

# HOMEWORK -1

①

Subject:- Information theory. (1 credit)

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Degree - PhD in communications & signal processing.

References :- Book by Papoulis named "Random Variables & Stochastic process"

→ lectures - MIT lecture series on probability

Exercise 1.5 1. Show that for all (+)ve integers  $(n, k, l)$  with  $(n > k > l)$ , prove that  $\frac{n-l}{k-l} \geq \frac{n}{k}$

Solution - Starting with the initial condition -

$$n > k > l$$

$$\Rightarrow n > k \quad \text{①} \quad k > l \quad \text{②}$$

From ①

$$n > k$$

Multiplying by  $\frac{1}{e}$  both the sides -

$$\frac{n}{l} > \frac{k}{l} \quad (l \text{ is } (+)\text{ve} \Rightarrow \frac{1}{l} \text{ is also } (+)\text{ve} \text{ integer} \therefore \text{multiplication of } (+)\text{ve integer won't change inequality})$$

$$\Rightarrow \frac{l}{n} < \frac{l}{k} \Rightarrow \frac{l}{n}-1 < \frac{l}{k}-1$$

Subtracting + both the sides

$$-\left(\frac{l}{n}-1\right) > -\left(\frac{l}{k}-1\right)$$

$$\Rightarrow -\left(\frac{l-n}{n}\right) > -\left(\frac{l-k}{k}\right) \quad \left\{ \begin{array}{l} \therefore n > k > l \\ \text{Ex. } n=3, k=2 \& l=1 \\ \frac{1}{3} < \frac{1}{2} \end{array} \right.$$

$$\Rightarrow \frac{n-l}{n} > \frac{k-l}{k} \quad \left. \begin{array}{l} \\ \text{but } -\left(\frac{1}{3}-1\right) > -\left(\frac{1}{2}-1\right) \end{array} \right\}$$

$$\Rightarrow (n-l)k > (k-l)n \quad \left\{ \text{cross multiplication} \right\}$$

$$\Rightarrow \frac{n-l}{k-l} > \frac{n}{k}$$

$$\Rightarrow \boxed{\frac{n}{k} \leq \frac{n-l}{k-l}} \quad \text{QED //}$$

1.5.2 Prove the following bounds for the binomial coefficient : ②

$$\left(\frac{n}{k}\right)^k \leq \binom{n}{k} \leq \left(\frac{n^k}{k!}\right)$$

Proof - As we know -

$$\begin{aligned} \binom{n}{k} &= \frac{n!}{k!(n-k)!} \\ \Rightarrow \binom{n}{k} &= \frac{n(n-1)(n-2)\dots(n-(k-1))(n-k)!}{k!(n-k)!} \\ &= \frac{n(n-1)(n-2)(n-3)\dots(n-(k-1))}{k!} = \prod_{i=0}^{k-1} \frac{n-i}{k-i} \\ \binom{n}{k} &= \frac{\prod_{i=0}^{k-1} \frac{n-i}{k(k-1)\dots(k-(k-1))}}{\prod_{i=0}^{k-1} \frac{n-i}{k-i}} = \prod_{i=0}^{k-1} \frac{\frac{n-i}{k}}{\frac{n-i}{k-i}} \quad \text{--- (1)} \end{aligned}$$

$$\therefore \binom{n}{k} = \frac{n(n-1)(n-2)\dots(n-(k-1))}{k(k-1)(k-2)\dots(k-k+1)} \quad \text{--- (2)}$$

from Above eqn (1)  $\rightarrow$  we can see each factor in the product (numerator as well as denominator) is at least  $\frac{n}{k}$  & there are total  $(k-1+1=k)$  k factors.

$$\therefore \text{lower bound, } \frac{n \times n \dots (k \text{ times})}{k \times k \dots (k \text{ times})} = \left(\frac{n}{k}\right)^k$$

$$\because k^k > k! \Rightarrow \left(\frac{n}{k}\right)^k \leq \frac{n(n-1)(n-2)\dots(n-(k-1))}{k(k-1)(k-2)\dots(k-k+1)}$$

$$\Rightarrow \boxed{\left(\frac{n}{k}\right)^k \leq \binom{n}{k}} \quad [\text{From (2)}] \quad \text{--- (3)}$$

$$\text{for upper bound: } \binom{n}{k} = \prod_{i=0}^{k-1} \frac{n-i}{k-i} = \prod_{i=0}^{k-1} \frac{n-i}{k}$$

$$\begin{aligned} \Rightarrow \binom{n}{k} &= \frac{n(n-1)(n-2)\dots(n-(k-1))}{k!} \\ &\leq \frac{n \cdot n \dots (k \text{ times})}{k!} \end{aligned}$$

$$\therefore \boxed{\binom{n}{k} \leq \frac{n^k}{k!}} \quad \text{--- (4)} \rightarrow \boxed{\left(\frac{n}{k}\right)^k \leq \binom{n}{k} \leq \left(\frac{n^k}{k!}\right)} \quad \text{QED}$$

