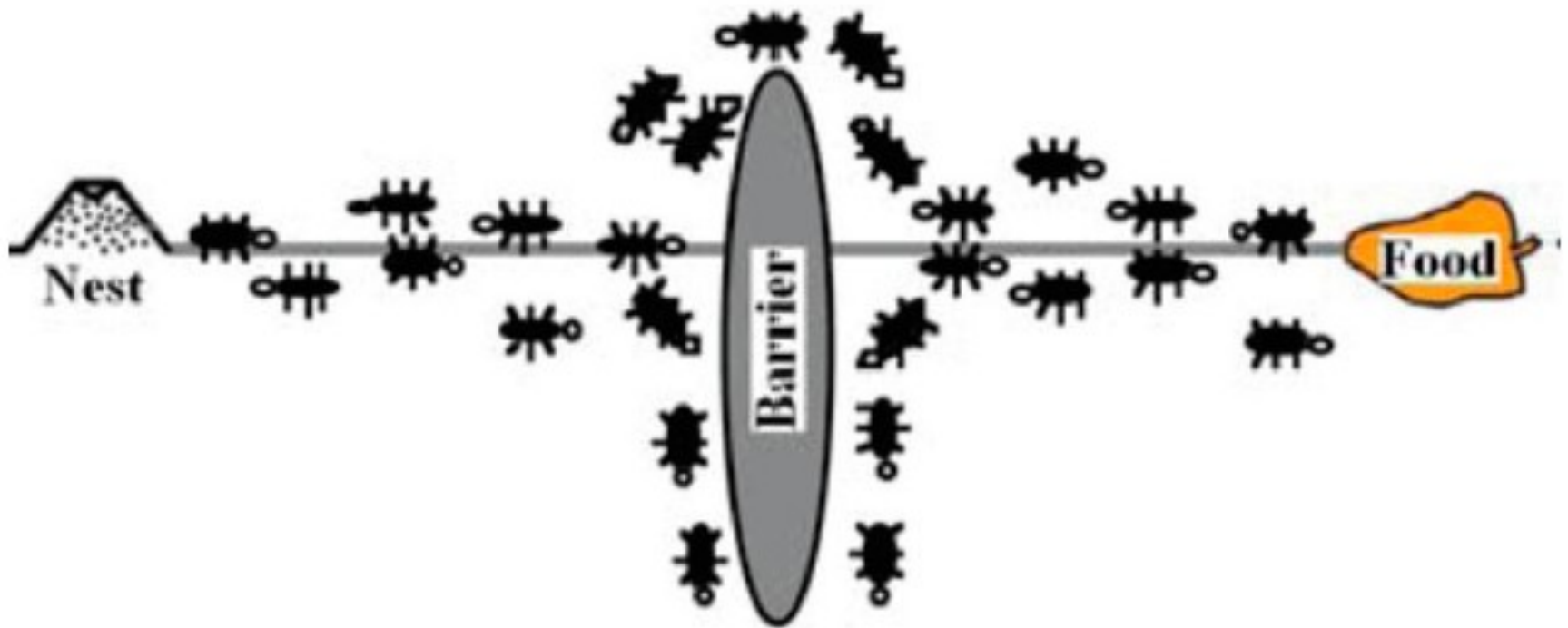


Ant Algorithm





Initialize the pheromone matrix τ for each pair of cities
Place the m ants on n random cities
for $t = 1$ to nc **do**
 for $i = 1$ to n **do**
 for $k = 1$ to m **do**
 Choose next city j according to the transition rule
 for $k = 1$ to m **do**
 Calculate tour distance L_k for ant k
 if an improved tour is found **then**
 Update T^* and L^*
Update the pheromone matrix τ



Initialize T_{Global}^* {this data is shared, everything else is private}

parallel region with $nColonies$ threads

Initialize the pheromone matrix τ for each pair of cities

Place the m ants on n random cities

for $t = 1$ to nc **do**

for $i = 1$ to n **do**

for $k = 1$ to m **do**

 Choose next city j according to the transition rule

for $k = 1$ to m **do**

 Calculate tour distance L_k for ant k

if an improved tour is found **then**

 Update T^* and L^*

if this is an exchange cycle **then**

if $L^* < L_{Global}^*$ **then**

 Critical section

if $L^* < L_{Global}^*$ **then**

$T_{Global}^* = T^*$

 End critical section

 Synchronization barrier

$T^* = T_{Global}^*$

 Update the pheromone matrix τ



visited using a stochastic mechanism. An ant k at city i has not visited set of cities S_p then P_{ij}^k be the probability to visit edge k after edge i .

$$P_{ij}^k = \begin{cases} \frac{\tau_{ij}^\alpha \eta_{ij}^\beta}{\sum_{j \in S_p} \tau_{ij}^\alpha \eta_{ij}^\beta} & \text{if } j \in S_p \\ 0 & \text{otherwise} \end{cases} \quad (1)$$

S_p represents the set of cities which has not been visited yet and to be visited again so that the probability of the ant visiting a city which has already visited becomes 0. Where τ_{ij} is the pheromone content on the edge joining node i to j . η_{ij} represents the heuristic value which is inverse of the distance between the city i to j , which is given by:



$$\eta_{ij} = \frac{1}{d_{ij}}$$

Where d_{ij} is the distance between the city i to j . α and β represents the dependency of probability on the pheromone content or the heuristic value respectively. Increasing the value of α and β may vary the convergence of ACO.



