

HOMEWORK -1

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Subject:- Information theory. (1 credit)

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Degree - PhD in communications & signal processing.

References :- Book by Papoulis named "Random Variables & Stochastic process"

→ lectures - MIT lecture series on probability

Exercise 1.5 1. Show that for all (+)ve integers (n, k, l) with ($n > k > l$), prove that $\frac{n-l}{k-l} \geq \frac{n}{k}$

Solution - Starting with the initial condition -

$$n > k > l$$

$$\Rightarrow n > k \quad \text{①} \quad k > l \quad \text{②}$$

From ①

$$n > k$$

Multiplying by $\frac{1}{e}$ both the sides -

$$\frac{n}{l} > \frac{k}{l} \quad (\text{ } l \text{ is (+)ve} \Rightarrow \frac{1}{l} \text{ is also (+)ve} \text{ and multiplication of (+)ve integer won't change inequality})$$

$$\Rightarrow \frac{l}{n} < \frac{l}{k} \Rightarrow \frac{l}{n}-1 < \frac{l}{k}-1$$

Subtracting + both the sides

$$-\left(\frac{l}{n}-1\right) > -\left(\frac{l}{k}-1\right)$$

$$\Rightarrow -\left(\frac{l-n}{n}\right) > -\left(\frac{l-k}{k}\right) \quad \left\{ \begin{array}{l} \therefore n > k > l \\ \text{Ex. } n=3, k=2 \& l=1 \\ \frac{1}{3} < \frac{1}{2} \end{array} \right.$$

$$\Rightarrow \frac{n-l}{n} > \frac{k-l}{k} \quad \left. \begin{array}{l} \text{but } -\left(\frac{1}{3}-1\right) > -\left(\frac{1}{2}-1\right) \end{array} \right\}$$

$$\Rightarrow (n-l)k > (k-l)n \quad \left\{ \text{cross multiplication} \right\}$$

$$\Rightarrow \frac{n-l}{k-l} > \frac{n}{k}$$

$$\Rightarrow \boxed{\frac{n}{k} \leq \frac{n-l}{k-l}} \quad \text{QED //}$$

1.5.2 Prove the following bounds for the binomial coefficient : ②

As we know -

$$\left(\frac{n}{k}\right)^k \leq \binom{n}{k} \leq \left(\frac{n^k}{k!}\right)$$

Proof - As we know -

$$n \binom{n}{k} = \binom{n}{k} = \frac{n!}{k!(n-k)!}$$

$$\Rightarrow \binom{n}{k} = \frac{n(n-1)(n-2) \dots (n-(k-1))(n-k)!}{k!(n-k)!}$$

$$= \frac{n(n-1)(n-2)(n-3) \dots (n-(k-1))}{k!} = \prod_{i=0}^{k-1} \frac{n-i}{k-i}$$

$$\binom{n}{k} = \frac{\prod_{i=0}^{k-1} \frac{n-i}{k-(k-i)}}{k(k-1)\dots(k-(k-1))} = \prod_{i=0}^{k-1} \frac{n-i}{k-i} \quad \text{--- (1)}$$

$$\therefore \binom{n}{k} = \frac{n(n-1)(n-2) \dots (n-(k-1))}{k(k-1)(k-2) \dots (k-(k-1))} \quad \text{--- (2)}$$

from Above eqn (1) \rightarrow we can see each factor in the product (numerator as well as denominator) is at least $\frac{n}{k}$ & there are total $(k-1+1=k)$ k factors.

$$\therefore \text{lower bound, } \frac{n \times n \dots (k \text{ times})}{k \times k \dots (k \text{ times})} = \left(\frac{n}{k}\right)^k$$

$$\because k^k > k! \Rightarrow \left(\frac{n}{k}\right)^k \leq \frac{n(n-1)(n-2) \dots (n-(k-1))}{k(k-1)(k-2) \dots (k-(k-1))}$$

$$\Rightarrow \boxed{\left(\frac{n}{k}\right)^k \leq \binom{n}{k}} \quad [\text{From (2)}] \quad \text{--- (3)}$$

$$\text{for upper bound: } \binom{n}{k} = \prod_{i=0}^{k-1} \frac{n-i}{k-i} = \prod_{i=0}^{k-1} \frac{n-i}{k-i}$$

$$\Rightarrow \binom{n}{k} = \frac{n(n-1)(n-2) \dots (n-(k-1))}{k!} \leq \frac{n \cdot n \dots (k \text{ times})}{k!}$$

$$= \frac{n^k}{k!} \quad \text{from (3) \& (1) ---}$$

$$\therefore \boxed{\binom{n}{k} \leq \frac{n^k}{k!}} \quad \text{--- (4)} \rightarrow \boxed{\left(\frac{n}{k}\right)^k \leq \binom{n}{k} \leq \left(\frac{n^k}{k!}\right)} \quad \text{QED}$$