Q.5.3 Prove that for all integers ( $+ve$  integers),  $n > k > 0$   
 we have  $\frac{\sqrt{2\pi}}{e^2} \frac{1}{\sqrt{np(1-p)}} 2^{nH_2(p)} \leq \binom{n}{k} \leq \frac{e}{2\pi} \frac{1}{\sqrt{np(1-p)}} 2^{nH_2(p)}$

where  $p = k/n$  &  $H_2(p) = -p \log_2 p + -(1-p) \log_2(1-p)$

Stirling's approximation -

$$\sqrt{2\pi} n^{n+1/2} e^{-n} \leq n! \leq e^{n(n+1/2)} e^{-n} \quad \text{--- (1)}$$

Proof -

as we know -

$$\binom{n}{k} = {}^n C_k = \frac{n!}{k!(n-k)!} = \frac{n!}{(np)!(n-np)!},$$

$$\Rightarrow \binom{n}{k} = \frac{n!}{(np)!(n(1-p))!} = \frac{n!}{(np)!(nq)!} \quad \begin{cases} \because k = np \\ \text{let } q = 1-p \end{cases}$$

Using (1) -

$$n! \leq e^{n(n+1/2)} e^{-n}$$

$$\Rightarrow \frac{n!}{(np)!(nq)!} \leq \frac{e^{n(n+1/2)} e^{-n}}{e^{(np)(np+1/2)} e^{-np} e^{(nq)(nq+1/2)} e^{-nq}}$$

$$\Rightarrow \binom{n}{k} \leq \frac{n^n \cdot n^{1/2} e^{-n}}{(np)^{np} \cdot (np)^{1/2} \cdot e^{-np} \cdot e \cdot (nq)^{nq} \cdot (nq)^{1/2} e^{-nq}} \quad \begin{cases} \because a^{(b+c)} = a^b \cdot a^c \end{cases}$$

$$= \frac{n^n \cdot n^{1/2} \cdot e^{-n}}{n^{np} \cdot p^{np} \cdot n^{1/2} \cdot p^{1/2} e^{-np} \cdot e \cdot (nq)^{nq} \cdot (q)^{nq} \cdot n^{1/2} q^{1/2} e^{-nq}} \quad \begin{cases} \because (ab)^c = a^c \cdot b^c \end{cases}$$

$$\binom{n}{k} \leq \frac{n^n \cdot n^{-np} \cdot n^{-nq} \cdot e^{-n} \cdot e^{np} \cdot e^{nq}}{e^{(n^{1/2} p^{1/2} q^{1/2})} (p^{np} \cdot q^{nq})} = \frac{n^{n(1-p-q)} \cdot e^{n(-1+p+q)}}{e \sqrt{npq} (p^{np} \cdot q^{nq})} \quad \text{--- (2)}$$

$$\therefore 1-p-q = 1-p-(1-p) = 1-p-1+p = 0$$

$$\Rightarrow -1+p+q = 0$$

from eqn ② -

④

$$\therefore \binom{n}{k} \leq \frac{n^{n \cdot 0} \cdot e^{n \cdot 0}}{e^{\sqrt{np(1-p)}} p^{np} q^{np}} = \frac{1}{e^{\sqrt{np(1-p)}} p^{np} q^{np}} \quad (4)$$

Now since entropy  $H_2(p)$  is given by -

$$H_2(p) = -[p \log_2 p + [ \log_2 (1-p)] \cdot [1-p]] \quad (\text{given})$$

$$\Rightarrow H_2(p) = -[\log_2 p^p + \log_2 (1-p)^{(1-p)}] \\ = -\log_2 (p^p \cdot (1-p)^{(1-p)}) \quad [\because \log a + \log b = \log ab]$$

$$\Rightarrow H_2(p) = -\log_2 (p^p \cdot q^q)$$

$$\Rightarrow p^p \cdot q^q = 2^{-H_2(p)}$$

taking  $n^{\text{th}}$  power both sides -

$$\Rightarrow p^{np} \cdot q^{nq} = 2^{-nH_2(p)} \quad (3)$$

Now, using equation (3) in eqn (4) -

- (5)

$$\binom{n}{k} \leq \frac{1}{e^{\sqrt{np(1-p)}} 2^{-nH_2(p)}}$$

$$\text{Now, } \frac{1}{e} = 0.3678$$

$$\& \frac{e}{2\pi} = 0.432$$

$$\therefore \frac{1}{e} < \frac{e}{2\pi}$$

$$\Rightarrow \frac{1}{e^{\sqrt{np(1-p)}} 2^{-nH_2(p)}} < \frac{e}{2\pi} \frac{1}{\sqrt{np(1-p)} 2^{-nH_2(p)}}$$

