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MODULE *HPaxos\_2*

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EXTENDS *HQuorum*, *HLearnerGraph*, *HMessage*, *TLAPS*

$\text{Assert}(P, \text{str}) \triangleq P$

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Algorithm specification

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```
--algorithm HPaxos2{
  variables msgs = {},
            known_msgs = [ $x \in \text{Acceptor} \cup \text{Learner} \mapsto \{\}$ ],
            recent_msgs = [ $a \in \text{Acceptor} \mapsto \{\}$ ],
            prev_msg = [ $a \in \text{Acceptor} \mapsto \text{NoMessage}$ ],
            decision = [ $lb \in \text{Learner} \times \text{Ballot} \mapsto \{\}$ ],
            BVal  $\in [\text{Ballot} \rightarrow \text{Value}]$ ;
```

```
  define {
    Get1a(m)  $\triangleq$ 
      { $x \in \text{Tran}(m) :$ 
        $\wedge x.type = \text{"1a"}$ 
        $\wedge \forall y \in \text{Tran}(m) :$ 
          $y.type = \text{"1a"} \Rightarrow y.bal \leq x.bal$ }
```

```
    B(m, bal)  $\triangleq \exists x \in \text{Get1a}(m) : bal = x.bal$ 
```

```
    V(m, val)  $\triangleq \exists x \in \text{Get1a}(m) : val = \text{BVal}[x.bal]$ 
```

```
    SameBallot(x, y)  $\triangleq$ 
       $\forall b \in \text{Ballot} : B(x, b) \equiv B(y, b)$ 
```

```
    Maximal ballot number of any messages known to acceptor a
    MaxBal(a, mbal)  $\triangleq$ 
       $\wedge \exists m \in \text{known\_msgs}[a] : B(m, \text{mbal})$ 
       $\wedge \forall x \in \text{known\_msgs}[a] :$ 
         $\forall b \in \text{Ballot} : B(x, b) \Rightarrow b \leq \text{mbal}$ 
```

```
    KnownRefs(a, m)  $\triangleq \forall r \in m.ref : r \in \text{known\_msgs}[a]$ 
```

```
    The acceptor is _caught_ in a message x if the transitive references of x
    include evidence such as two different messages both signed by the acceptor,
    which have equal previous messages.
```

```
    CaughtMsg(x)  $\triangleq$ 
      { $m \in \text{Tran}(x) :$ 
        $\wedge m.type \neq \text{"1a"}$ 
        $\wedge \exists m1 \in \text{Tran}(x) :$ 
          $\wedge m1.type \neq \text{"1a"}$ 
          $\wedge m.acc = m1.acc$ }
```

$$\begin{aligned}
& \wedge m \neq m1 \\
& \wedge m \notin \text{PrevTran}(m1) \\
& \wedge m1 \notin \text{PrevTran}(m) \} \\
\text{Caught}(x) & \triangleq \{m.\text{acc} : m \in \text{CaughtMsg}(x)\} \\
\text{CaughtMsg0}(x) & \triangleq \\
& \{m \in \text{Tran}(x) : \\
& \quad \wedge m.\text{type} \neq \text{"1a"} \\
& \quad \wedge \exists m1 \in \text{Tran}(x) : \\
& \quad \quad \wedge m1.\text{type} \neq \text{"1a"} \\
& \quad \quad \wedge m.\text{acc} = m1.\text{acc} \\
& \quad \quad \wedge m \neq m1 \\
& \quad \quad \wedge m.\text{prev} = m1.\text{prev} \} \\
\text{Caught0}(x) & \triangleq \{m.\text{acc} : m \in \text{CaughtMsg0}(x)\} \\
\text{Connected} \\
\text{ConByQuorum}(a, b, x, S) & \triangleq \quad a : \text{Learner}, b : \text{Learner}, x : 1b, S \in \text{ByzQuorum} \\
& \wedge [\text{from} \mapsto a, \text{to} \mapsto b, q \mapsto S] \in \text{TrustSafe} \\
& \wedge S \cap \text{Caught}(x) = \{\} \\
\text{Con}(a, x) & \triangleq \quad a : \text{Learner}, x : 1b \\
& \{b \in \text{Learner} : \\
& \quad \exists S \in \text{ByzQuorum} : \text{ConByQuorum}(a, b, x, S)\} \\
2a\text{-message is } \textit{buried} & \text{ if there exists a quorum of acceptors that have seen} \\
& 2a\text{-messages with different values, the same learner, and higher ballot} \\
& \text{numbers.} \\
\text{Buried}(l, x, y) & \triangleq \quad x : 2a, y : 1b \\
\text{LET } Q & \triangleq \{m \in \text{Tran}(y) : \\
& \quad \exists z \in \text{Tran}(m) : \\
& \quad \quad \wedge z.\text{type} = \text{"2a"} \\
& \quad \quad \wedge l \in x.\text{lrn} \\
& \quad \quad \wedge l \in z.\text{lrn} \\
& \quad \quad \wedge \forall bx, bz \in \text{Ballot} : \\
& \quad \quad \quad B(x, bx) \wedge B(z, bz) \Rightarrow bx < bz \\
& \quad \quad \wedge \forall vx, vz \in \text{Value} : \\
& \quad \quad \quad V(x, vx) \wedge V(z, vz) \Rightarrow vx \neq vz \} \\
\text{IN } [lr \mapsto l, q \mapsto \{m.\text{acc} : m \in Q\}] & \in \text{TrustLive} \\
\text{Connected } 2a \text{ messages and learners} \\
\text{Con2as}(l, x) & \triangleq \quad l : \text{Learner}, x : 1b \\
& \{m \in \text{Tran}(x) : \\
& \quad \wedge m.\text{type} = \text{"2a"} \\
& \quad \wedge m.\text{acc} = x.\text{acc} \\
& \quad \wedge \exists \text{beta} \in \text{Con}(l, x) : \neg \text{Buried}(\text{beta}, m, x)\}
\end{aligned}$$