Q.5.3 Prove that for all integers ($+ve$ integers), $n > k > 0$
 we have $\frac{\sqrt{2\pi}}{e^2} \frac{1}{\sqrt{np(1-p)}} 2^{nH_2(p)} \leq \binom{n}{k} \leq \frac{e}{2\pi} \frac{1}{\sqrt{np(1-p)}} 2^{nH_2(p)}$

where $p = k/n$ & $H_2(p) = -p \log_2 p + -(1-p) \log_2(1-p)$

Stirling's approximation -

$$\sqrt{2\pi} n^{n+1/2} e^{-n} \leq n! \leq e^{n(n+1/2)} e^{-n} \quad \text{--- (1)}$$

Proof -

as we know -

$$\binom{n}{k} = {}^n C_k = \frac{n!}{k!(n-k)!} = \frac{n!}{(np)!(n-np)!},$$

$$\Rightarrow \binom{n}{k} = \frac{n!}{(np)!(n(1-p))!} = \frac{n!}{(np)!(nq)!} \quad \begin{cases} \because k = np \\ \text{let } q = 1-p \end{cases}$$

Using (1) -

$$n! \leq e^{n(n+1/2)} e^{-n}$$

$$\Rightarrow \frac{n!}{(np)!(nq)!} \leq \frac{e^{n(n+1/2)} e^{-n}}{e^{(np)(np+1/2)} e^{-np} e^{(nq)(nq+1/2)} e^{-nq}}$$

$$\Rightarrow \binom{n}{k} \leq \frac{n^n \cdot n^{1/2} e^{-n}}{(np)^{np} \cdot (np)^{1/2} \cdot e^{-np} \cdot e \cdot (nq)^{nq} \cdot (nq)^{1/2} e^{-nq}} \quad \begin{cases} \because a^{(b+c)} = a^b \cdot a^c \end{cases}$$

$$= \frac{n^n \cdot n^{1/2} \cdot e^{-n}}{n^{np} \cdot p^{np} \cdot n^{1/2} \cdot p^{1/2} e^{-np} \cdot e \cdot (nq)^{nq} (q)^{nq} n^{1/2} q^{1/2} e^{-nq}} \quad \begin{cases} \because (ab)^c = a^c \cdot b^c \end{cases}$$

$$\binom{n}{k} \leq \frac{n^n \cdot n^{-np} \cdot n^{-nq} \cdot e^{-n} \cdot e^{np} \cdot e^{nq}}{e^{(n^{1/2} p^{1/2} q^{1/2})} (p^{np} \cdot q^{nq})} = \frac{n^{n(1-p-q)} \cdot e^{n(-1+p+q)}}{e \sqrt{npq} (p^{np} \cdot q^{nq})} \quad \text{--- (2)}$$

$$\therefore 1-p-q = 1-p-(1-p) = 1-p-1+p = 0$$

$$\Rightarrow -1+p+q = 0$$

from eqn ② -

④

$$\therefore \binom{n}{k} \leq \frac{n^{n \cdot 0} \cdot e^{n \cdot 0}}{e^{\sqrt{np(1-p)}} p^{np} q^{np}} = \frac{1}{e^{\sqrt{np(1-p)}} p^{np} q^{np}} \quad (4)$$

Now since entropy $H_2(p)$ is given by -

$$H_2(p) = -[p \log_2 p + [\log_2 (1-p)] \cdot [1-p]] \quad (\text{given})$$

$$\Rightarrow H_2(p) = -[\log_2 p^p + \log_2 (1-p)^{1-p}] \\ = -\log_2 (p^p, (1-p)^{1-p}) \quad [\because \log a + \log b = \log ab]$$

$$\Rightarrow H_2(p) = -\log_2 (p^p, q^q)$$

$$\Rightarrow p^p \cdot q^q = 2^{-H_2(p)}$$

taking n^{th} power both sides -

$$\Rightarrow p^{np}, q^{nq} = 2^{-nH_2(p)} \quad (3)$$

Now, using equation (3) in eqn (4) -

- (5)

$$\binom{n}{k} \leq \frac{1}{e^{\sqrt{np(1-p)}} 2^{-nH_2(p)}}$$

$$\text{Now, } \frac{1}{e} = 0.3678$$

$$\& \frac{e}{2^\pi} = 0.432$$

$$\therefore \frac{1}{e} < \frac{e}{2^\pi}$$

$$\Rightarrow \frac{1}{e^{\sqrt{np(1-p)}} 2^{-nH_2(p)}} < \frac{e}{2^\pi} \frac{1}{\sqrt{np(1-p)} 2^{-nH_2(p)}}$$