∴ from (5) -

$$\binom{n}{k} \leq \frac{e}{2\pi} \frac{1}{\sqrt{np(1-p)} 2^{-nH_2(p)}}$$

$$= \frac{e}{2\pi \sqrt{np(1-p)}} 2^{nH_2(p)}$$

$$\boxed{\binom{n}{k} \leq \frac{e}{2\pi \sqrt{np(1-p)}} 2^{nH_2(p)}} \quad \begin{array}{l} \text{Upper bound} \\ \text{-(1) proved} \end{array}$$

for lower bound - from lower bound of Stirling's eqn - (5)

$$n! \geq \sqrt{2\pi} n^{n+1/2} e^{-n}$$

$$\begin{aligned} \binom{n}{k} &= \frac{n!}{(np)!(nq)!} \geq \frac{\sqrt{2\pi} n^{n+1/2} e^{-n}}{\sqrt{2\pi} (np)^{nq+1/2} e^{-nq} \sqrt{2\pi} (np)^{np+1/2} e^{-np}} \\ &= \frac{n^n \cdot n^{1/2} e^{-n} \cdot e^{nq} \cdot e^{np}}{\sqrt{2\pi} (nq)^{nq} (nq)^{1/2} (np)^{np} (np)^{1/2}} \\ &= \frac{1}{\sqrt{2\pi}} \frac{n^n \cdot n^{-nq} \cdot n^{-np} \cdot n^{1/2} e^{n(-1+q+p)}}{q^{nq} \cdot p^{np} \cdot n^{1/2} \cdot n^{1/2} \cdot q^{1/2} \cdot p^{1/2}} \\ &= \frac{1}{\sqrt{2\pi}} \frac{1}{n^{1/2} q^{1/2} p^{1/2}} \frac{n^{-n(-1+q+p)} \cdot e^{n(-1+q+p)}}{p^{np} q^{nq}} \end{aligned} \quad (6)$$

$$\therefore p^{np} \cdot q^{nq} = 2^{-nH_2(p)}$$

$$\& -1+q+p = 0 \Rightarrow n^{-n+0} = 1 \& e^{n \times 0} = 1$$

$$\therefore \text{from (6)} - \binom{n}{k} \geq \frac{1}{\sqrt{2\pi}} \frac{1}{\sqrt{npq}} \times \frac{1}{2^{-nH_2(p)}} \quad (7)$$

$$\binom{n}{k} \geq \frac{1}{\sqrt{2\pi}} \frac{1}{\sqrt{npq}} 2^{nH_2(p)}$$

$$\therefore \frac{1}{\sqrt{2\pi}} = 0.398 \quad \& \frac{\sqrt{2\pi}}{e^2} = 0.339$$

$$\Rightarrow \frac{1}{\sqrt{2\pi}} \geq \frac{\sqrt{2\pi}}{e^2} \Rightarrow \frac{1}{\sqrt{2\pi}} \frac{1}{\sqrt{npq}} 2^{nH_2(p)} \geq \frac{\sqrt{2\pi}}{e^2} \frac{1}{\sqrt{npq}} 2^{nH_2(p)}$$

from (a) -

$$\binom{n}{k} \geq \frac{\sqrt{2\pi}}{e^2} \frac{1}{\sqrt{npq}} 2^{nH_2(p)}$$

$$\Rightarrow \boxed{\binom{n}{k} \geq \frac{\sqrt{2\pi}}{e^2} \frac{1}{\sqrt{np(1-p)}} 2^{nH_2(p)}} \quad -(b) \text{ upper bound proved}$$

From (a) & (b) -

$$\boxed{\frac{\sqrt{2\pi}}{e^2} \frac{2^{nH_2(p)}}{\sqrt{np(1-p)}} \leq \binom{n}{k} \leq \frac{e}{2\pi} \frac{1}{\sqrt{np(1-p)}} 2^{nH_2(p)}}$$

QED

Exercise 1.5. 4 use the above to conclude for any (6)

$$0 < p \leq \frac{1}{2}, \lim_{n \rightarrow \infty} \frac{1}{n} \log_2 \left( \begin{matrix} n \\ \lfloor np \rfloor \end{matrix} \right) = H_2(p)$$

$$\text{LHS} = \lim_{n \rightarrow \infty} \frac{1}{n} \log_2 \left( \begin{matrix} n \\ \lfloor np \rfloor \end{matrix} \right)$$

consider  $n =$  integer multiple of  $p$

$$n = mp, \quad m \in \mathbb{Z}^+$$

$$\Rightarrow p = \frac{m}{n} \quad \text{--- (1)}$$

$$\Rightarrow \lfloor np \rfloor = \lfloor n \times \frac{m}{n} \rfloor = \lfloor m \rfloor = m \quad (\because m \in \mathbb{Z}^+)$$

Now consider LHS -

$$\lim_{n \rightarrow \infty} \frac{1}{n} \log_2 \left( \begin{matrix} n \\ \lfloor np \rfloor \end{matrix} \right) \stackrel{2^{nH_2(p)}}{=} \quad \text{--- (2)}$$

$$\leq \lim_{n \rightarrow \infty} \frac{1}{n} \log_2 \left[ \frac{e}{2\pi} \frac{1}{\sqrt{np(1-p)}} \right]$$

$$\left\{ \because \binom{n}{k} \leq \frac{e}{2\pi} \frac{1}{\sqrt{np(1-p)}}^2, k = np \quad \begin{matrix} \text{upper bound} \\ \text{of } \text{at previous} \\ \text{expt.} \end{matrix} \right\}$$

$$\Rightarrow \log_2 \left( \begin{matrix} n \\ np \end{matrix} \right) \leq \log_2 \left[ \frac{e}{2\pi} \frac{1}{\sqrt{np(1-p)}} \right]^{2^{nH_2(p)}}$$