```

Fresh 1b messages
Fresh( $l, x$ )  $\triangleq$   $l : \text{Learner}, x : 1b$ 
 $\forall m \in \text{Con2as}(l, x) : \forall v \in \text{Value} : V(x, v) \equiv V(m, v)$ 

Quorum of messages referenced by 2a for a learner instance
 $q(l, x) \triangleq x : 2a$ 
LET  $Q \triangleq \{m \in \text{Tran}(x) :$ 
 $\quad \wedge m.type = \text{"1b"}$ 
 $\quad \wedge \text{Fresh}(l, m)$ 
 $\quad \wedge \forall b \in \text{Ballot} : B(m, b) \equiv B(x, b)\}$ 
IN  $\{m.acc : m \in Q\}$ 

ChainRef( $m$ )  $\triangleq$ 
 $\vee m.prev = \text{NoMessage}$ 
 $\vee m.prev \in m.ref \wedge m.prev.acc = m.acc$ 

WellFormed1b( $m$ )  $\triangleq$ 
 $\forall y \in \text{Tran}(m) :$ 
 $\quad m \neq y \wedge \text{SameBallot}(m, y) \Rightarrow y.type = \text{"1a"}$ 

WellFormed2a( $m$ )  $\triangleq$ 
 $m.lrn = \{l \in \text{Learner} : [lr \mapsto l, q \mapsto q(l, m)] \in \text{TrustLive}\}$ 

WellFormed( $m$ )  $\triangleq$ 
 $\wedge m \in \text{Message}$ 
 $\wedge \exists b \in \text{Ballot} : B(m, b)$  TODO prove it
 $\wedge \text{ChainRef}(m)$ 
 $\wedge m.type = \text{"1b"} \Rightarrow$ 
 $\quad \wedge (\exists r \in m.refs : r.type = \text{"1a"})$ 
 $\quad \wedge \text{WellFormed1b}(m)$ 
 $\wedge m.type = \text{"2a"} \Rightarrow$ 
 $\quad \wedge m.ref \neq \{\}$ 
 $\quad \wedge \text{WellFormed2a}(m)$ 

Known2a( $l, b, v$ )  $\triangleq$ 
 $\{x \in \text{known\_msgs}[l] :$ 
 $\quad \wedge x.type = \text{"2a"}$ 
 $\quad \wedge l \in x.lrn$ 
 $\quad \wedge B(x, b)$ 
 $\quad \wedge V(x, v)\}$ 

ChosenIn( $l, b, v$ )  $\triangleq$ 
 $\exists S \in \text{SUBSET } \text{Known2a}(l, b, v) :$ 
 $\quad [lr \mapsto l, q \mapsto \{m.acc : m \in S\}] \in \text{TrustLive}$ 
}

macro Send(  $m$  ) {  $msgs := msgs \cup \{m\}$  }

```

```

macro Send1a( b ) {
  Send([type  $\mapsto$  "1a", bal  $\mapsto$  b, prev  $\mapsto$  NoMessage, ref  $\mapsto$  {}])
}

macro Recv( m ) {
  when  $\wedge m \notin \text{known\_msgs}[\text{self}]$ 
     $\wedge \text{KnownRefs}(\text{self}, m)$ ;
  known_msgs[self] := known_msgs[self]  $\cup$  {m}
}

macro ProposerSendAction( b ) { Send1a(b) }

macro Process1a( m ) {
  when m.type = "1a";
  with ( new1b = [type  $\mapsto$  "1b",
    acc  $\mapsto$  self,
    prev  $\mapsto$  prev_msg[self],
    ref  $\mapsto$  recent_msgs[self]  $\cup$  {m} ] )
  {
    assert new1b  $\in$  Message;
    either {
      when WellFormed1b(new1b);
      prev_msg[self] := new1b;
      recent_msgs[self] := {new1b};
      Send(new1b)
    }
    or {
      when  $\neg \text{WellFormed1b}(\text{new1b})$ ;
      skip
    }
  }
}

macro Process1b( m ) {
  when m.type = "1b";
  with ( LL  $\in$  SUBSET Learner,
    new2a = [type  $\mapsto$  "2a",
      lrn  $\mapsto$  LL,
      acc  $\mapsto$  self,
      prev  $\mapsto$  prev_msg[self],
      ref  $\mapsto$  recent_msgs[self]  $\cup$  {m} ] )
  {
    assert new2a  $\in$  Message ;
    when WellFormed2a(new2a);
    prev_msg[self] := new2a;
    recent_msgs[self] := {new2a};
  }
}

```

