Graph Coloring Problem

Problem Statement: The Graph Coloring Problem is a classic challenge in graph theory. The goal is to color the vertices of a graph in such a way that no two adjacent vertices (those connected directly by an edge) have the same color.

Given that, an undirected graph G = (V, E); where V represents the set of vertices and E represents the set of edges, the objective is to find a minimum k-coloring, where k represents the smallest number of colors needed to color all vertices like any edge $(u, v) \in E$, the assigned colors satisfy $C(u) \neq C(v)$.

Algorithms for Solving Graph Coloring

- 1- Greedy Coloring Algorithm
- 2- Backtracking Algorithm
- 3- DSATUR Algorithm

Explain the algorithms for the problem

1. Backtracking Algorithm: This algorithm tries to color the graph by assigning colors to vertices and backtracking when it encounters a conflict.

Sketch of the Algorithm:

- 1. Start with the first vertex.
- 2. Assign it the first available color.
- 3. Move to the next vertex and assign a color that does not conflict with adjacent vertices.
- 4. If no valid color is found, backtrack and try a different color.
- 5. Continue until all vertices are colored or backtracking exhausts all options.
- 2. **Greedy Coloring Algorithm:** This algorithm assigns colors to vertices one by one, following a specific order. It assigns the smallest available color that hasn't been used by adjacent vertices. Example: For a graph with vertices $V = \{1,2,3,4\}$ and edges $E = \{(1,2),(2,3),(3,4)\}$, the greedy algorithm might color the vertices as follows: 1 (color 1), 2 (color 2), 3 (color 1), 4 (color 2).

Sketch of the Algorithm:

- 1. Sort vertices in a predefined order (e.g., descending degree order).
- 2. Assign the smallest possible color to the first vertex.
- 3. Move to the next vertex and assign the lowest available color that does not conflict with adjacent vertices.

- 4. Repeat until all vertices are colored.
- 3. **DSATUR Algorithm:** This algorithm selects the next vertex to color based on the saturation degree, which is the number of different colors to which it is adjacent.

Sketch of the Algorithm:

- 1. Start with the vertex to the highest degree.
- 2. Assign it the first available color.
- 3. Select the next vertex with the highest **saturation degree**.
- 4. Assign the smallest valid color.
- 5. Repeat until all vertices are colored.