

1. CONVERSION OF REGULAR EXPRESSION TO NFA

AIM:

To write a C program to convert the regular expression to NFA.

ALGORITHM:

1. Start the program.
2. Declare all necessary header files.
3. Define the main function.
4. Declare the variables and initialize variables r & c to ‘0’.
5. Use a for loop within another for loop to initialize the matrix for NFA states.
6. Get a regular expression from the user & store it in ‘m’.
7. Obtain the length of the expression using strlen() function and store it in ‘n’.
8. Use for loop upto the string length and follow steps 8 to 12.
9. Use switch case to check each character of the expression
- 10.If case is ‘*’, set the links as ‘E’ or suitable inputs as per rules.
- 11.If case is ‘+’, set the links as ‘E’ or suitable inputs as per rules.
- 12.Check the default case, i.e.,for single alphabet or 2 consecutive alphabets and set the links to respective alphabet.
- 13.End the switch case.
- 14.Use for loop to print the states along the matrix.
- 15.Use a for loop within another for lop and print the value of respective links.
- 16.Print the states start state as ‘0’ and final state.
- 17.End the program.

PROGRAM:

```
#include<stdio.h>
#include<conio.h>

void main()
{
char m[20],t[10][10];
intn,i,j,r=0,c=0;
clrscr();
printf("\n\t\tSIMULATION OF NFA");
printf("\n\t\t*****");
for(i=0;i<10;i++)
{
for(j=0;j<10;j++)
{
t[i][j]=' ';
}
}
printf("\n\nEnter a regular expression:");
scanf("%s",m);
n=strlen(m);
for(i=0;i<n;i++)
{
switch(m[i])
{
case '|': {
    t[r][r+1]='E';
    t[r+1][r+2]=m[i-1];
    t[r+2][r+5]='E';

    t[r][r+3]='E';
    t[r+4][r+5]='E';
    t[r+3][r+4]=m[i+1];
    r=r+5;
    break;
}
case '*':{
```

```

t[r-1][r]='E';
t[r][r+1]='E';
t[r][r+3]='E';
t[r+1][r+2]=m[i-1];
t[r+2][r+1]='E';
t[r+2][r+3]='E';
r=r+3;
break;
}
case '+': {
t[r][r+1]=m[i-1];
t[r+1][r]='E';
r=r+1;
break;
}
default:
{
    if(c==0)
    {
        if((isalpha(m[i]))&&(isalpha(m[i+1])))
        {
            t[r][r+1]=m[i];
            t[r+1][r+2]=m[i+1];
            r=r+2;
            c=1;
        }
        c=1;
    }
    else if(c==1)
    {
        if(isalpha(m[i+1]))
        {
            t[r][r+1]=m[i+1];
            r=r+1;
            c=2;
        }
    }
    else
    {

```

```

        if(isalpha(m[i+1]))
        {
            t[r][r+1]=m[i+1];
            r=r+1;
            c=3;
        }
    }
break;
}
}
printf("\n");
for(j=0;j<=r;j++)
printf(" %d",j);
printf("\n");
printf("\n");
for(i=0;i<=r;i++)
{
for(j=0;j<=r;j++)
{
printf(" %c",t[i][j]);
}
printf(" | %d",i);
printf("\n");
}
printf("\nStart state: 0\nFinal state: %d",i-1);
getch();
}

```

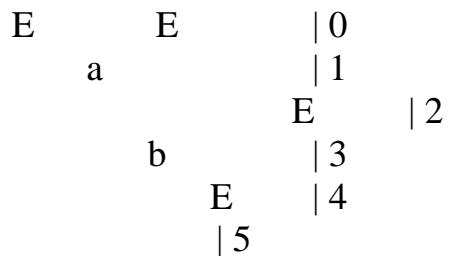
OUTPUT:

Enter a regular Expression: a|b

SIMULATION OF NFA *****

Enter a regular expression:a|b

0 1 2 3 4 5



Start state: 0

Final state: 5

RESULT:

Thus the C program to convert regular expression to NFA has been executed and the output has been verified successfully.