Since log function is monotonically increasing function.

from (2) -

$$\lim_{n \rightarrow \infty} \frac{1}{n} \log_2 \left( \begin{matrix} n \\ \lfloor np \rfloor \end{matrix} \right) \leq \lim_{n \rightarrow \infty} \frac{1}{n} \left[ \log_2 \left( \frac{e}{2\pi} \frac{1}{\sqrt{np(1-p)}} \right) + \log_2 (2^{nH_2(p)}) \right]$$

$$= \lim_{n \rightarrow \infty} \left[ \frac{1}{n} \log_2 \left( \frac{e}{2\pi} \frac{1}{\sqrt{np(1-p)}} \right) + \frac{1}{n} \log_2 (2^{nH_2(p)}) \right].$$

$$= \lim_{n \rightarrow \infty} \left[ \frac{1}{n} \log_2 \left( \frac{e}{2\pi} \frac{1}{\sqrt{np(1-p)}} \right) + \frac{1}{n} \times n H_2(p) \right]$$

$$= \lim_{n \rightarrow \infty} \left[ \underbrace{\frac{1}{n} \log_2 \left( \frac{e}{2\pi} \frac{1}{\sqrt{np(1-p)}} \right)}_{E} + H_2(p) \right]$$

When  $n \rightarrow \infty \Rightarrow E \rightarrow 0$

$$\therefore \lim_{n \rightarrow \infty} \frac{1}{n} \log_2 \left( \begin{matrix} n \\ \lfloor np \rfloor \end{matrix} \right) \rightarrow H_2(p)$$

since  $\lfloor np \rfloor$  is used  
in int leads to sign.

$$\boxed{\lim_{n \rightarrow \infty} \frac{1}{n} \log_2 \left( \begin{matrix} n \\ \lfloor np \rfloor \end{matrix} \right) = H_2(p)}$$

QED

Exercise 1.4 (as the empirical frequency is close to the true probability if the number of samples is large). Given a sequence  $x^n = (x_1, x_2, \dots, x_n)$ , where each  $x_i \in X$ , the relative frequency / Empirical freq. of occurrence of  $a \in X$  is defined as

$$\eta_a(x^n) = \frac{\text{No. of times } a \text{ occurs in } x^n}{n}$$

Consider a sequence of  $n$  i.i.d random variables  $x_1, x_2, \dots, x_n$  over the alphabet  $X$  & having distribution  $P_X$ .

1. prove that for each  $a \in X$  &  $\epsilon > 0$ ,

$$\lim_{n \rightarrow \infty} \Pr [ |\eta_a(x^n) - P_X(a)| > \epsilon ] = 0 \quad \text{--- (1)}$$

Also show that the probability above decays exponentially in  $n$ .

Solution since  $\eta_a(x^n) = \frac{\text{No. of times } a \text{ occurs in } x^n}{\text{length of sequence}} = \frac{k}{n}$  (let)

= empirical frequency

$\therefore (\eta_a(x^n) : a \in X)$  forms a valid probability mass function. &  $P_X(a) = \text{probability of occurrence of } a \in X \text{ in a single trial}$

Now, considering a sequence of  $n$  i.i.d random variables  $x_1, x_2, \dots, x_n$  over alphabet  $X$  & having

probability distribution  $P_X$ .  $\Rightarrow$  it is same as Bernoulli R.V. with  $n$  trials &  $k = \text{no. of successes}$  &  $n = \text{length of sequence}$

No. of times  $a$  occurs in  $x^n = \text{No. of successes}$  &  $a \in X$  —

Now, since  $x_1, x_2, \dots, x_n$  are independent  $\Rightarrow k = \text{No. of times } a \text{ occurs in } x^n$ .

$\therefore \text{No. of "successes" in } n \text{-trials} = k = \text{No. of times } a \text{ occurs in } x^n$ .

