Numerical Introductory Seminar - Prime Factorization Anna Friesen

Motivation: The cryptosystem RSA multiplies two large primes

to create one part of the keys

Example: Trial Division (inefficient for large prime factors)

Modification: Test in Trial Division only primes Sieve of Eratosthenes: Creates list of primes (and zeros)

Code in Mathematica (Sieve of Eratosthenes)

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\begin{split} \mathbf{SoE}[\mathbf{n}] &:= \mathbf{Module}[\{v = \mathrm{Table}[t, \{t, 1, n\}], y = 2, k\}, \\ v = \mathbf{ReplacePart}[v, 0, 1]; \\ \mathbf{While}[y * y \leq n, k = y + y; \\ \mathbf{While}[k \leq n, \\ v = \mathbf{ReplacePart}[v, 0, k]; \\ k = k + y]; \\ y = y + 1]; \\ \mathbf{Return}[\mathbf{v}] \end{split}
```

Pollard (p-1) factorization method

Theoretical base: Fermat's little theorem
Previous algorithm: Pollard rho method

Algorithm (Pollard (p-1) factorization method)

- 1. Input: $n \geq 2$
- 2. choose a with $1 \leq a \leq n-1$ and arbitrary $B \in \mathbb{N}$
- 3. $\forall q \in \mathbb{P}, q \leq B$:
 - $a := a^q \pmod{n}$
 - p := gcd(a-1,n)
 - if $p \mid n$ break
 - ullet else select new a and go to step 3
- 4. return p