∴ from (5) -

$$\binom{n}{k} \leq \frac{e}{2^\pi} \frac{1}{\sqrt{np(1-p)} 2^{-nH_2(p)}}$$

$$= \frac{e}{2^\pi \sqrt{np(1-p)}} 2^{nH_2(p)}$$

$$\boxed{\binom{n}{k} \leq \frac{e}{2^\pi \sqrt{np(1-p)}} 2^{nH_2(p)}} \quad \begin{array}{l} \text{Upper bound} \\ -(1) \text{ proved} \end{array}$$

for lower bound - from lower bound of Stirling's eqn - (5)

$$n! \geq \sqrt{2\pi} n^{n+1/2} e^{-n}$$

$$\begin{aligned} \binom{n}{k} &= \frac{n!}{(np)!(nq)!} \geq \frac{\sqrt{2\pi} n^{n+1/2} e^{-n}}{\sqrt{2\pi} (np)^{nq+1/2} e^{-nq} \sqrt{2\pi} (np)^{np+1/2} e^{-np}} \\ &= \frac{n^n \cdot n^{1/2} e^{-n} \cdot e^{nq} \cdot e^{np}}{\sqrt{2\pi} (nq)^{nq} (nq)^{1/2} (np)^{np} (np)^{1/2}} \\ &= \frac{1}{\sqrt{2\pi}} \frac{n^n \cdot n^{-nq} \cdot n^{-np} \cdot n^{1/2} e^{n(-1+q+p)}}{q^{nq} \cdot p^{np} \cdot n^{1/2} \cdot n^{1/2} \cdot q^{1/2} \cdot p^{1/2}} \\ &= \frac{1}{\sqrt{2\pi}} \frac{1}{n^{1/2} q^{1/2} p^{1/2}} \frac{n^{-n(-1+q+p)} \cdot e^{n(-1+q+p)}}{p^{np} q^{nq}} \end{aligned} \quad (6)$$

$$\therefore p^{np} \cdot q^{nq} = 2^{-nH_2(p)}$$

$$\& -1+q+p = 0 \Rightarrow n^{-n+0} = 1 \& e^{n \times 0} = 1$$

$$\therefore \text{from (6)} - \binom{n}{k} \geq \frac{1}{\sqrt{2\pi}} \frac{1}{\sqrt{npq}} \times \frac{1}{2^{-nH_2(p)}} \quad (7)$$

$$\binom{n}{k} \geq \frac{1}{\sqrt{2\pi}} \frac{1}{\sqrt{npq}} 2^{nH_2(p)}$$

$$\therefore \frac{1}{\sqrt{2\pi}} = 0.398 \quad \& \frac{\sqrt{2\pi}}{e^2} = 0.339$$

$$\Rightarrow \frac{1}{\sqrt{2\pi}} \geq \frac{\sqrt{2\pi}}{e^2} \Rightarrow \frac{1}{\sqrt{2\pi}} \frac{1}{\sqrt{npq}} 2^{nH_2(p)} \geq \frac{\sqrt{2\pi}}{e^2} \frac{1}{\sqrt{npq}} 2^{nH_2(p)}$$

from (a) -

$$\binom{n}{k} \geq \frac{\sqrt{2\pi}}{e^2} \frac{1}{\sqrt{npq}} 2^{nH_2(p)}$$

$$\Rightarrow \boxed{\binom{n}{k} \geq \frac{\sqrt{2\pi}}{e^2} \frac{1}{\sqrt{np(1-p)}} 2^{nH_2(p)}} \quad -(b) \text{ upper bound proved}$$

From (a) & (b) -

$$\boxed{\frac{\sqrt{2\pi}}{e^2} \frac{2^{nH_2(p)}}{\sqrt{np(1-p)}} \leq \binom{n}{k} \leq \frac{e}{2\pi} \frac{1}{\sqrt{np(1-p)}} 2^{nH_2(p)}}$$

QED

Exercise 1.5. 4 use the above to conclude for any (6)

$$0 < p \leq \frac{1}{2}, \lim_{n \rightarrow \infty} \frac{1}{n} \log_2 (\lfloor np \rfloor) = H_2(p)$$

$$\text{LHS} = \lim_{n \rightarrow \infty} \frac{1}{n} \log_2 \left(\frac{n}{\lfloor np \rfloor} \right)$$

consider $n = \text{integer multiple of } p$

$$n = mp, \quad m \in \mathbb{Z}^+$$

$$\Rightarrow p = \frac{m}{n} \quad \text{--- (1)}$$

$$\Rightarrow \lfloor np \rfloor = \lfloor n \times \frac{m}{n} \rfloor = \lfloor m \rfloor = m \quad (\because m \in \mathbb{Z}^+)$$