```

    Send(new2a)
  }
}

macro Process2a( m ) {
  when m.type = "2a" ;
  recent_msgs[self] := recent_msgs[self] ∪ {m}
}

macro FakeSend1b( ) {
  with ( fin ∈ FINSUBSET(msgs, RefCardinality),
        new1b = [type ↦ "1b", acc ↦ self, ref ↦ fin] )
  {
    when WellFormed(new1b);
    Send(new1b)
  }
}

macro FakeSend2a( ) {
  with ( fin ∈ FINSUBSET(msgs, RefCardinality),
        LL ∈ SUBSET Learner,
        new2a = [type ↦ "2a", lrn ↦ LL, acc ↦ self, ref ↦ fin] )
  {
    when WellFormed(new2a);
    Send(new2a)
  }
}

macro LearnerRecv( m ) {
  when WellFormed(m);
  Recv(m)
}

macro LearnerDecide( b, v ) {
  when ChosenIn(self, b, v);
  decision[⟨self, b⟩] := decision[self, b] ∪ {v}
}

process ( proposer ∈ Proposer ) {
  propose: while ( TRUE ) {
    with ( b ∈ Ballot ) { ProposerSendAction(b) }
  }
}

process ( safe_acceptor ∈ SafeAcceptor ) {
  safe: while ( TRUE ) {
    with ( m ∈ msgs ) {

```

```

    Recv(m);
    when WellFormed(m);
    either Process1a(m)
    or     Process1b(m)
    or     Process2a(m)
    }
  }
}

process ( learner ∈ Learner ) {
  learn: while ( TRUE ) {
    either with ( m ∈ msgs ) LearnerRecv(m)
    or     with ( b ∈ Ballot, v ∈ Value ) LearnerDecide(b, v)
  }
}

process ( fake_acceptor ∈ FakeAcceptor ) {
  fake: while ( TRUE ) {
    either FakeSend1b()
    or     FakeSend2a()
  }
}
}

*****
BEGIN TRANSLATION (chksum(pcal) = "e34bd140" ∧ chksum(tla) = "89878241")
VARIABLES msgs, known_msgs, recent_msgs, prev_msg, decision, BVal

define statement
Get1a(m)  $\triangleq$ 
  {x ∈ Tran(m) :
    ∧ x.type = "1a"
    ∧ ∀ y ∈ Tran(m) :
      y.type = "1a" ⇒ y.bal ≤ x.bal}

B(m, bal)  $\triangleq$  ∃ x ∈ Get1a(m) : bal = x.bal

V(m, val)  $\triangleq$  ∃ x ∈ Get1a(m) : val = BVal[x.bal]

SameBallot(x, y)  $\triangleq$ 
  ∀ b ∈ Ballot : B(x, b) ≡ B(y, b)

MaxBal(a, mbal)  $\triangleq$ 
  ∧ ∃ m ∈ known_msgs[a] : B(m, mbal)
  ∧ ∀ x ∈ known_msgs[a] :
    ∀ b ∈ Ballot : B(x, b) ⇒ b ≤ mbal

