

Disclaimer:

1. The solution is just for your reference. They may contain some mistakes. DO TRY to solve the problems by yourself. Please also pay attentions to the course website for the updates.
2. Try not to use pseudoinstructions for any exercises that ask you to produce MIPS code. Your goal should be to learn the real MIPS instruction set, and if you are asked to count instructions, your count should reflect the actual instructions that will be executed and not the pseudoinstructions.

2.23 \$t2 = 3

2.24

For jump instruction, it obtains bit 31-28 from PC (which is 0010 in this example). Therefore, it is impossible to jump to 0x40000000.

For more details, please refer https://chortle.ccsu.edu/AssemblyTutorial/Chapter-17/ass17_5.html

For beq instruction, the address range is just 16-bit. Therefore, it is also impossible.

2.26.1 20

2.26.2

while (i>0)

```
{  
    i = i - 1;  
    B += 2;  
}
```

2.26.3

For any value of $N > 0$, (N full loops + 2 lines) of MIPS instructions must be executed; each full loop contains 5 MIPS instructions. Therefore, if the register \$t1 is initialized to the value N , a total of $(5*N)+2$ MIPS instructions are executed. The additional 2 instruction is because of the need to exit the loop

2.27

```
    add $t0, $0, $0      # i = 0  
L1: slt  $t2, $t0, $s0    # i < a  
    beq $t2, $0, Exit     # $t2 == 0, go to Exit  
    add $t1, $0, $0      # j = 0
```

```

L2: slt  $t2, $t1, $s1      # j < b
    beq  $t2, $0, L3        # if $t2 == 0, go to L3
    add  $t2, $t0, $t1      # i+j
    sll  $t4, $t1, 4        # $t4 = 4*j    (sll was written instead of mul.)
    add  $t3, $t4, $s2      # $t3 = &D[4*j]
    sw   $t2, 0($t3)        # D[4*j] = i+j
    addi $t1, $t1, 1        # j = j+1
    j    L2
L3: addi $t0, $t0, 1        # i = i+1
    j    L1
Exit:

```

fib:

```

    addi $sp, $sp, -12      # allocate stack frame of 12 bytes
    sw   $a0, 8($sp)       # save n
    sw   $ra, 4($sp)       # save return address
    sw   $s0, 0($sp)       # save $s0

```

```

    slti  $t0, $a0, 2      # fib(i) = i for i = 0, 1
    beq  $t0, $0, else
    add  $v0, $a0, $0      # $v0 = 0 or 1
    j    exit             # go to exit

```

else:

```

    addi $a0, $a0, -1      # fib(n-1)
    jal  fib               # recursive call
    add  $s0, $v0, $0
    addi $a0, $a0, -1      # fib(n-2)
    jal  fib               # recursive call
    add  $v0, $v0, $s0

```

exit:

```

    lw   $a0, 8($sp)       # restore $a0
    lw   $ra, 4($sp)       # restore return address
    lw   $s0, 0($sp)       # restore $s0
    addi $sp, $sp, 12      # free stack frame
    jr   $ra              # return to caller

```

2.39 Generally, all solutions are similar:

```

lui $t1, top_16_bits
ori $t1, $t1, bottom_16_bits

```

Therefore, the answer is

```

lui $t1, 0x2001
ori $t1, $t1, 0x4924

```

2.40

Address in Exercise 2.39 is 20014924_{16}

No, jump can go up to $0x0FFFFFFC$.

2.41

Address in Exercise 2.39 is 20014924_{16}

No, range is $0x604 + 0x1FFFC = 0x0002\ 0600$ to $0x604 - 0x20000$
 $= 0 \times FFFE\ 0604$.

2.42

Address in Exercise 2.39 is 20014924_{16}

Yes, range is $0x1FFFF004 + 0x1FFFC = 0x2001F000$ to $0x1FFFF004 - 0x20000 = 1FFDF004$