# Lesson 44: Repeated Roots - Why We Get $t^k Terms$

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#### 1 Introduction: The Problem of Repeated Roots

When the characteristic equation of a linear ODE with constant coefficients has repeated roots, the naive approach of using only exponential solutions fails to generate enough linearly independent solutions.

**Definition 1** (Multiplicity of a Root). If the characteristic polynomial p(r) can be factored as  $p(r) = (r - r_0)^m \ q(r)$  where  $q(r_0) \neq 0$ , then  $r_0$  is a root of multiplicity m.

**Theorem 1** (Dimension of Solution Space). An n-th order linear homogeneous ODE has an n-dimensional solution space. Therefore, we need exactly n linearly independent solutions.

#### 2 Failure of Simple Exponentials

**Example 1** (The Problem). Consider y'' - 4y' + 4y = 0 with characteristic equation  $(r-2)^2 = 0$ .

 $\label{eq:constraint} \mbox{\it If we only use $y=e^{2t}$, $we have just one solution. } \mbox{\it Butwe need two linearly independent solutions for this second erequation.}$ 

When a root r has multiplicity m > 1, the solution  $e^{rt}$  alone cannot span the m-dimensional subspace of solutions corresponding to that root.

#### 3 Derivation via Reduction of Order

**Method 1** (Reduction of Order for Repeated Roots). Given one solution  $y_1 = e^{rt}$  where risar epeated root, we  $v(t)y_1 = v(t)e^{rt}$ 

**Theorem 2** (Second Solution for Double Root). If r is a double root of the characteristic equation for y''+py'+qy=0, and  $y_1=e^{rt}isonesolution$ , then  $y_2=te^{rt}isasecond$ , linearly independent solution

Proof. Let  $y_2 = v(t)e^{rt}$ . Then:

Substituting into the differential equation:

v''

$$e^{rt} + 2rv'e^{rt} + r^{2vert} + p(v'e^{rt} + rve^{rt}) + qve^{rt} = 0$$

Factoring out  $e^{rt}$ :

$$v'' + v'(2r + p) + v(r^2 + pr + q) = 0$$

Since r is a double root of  $\lambda^2 + p\lambda + q = 0$ :

- 1.  $r^2 + pr + q = 0$  (r is a root)
- 2. 2r + p = 0 (derivative of characteristic polynomial at r equals 0)

Therefore: v'' = 0, giving  $v = c_{1t} + c_2$ . Taking  $c_1 = 1$ ,  $c_2 = 0$  yields  $y_2 = te^{rt}$ .

### 4 The Differential Operator Approach

The differential equation  $(D-r)^m[y] = 0$  naturally produces solutions  $t^k e^{rt}$  for  $k=0, 1, \ldots, m-1$ .

**Theorem 3** (Operator Factorization). If the characteristic polynomial factors as  $p(\lambda) = (\lambda - r)^m$ , then the differential operator factors as  $z = (D - r)^m$ 

where

$$D = d_{\overline{dt}}$$
.

**Lemma 1** (Kernel of  $(D-r)^m$ ). The kernel (null space) of the operator  $(D-r)^m$  is  $:ker((D-r)^m) = span\{e^{rt}, te^{rt}, t^{2ert}, \dots, t^{m-1}e^{rt}\}$ 

Proof by Induction. Base case (m = 1):  $(D - r)[y] = 0 \Rightarrow y = ce^{rt}$ .

Inductive step: Assume true for m = k. For m = k + 1: If  $(D - r)^{k+1}[y] = 0$ , let w = (D-r)[y]. Then  $(D-r)^k[w] = 0$ , sow =  $p_k(t)e^{rt}wherep_kisapolynomial of degree <math>\leq k-1$ .

Solving  $(D-r)[y] = p_k(t)e^{rt}$ : Using integrating factor  $e^{-rt}$ :  $\frac{d}{dt}[e^{-rt}y] = p_k(t)$ 

Integrating:  $e^{-rt}y = P_{k+1}(t) + C$  where  $P_{k+1}$  has degree  $\leq k$ .

Therefore:  $y = [P_{k+1}(t) + C]e^{rt} = p_{k+1}(t)e^{rt}wheredeg(p_{k+1}) \le k.$ 

## 5 The Limit Perspective

**Theorem 4** (Repeated Roots as Limits). The solution  $te^{rt}$  can be understood as the limit of solutions for near the Heuristic Derivation. Consider two roots  $r_1 = r$  and  $r_2 = r + \epsilon$  with solutions:

$$y = c_1$$

$$e^{rt} + c_2 e^{(r+\epsilon)t}$$

Rewriting:

$$y = c_1$$

$$e^{rt} + c_2 e^{rt} e^{\epsilon} t$$

For small  $\epsilon$ :  $e^{\epsilon t} \approx 1 + \epsilon t$ Thus:

$$y \approx (c_1 + c_2)$$

$$e^{rt} + c_2 \epsilon t e^{rt}$$

As  $\epsilon \to 0$ , setting  $A = c_1 + c_2$  and  $B = \lim_{\epsilon \to 0} c_2 \epsilon$ :

$$y = A$$

$$e^{rt} + Bte^{rt}$$

### 6 Higher Multiplicities

**Theorem 5** (General Multiplicity Case). If r is a root of multiplicity m, then the m linearly independent solutions are:

$$e^{rt}, t$$

$$e^{rt}$$
,  $t^{2ert}$ , ...,  $t^{m-1}e^{rt}$ 

For a root r of multiplicity m:

Solutions = 
$$\{t$$

$$e^{rt}$$
: k = 0, 1, 2, ..., m-1

#### 7 Complex Repeated Roots

**Theorem 6** (Complex Repeated Roots). If  $\alpha \pm i\beta$  are complex conjugate roots each of multiplicity m, the 2m real-valued linearly independent solutions are:

t

$$^k e^{\alpha} t \cos(\beta t), t^k e^{\alpha} t \sin(\beta t)$$
 for  $k = 0, 1, \dots, m-1$ 

For complex repeated roots, both the sine AND cosine terms need all powers of t up to  $t^{m-1}$ 

#### 8 The Wronskian for Repeated Roots

**Theorem 7** (Linear Independence via Wronskian). The functions  $\{t^k e^{rt} : k = 0, 1, \ldots, m-1\}$  are linearly independent.

Wronskian Calculation. For simplicity, consider m=2 with solutions  $e^{rt}$  and  $te^{rt}$ :

$$W(t) = \begin{vmatrix} e^{rt} & te^{rt} \\ re^{rt} & (1+rt)e^{rt} \end{vmatrix}$$

$$W(t) =$$

$$e^{rt} \cdot (1+rt)e^{rt} - te^{rt} \cdot re^{rt} = e^{2rt}(1+rt-rt) = e^{2rt} \neq 0$$

For general m, the Wronskian is:

$$W(t) =$$

$$e^{mrt} \cdot \prod_{0 \le i < j < m} (j - i) \ne 0$$
  $\square$ 

#### 9 Connection to Jordan Normal Form

 $\label{thm:contraction} The appearance of \it t^k terms in solutions corresponds exactly to the \it Jordan block structure of \it t^k terms in solutions corresponds exactly to the \it Jordan block structure of \it t^k terms in solutions corresponds exactly to the \it Jordan block structure of \it t^k terms in solutions corresponds exactly to the \it Jordan block structure of \it t^k terms in solutions corresponds exactly to the \it Jordan block structure of \it t^k terms in solutions corresponds exactly to the \it Jordan block structure of \it t^k terms in solutions corresponds exactly to the \it Jordan block structure of \it t^k terms in solutions corresponds exactly to the \it Jordan block structure of \it t^k terms in \it solutions corresponds exactly to the \it Jordan block structure of \it t^k terms in \it solutions of \it solutions of \it t^k$ 

**Theorem 8** (Jordan Form and Solutions). For a system  $\mathbf{x}' = A\mathbf{x}$  where A has a Jordan block:

$$J = \begin{pmatrix} \lambda & 1 & 0 & \cdots & 0 \\ 0 & \lambda & 1 & \cdots & 0 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ 0 & 0 & 0 & \lambda & 1 \\ 0 & 0 & 0 & 0 & \lambda \end{pmatrix}$$

The matrix exponential is:

$$e^{Jt} =$$

$$e^{\lambda} t \begin{pmatrix} 1 & t & \frac{t^2}{2!} & \cdots & \frac{t^{m-1}}{(m-1)!} \\ 0 & 1 & t & \cdots & \frac{t^{m-2}}{(m-2)!} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ 0 & 0 & 0 & 1 & t \\ 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

#### 10 Algorithm for Repeated Roots

Method 2 (Complete Solution Construction). 1. Factor the characteristic polynomial completely.

2. For each distinct root  $r_i$  with multiplicity  $m_i$ :

- 3. If  $r_i is real : Include solutions t^k e r_i t for k = 0, 1, ..., m_i 1 If r_i = \alpha + i\beta$  is complex (with conjugate  $\alpha i\beta$ ): Include solutions  $t^k e^{\alpha} t \cos(\beta t)$  and  $t^k e^{\alpha} t \sin(\beta t)$  for  $k = 0, 1, ..., m_i 1$
- 4. Form the general solution as a linear combination of all these functions.

# 11 Worked Examples

**Example 2** (Triple Root). Solve: y''' - 3y'' + 3y' - y = 0 **Solution:** Characteristic equation:  $r^3 - 3r^2 + 3r - 1 = 0$ Recognize this as  $(r-1)^3 = 0$  (expand to verify). Root: r = 1 with multiplicity 3. Solutions:  $e^t$ ,  $te^t$ ,  $t^{2et}$ General solution:  $y(t) = (c_1 + c_2t + c_3t^2)e^t$ 

**Example 3** (Complex Double Roots). Solve:  $y^{(4)} + 8y'' + 16y = 0$  Solution: Characteristic equation:  $r^4 + 8r^2 + 16 = 0$  This is  $(r^2 + 4)^2 = 0$ , giving  $(r - 2i)^2(r + 2i)^2 = 0$  Roots:  $\pm 2i$ , each with multiplicity 2. Real solutions:

• From 2i (mult. 2):  $\cos(2t)$ ,  $t\cos(2t)$ ,  $\sin(2t)$ ,  $t\sin(2t)$ General solution:  $y(t) = (c_1 + c_2t)\cos(2t) + (c_3 + c_4t)\sin(2t)$ 

Prof. Ditkowski often presents the characteristic polynomial already factored. Don't waste time trying to factor what's already factored - just identify multiplicities and write solutions systematically.