KnownRefs(a, m)  $\triangleq$  ∀ r ∈ m.ref : r ∈ known_msgs[a]

```

$$\begin{aligned}
CaughtMsg(x) &\triangleq \\
&\{m \in Tran(x) : \\
&\quad \wedge m.type \neq "1a" \\
&\quad \wedge \exists m1 \in Tran(x) : \\
&\quad \quad \wedge m1.type \neq "1a" \\
&\quad \quad \wedge m.acc = m1.acc \\
&\quad \quad \wedge m \neq m1 \\
&\quad \quad \wedge m \notin PrevTran(m1) \\
&\quad \quad \wedge m1 \notin PrevTran(m)\} \\
Caught(x) &\triangleq \{m.acc : m \in CaughtMsg(x)\}
\end{aligned}$$

$$\begin{aligned}
CaughtMsg0(x) &\triangleq \\
&\{m \in Tran(x) : \\
&\quad \wedge m.type \neq "1a" \\
&\quad \wedge \exists m1 \in Tran(x) : \\
&\quad \quad \wedge m1.type \neq "1a" \\
&\quad \quad \wedge m.acc = m1.acc \\
&\quad \quad \wedge m \neq m1 \\
&\quad \quad \wedge m.prev = m1.prev\} \\
Caught0(x) &\triangleq \{m.acc : m \in CaughtMsg0(x)\}
\end{aligned}$$

$$\begin{aligned}
ConByQuorum(a, b, x, S) &\triangleq \\
&\wedge [from \mapsto a, to \mapsto b, q \mapsto S] \in TrustSafe \\
&\wedge S \cap Caught(x) = \{\}
\end{aligned}$$

$$\begin{aligned}
Con(a, x) &\triangleq \\
&\{b \in Learner : \\
&\quad \exists S \in ByzQuorum : ConByQuorum(a, b, x, S)\}
\end{aligned}$$

$$\begin{aligned}
Buried(l, x, y) &\triangleq \\
LET \ Q &\triangleq \{m \in Tran(y) : \\
&\quad \exists z \in Tran(m) : \\
&\quad \quad \wedge z.type = "2a" \\
&\quad \quad \wedge l \in x.lrn \\
&\quad \quad \wedge l \in z.lrn \\
&\quad \quad \wedge \forall bx, bz \in Ballot : \\
&\quad \quad \quad B(x, bx) \wedge B(z, bz) \Rightarrow bx < bz \\
&\quad \quad \wedge \forall vx, vz \in Value : \\
&\quad \quad \quad V(x, vx) \wedge V(z, vz) \Rightarrow vx \neq vz\} \\
IN \ [lr \mapsto l, q \mapsto \{m.acc : m \in Q\}] &\in TrustLive
\end{aligned}$$

$$Con2as(l, x) \triangleq$$

$$\begin{aligned}
& \{m \in \text{Tran}(x) : \\
& \quad \wedge m.type = \text{"2a"} \\
& \quad \wedge m.acc = x.acc \\
& \quad \wedge \exists beta \in \text{Con}(l, x) : \neg \text{Buried}(beta, m, x)\}
\end{aligned}$$

$$\begin{aligned}
\text{Fresh}(l, x) & \triangleq \\
& \forall m \in \text{Con2as}(l, x) : \forall v \in \text{Value} : V(x, v) \equiv V(m, v)
\end{aligned}$$

$$\begin{aligned}
q(l, x) & \triangleq \\
& \text{LET } Q \triangleq \{m \in \text{Tran}(x) : \\
& \quad \wedge m.type = \text{"1b"} \\
& \quad \wedge \text{Fresh}(l, m) \\
& \quad \wedge \forall b \in \text{Ballot} : B(m, b) \equiv B(x, b)\} \\
& \text{IN } \{m.acc : m \in Q\}
\end{aligned}$$

$$\begin{aligned}
\text{ChainRef}(m) & \triangleq \\
& \vee m.prev = \text{NoMessage} \\
& \vee m.prev \in m.ref \wedge m.prev.acc = m.acc
\end{aligned}$$

$$\begin{aligned}
\text{WellFormed1b}(m) & \triangleq \\
& \forall y \in \text{Tran}(m) : \\
& \quad m \neq y \wedge \text{SameBallot}(m, y) \Rightarrow y.type = \text{"1a"}
\end{aligned}$$

$$\begin{aligned}
\text{WellFormed2a}(m) & \triangleq \\
& m.lrn = \{l \in \text{Learner} : [lr \mapsto l, q \mapsto q(l, m)] \in \text{TrustLive}\}
\end{aligned}$$

$$\begin{aligned}
\text{WellFormed}(m) & \triangleq \\
& \wedge m \in \text{Message} \\
& \wedge \exists b \in \text{Ballot} : B(m, b) \\
& \wedge \text{ChainRef}(m) \\
& \wedge m.type = \text{"1b"} \Rightarrow \\
& \quad \wedge (\exists r \in m.refs : r.type = \text{"1a"}) \\
& \quad \wedge \text{WellFormed1b}(m) \\
& \wedge m.type = \text{"2a"} \Rightarrow \\
& \quad \wedge m.ref \neq \{\} \\
& \quad \wedge \text{WellFormed2a}(m)
\end{aligned}$$

$$\begin{aligned}
\text{Known2a}(l, b, v) & \triangleq \\
& \{x \in \text{known\_msgs}[l] : \\
& \quad \wedge x.type = \text{"2a"} \\
& \quad \wedge l \in x.lrn \\
& \quad \wedge B(x, b) \\
& \quad \wedge V(x, v)\}
\end{aligned}$$

$$\begin{aligned}
\text{ChosenIn}(l, b, v) & \triangleq \\