Now consider LHS -

$$\lim_{n \rightarrow \infty} \frac{1}{n} \log_2 \left(\frac{n}{\lfloor np \rfloor} \right) \xrightarrow{\text{for } 0 < p \leq \frac{1}{2}} 2^{nH_2(p)} \quad \text{--- (2)}$$

$$\leq \lim_{n \rightarrow \infty} \frac{1}{n} \log_2 \left[\frac{e}{2^n} \frac{1}{\sqrt{np(1-p)}} \right]$$

$$\left\{ \because \binom{n}{k} \leq \frac{e}{2^n} \frac{1}{\sqrt{np(1-p)}} \right. ; \quad k = np \quad \begin{matrix} \text{(Upper bound} \\ \text{of previous} \\ \text{step)} \end{matrix}$$

$$\Rightarrow \log_2 \left(\frac{n}{\lfloor np \rfloor} \right) \leq \log_2 \left[\frac{e}{2^n} \frac{1}{\sqrt{np(1-p)}} \right] \quad \text{--- (3)}$$

since log function is monotonically increasing function.

from (2) -

$$\lim_{n \rightarrow \infty} \frac{1}{n} \log_2 \left(\frac{n}{\lfloor np \rfloor} \right) \leq \lim_{n \rightarrow \infty} \frac{1}{n} \left[\log_2 \left(\frac{e}{2^n} \frac{1}{\sqrt{np(1-p)}} \right) + \log_2 (2^{nH_2(p)}) \right]$$

$$= \lim_{n \rightarrow \infty} \left[\frac{1}{n} \log_2 \left(\frac{e}{2^n} \frac{1}{\sqrt{np(1-p)}} \right) + \frac{1}{n} \log_2 (2^{nH_2(p)}) \right]$$

$$= \lim_{n \rightarrow \infty} \left[\frac{1}{n} \log_2 \left(\frac{e}{2^n} \frac{1}{\sqrt{np(1-p)}} \right) + \frac{1}{n} \times nH_2(p) \right]$$

$$= \lim_{n \rightarrow \infty} \left[\underbrace{\frac{1}{n} \log_2 \left(\frac{e}{2^n} \frac{1}{\sqrt{np(1-p)}} \right)}_{E} + H_2(p) \right]$$

when $n \rightarrow \infty \Rightarrow E \rightarrow 0$

$$\therefore \lim_{n \rightarrow \infty} \frac{1}{n} \log_2 \left(\frac{n}{\lfloor np \rfloor} \right) \rightarrow H_2(p)$$

$$\boxed{\lim_{n \rightarrow \infty} \frac{1}{n} \log_2 \left(\frac{n}{\lfloor np \rfloor} \right) = H_2(p)}$$

since $\lfloor np \rfloor$ is used
limit leads to sign.

QED

Exercise 1.4 (as the empirical frequency is close to the true probability if the number of samples is large). Given a sequence $x^n = (x_1, x_2, \dots, x_n)$, where each $x_i \in X$, the relative frequency / Empirical freq. of occurrence of $a \in X$ is defined as

$$\eta_a(x^n) = \frac{\text{No. of times } a \text{ occurs in } x^n}{n}$$

Consider a sequence of n i.i.d random variables x_1, x_2, \dots, x_n over the alphabet X & having distribution P_X .

1. prove that for each $a \in X$ & $\epsilon > 0$,

$$\lim_{n \rightarrow \infty} \Pr [|\eta_a(x^n) - P_X(a)| > \epsilon] = 0 \quad \text{--- (1)}$$

Also show that the probability above decays exponentially in n .

Solution since $\eta_a(x^n) = \frac{\text{No. of times } a \text{ occurs in } x^n}{\text{length of sequence}} = \frac{k}{n}$ (let)

= empirical frequency

$\therefore (\eta_a(x^n) : a \in X)$ forms a valid probability mass function. & $P_X(a) = \text{probability of occurrence of } a \in X \text{ in a single trial}$

Now, considering a sequence of n i.i.d random variables x_1, x_2, \dots, x_n over alphabet X & having

probability distribution P_X . \Rightarrow it is same as Bernoulli R.V. with n trials & $k = \text{no. of successes}$ & $n = \text{length of sequence}$

No. of times a occurs in $x^n = \text{No. of successes}$ & $a \in X$ —

Now, since x_1, x_2, \dots, x_n are independent $\Rightarrow k = \text{No. of times } a \text{ occurs in } x^n$.

$\therefore \text{No. of "successes" in } n \text{-trials} = k = \text{No. of times } a \text{ occurs in } x^n$.

