arrival at customer node  $c$ .

Deep-dive into constraints for  $z_c^v$ .

$\rightarrow$  if  $x_{ck}^v = 1$  i.e. if vehicle  $v$  visits node  $k$  from customer node  $c$ , then,

$z_k^v$  will reduce in comparison to  $z_c^v$  by at least as much as the demand at customer node  $c \Rightarrow q_c$ .

$$\therefore z_k^v \leq z_c^v - q_c$$

$\rightarrow$  if  $x_{ck}^v = 0$  i.e. if vehicle  $v$  does not visit node  $k$  from customer node  $c$ , then,

$z_k^v$  can take any value. There is no relation between  $z_k^v$  and  $z_c^v$ .

$$\therefore z_k^v \in \mathbb{R}_+$$

We re-write this as,

$$z_k^v \leq M$$

where  $M$  is large enough.

#2 Combining the two if conditions using  $x_{ck}^v$

$$Z_k^v \leq \underbrace{(x_{ck}^v)(Z_c^v - q_v)}_{\text{if } x_{ck}^v = 1 \text{ this part of RHS will be active}} + \underbrace{(1 - x_{ck}^v)(M)}_{\text{if } x_{ck}^v = 0 \text{ then this section of RHS will be active.}}$$

if  $x_{ck}^v = 1$  this part of RHS will be active

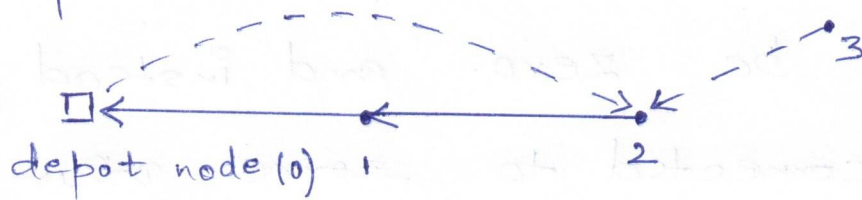
if  $x_{ck}^v = 0$  then this section of RHS will be active.

In addition to this constraint, we have

$$Z_c^v \leq q_v \longrightarrow \text{Vehicle Capacity Constraint}$$

$$Z_d^v = 0 \longrightarrow \text{Boundary Condition.}$$

Examining constraints on  $Z_c^v$  through a small example,



$$\Rightarrow \boxed{Z_0^v = 0} \text{ (Boundary Condition)}$$

also, for connection between 1 and 0.

$$- Z_0^v \leq x_{10}^v (Z_1^v - q_1) + (1 - x_{10}^v) M$$

$$\text{Since } x_{10}^v = 1$$

$$Z_0^v \leq Z_1^v - q_1$$

$$\Rightarrow Z_1^v \geq \cancel{Z_0^v} + q_1 \Rightarrow \boxed{Z_1^v \geq q_1}$$

also, for connection between 2 and 1.

$$- Z_1^v \leq x_{21}^v (Z_2^v - q_2) + (1 - x_{21}^v) M$$

$$\text{Since } x_{21}^v = 1$$

$$\cancel{Z_1^v} \leq Z_2^v - q_2$$

$$\Rightarrow \boxed{Z_2^v \geq q_1 + q_2}$$

Likewise, we can write,

$$Z_3^v \geq q_1 + q_2 + q_3$$

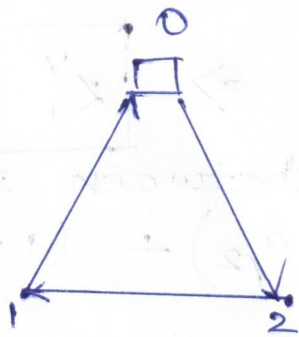
$$\text{if } x_{32}^v = 1.$$

Let's now assume  $q_1 + q_2 + q_3 > q_v$   
vehicle capacity  $\leftarrow$



#11. Because capacity constraint is violated  
connecting  $3 \rightarrow 2 \rightarrow 1 \rightarrow 0$ ,  
 $x_{32}^v$  has to be zero and instead, node 2  
should be connected to some other node

If  $T(2) = \{0, 3\}$  then the only option  
left is to connect node 2 back to  
the depot node (0). Consequently, the  
tour should be,



Hence, the load balance, vehicle capacity, <sup>boundary condition</sup> and customer service constraints together  
enforce that a delivery tour starts  
and ends at the depot node.