& \exists S \in \text{SUBSET } \text{Known2a}(l, b, v) :
\end{aligned}$$



$$[lr \mapsto l, q \mapsto \{m.acc : m \in S\}] \in TrustLive$$

$$vars \triangleq \langle msgs, known\_msgs, recent\_msgs, prev\_msg, decision, BVal \rangle$$

$$ProcSet \triangleq (Proposer) \cup (SafeAcceptor) \cup (Learner) \cup (FakeAcceptor)$$

$$\begin{aligned} Init &\triangleq \text{Global variables} \\ &\wedge msgs = \{\} \\ &\wedge known\_msgs = [x \in Acceptor \cup Learner \mapsto \{\}] \\ &\wedge recent\_msgs = [a \in Acceptor \mapsto \{\}] \\ &\wedge prev\_msg = [a \in Acceptor \mapsto NoMessage] \\ &\wedge decision = [lb \in Learner \times Ballot \mapsto \{\}] \\ &\wedge BVal \in [Ballot \rightarrow Value] \end{aligned}$$

$$\begin{aligned} proposer(self) &\triangleq \wedge \exists b \in Ballot : \\ &\quad msgs' = (msgs \cup \{([type \mapsto "1a", bal \mapsto b, prev \mapsto NoMessage, ref \mapsto \{\}])\}) \\ &\quad \wedge UNCHANGED \langle known\_msgs, recent\_msgs, prev\_msg, decision, \\ &\quad \quad BVal \rangle \end{aligned}$$

$$\begin{aligned} safe\_acceptor(self) &\triangleq \wedge \exists m \in msgs : \\ &\quad \wedge \wedge m \notin known\_msgs[self] \\ &\quad \wedge KnownRefs(self, m) \\ &\quad \wedge known\_msgs' = [known\_msgs \text{ EXCEPT } ![self] = known\_msgs[self] \cup \{m\}] \\ &\quad \wedge WellFormed(m) \\ &\quad \wedge \vee \wedge m.type = "1a" \\ &\quad \quad \wedge LET new1b \triangleq [type \mapsto "1b", \\ &\quad \quad \quad acc \mapsto self, \\ &\quad \quad \quad prev \mapsto prev\_msg[self], \\ &\quad \quad \quad ref \mapsto recent\_msgs[self] \cup \{m\}] IN \\ &\quad \quad \wedge Assert(new1b \in Message, \\ &\quad \quad \quad \text{"Failure of assertion at line 164, column 7 of macro called at line"} \\ &\quad \quad \wedge \vee \wedge WellFormed1b(new1b) \\ &\quad \quad \quad \wedge prev\_msg' = [prev\_msg \text{ EXCEPT } ![self] = new1b] \\ &\quad \quad \quad \wedge recent\_msgs' = [recent\_msgs \text{ EXCEPT } ![self] = \{new1b\}] \\ &\quad \quad \quad \wedge msgs' = (msgs \cup \{new1b\}) \\ &\quad \quad \vee \wedge \neg WellFormed1b(new1b) \\ &\quad \quad \wedge TRUE \\ &\quad \quad \wedge UNCHANGED \langle msgs, recent\_msgs, prev\_msg \rangle \\ &\quad \vee \wedge m.type = "1b" \\ &\quad \quad \wedge \exists LL \in SUBSET Learner : \\ &\quad \quad \quad LET new2a \triangleq [type \mapsto "2a", \\ &\quad \quad \quad \quad lrn \mapsto LL, \\ &\quad \quad \quad \quad acc \mapsto self, \\ &\quad \quad \quad \quad prev \mapsto prev\_msg[self], \\ &\quad \quad \quad \quad ref \mapsto recent\_msgs[self] \cup \{m\}] IN \\ &\quad \quad \quad \wedge Assert(new2a \in Message, \end{aligned}$$

“Failure of assertion at line 187, column 7 of macro called at line

$$\begin{aligned}
& \wedge \text{WellFormed2a}(\text{new2a}) \\
& \wedge \text{prev\_msg}' = [\text{prev\_msg} \text{ EXCEPT } ![\text{self}] = \text{new2a}] \\
& \wedge \text{recent\_msgs}' = [\text{recent\_msgs} \text{ EXCEPT } ![\text{self}] = \{\text{new2a}\}] \\
& \wedge \text{msgs}' = (\text{msgs} \cup \{\text{new2a}\}) \\
\vee & \wedge m.\text{type} = \text{"2a"} \\
& \wedge \text{recent\_msgs}' = [\text{recent\_msgs} \text{ EXCEPT } ![\text{self}] = \text{recent\_msgs}[\text{self}] \cup \{m\}] \\
& \wedge \text{UNCHANGED } \langle \text{msgs}, \text{prev\_msg} \rangle \\
& \wedge \text{UNCHANGED } \langle \text{decision}, \text{BVal} \rangle
\end{aligned}$$

$$\begin{aligned}
\text{learner}(\text{self}) \triangleq & \wedge \vee \wedge \exists m \in \text{msgs} : \\
& \wedge \text{WellFormed}(m) \\
& \wedge \wedge m \notin \text{known\_msgs}[\text{self}] \\
& \wedge \text{KnownRefs}(\text{self}, m) \\