$\therefore P_X(a) = p \text{ & } 1-p=q$. (let) ; $\epsilon > 0$

considering $\left| \frac{k}{n} - p \right| > \epsilon$ $\Rightarrow \left(\frac{k}{n} - p \right) > \epsilon \quad \text{--- (1)} \quad \& \quad \left(\frac{k}{n} - p \right) < -\epsilon \quad \text{--- (2)}$

By (1) $\left(\frac{k}{n} - p \right) > \epsilon$
 $\Rightarrow (k - np) > n\epsilon$

$$\Rightarrow k - n(p + \varepsilon) > 0$$

(8)

taking exponential both sides -

$$e^{mx} ; m > 0$$

$$\text{for any } (t) \text{ we have } e^{m(k-n(p+\varepsilon))} > e^0 \quad \text{--- (1)} \quad (\because e^{mx} \text{ is increasing for } m > 0)$$

$$\Rightarrow e^{m(k-n(p+\varepsilon))} > 1 \quad \text{--- (1)}$$

$$P\left\{\frac{k}{n} - p > \varepsilon\right\} = P\{k > n\varepsilon + np\} = P\{k > n(p+\varepsilon)\}$$

$$= \sum_{k=\lceil n(p+\varepsilon) \rceil}^n \binom{n}{k} p^k q^{n-k} \quad [\because k \rightarrow \text{No. of success in } n\text{-Bernoulli trials}]$$

$$< \sum_{k=\lceil n(p+\varepsilon) \rceil}^n e^{m(k-n(p+\varepsilon))} \binom{n}{k} p^k q^{n-k} \quad \text{from (1)} \quad \text{--- (2)}$$

~~$$\sum_{k=0}^n e^{m(k-n(p+\varepsilon))} \binom{n}{k} p^k q^{n-k}$$~~

$$= \sum_{k=0}^{\lceil n(p+\varepsilon) \rceil} e^{m(k-n(p+\varepsilon))} \binom{n}{k} p^k q^{n-k}$$

$$+ \sum_{k=0}^{\lceil n(p+\varepsilon) \rceil} e^{m(k-n(p+\varepsilon))} \binom{n}{k} p^k q^{n-k}$$

$$\Rightarrow \sum_{k=0}^{\lceil n(p+\varepsilon) \rceil} e^{m(k-n(p+\varepsilon))} \binom{n}{k} p^k q^{n-k} \leq \sum_{k=0}^{\lceil n(p+\varepsilon) \rceil} e^{m(k-n(p+\varepsilon))} \binom{n}{k} p^k q^{n-k}$$

$$\text{From (2)} - \quad P\left\{\frac{k}{n} - p > \varepsilon\right\} \leq \sum_{k=0}^n e^{m(k-n(p+\varepsilon))} \binom{n}{k} p^k q^{n-k} \quad \text{--- (3)}$$

$$\therefore (x+y)^n = \sum_{r=0}^n nCr x^r y^{n-r} \quad \text{--- (4)}$$

$$\text{From (3)} - \quad P\left\{\frac{k}{n} - p > \varepsilon\right\} \leq \sum_{k=0}^n e^{-mn\varepsilon} \binom{n}{k} (pe^{mq})^k (qe^{-mp})^{n-k}$$

$$= e^{-mn\varepsilon} (pe^{mq} + qe^{-mp})^n \quad \text{[using (4)]} \quad \text{--- (5)}$$

$$\text{Consider} - pe^{mq} + qe^{-mp}$$

$$< p(mq + e^{m^2q^2}) + q(-mp + e^{m^2p^2})$$

$$= pe^{m^2q^2} + qe^{m^2p^2} \quad \left[\because e^x \leq x + e^{x^2} \text{ for } x > 0\right]$$

$$= p e^{M^2 q^2} + q e^{M^2 p^2} \leq e^{M^2} \quad [\because e^x \leq x + e^{x^2} \text{ for } x > 0] \\ \Rightarrow e^{-mn\varepsilon} \left[p e^{M^2 q^2} + q e^{M^2 p^2} \right]^n \leq e^{M^2 n} e^{-mn\varepsilon}$$

\Rightarrow from eqn (5) —

$$P\left\{\frac{k}{n} - p > \varepsilon\right\} \leq e^{M^2 n} e^{-mn\varepsilon} \\ \Rightarrow P\left\{\frac{k}{n} - p > \varepsilon\right\} \leq e^{m^2 - mn\varepsilon}, \varepsilon > 0 \quad (6)$$

let $g(m) = e^{m^2 - mn\varepsilon}$

let's find min. value of $g(m) \Rightarrow$ min. value of $m^2 - mn\varepsilon$

$$\Rightarrow f(m) = m^2 - mn\varepsilon$$

$$\frac{df(m)}{dm} = 2m - n\varepsilon = 0 \\ \Rightarrow m = \frac{n\varepsilon}{2}$$

$\frac{d^2 f(m)}{dm^2} = -n\varepsilon < 0 \Rightarrow m = \frac{n\varepsilon}{2}$ corresponds to min. value of $g(m)$.