EX.NO : 3

Converting NFA to DFA

AIM: To write a program to convert NFA to DFA

ALGORITHM:

1. Start the program
2. Assign an input string terminated by end of file, DFA with start
3. The final state is assigned to F
4. Assign the state to S
5. Assign the input string to variable C
6. While C!=e of do
 S=move(s,c)
 C=next char
7. If it is in ten return yes else no
8. Stop the program

SOURCE CODE:

```
#include<conio.h>
#include<string.h>
#include<process.h>
#include<math.h>
Int n[11],I,j,c,k,h,l,h1,f,h2,temp1[12],temp2[12],count=0,ptr=0;
Char a[20][20],s[5][8];
Intr[5][2],ecl[5][8],str[5],flag;
Intr[5][2],ecl[5][8],st[5],flag;
Void ecls(int b[10],int x)
{
I=0;
K=-1;flag=0;
While(l<x)
```

```

{
N[++k]=b[l];
I=b[l];
h=k+1;
a:
for(j=I;j<=11;j++)
{
If(a[i][j]=='e')
{n[++k]=j;
}
If(j==11&&h<=k)
{
I=n[h];
H++;
Goto a;
}}
L++;
}for(i=0;i<k;i++)
for(j=i+1;j<k;j++)
if(n[i]>n[j])
{
C=n[i];
N[i]=n[j];
N[j]=c;
}for(i=0;i<ptr;i++)
for(j=0;j<k;j++)
{
If(ecl[i][j]!=n[j])
{
If(i<count)
{i++;
J=0;
}
Else
Goto b;
}
Else if((ecl[i][j]==n[j])&&(j==k))
{tr[ptr][f]=st[i];
Flag=1;
Break;
}}
B:
If(flag==0)
{for(i=0;i<=k;i++)
Ecl[count][i]=n[i];
st[count]=count+65;
}

```

```

tr[ptr][f]=st[count];
count++;
}
}
Void mova(int g)
{h1=0;
for(i=0;i<7;i++)
{if(ecl[g][i]==3)
Temp1[h1++]=4;
if(ecl[g][i]==8)
Temp1[h1++]=9;
}
Printf("\n move(%c,a):",st[g]);
For(i=0;i<h1;i++)
Printf("%d",temp1[i]);
F=0;
Ecls(temp1,h1);
}
Void movb(int g)
{
H2=0;
For(i=0;i<7;i++)
{
If(ecl[g][i]==5)
Temp2[h2++]=6;
If(ecl[g][i]==9)
Temp2[h2++]=10;
If(ecl[g][i]==10)
Temp2[h2++]=11;
Printf("move(%c,b):",st[g]);
For(i=0;i<h2;i++)
Printf("%d",temp2[i]);
F=1;
Ecls(temp2,h2);
}
Void main()
{
Clrscr();
Printf("\n the no. of states in NFA (a/b)*abb are:11");
For(i=0;i<=11;i++)
For(j=0;j<=11;j++)
A[i][j]='\0';
A[1][2]='e';
A[1][8]='e';
A[2][3]='e';
A[2][5]='e';
A[3][4]='a';

```

```

A[5][6]=’b’;
A[4][7]=’e’;
A[6][7]=’e’;
A[7][8]=’e’;
A[7][2]=’e’;
A[8][9]=’a’;
A[9][10]=’b’;
A[10][11]=’b’;
Printf(“\n the transmission table is as Follows”);
Printf(“\n states 1 2 3 4 5 6 7 8 9 10 11”);
Getch();
For(i=1;i<=11;i++)
{
Printf(“\n %d \t”,i);
For(j=1;j<=11;j++)
Printf(“%c”,a[i][j]);
}
Getch();
Printf(“\n \n press any key to continue”);
Clrscr();
I=1;k=1;h=1;
N[0]=I;
Printf(“\n”);
A:
For(j=1;j<=11;j++)
{if(a[i][j]==’e’)
{
N[k++]=j;
}
If(j==11&&h<k)
{
I=n[h];
H++;
Goto a;
}}
For(i=1;j<k;i++)
For(j=i+1;j<k;j++)
If(n[i]>n[j])
{c=n[i];
N[i]=n[j];
N[j]=c;
}
Count++;
St[0]=65;
For(i=0;i<k;i++)
Ecl[0][i]=n[i];