$\therefore P_X(a) = p \text{ & } 1-p=q$ . (let) ;  $\epsilon > 0$

considering  $\left| \frac{k}{n} - p \right| > \epsilon$   $\Rightarrow \left( \frac{k}{n} - p \right) > \epsilon \quad \text{--- (1)} \quad \& \quad \left( \frac{k}{n} - p \right) < -\epsilon \quad \text{--- (2)}$

By (1)  $\left( \frac{k}{n} - p \right) > \epsilon$   
 $\Rightarrow (k - np) > n\epsilon$

$$\Rightarrow k - n(p + \varepsilon) > 0$$

(8)

taking exponential both sides -

$$e^{mx} ; m > 0$$

$$\text{for any } (t) \text{ we have } e^{m(k-n(p+\varepsilon))} > e^0 \quad \text{--- (1)} \quad (\because e^{mx} \text{ is increasing for } m > 0)$$

$$\Rightarrow e^{m(k-n(p+\varepsilon))} > 1 \quad \text{--- (1)}$$

$$P\left\{\frac{k}{n} - p > \varepsilon\right\} = P\{k > n\varepsilon + np\} = P\{k > n(p+\varepsilon)\}$$

$$= \sum_{k=\lceil n(p+\varepsilon) \rceil}^n \binom{n}{k} p^k q^{n-k} \quad [\because k \rightarrow \text{No. of success in } n\text{-Bernoulli trials}]$$

$$< \sum_{k=\lceil n(p+\varepsilon) \rceil}^n e^{m(k-n(p+\varepsilon))} \binom{n}{k} p^k q^{n-k} \quad \text{from (1)} \quad \text{--- (2)}$$

~~$$\sum_{k=0}^n e^{m(k-n(p+\varepsilon))} \binom{n}{k} p^k q^{n-k}$$~~

$$= \sum_{k=0}^{\lceil n(p+\varepsilon) \rceil} e^{m(k-n(p+\varepsilon))} \binom{n}{k} p^k q^{n-k}$$

$$+ \sum_{k=0}^{\lceil n(p+\varepsilon) \rceil} e^{m(k-n(p+\varepsilon))} \binom{n}{k} p^k q^{n-k}$$

$$\Rightarrow \sum_{k=0}^{\lceil n(p+\varepsilon) \rceil} e^{m(k-n(p+\varepsilon))} \binom{n}{k} p^k q^{n-k} \leq \sum_{k=0}^{\lceil n(p+\varepsilon) \rceil} e^{m(k-n(p+\varepsilon))} \binom{n}{k} p^k q^{n-k}$$

$$\text{From (2)} - \quad P\left\{\frac{k}{n} - p > \varepsilon\right\} \leq \sum_{k=0}^n e^{m(k-n(p+\varepsilon))} \binom{n}{k} p^k q^{n-k} \quad \text{--- (3)}$$

$$\therefore (x+y)^n = \sum_{r=0}^n nCr x^r y^{n-r} \quad \text{--- (4)}$$

$$\text{From (3)} - \quad P\left\{\frac{k}{n} - p > \varepsilon\right\} \leq \sum_{k=0}^n e^{-mn\varepsilon} \binom{n}{k} (pe^{mq})^k (qe^{-mp})^{n-k}$$

$$= e^{-mn\varepsilon} (pe^{mq} + qe^{-mp})^n \quad \text{[using (4)]} \quad \text{--- (5)}$$

$$\text{Consider} - pe^{mq} + qe^{-mp}$$

$$< p(mq + e^{m^2q^2}) + q(-mp + e^{m^2p^2})$$

$$= pe^{m^2q^2} + qe^{m^2p^2} \quad \left[\because e^x \leq x + e^{x^2} \text{ for } x > 0\right]$$

$$= p e^{M^2 q^2} + q e^{M^2 p^2} \leq e^{M^2} \quad [\because e^x \leq x + e^{x^2} \text{ for } x > 0] \\ \Rightarrow e^{-m n \varepsilon} \left[ p e^{M^2 q^2} + q e^{M^2 p^2} \right]^n \leq e^{M^2 n} e^{-m n \varepsilon} \\ \Rightarrow \text{from eqn (5)} -$$

$$P \left\{ \frac{k}{n} - p > \varepsilon \right\} \leq e^{M^2 n} e^{-m n \varepsilon} \\ \Rightarrow P \left\{ \frac{k}{n} - p > \varepsilon \right\} \leq e^{m^2 - m n \varepsilon}, \varepsilon > 0 \quad (6)$$

let  $g(m) = e^{m^2 - m n \varepsilon}$   
 let's find min. value of  $g(m) \Rightarrow$  min. value of  $m^2 - m n \varepsilon$

$$\Rightarrow f(m) = m^2 - m n \varepsilon$$

$$\frac{df(m)}{dm} = 2m - n \varepsilon = 0 \\ \Rightarrow m = \frac{n \varepsilon}{2}$$

$\frac{d^2 f(m)}{dm^2} = -n \varepsilon < 0 \Rightarrow m = \frac{n \varepsilon}{2}$  corresponds to  
 min. value of  $g(m)$ .