\Rightarrow at $m = \frac{n\varepsilon}{2}$, $e^{m^2 - mn\varepsilon}$ is minimum at $m = \frac{n\varepsilon}{2}, \varepsilon > 0$

\therefore from (6) —

$$P\left\{\frac{k}{n} - p > \varepsilon\right\} \leq e^{\left(\frac{n\varepsilon}{2}\right)^2 - \left(\frac{n\varepsilon}{2}\right)n\varepsilon} \\ = e^{-\frac{n\varepsilon^2}{4}}$$

when $m = \varepsilon/2$ ($\sqrt{mn} \Rightarrow$ normalized)

$$\Rightarrow P\left\{\frac{k}{n} - p > \varepsilon\right\} \leq e^{-n\varepsilon^2/4} \quad (7)$$

Similarly, $P\left\{\frac{k}{n} - p < -\varepsilon\right\} \leq e^{-n\varepsilon^2/4} \quad (8)$

Adding (7) & (8) —

$$\Rightarrow P\left\{\left|\frac{k}{n} - p\right| > \varepsilon\right\} \leq 2e^{-n\varepsilon^2/4}$$

$$\Rightarrow P\left\{\left|n_a(x^n) - P_X(a)\right| > \varepsilon\right\} \leq 2e^{-n\varepsilon^2/4} \quad \forall a \in X \text{ and } \varepsilon > 0$$

when $n \rightarrow \infty$ $\Rightarrow P\left\{\left|n_a(x^n) - P_X(a)\right| > \varepsilon\right\} = 0 \quad [\because e^{-\infty} = 0]$

this shows probability $P\left\{\left|n_a(x^n) - P_X(a)\right| > \varepsilon\right\}$ decays exponentially in n .

 Exercise 1.1.3 - Use this to prove that (16)
 $\lim_{n \rightarrow \infty} \Pr \left[|\eta_a(x^n) - P_X(a)| > \varepsilon \text{ for some } a \in X \right] = 0$

This equation does not enforce $a \in X$, but any $a \in X$.

Let $k = \text{No. of times for some } a \text{ occurs in } x^n$, $n = \text{length of sequence.}$

$$\Rightarrow \eta_a(x^n) = \frac{k}{n} \quad \text{and} \quad q = 1-p \quad (\text{let})$$

$$\Pr \left[|\eta_a(x^n) - P_X(a)| > \varepsilon \right] = \Pr \left[\left| \frac{k}{n} - p \right| > \varepsilon \right]$$

$$= \Pr \left[\left(\frac{k}{n} - p \right)^2 > \varepsilon^2 \right]$$

$$= \Pr \left[(k - np)^2 > n^2 \varepsilon^2 \right]$$

$$\stackrel{\text{consider}}{\leq} \Pr \left[|\eta_a(x^n) - P_X(a)| > \varepsilon \right] = \Pr \left[(k - np)^2 > n^2 \varepsilon^2 \right]$$

$$(k - np)^2 > n^2 \varepsilon^2$$

$$\Rightarrow \sum_{k=0}^n (k - np)^2 P_n(k)$$

$$> \sum_{k=0}^n n^2 \varepsilon^2 P_n(k)$$

$$\left\{ \begin{array}{l} \text{where } P_n(k) = P_X \left\{ \begin{array}{l} X = \text{No. of times} \\ a \text{ occurs in } x^n \end{array} \right. \right.$$

$$= \binom{n}{k} p^k q^{n-k}$$

& $P_X(a) = \text{probability of occurrence of any } a \in X \text{ in a single trial} = p \quad (\text{let})$

$$q = 1-p$$

$$\Rightarrow \sum_{k=0}^n (k - np)^2 P_n(k) > n^2 \varepsilon^2 \sum_{k=0}^n P_n(k)$$

$$\Rightarrow \sum_{k=0}^n (k - np)^2 P_n(k) > n^2 \varepsilon^2 \quad \left[\because \sum_{k=0}^n P_n(k) = \text{sum of pmf} = 1 \right]$$

$$\Rightarrow \boxed{\sum_{k=0}^n k^2 P_n(k) - 2np \sum_{k=0}^n k P_n(k) + n^2 p^2 > n^2 \varepsilon^2} \quad (1)$$

$$\text{calculating } \sum_{k=0}^n k P_n(k) = \sum_{k=1}^n k \binom{n}{k} p^k (1-p)^{n-k} \quad [\because k \geq 0]$$

$$= \sum_{k=1}^n k \frac{n!}{(n-k)! k!} p^k (1-p)^{n-k}$$

$$= \sum_{k=1}^n \frac{n!}{(n-k)! (k-1)!} p^k (1-p)^{n-k}$$

$$= np \sum_{k=1}^n \frac{(n-1)!}{(n-k)! (k-1)!} p^{k-1} q^{n-k}$$

$$\hat{=} np \leq \sum_{k=0}^{n-k} \binom{n-1}{k-1} p^{(x-1)} (1-p)^{(n-1)-(x-1)} \quad (1)$$