```

```

Printf("the moves are of the Following manner");
Mova(ptr);
Movb(ptr);
Ptr++;
}
Printf("\n the new states of DFA are as Follows");
For(i=0;i<5;i++)
{printf("\n %c",st[i]);
For(j=0;j<7;j++)
Printf("%d",ecl[i][j]);
}
Printf(" the transition table are as Follows");
Printf("\n a \n b \n");
For(i=0;i<5;i++)
{
Printf("%c",st[i]);
For(j=0;j<2;j++)
Printf("%c \t",tr[i][j]);
}
Getch();
}

```

OUTPUT:

The no. of states in NFA(a/b)*abb are:11

The transition table is as Follows

	1	2	3	4	5	6	7	8	9	10	11
1	e		e								
2		e		e							
3		a									
4			e								
5			b	e							
6				e							
7		e			e						
8											
9					a						
10					b						
11					b						

RESULT:

Thus the above conversion program is successfully executed

EX.NO : 4

Computation of FIRST and FOLLOW sets

AIM: To calculate the first and Follow of the given expression

ALGORITHM:

1. Start the program
2. In the production the first terminal on R.H.S becomes the first of it
3. If the first character is non-terminal then its first is taken else Follow of left is taken
4. To find Follow find where all the non terminals appear. the first of its Follows is its Follow
5. If the Follow is t then the Follow of left is taken
6. Finally print first and its Follow
7. Stop the program

SOURCE CODE:

```
#include<stdio.h>
#include<conio.h>
#include<process.h>
Void main()
{
Intnop,x=0,y=0,l=0,k=0,c=0,s=0,z=0;
Char p[10][10],o,fi[10][10],fo[10][10];
Clrscr();
For(x=0;x<10;x++)
{
For(y=0;y<10;y++)
{
P[x][y]=’0’;
Fi[x][y]=’0’;
Fo[x][y]=’0’;
```

```

    }
    Printf("enter the no of productions:");
    Scanf("%d",&nop);
    Printf("\n enter the number of production");
    For(x=0;x<nop;x++)
    {
        Scanf("%s",&p[x]);
        If(p[x][0]==p[x][2])
        {
            Printf("\a production is not free from left recursion");
            Printf("\a program has to be terminated");
            Printf("\a press any key");
            Getch();
            Exit(0);
        }
        Printf("\n\n")
        Printf("first \n");
        For(y=0;y<nop;y++)
        {
            Printf("first(%c)=",p[y][0]);
            If(p[y][2]>='A' &&p[y][2]<='z')
            {
                O=p[y][2];
                For(x=y+1;x<nop;x++)
                {
                    If(p[x][0]==o)
                    {
                        If(p[x][2]>='A'&&p[x][2]<='z')
                        O=p[x][2];
                        Else if(p[x][2]<'A'||p[x][2]>'z')
                        {
                            Printf("%c",p[x][2]);
                            Fi[y][k++]=p[x][2];
                            For(l=0;l<strlen(p[x]);l++)
                            {
                                If(p[x][l]== '/')
                                {
                                    Printf("%c",p[x][l+1]);
                                    Fi[y][k++]=p[x][l+1];
                                    Break;
                                }
                            }
                        }
                    }
                }
            }
        }
    }
}
Else if(p[y][2]<'A' || p[y][2]>'Z')
{
    Printf("%c",p[y][2]);
    Fi[y][k++]=p[y][2];
}

```

```

L=strlen(p[y]);
For(c=0;c<l;c++)
{
If(p[y][c]== '/')
{
Printf("%c",p[y][c+1]);
Fi[y][k++]=p[y][c+1];
}
Printf("\n");
K=0;
Printf("Follow \n");
For(x=0;x<nop;x++)
{
For(y=0;y<nop;y++)
For(l=2;l<strlen(p[x]);l++)
If(p[y][l]==p[x][0])
{
If(p[y][l+1]<'A'||p[y][l+1]>'Z'||p[y][l+1]!='/'||p[y][l+1]!='0')
{
Fo[x][k++]=p[y][l+1];
If(x==0)
{
Fo[x][k++]='$';
}
C=k;
If(p[y][l+1]=='0'||p[y][l+1]=='/')
{
For(s=0;s<c+10;s++)
{
Fo[x][k++]=fo[x-1][s];
}
C=k;
If(p[y][l+1]>='A'&&p[y][l+1]<='z')
{
For(s=0;s<=c;s++)
{
Fo[x][k++]=fo[x-1][s];
Fo[x][k++]=fo[x-2][s];
Fo[x][k++]=fi[x-1][s];
}
C=k;
}
Printf("Follow(%c)=",p[x][0]);
For(z=0;z<=k+10;z++)
If(fo[x][z]=='*'||fo[x][z]=='+'||fo[x][z]=='$'||fo[x][z]==')'||fo[x][z]==('||fo[x][z]==-
'||fo[x][z]=='%')