$\Rightarrow$  at  $m = \frac{n \varepsilon}{2}$ ,  $e^{m^2 - m n \varepsilon}$  is minimum at  $m = \frac{n \varepsilon}{2}, \varepsilon > 0$

$\therefore$  from (6) -

$$P \left\{ \frac{k}{n} - p > \varepsilon \right\} \leq e^{\left(\frac{n \varepsilon}{2}\right)^2 - \left(\frac{n \varepsilon}{2}\right) n \varepsilon} \\ = e^{-\frac{n \varepsilon^2}{4}}$$

when  $m = \varepsilon/2$  ( $\sqrt{2} \approx 1.414$  normalized)

$$\Rightarrow P \left\{ \frac{k}{n} - p > \varepsilon \right\} \leq e^{-n \varepsilon^2/4} \quad (7)$$

Similarly,  $P \left\{ \frac{k}{n} - p < -\varepsilon \right\} \leq e^{-n \varepsilon^2/4} \quad (8)$

Adding (7) & (8) -

$$\Rightarrow P \left\{ \left| \frac{k}{n} - p \right| > \varepsilon \right\} \leq 2e^{-n \varepsilon^2/4}$$

$$\Rightarrow P \left\{ \left| \eta_a(x^n) - P_X(a) \right| > \varepsilon \right\} \leq 2e^{-n \varepsilon^2/4} \quad \forall a \in X \quad \& \varepsilon > 0$$

When  $n \rightarrow \infty$   $\Rightarrow P \left\{ \left| \eta_a(x^n) - P_X(a) \right| > \varepsilon \right\} = 0 \quad [\because e^{-\infty} = 0]$   
 this shows probability  $P \left\{ \left| \eta_a(x^n) - P_X(a) \right| > \varepsilon \right\}$  decays exponentially in  $n$ .

 Exercise 1.1.3 - Use this to prove that (16)  
 $\lim_{n \rightarrow \infty} \Pr \left[ |\eta_a(x^n) - P_X(a)| > \varepsilon \text{ for some } a \in X \right] = 0$

This equation does not enforce  $a \in X$ , but any  $a \in X$ .

Let  $k = \text{No. of times for some } a \text{ occurs in } x^n$ ,  $n = \text{length of sequence.}$

$$\Rightarrow \eta_a(x^n) = \frac{k}{n} \quad \text{and} \quad q = 1-p \quad (\text{let})$$

$$\Pr \left[ |\eta_a(x^n) - P_X(a)| > \varepsilon \right] = \Pr \left[ \left| \frac{k}{n} - p \right| > \varepsilon \right]$$

$$= \Pr \left[ \left( \frac{k}{n} - p \right)^2 > \varepsilon^2 \right]$$

$$= \Pr \left[ (k - np)^2 > n^2 \varepsilon^2 \right]$$

$$\stackrel{\text{consider}}{\leq} \Pr \left[ |\eta_a(x^n) - P_X(a)| > \varepsilon \right] = \Pr \left[ (k - np)^2 > n^2 \varepsilon^2 \right]$$

$$(k - np)^2 > n^2 \varepsilon^2$$

$$\Rightarrow \sum_{k=0}^n (k - np)^2 P_n(k)$$

$$> \sum_{k=0}^n n^2 \varepsilon^2 P_n(k)$$

$$\left\{ \begin{array}{l} \text{where } P_n(k) = P_X \left\{ \begin{array}{l} X = \text{No. of times} \\ a \text{ occurs in } x^n \end{array} \right. \right.$$

$$= \binom{n}{k} p^k q^{n-k}$$

$$\left. \begin{array}{l} \text{& } P_X(a) = \text{probability of} \\ \text{occurrence of any } a \in X \text{ in} \\ \text{a single trial} = p \quad (\text{let}) \\ q = 1-p \end{array} \right.$$

$$\Rightarrow \sum_{k=0}^n (k - np)^2 P_n(k) > n^2 \varepsilon^2 \sum_{k=0}^n P_n(k)$$

$$\Rightarrow \sum_{k=0}^n (k - np)^2 P_n(k) > n^2 \varepsilon^2 \quad \left[ \because \sum_{k=0}^n P_n(k) = \text{sum of pmf} = 1 \right]$$

$$\Rightarrow \boxed{\sum_{k=0}^n k^2 P_n(k) - 2np \sum_{k=0}^n k P_n(k) + n^2 p^2 > n^2 \varepsilon^2} \quad (1)$$

$$\text{calculating } \sum_{k=0}^n k P_n(k) = \sum_{k=1}^n k \binom{n}{k} p^k (1-p)^{n-k} \quad [\because k \geq 0]$$

$$= \sum_{k=1}^n k \frac{n!}{(n-k)! k!} p^k (1-p)^{n-k}$$

$$= \sum_{k=1}^n \frac{n!}{(n-k)! (k-1)!} p^k (1-p)^{n-k}$$

$$= np \sum_{k=1}^n \frac{(n-1)!}{(n-k)! (k-1)!} p^{k-1} q^{n-k}$$

$$\hat{=} np \leq \sum_{k=0}^{n-k} \binom{n-1}{k-1} p^{(x-1)} (1-p)^{(n-1)-(x-1)} \quad (1)$$

$$\sum_{k=1}^n k P_n(k) = np [p + 1-p]^{n-1} = np \times 1 = np \quad - (2)$$

$$\therefore \sum_{n=0}^n n x p^x (1-p)^{n-x} = (p+1-p)^n$$

$$\text{or } (x+y)^n = \sum_{x=0}^n \binom{n}{x} x^x y^{n-x}$$

$$\sum_{k=0}^n k^2 P_n(k) = \sum_{k=1}^n k \frac{n!}{(n-k)! (k-1)!} p^k q^{n-k}$$

