

Day 2

Big O notation and time Complexity
(time to takes to run your function)

Types of complexity:-

* $O(1)$ - constant running time - Best
(linear search)

* $O(\log n)$ - logarithmic running time -
(Binary search) Good

* $O(n)$ - linear running time - Fair

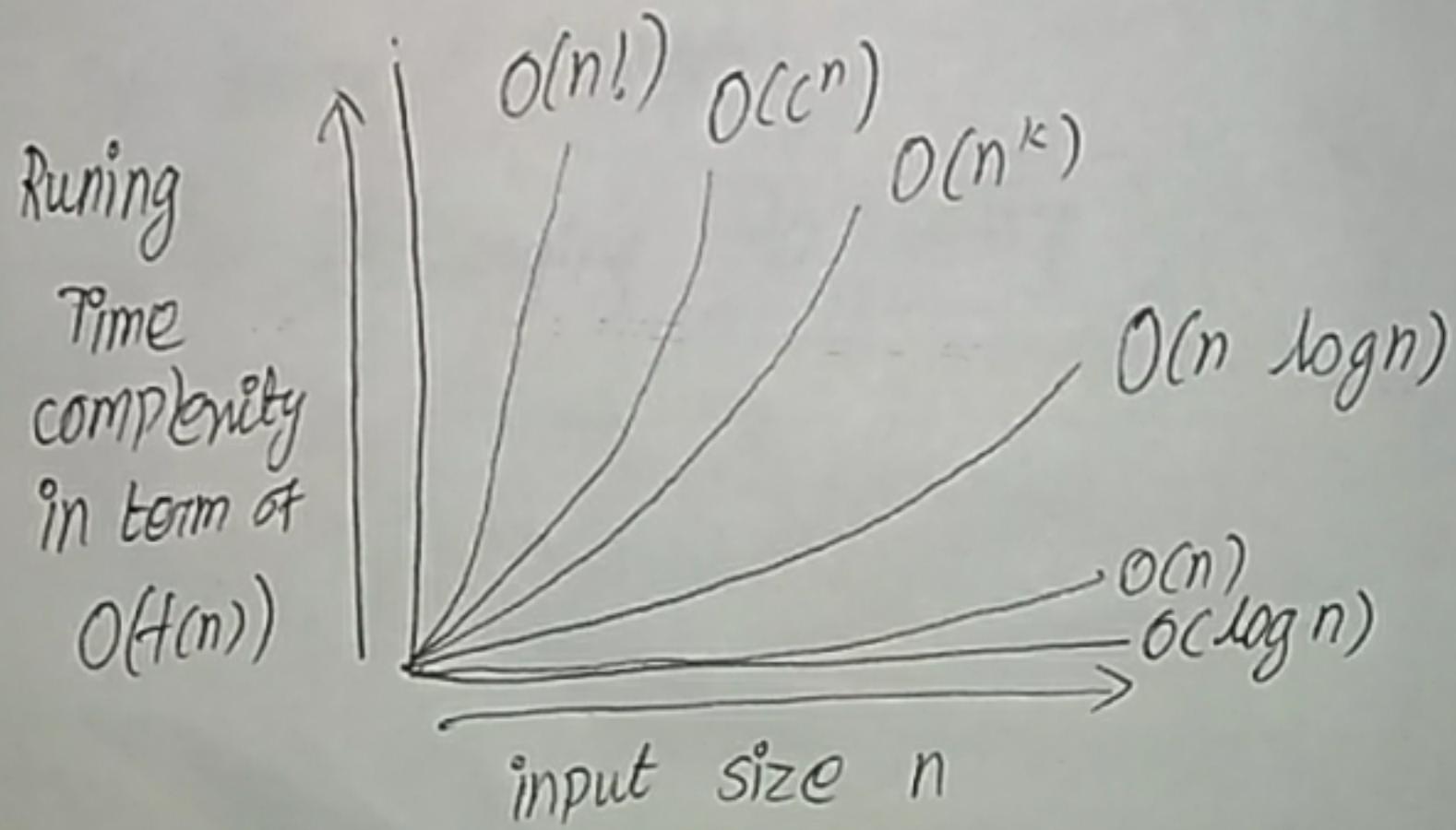
Bad ← * $O(n \log n)$ - log-linear running time
(merge sort)

{ * $O(n^k)$ - polynomial running time

* $O(2^n)$ - Exponential running time
(Bubble sort)
(Tower of Hanoi)

* $O(n!)$ - Factorial running time

(Brute force search algorithm)



* Time complexity : A way of showing how the runtime of a function increases as the size of input increases.