$$\sum_{k=1}^n k P_n(k) = np [p + 1-p]^{n-1} = np \times 1 = np \quad - (2)$$

$$\therefore \sum_{n=0}^n n x p^x (1-p)^{n-x} = (p+1-p)^n$$

$$\text{or } (x+y)^n = \sum_{x=0}^n \binom{n}{x} x^x y^{n-x}$$

$$\sum_{k=0}^n k^2 P_n(k) = \sum_{k=1}^n k \frac{n!}{(n-k)! (k-1)!} p^k q^{n-k}$$

$$= n^2 p^2 \sum_{k=2}^n \frac{(n-2)!}{(n-k)! (k-2)!} p^{k-2} q^{n-k}$$

$$+ \sum_{k=1}^n \frac{n!}{(n-k)! (k-1)!} p^k q^{n-k}$$

$$= n^2 p^2 [p+1-p]^{n-2} + npq$$

$$\left| \sum_{k=0}^n k^2 P_n(k) = n^2 p^2 + npq \right| \quad - (3)$$

Now putting values from eqn (2) & (3) into eqn (1) -

$$\sum_{k=0}^n (k-np)^2 P_n(k) = \sum_{k=0}^n k^2 P_n(k) - 2np \sum_{k=0}^n k P_n(k) + n^2 p^2$$

$$= n^2 p^2 + npq - 2np \cdot np + n^2 p^2 = npq.$$

$$\sum_{k=0}^n (k-np)^2 P_n(k) = \sum_{|k-np| \leq n\epsilon} (k-np)^2 P_n(k)$$

$$+ \sum_{|k-np| > n\epsilon} (k-np)^2 P_n(k)$$

$$\geq \sum_{|k-np| > n\epsilon} (k-np)^2 P_n(k) > n^2 \epsilon^2 \sum_{|k-np| > n\epsilon} P_n(k)$$

$$= n^2 \epsilon^2 p \{ |k-np| > n\epsilon \}$$

$$\therefore P(|\frac{k}{n} - p| > \epsilon, \text{ for some } \alpha \in X) < \frac{p}{n\epsilon^2}$$

$$\Rightarrow P(|\eta_a(x^n) - P_X(a)| > \epsilon, \text{ for some } a \in X) < \frac{P_X(a)(1-P_X(a))}{n\epsilon^2}$$

for some $a \in X$ & $\epsilon > 0$

When $n \rightarrow \infty$ $P(|\eta_a(x^n) - P_X(a)| > \epsilon, \text{ for some } a \in X) = 0$ QED.

Exercise 1.4.3

(12)

Using program - $n = 10000$

$$\sum_{a \in X} |\eta_a(x^n) - p_a(a)| = 0.0176667\%$$

for $\forall a$ $a \in X$ $\sum_{a \in X} |\eta_a(x^n) - p_a(a)| = 0.0033333 [x \in \{\frac{1}{3}, \frac{2}{3}\}]$

Exercise 1.2.0 Consider n independent tosses of a biased coin which lands heads with probability p .

- Find an expression for the probability that the number of heads is exactly k . forms bernoulli q. q. d.

Probability for getting head = p

Probability for getting tail = $1-p$

$$P(X=k) = P(\text{No. of heads is } k) = p(x=k)$$

$$\therefore P(X=k) = \sum_{k=0}^n nC_k p^k (1-p)^{n-k} \quad \text{required expression}$$

Exercise 1.2.2
Find an expression for the probability that the number of heads is atmost m .

$$P_r(\text{No. of heads is atmost } m) = P(x \leq m)$$

$$f_{X(m)} = \sum_{k=0}^m nC_k p^k (1-p)^{n-k} \quad \text{= required } f_X(m) \text{ expression}$$

Exercise 1.2.3

$$n = 17, m = 6, p = 0.3824772289$$

$$P(X \leq m) = f_X(m) = \sum_{k=0}^6 17C_k \cdot p^k \cdot (1-p)^{17-k}$$

$$P(X \leq 6) = 0.507590366791487$$

using computer program

Exercise 1.2.4 - finding bounds for the probability

calculated in 1.2.3.

Expectation for bernoulli R.V -

$$\mu = nE[x_i] = n[1 \cdot p + (1-p) \cdot 0] = np \\ = 17 \times 0.3824772289 \\ \approx 6.502$$

Now from Markov's inequality :-

(3)

$$\Pr[X \geq m] \leq \frac{\mu}{m+1} \Rightarrow \cancel{\Pr[X \geq m]} \cdot \Pr[X \geq m+1] \leq \frac{\mu}{m+1}$$

$$\Pr[X \leq m] = 1 - \Pr[X \geq m+1]$$

$$\Pr[X \leq m] \geq 1 - \frac{\mu}{m+1} = 1 - \frac{np}{m+1}$$

$$\Pr[X \leq 6] \geq 1 - \frac{6 \cdot 502}{7} = 0.0711267298$$

$$\therefore \boxed{\Pr[X \leq 6] \geq 0.07112} \text{ (lower bound)}$$

∴ $0.50759 > 0.07112 \Rightarrow$ Markov inequality satisfied.