```

```
Printf("%c",fo[x][z]);
K=0;
Printf("\n");
}
Getch();
}
```

OUTPUT:

Enter the no. of production:5

E->TE'
E'->+TE'/e
T->FT'
T'->*FT'/e
F->(E) /id

First

First(E)={(.id}
First(E')={+,e}
First(T)={(.id}
First(T')={*,e}
First(F)={(.id}

Follow

Follow(E)={),\$}
Follow(E')={),+\$}
Follow(T)={),+\$}
Follow(T')={*,+,),\$}
Follow(F)={*,+,),\$}

RESULT:

Thus the above computation of FIRST & FOLLOW program is successfully executed.

EX.NO :6

Construction of Predictive Parsing Table

AIM: To write a C program for the implementation of predictive parsing table

ALGORITHM:

1. start the program
2. Assign an input string and parsing table in for then G.
3. Set ip to point to the first symbol of to \$
4. Repeat if X is a terminal of \$ then if n=a,then pop X from the stack
5. Push Y into the stack with Y,on top
6. Output the production x->x,y
7. End else error until x=\$
8. Terminate the program

SOURCE CODE:

```
#include<stdio.h>
#include<conio.h>
#include<string.h>
char str[10],out,in,output[10],input[10],temp;
char tl[10]={‘x’,’+’,’*’,’(’,’)’,$’,@’};
char ntl[10]={‘e’,’e’,’t’,’t’,’f’};
int err=0,flag=0,i,j,k,l,m;
char
c[10][10][7]={{{{“te”},{“error!”},{{“error!”},{“te”},{“error!”}},{{“error!”},{“error!”}},{{“error!”},“te”,“error!”},{{“error!”},“@”,“@”}},{{{“ft”},“error!”},{{“error!”},“ft”,“error!”},{{“error!”},“ft”,“error!”},{{“error!”},“@”,“*ft”,“error!”},{{“error!”},“@”,“@”}},{{{“x”},“error!”},{{“error!”},“(e)”},{{“error!”},“error!”},{{“error!”},“error!”}}};
struct stack
```

```

{
char sic[10];
int top;
};
void push(struct stack *s,char p)
{
s->sic[++s->top]=p;
s->sic[s->top+1]='\0';
}
char pop(struct stack *s)
{
char a;
a=s->sic[s->top];
s->sic[s->top--]='\0';
return(a);
}
char stop(struct stack *s)
{
return(s->sic[s->top]);
}
voidpobo(struct stack *s)
{
m=o;
while(str[m]!='\0')
m++;
m--;
while(m!=-1)
{
if(str[m]!='@')
push(s,str[m]);
m--;
}}
void search(int l)
{
for(k=0;k<7;k++)
if(in==tl[k])
break;
if(l==0)
strcpy(str,c[l][k]);
else if(l==1)
strcpy(str,c[l][k]);
else if(l==2)
strcpy(str,c[l][k]);
else if(l==3)
strcpy(str,c[l][k]);
else strcpy(str,c[l][k]);}

```

```

void main()
{
struct stack s1;
struct stack *s;
s=&s1;
s->top=-1;
clrscr();
printf("\t\t parsing table \t\t");
for(i=0;i<5;i++)
{
printf("%c\t",ntl[i]);
for(j=0;j<6;j++)
if(strcmp(c[i][j],"error!")==0)
printf("error!\t");
else
printf("%c->%s" \t",ntl[i],c[i][j]);
}
push(s,'$');
push(s,'e');
printf("enter the input string");
scanf("%s",input);

printf("\n\n the behaviour of the parser for given input string is: \n\ ");
printf("\n stack\n input\n output");
i=0;
in=input[i];
printf("%s\t",s->sic);
for(k=i;k<strlen(input);k++)
printf("%c",input[k]);
if(strcmp(str,' ')!=0)
printf("\t%c->%s" ntl[j],str);
while((s->sic[s->top]!='$')&&err!=1&&strcmp(str,"error")!=0)
{
strcpy(str," ");
flag=0;
for(j=0;j<7;j++)
if(in==tl[j])
{
flag=1;
break;
}
if(flag==0)
in='x';
flag=0;
out=stop(s);
for(j=0;j<7;j++)