& \wedge \text{known\_msgs}' = [\text{known\_msgs} \text{ EXCEPT } ![\text{self}] = \text{known\_msgs}[\text{self}] \cup \{m\}] \\
& \wedge \text{UNCHANGED } \text{decision} \\
\vee & \wedge \exists b \in \text{Ballot} : \\
& \exists v \in \text{Value} : \\
& \wedge \text{ChosenIn}(\text{self}, b, v) \\
& \wedge \text{decision}' = [\text{decision} \text{ EXCEPT } ![\langle \text{self}, b \rangle] = \text{decision}[\text{self}, b] \cup \{v\}] \\
& \wedge \text{UNCHANGED } \text{known\_msgs} \\
& \wedge \text{UNCHANGED } \langle \text{msgs}, \text{recent\_msgs}, \text{prev\_msg}, \text{BVal} \rangle
\end{aligned}$$

$$\begin{aligned}
\text{fake\_acceptor}(\text{self}) \triangleq & \wedge \vee \wedge \exists \text{fin} \in \text{FINSUBSET}(\text{msgs}, \text{RefCardinality}) : \\
& \text{LET } \text{new1b} \triangleq [\text{type} \mapsto \text{"1b"}, \text{acc} \mapsto \text{self}, \text{ref} \mapsto \text{fin}] \text{IN} \\
& \wedge \text{WellFormed}(\text{new1b}) \\
& \wedge \text{msgs}' = (\text{msgs} \cup \{\text{new1b}\}) \\
\vee & \wedge \exists \text{fin} \in \text{FINSUBSET}(\text{msgs}, \text{RefCardinality}) : \\
& \exists \text{LL} \in \text{SUBSET } \text{Learner} : \\
& \text{LET } \text{new2a} \triangleq [\text{type} \mapsto \text{"2a"}, \text{lrn} \mapsto \text{LL}, \text{acc} \mapsto \text{self}, \text{ref} \mapsto \text{fin}] \text{IN} \\
& \wedge \text{WellFormed}(\text{new2a}) \\
& \wedge \text{msgs}' = (\text{msgs} \cup \{\text{new2a}\}) \\
& \wedge \text{UNCHANGED } \langle \text{known\_msgs}, \text{recent\_msgs}, \text{prev\_msg}, \\
& \quad \text{decision}, \text{BVal} \rangle
\end{aligned}$$

$$\begin{aligned}
\text{Next} \triangleq & (\exists \text{self} \in \text{Proposer} : \text{proposer}(\text{self})) \\
& \vee (\exists \text{self} \in \text{SafeAcceptor} : \text{safe\_acceptor}(\text{self})) \\
& \vee (\exists \text{self} \in \text{Learner} : \text{learner}(\text{self})) \\
& \vee (\exists \text{self} \in \text{FakeAcceptor} : \text{fake\_acceptor}(\text{self}))
\end{aligned}$$

$$\text{Spec} \triangleq \text{Init} \wedge \Box[\text{Next}]_{\text{vars}}$$

END TRANSLATION

$$\text{Send}(m) \triangleq \text{msgs}' = \text{msgs} \cup \{m\}$$

$$\begin{aligned}
\text{Recv}(a, m) &\triangleq \\
&\wedge m \notin \text{known\_msgs}[a] \text{ ignore known messages} \\
&\wedge \text{KnownRefs}(a, m) \\
&\wedge \text{known\_msgs}' = [\text{known\_msgs} \text{ EXCEPT } ![a] = \text{known\_msgs}[a] \cup \{m\}] \\
\\
\text{Send1a}(b) &\triangleq \\
&\wedge \text{Send}([type \mapsto \text{"1a"}, bal \mapsto b, prev \mapsto \text{NoMessage}, ref \mapsto \{\}]) \\
&\wedge \text{UNCHANGED } \langle \text{known\_msgs}, \text{recent\_msgs}, \text{prev\_msg} \rangle \\
&\wedge \text{UNCHANGED } \text{decision} \\
&\wedge \text{UNCHANGED } BVal \\
\\
\text{Process1a}(a, m) &\triangleq \\
&\text{LET } new1b \triangleq [type \mapsto \text{"1b"}, acc \mapsto a, \\
&\quad prev \mapsto \text{prev\_msg}[a], \\
&\quad ref \mapsto \text{recent\_msgs}[a] \cup \{m\}] \text{IN} \\
&\wedge m.type = \text{"1a"} \\
&\wedge \text{Recv}(a, m) \\
&\wedge \text{WellFormed}(m) \\
&\wedge new1b \in \text{Message} \\
&\wedge \vee \wedge \text{WellFormed1b}(new1b) \\
&\quad \wedge \text{Send}(new1b) \\
&\quad \wedge \text{recent\_msgs}' = [\text{recent\_msgs} \text{ EXCEPT } ![a] = \{new1b\}] \\
&\quad \wedge \text{prev\_msg}' = [\text{prev\_msg} \text{ EXCEPT } ![a] = new1b] \\
&\vee \wedge \neg \text{WellFormed1b}(new1b) \\
&\quad \wedge \text{UNCHANGED } \langle \text{msgs}, \text{prev\_msg}, \text{recent\_msgs} \rangle \\
&\wedge \text{UNCHANGED } \text{decision} \\
&\wedge \text{UNCHANGED } BVal \\
\\
\text{Process1b}(a, m) &\triangleq \\
&\wedge m.type = \text{"1b"} \\