$$= n^2 p^2 \sum_{k=2}^n \frac{(n-2)!}{(n-k)! (k-2)!} p^{k-2} q^{n-k}$$

$$+ \sum_{k=1}^n \frac{n!}{(n-k)! (k-1)!} p^k q^{n-k}$$

$$= n^2 p^2 [p+1-p]^{n-2} + npq$$

$$\left| \sum_{k=0}^n k^2 P_n(k) = n^2 p^2 + npq \right| \quad - (3)$$

Now putting values from eqn (2) & (3) into eqn (1) -

$$\sum_{k=0}^n (k-np)^2 P_n(k) = \sum_{k=0}^n k^2 P_n(k) - 2np \sum_{k=0}^n k P_n(k) + n^2 p^2$$

$$= n^2 p^2 + npq - 2np \cdot np + n^2 p^2 = npq.$$

$$\sum_{k=0}^n (k-np)^2 P_n(k) = \sum_{|k-np| \leq n\epsilon} (k-np)^2 P_n(k)$$

$$+ \sum_{|k-np| > n\epsilon} (k-np)^2 P_n(k)$$

$$\geq \sum_{|k-np| > n\epsilon} (k-np)^2 P_n(k) > n^2 \epsilon^2 \sum_{|k-np| > n\epsilon} P_n(k)$$

$$= n^2 \epsilon^2 p \{ |k-np| > n\epsilon \}$$

$$\therefore P(|\frac{k}{n} - p| > \epsilon, \text{ for some } \alpha \in X) < \frac{p}{n\epsilon^2}$$

$$\Rightarrow P(|\eta_a(x^n) - P_X(a)| > \epsilon, \text{ for some } a \in X) < \frac{P_X(a)(1-P_X(a))}{n\epsilon^2}$$

for some  $a \in X$  &  $\epsilon > 0$

When  $n \rightarrow \infty$   $P(|\eta_a(x^n) - P_X(a)| > \epsilon, \text{ for some } a \in X) = 0$  QED.

Exercise 1.4.3

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Using program -  $n = 1000$ 

$$\sum_{x \in X} |y_a(x^n) - p_y(a)| = \underline{0.0176667}.$$

Exercise 1.2.0 Consider  $n$  independent tosses of a biased coin which lands heads with probability  $p$ .

1. Find an expression for the probability that the number of heads is exactly  $k$ . forms bernoulli g.g.d.

probability for getting head =  $p$

probability for getting tail =  $1-p$

$$P(x=k) = P(\text{No. of heads is } k) = p(x=k)$$

$$\therefore P(x=k) = \binom{n}{k} p^k (1-p)^{n-k} \quad \text{required expression}$$

Exercise 1.2.2  
find an expression for the probability that the number of heads is atmost  $m$ .

$$P(\text{No. of heads is atmost } m) = p(x \leq m)$$

$$f_x(m) = P(x \leq m) = \sum_{k=0}^m \binom{n}{k} p^k (1-p)^{n-k} \quad \begin{matrix} \text{required} \\ = f_x(m) \end{matrix} \text{ expression}$$

Exercise 1.2.3

$$n = 17, m = 6, p = 0.3829772289$$

$$P(x \leq m) = f_x(m) = \sum_{k=0}^6 \binom{17}{k} p^k (1-p)^{17-k}$$

$$P(x \leq 6) = 0.507590366791487 \quad \begin{matrix} \text{using} \\ \text{computer} \\ \text{program} \end{matrix}$$

Exercise 1.2.4 - finding bounds for the probability

calculated in 1.2.3.

Expectation for bernoulli R.V -

$$\mu = n E[x_1] = n [1 \cdot p + (1-p) \cdot 0] = np = 17 \times 0.3829772289 \approx 6.502$$

Now, from Markov's Inequality -

$$\Pr [X \geq m] \leq \frac{\mu}{m} \Rightarrow \Pr [X > m] \leq \frac{\mu}{m}$$

$$\Rightarrow \Pr [X \leq m] = 1 - \Pr [X > m] \geq 1 - \frac{\mu}{m}$$

$$\Pr [X \leq 6] \geq 1 - \frac{\mu}{m} = 1 - \frac{6}{6} = 0.6836854818833311$$

Probability of getting atmost 6 heads -

$$\boxed{\Pr [X \leq 6] \geq 0.6836854818833311}$$

(lower bound)  $0.50759 > 0.68$   
⇒ satisfies

Using Chernoff bound - for Bernoulli's trial

Cover tail  $\Pr [X \leq (1-\delta)\mu] \leq e^{-\frac{\mu\delta^2}{2}} \quad \forall 0 < \delta < 1$

$$\Pr [X \leq m] = ?$$

$$m = (1-\delta)\mu \Rightarrow \begin{aligned} \mu &= (1-\delta)\mu \\ 6 &= (1-\delta) \times np \\ 6 &= (1-\delta) 6.502 \end{aligned}$$