Using Chernoff bound for binomial's trial -

$$\Pr[X \leq (1-s) \mu] \leq e^{-\frac{\mu s^2}{2}} \quad \text{if } 0 < s < 1$$

$$\text{where } X = \sum_{i=1}^n y_i$$

$$\Pr[X \leq m] = ?$$

$$m = (1-s)\mu \Rightarrow f = (1-s)np$$

$$\Rightarrow f = 6 \cdot 502 - 6 \cdot 502s$$

$$\Rightarrow s = \frac{0.502}{6.502} = 0.077$$

$$\therefore \Pr[X \leq 6] \leq e^{-\frac{\mu s^2}{2}} = 0.980111815$$

$$\boxed{\Pr[X \leq 6] \leq 0.98111815}$$

$$0.50759 < 0.98111815$$

Chernoff bound satisfies.

$$\text{Chebyshev bound} \quad \Pr[|X-\mu| \geq a] \leq \frac{\text{Var}(X)}{a^2} = \frac{E[(X-\mu)^2]}{a^2}$$

$$\Pr[X \leq m] = 1 - \Pr[X \geq m+1] = 1 - \Pr[X \geq \mu + a^2]$$

$$= 1 - \frac{E[X^2]}{(m+1)^2} = 1 - \frac{\mu}{7^2} = 1 - \frac{6.502}{49} \quad (14)$$

$$\begin{aligned} P(X \geq m+1) &= P(|X-np| \geq m+1-np) \\ &\leq \frac{\text{Var}(X)}{(m+1-np)} \quad \left. \begin{array}{l} X-np \geq m+1-np \\ X-np \leq -m-1+nnp \\ X \leq 2np-m-1 \\ X \geq 0 \end{array} \right\} \\ &= \frac{np(1-p)}{(m+1-np)} \end{aligned}$$

$$\therefore P(X \leq m) = 1 - P(X \geq m+1) \geq 1 - \frac{np(1-p)}{(m+1-np)} = 1 - 15.1974 \quad (\text{Using program})$$

$$\therefore P(X \leq 6) \geq -15.1974 \quad (\text{Using Chebysev})$$

$\therefore 0.50759 > -15.1974 \Rightarrow$ this Chebysev inequality follows

$$\therefore 0.07112 \leq P[X \leq 6] \leq 0.98111$$

$$(1.2.5) \quad n=N=1002 \quad \& \quad m=M=408$$

$$\& \quad p = 0.3624772289$$

$P(X \leq m)$ = probability of getting at most m heads -

$$= \sum_{k=0}^m {}^n C_k p^k (1-p)^{n-k}$$

$$= f_X(m) = f_X(408)$$

$$\boxed{P(X \leq 408) = 0.9492697985} \quad (\text{Using program}) \quad (1)$$

Using Markov's Inequality $\mu = np = 3 \times 3 - 23$ (B)

$$\mu = np = 3 \times 3 - 23$$

$$\Pr [x \geq m] \leq \frac{\mu}{m} \Rightarrow \Pr [x \geq m+1] \leq \frac{\mu}{m+1}$$

$$\therefore P[x \leq m] = 1 - P[x \geq m+1]$$

$$\geq 1 - \frac{\mu}{m+1}$$

$$\Rightarrow [P[x \leq 408] \geq 0.0029] \quad (\text{using program})$$

Using Chernoff bound for binomial trial -

$$P(x \leq (1-s)\mu) \leq e^{-\frac{\mu s^2}{2}} \quad | \quad 0 < s < 1$$

$$P(x \leq m) = ?$$

$$m = (1-s)\mu \Rightarrow s = \frac{\mu - m}{\mu}$$

$$s = 0.646$$

$$[P(x \leq m) \leq 0.9894685]$$

Using Chebyshev - $P(|x-\mu| \geq a) \leq \frac{\text{var}(x)}{a^2} = \frac{E[(x-\mu)^2]}{a^2}$

$$P(x \leq m) = 1 - P(x \geq m+1) =$$

$$P(x \geq m+1) = P(x-np \geq m+1-np)$$

$$\leq \frac{\text{var}(x)}{m+1-np} = \frac{np(1-p)}{(m+1-np)}$$

$$\therefore P(x \leq m) \geq 1 - \frac{np(1-p)}{(m+1-np)}$$

$$[P[x \leq 408] \geq 0.99769]$$