&\wedge \text{Recv}(a, m) \\
&\wedge \text{WellFormed}(m) \\
&\wedge \exists LL \in \text{SUBSET } \text{Learner} : \\
&\quad \text{LET } new2a \triangleq [type \mapsto \text{"2a"}, lrn \mapsto LL, acc \mapsto a, \\
&\quad \quad prev \mapsto \text{prev\_msg}[a], \\
&\quad \quad ref \mapsto \text{recent\_msgs}[a] \cup \{m\}] \text{IN} \\
&\quad \wedge new2a \in \text{Message} \\
&\quad \wedge \text{WellFormed2a}(new2a) \\
&\quad \wedge \text{Send}(new2a) \\
&\quad \wedge \text{recent\_msgs}' = [\text{recent\_msgs} \text{ EXCEPT } ![a] = \{new2a\}] \\
&\quad \wedge \text{prev\_msg}' = [\text{prev\_msg} \text{ EXCEPT } ![a] = new2a] \\
&\wedge \text{UNCHANGED } \text{decision} \\
&\wedge \text{UNCHANGED } BVal \\
\\
\text{Process2a}(a, m) &\triangleq \\
&\wedge m.type = \text{"2a"}
\end{aligned}$$

$$\begin{aligned}
& \wedge \text{Recv}(a, m) \\
& \wedge \text{WellFormed}(m) \\
& \wedge \text{recent\_msgs}' = [\text{recent\_msgs} \text{ EXCEPT } ![a] = \text{recent\_msgs}[a] \cup \{m\}] \\
& \wedge \text{UNCHANGED } \langle \text{msgs}, \text{prev\_msg} \rangle \\
& \wedge \text{UNCHANGED } \text{decision} \\
& \wedge \text{UNCHANGED } B\text{Val} \\
\\
\text{ProposerSendAction}(p) & \triangleq \\
& \exists \text{bal} \in \text{Ballot} : \text{Send1a}(\text{bal}) \\
\\
\text{AcceptorProcessAction}(a) & \triangleq \\
& \exists m \in \text{msgs} : \\
& \quad \vee \text{Process1a}(a, m) \\
& \quad \vee \text{Process1b}(a, m) \\
& \quad \vee \text{Process2a}(a, m) \\
\\
\text{FakeSend1b}(a) & \triangleq \\
& \wedge \exists \text{fin} \in \text{FINSUBSET}(\text{msgs}, \text{RefCardinality}) : \\
& \quad \text{LET } \text{new1b} \triangleq [\text{type} \mapsto \text{"1b"}, \text{acc} \mapsto a, \text{ref} \mapsto \text{fin}] \text{IN} \\
& \quad \wedge \text{WellFormed}(\text{new1b}) \\
& \quad \wedge \text{Send}(\text{new1b}) \\
& \wedge \text{UNCHANGED } \langle \text{known\_msgs}, \text{recent\_msgs}, \text{prev\_msg} \rangle \\
& \wedge \text{UNCHANGED } \text{decision} \\
& \wedge \text{UNCHANGED } B\text{Val} \\
\\
\text{FakeSend2a}(a) & \triangleq \\
& \wedge \exists \text{fin} \in \text{FINSUBSET}(\text{msgs}, \text{RefCardinality}) : \\
& \quad \exists LL \in \text{SUBSET } \text{Learner} : \\
& \quad \text{LET } \text{new2a} \triangleq [\text{type} \mapsto \text{"2a"}, \text{lrn} \mapsto LL, \text{acc} \mapsto a, \text{ref} \mapsto \text{fin}] \text{IN} \\
& \quad \wedge \text{WellFormed}(\text{new2a}) \\
& \quad \wedge \text{Send}(\text{new2a}) \\
& \wedge \text{UNCHANGED } \langle \text{known\_msgs}, \text{recent\_msgs}, \text{prev\_msg} \rangle \\
& \wedge \text{UNCHANGED } \text{decision} \\
& \wedge \text{UNCHANGED } B\text{Val} \\
\\
\text{LearnerRecv}(l, m) & \triangleq \\
& \wedge \text{Recv}(l, m) \\
& \wedge \text{WellFormed}(m) \\
& \wedge \text{UNCHANGED } \langle \text{msgs}, \text{recent\_msgs}, \text{prev\_msg} \rangle \\
& \wedge \text{UNCHANGED } \text{decision} \\
& \wedge \text{UNCHANGED } B\text{Val} \\
\\
\text{LearnerDecide}(l, b, v) & \triangleq \\
& \wedge \text{ChosenIn}(l, b, v) \\
& \wedge \text{decision}' = [\text{decision} \text{ EXCEPT } ![\langle l, b \rangle] = \text{decision}[l, b] \cup \{v\}] \\
& \wedge \text{UNCHANGED } \langle \text{msgs}, \text{known\_msgs}, \text{recent\_msgs}, \text{prev\_msg} \rangle \\
& \wedge \text{UNCHANGED } B\text{Val}
\end{aligned}$$

$$\begin{aligned}
\text{LearnerAction}(\text{lrn}) &\triangleq \\
&\vee \exists m \in \text{msgs} : \\
&\quad \text{LearnerRecv}(\text{lrn}, m) \\
&\vee \exists \text{bal} \in \text{Ballot} : \\
&\quad \exists \text{val} \in \text{Value} : \\
&\quad \text{LearnerDecide}(\text{lrn}, \text{bal}, \text{val})