```

```

if(out==tl[j])
{
flag=1;
break;
}
if(flag==1)
{
if(out==in)
{
temp=pop(s);
in=input[++i];
if(str=='@')
temp=pop(s);
}
else
{
strcpy(sstr,"error!");
err=1;
}}
else
{
flag=0;
for(j=0;j<5;j++)
if(out==ntl[j])
{
flag=1;
break;
}
if(flag==1)
{
search(j);
temp=pop(s);
pobo(s);
}
else
{
strcpy(str,"error!");
err=1;
}}
if(strcmp(str,"error!")!=0)
{
printf("%s\t",s->sic);
for(k=i;k<strlen(input);k++)
printf("%c",input[k]);
if((strcmp(str," ")!=0)&&(strcmp(str,"error")!=0))
printf("\t %c->%s",ntl[j],str);
}

```

```

    }
if(strcmp(str,"error")==0)
printf("\n the string is not accepted!!");
else
printf("\t \t accept \n\n the string is accepted");
getch();
}

```

OUTPUT:

Parsing table

X	+	*	()	\$
E.E->Te	ERROR!	ERROR!	E->te	ERROR!	ERROR!
E	ERROR!	E->+te	ERROR!	ERROR!	E->@ e->@
T	T->Ft	ERROR!	ERROR!	T->ft	ERROR!
T	ERROR!	T->@	t->*ft	ERROR	t->@ t->@
F.F->x	ERROR!	ERROR!	F->	(E)	ERROR! ERROR!

Enter the input string: x+X\$

The behaviour of the parser for given input string is

Stack	input	output
SE	X+X\$	
SeT	X+X\$	E->Te
Set F	X+X\$	T->Ft
Set X	X+X\$	F->X
Set	+X\$	
Se	+X\$	t->@
SeT+	+X\$	e->+Te
seT	X\$	
SetF	X\$	T->Ft

RESULT:

Thus the Predictive Parser program is executed successfully.

EX.NO : 5

Implementation of Shift Reduce Parsing

AIM: To write a C program for shift reduce parsing

ALGORITHM:

1. start the program
2. read the expression and declare the variables
3. set \$ symbol to indicate the start of stack
4. Repeat for i=0 to n where n is the no. of productions
 - A. read prod(i)
- B.set problem[i]=strlen(prod[i])
5. check for the non-terminal which corresponds to the terminal
6. If it equals then replace the terminal with non-terminal
7. Repeat this until terminals are replaced by non-terminals

SOURCE CODE:

```
#include<iostream.h>
#include<conio.h>
#include<string.h>
#include<stdlib.h>
void push(char);
char pop();
void printstack();
struct grammar
{
    char lpr,rpr[10];
};char stack[20];
int top=-1;
void main()
```

```

{
grammar gr[10];
char buffer[10];
char ch,ch1,temp[10],start;
int i,j,k,s,t,len;
clrscr();
cout<<"enter the no of productions:";
cin>>n;
for(i=0;i<n;i++)
{
cout<<"\n enter the left side of productions"<<i+1<<":";
cin>>gr[i].lpr;
cout<<"\n enter the right side of productions"<<":";
cin>>gr[i].rpr;
}
cout<<"\n enter the input string:";
cin>>buffer;
cout<<"\n the grammar is:\n";
for(i=0;i<n;i++)
cout<<gr[i].lpr<<"-->"<<gr[i].rpr<<endl;
cout<<"\n input string is :"<<buffer;
push('$');
start=gr[0].lpr;
len=strlen(buffer);
buffer[len]='$';
buffer[len+1]='\0';
cout<<"\n \n stack \t\t buffer\t\t action\n";
cout<<stack<<"\t\t"<<buffer<<endl;
getch();
while(1)
{
ch=buffer[i];
lab:t=0;
for(k=top;k>0;k--)
{
temp[t++]=stack[t]='\0';
strrev(temp);
for(j=0;j<n;j++)
{
if(strcmp(temp,gr[j].rpr)==0)
{
for(s=0;s<t;s++)
ch1=pop();
push(gr[j].lpr);
printstack();
cout<<"\t\t" <<&buffer[i]<<"\t\tReduce"<<endl;
}
}
}
}
}