After solution construction we have to update the pheromone accordingly, as follows:

$$\tau_{ij} \leftarrow (1 - \rho) \cdot \tau_{ij} + \sum_{k=1}^m \Delta \tau_{ij}^k$$

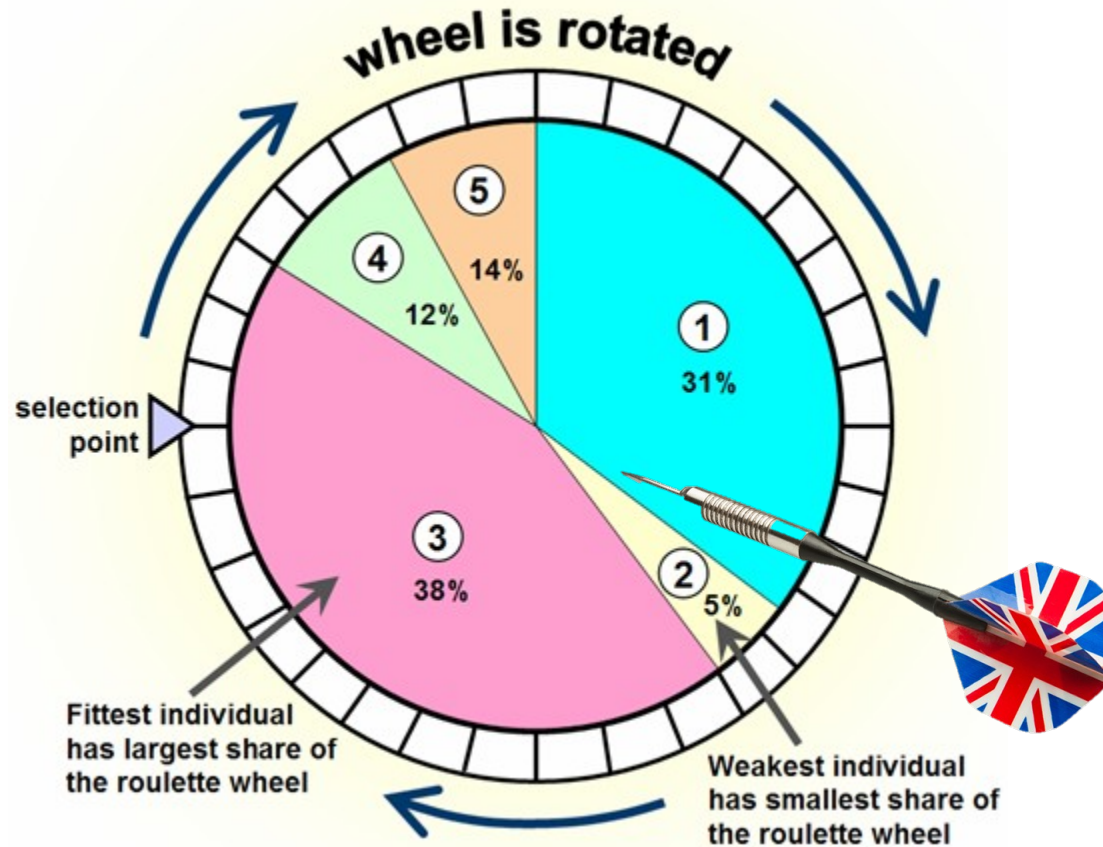
Where ρ is the evaporation rate, m is the number of ants, and $\Delta \tau_{ij}^k$ is the quantity of pheromone laid on edge(i,j) by an ant k :

$$\Delta \tau_{ij}^k = \begin{cases} Q/L_k & \text{if ant } k \text{ uses edge } (i,j) \text{ in its tour,} \\ 0 & \text{otherwise} \end{cases}$$

Where Q is a constant and L_k is the length of the tour constructed by an ant k .



Roulette Wheel Selection



Printing the best tour

```
struct {  
    int cost;  
    int rank;  
} loc_data, global_data;  
  
loc_data.cost = Tour_cost(loc_best_tour);  
loc_data.rank = my_rank;  
  
MPI_Allreduce(&loc_data, &global_data, 1, MPI_2INT, MPI_MINLOC, comm);  
if (global_data.rank == 0) return;  /* 0 already has the best tour */  
if (my_rank == 0)  
    Receive best tour from process global_data.rank;  
else if (my_rank == global_data.rank)  
    Send best tour to process 0;
```



Homework 6

- 使用 MPI+OpenMP 實作，每一台電腦各啟動一個 process，每個 process 再 fork 出 multi-thread