\end{aligned}$$

$$\begin{aligned}
\text{FakeAcceptorAction}(a) &\triangleq \\
&\vee \text{FakeSend1b}(a) \\
&\vee \text{FakeSend2a}(a)
\end{aligned}$$

$$\begin{aligned}
\text{NextTLA} &\triangleq \\
&\vee \exists p \in \text{Proposer} : \\
&\quad \text{ProposerSendAction}(p) \\
&\vee \exists \text{acc} \in \text{SafeAcceptor} : \\
&\quad \text{AcceptorProcessAction}(\text{acc}) \\
&\vee \exists \text{lrn} \in \text{Learner} : \\
&\quad \text{LearnerAction}(\text{lrn}) \\
&\vee \exists \text{acc} \in \text{FakeAcceptor} : \\
&\quad \text{FakeAcceptorAction}(\text{acc})
\end{aligned}$$

THEOREM  $\text{NextDef} \triangleq \text{Next} \equiv \text{NextTLA}$

- ⟨1⟩1. ASSUME NEW  $\text{self} \in \text{Proposer}$   
PROVE  $\text{proposer}(\text{self}) \equiv \text{ProposerSendAction}(\text{self})$   
BY DEF  $\text{proposer}$ ,  $\text{ProposerSendAction}$ ,  $\text{Send1a}$ ,  $\text{Send}$
- ⟨1⟩2. ASSUME NEW  $\text{self} \in \text{SafeAcceptor}$   
PROVE  $\text{safe\_acceptor}(\text{self}) \equiv \text{AcceptorProcessAction}(\text{self})$   
BY Zenon DEF  $\text{safe\_acceptor}$ ,  $\text{AcceptorProcessAction}$ ,  $\text{Process1a}$ ,  $\text{Process1b}$ ,  $\text{Process2a}$ ,  $\text{Recv}$ ,  $\text{Send}$ ,  $\text{Ass}$
- ⟨1⟩3. ASSUME NEW  $\text{self} \in \text{Learner}$   
PROVE  $\text{learner}(\text{self}) \equiv \text{LearnerAction}(\text{self})$   
BY Zenon DEF  $\text{learner}$ ,  $\text{LearnerAction}$ ,  $\text{LearnerRecv}$ ,  $\text{LearnerDecide}$ ,  $\text{Recv}$
- ⟨1⟩4. ASSUME NEW  $\text{self} \in \text{FakeAcceptor}$   
PROVE  $\text{fake\_acceptor}(\text{self}) \equiv \text{FakeAcceptorAction}(\text{self})$   
BY Zenon DEF  $\text{fake\_acceptor}$ ,  $\text{FakeAcceptorAction}$ ,  $\text{FakeSend1b}$ ,  $\text{FakeSend2a}$ ,  $\text{Send}$
- ⟨1⟩5. QED BY ⟨1⟩1, ⟨1⟩2, ⟨1⟩3, ⟨1⟩4 DEF  $\text{Next}$ ,  $\text{NextTLA}$

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$$\begin{aligned}
\text{Safety} &\triangleq \\
&\forall L1, L2 \in \text{Learner} : \forall B1, B2 \in \text{Ballot} : \forall V1, V2 \in \text{Value} : \\
&\quad \langle L1, L2 \rangle \in \text{Ent} \wedge \\
&\quad V1 \in \text{decision}[L1, B1] \wedge V2 \in \text{decision}[L2, B2] \Rightarrow \\
&\quad V1 = V2
\end{aligned}$$

THEOREM  $\text{SafetyResult} \triangleq \text{Spec} \Rightarrow \Box \text{Safety}$

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Sanity check propositions

$$\begin{aligned}
\text{SanityCheck0} &\triangleq \\
&\forall L \in \text{Learner} : \text{Cardinality}(\text{known\_msgs}[L]) = 0 \\
\text{SanityCheck1} &\triangleq \\
&\forall L \in \text{Learner} : \forall m1, m2 \in \text{known\_msgs}[L] : \\
&\quad \forall b1, b2 \in \text{Ballot} : \\
&\quad \quad B(m1, b1) \wedge B(m2, b2) \Rightarrow b1 = b2 \\
\text{2aNotSent} &\triangleq \\
&\forall M \in \text{msgs} : M.\text{type} \neq \text{"2a"} \\
\text{2aNotSentBySafeAcceptor} &\triangleq \\
&\forall M \in \text{msgs} : M.\text{type} = \text{"2a"} \Rightarrow M.\text{acc} \notin \text{SafeAcceptor} \\
\text{1bNotSentBySafeAcceptor} &\triangleq \\
&\forall M \in \text{msgs} : M.\text{type} = \text{"1b"} \Rightarrow M.\text{acc} \notin \text{SafeAcceptor} \\
\text{NoDecision} &\triangleq \\
&\forall L \in \text{Learner} : \forall BB \in \text{Ballot} : \forall VV \in \text{Value} : \\
&\quad VV \notin \text{decision}[L, BB] \\
\text{UniqueDecision} &\triangleq \\
&\forall L1, L2 \in \text{Learner} : \forall B1, B2 \in \text{Ballot} : \forall V1, V2 \in \text{Value} : \\
&\quad V1 \in \text{decision}[L1, B1] \wedge V2 \in \text{decision}[L2, B2] \Rightarrow \\
&\quad V1 = V2
\end{aligned}$$


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