$$\Rightarrow 6 = 6.502 - 6.502\delta$$

$$\Rightarrow 6.502\delta = 6.502 - 6 \Rightarrow \delta = \frac{6.502 - 6}{6.502} = 0.077$$

$$\therefore \Pr [X \leq 6] \leq e^{-\frac{\mu\delta^2}{2}} = 0.98079993851502398$$

$$\boxed{\Pr [X \leq 6] \leq 0.98079993851502398} \quad (\text{Upper bound})$$

Using Chebysev bound -  $\Rightarrow 0.50759 \leq 0.98079993851502398 \Rightarrow$  satisfies

$$\Pr [X \leq m] = 1 - \Pr [X < m] = 1 - \Pr [X^2 < m^2]$$

$$\Rightarrow 1 - \frac{\mathbb{E}[X^2]}{m^2} = 1 - \Pr [X^2 < m^2] = \Pr [X^2 \geq m^2]$$

$$\Rightarrow 1 - \Pr [X^2 > m^2] = \Pr [|X - \mu| > m - \mu] \quad \left[ \begin{array}{l} |X| < m \\ \Rightarrow X < m \wedge X > -m \\ \Rightarrow X < m \wedge X > \mu \quad (\because X \in \mathbb{Z}) \end{array} \right]$$

[∴ Using Chebysev]

$$\Pr [|X - \mu| > a] \leq \frac{\mathbb{E}[(X - \mu)^2]}{a^2}$$

$$\therefore 1 - 0.1806111 = 0.8193888$$

$$\boxed{\Pr [X \leq 6] \geq 0.8193888}$$

$$0.50759 > 0.8193888$$

Actual probability  $\approx 0.507$   
of satisfying all bounds.

1.2.5 Exercise 1.2.5

(14)

$$n = N = 1002 \text{ & } m = M = 408$$

$$\therefore p = 0.3824772289.$$

$P(X \leq m) = \text{Probability of getting at most } m \text{ heads}$

$$= \sum_{k=0}^m nC_k p^k (1-p)^{n-k}$$

$$= f_X(m) = f_X(408) = \sum_{k=0}^{408} 1002C_{402} (p)^k (1-p)^{1002-k}$$

$$P(X \leq 408) = 0.9492697985 \quad (1) \quad (\text{using program})$$

$$\text{Using Markov's Inequality: } \mu = np = 383.23$$

$$Pr[X \geq m] \leq \frac{\mu}{m} \Rightarrow Pr[X > m] < \frac{\mu}{m}$$

$$\Rightarrow Pr[X \leq m] = 1 - Pr[X > m] \geq 1 - \frac{\mu}{m}$$

$$\Rightarrow P(X \leq 408) \geq 1 - \frac{\mu}{m} = 0.06068092314264717$$

$$\therefore P(X \leq 408) \geq 0.0606809 \quad (\text{lower bound})$$

$$\therefore 0.9492697985 > 0.0606809 \quad (\text{using 1})$$

$\Rightarrow$  It satisfies this bound

Using Chernoff bound for Bernoulli's trial -

$$P(X \leq (1-s)\mu) \leq e^{-\frac{\mu s^2}{2}} \quad \text{if } 0 < s < 1$$

$$P(X \leq m) = ? \\ m = (1-s)\mu \Rightarrow s\mu = \mu - m \Rightarrow s = \frac{\mu - m}{\mu}$$

$$s = 0.646$$

$$\therefore P(X \leq m) \leq e^{-\frac{383.23 \times (0.646)^2}{2}} \\ = 0.4994685112703179$$

$$\therefore P(X \leq 408) \leq 0.4994685112$$

$$\therefore 0.949269 > 0.4994685112 \quad (\text{using})$$

$\Rightarrow$  Chernoff's bound satisfied //

Using Chebyshev bound -

(15)

$$P[X \leq m] = 1 - P[X < m] = 1 - P[X^2 < m^2]$$
$$\cancel{> C \neq \infty} = 1 - P[|X| < m]$$

$$\geq 1 - \frac{E[X^2]}{m^2}$$

$\left\{ \begin{array}{l} \because |X| < m \\ \Rightarrow X < m \text{ & } X > -m \\ \Rightarrow X < m \text{ & } X > 0 \quad (\because X \in \mathbb{Z}^+) \end{array} \right.$

{ Using chebyshev -

$$P[|X - \mu| > a] \leq \frac{E[(X - \mu)^2]}{a^2}$$

$$\therefore P[X \leq 408] \geq 1 - \frac{E[X]}{m^2} \quad [\because E[X^2] = E[X]]$$
$$= 1 - \frac{n p}{m^2}$$
$$= 1 - \frac{363.23}{408 \times 408} = -0.99769$$

$$\Rightarrow P[X \leq 408] \geq -0.99769$$

$\therefore 0.94926 \geq -0.99769$

$\Rightarrow$  This chebyshev bound also satisfies.