Constant Running Time - $O(1)$

An algorithm is said to have a constant time when it is not dependent on the input data (n). no matter the size of input data , the running time will always be the same.

Example :-

if $a > b$:

 return True

else :

 return False

don't need
to worry about
input data
size

logarithmic Running Time - $O(\log n)$

A logarithmic time complexity when it reduce the size of the input data in each step.

Example :-

```
for index in range(0, len(data), 3):  
    print(data[index])
```

"don't need to look at all
values of the input data"

* logarithmic time complexity are
commonly found in operations on
Binary tree or when using binary Search.

Linear Running Time - $O(n)$

A linear Time Complexity when the
running time increases at most linearly
with the size of the input data.

This is the best possible time complexity
when the algorithm must examine all
values in the input data.

"need to ~~not~~ look
at all values in the list. to
find the value we are looking for"

Example:-

```
for value in data:  
    print(value)
```

* linear time complexity are commonly found in linear search operations.

log-linear Running Time - $O(n \log n)$

* A quasilinear (or) log linear Running time Complexity when each operations in the input data have ~~been~~ a logarithm time complexity.

* It is commonly seen in sorting algorithms. (merge sort, timsort, heapsort)

Example:-

```
for each value in the data1 ( $O(n)$ )  
    use the binary search ( $O(\log n)$ ) to search  
    the same value in data2.
```

```
for value in data1:  
    result.append(binary search(data2, value))
```

Polynomial (or) Quadratic Runing time - O(n^k)

- * Quadratic Runing time complexity when it needs to perform a linear time operation for each value in the input data.
- * "Bubble sort" is a great example of quadratic time complexity since for each value it need to compare to all other values in the list.

Example :-

```
for x in data:
```

```
    for y in data:
```

```
        print(x,y)
```

"It ~~prob~~ operation have ~~a~~ k "for" loop Then time complexity is $O(n^k)$ "

k - "for" loop count

- * If the input size doubles the runtime quadruples.

Exponential Running Time = $O(2^n)$

* An exponential time complexity when the growth doubles with each addition to the input data set. This kind of time complexity is usually seen in brute-force algorithm.

* Using a exponential algorithm to do this, it becomes incredibly resource-expensive to brute-force crack a long password versus a short one. This is one reason that a long password is considered more secure than the shorter one.

Example :-

```
def fibonacci(n):
```

```
    if n <= 1
```

```
        return n
```

```
    return fibonacci(n-1) + fibonacci(n-2)
```

~~* As the input size doubles the runtime~~

* If the input size increases by one,

The runtime doubles

Factorial Runing Time Complexity - $O(n!)$

- * A factorial Running time complexity grows in a factorial way when based on the size of the input data.
- * Heap's algorithm is best example of factorial Running Complexity. which is used for generating all possible permutations of n object.

```
def Heap_permutation(data, n):  
    if n == 1:  
        print(data)  
        return  
  
    for i in range(n):  
        Heap_permutation(data, n-1)  
        if n%2 == 0:  
            data[i], data[n-1] = data[n-1], data[i]  
        else:  
            data[0], data[n-1] = data[n-1], data[0]
```

"it will grow in a factorial way, based on size of the input data"

How to analyzing the time complexity of an algorithm with several operation?

* we need to describe the algorithm based on the largest complexity among all operations.

example:-

def func(data):

 first_ele = data[0] $\rightarrow O(1)$

 for value in data:
 print (value) $\rightarrow O(n)$

 for x in data:
 for y in data:
 print (x,y) $\rightarrow O(n^2)$

* We can see that it has multiple time complexities: $O(1) + O(n) + O(n^2)$.

Based on this, we can describe the time complexity of ^{this} algorithm as worst case ($O(n^2)$).

* If operation has multiple time complexities, then the time complexity is worst case of the algorithm.