```

```

getch();
goto lab;
}
}
strrev(temp);
}
ch1=pop();
if(ch!='$')
{push(ch1);
push(ch);
printstack();
cout<<"\t\t\t"<<&buffer[i+1]<<"\t\tshift"<<endl;
getch();
i++;
}
else if(ch=='$' && ch1==start && top==0)
{
cout<<"\n string is accepted";
getch();
exit(0);
}
else
{
cout<<"\n string is not accepted";
getch();
exit(0);
}
}
}
}

void push(char a)
{
stack[++top]=a;
}
char pop(){
return stack[top--];
}
void printstack()
for(int i=0;i<top;i++)
cout<<stack[i];
}

```

OUTPUT:

Enter the no. of production: 3
 Enter the production: S->aABC
 A->Abc/b
 B->d

The production are

S>>aABC

A>>Abc/b

B>>d

ILR

S>>aABC

u>>bc/bu

B>>d

u>x

the first symbols are X

RESULT:

Thus the shift reduce parser program is successfully executed

Computation of Follow and Trailing Sets

AIM: To write a C program to Compute of Follow and Trailing Sets

ALGORITHM:

Step1 : Start the Program

Step2: Get the no of production and calculate the length of each production.

Step3: With a variable val for checking the valid non terminals if they are duplicated get all Non terminals in an array.

Step4: In each production check the first accurate of terminals and take that terminal and add it to Follow of non terminal in array and exit the loop.

Step5: Scan the production and find the last terminal and add it to respective trailing array of associated non terminal &exit the loop.

Step6: Consider production with a non terminal on right side and check the Follow of that non terminal associated production and also trailing and add it to the Follow and trailing of left side non terminal.

Step7: Write the Follow and trailing terminals for each non terminal.

SOURCE CODE:

```
#include<iostream.h>
#include<stdio.h>
#include<conio.h>
#include<string.h>
#include<ctype.h>
```

```

char nt[5],p[5],q[5],a[5][5][5],b[5][5][5],fi[5][5],left[5],right[5],lead[5][10],trail[5][10];
int n1,n[5],c[5][5],m,l[5],f[5],k,a1;

void leading(char,int);
void trailing(char,int);

void main()
{
clrscr();
cout<<"Enter the number of non-terminals ";
cin>>n1;
cout<<"Enter the set of non-terminals ";
for(int i=0;i<n1;i++)
    cin>>nt[i];
for(i=0;i<n1;i++)
{
    cout<<endl;
cout<<"Enter the number of productions for "<<nt[i]<<" ";
cin>>n[i];
    for(int j=0;j<n[i];j++)
{
        cout<<"\nEnter the productions ";
        gets(p);
        //q=strrev(p);
        c[i][j]=strlen(p);
        for(int k=0;k<c[i][j];k++)
        {
            a[i][j][k]=p[k];
            b[i][j][k]=p[c[i][j]-k-1];
        }
    }
}
for(i=0;i<n1;i++)
{
    cout<<endl;
    cout<<nt[i]<<"--->";
    for(int j=0;j<n[i];j++)
    {
        for(int k=0;k<c[i][j];k++)
        cout<<a[i][j][k];
        cout<<"/";
    }
}
cout<<"\n\n";
char x;
for(i=0;i<n1;i++)
{

```

```

l[i]=0;
f[i]=0;
x=nt[i];
k=i;
leading(x,k);
trailing(x,k);
}
/*for(intmn=0;mn<=n1;mn++)
{
cout<<right[mn]<<" ";
} */
int count=0;
char z;
for(intmn=0;mn<n1;mn++)
{
for(i=0;i<n1;i++)
{
z=right[i];
for(int k=0;k<n1;k++)
{
if(z==nt[k])
{
for(int d=0;d<l[k];d++)
{
for(int g=0;g<l[i];g++)
{
if(lead[i][g]==lead[k][d])
count=1;
}
if(count==0)
{
lead[i][l[i]-1]=lead[k][d];
l[i]++;
}
}
count=0;
}
for(d=0;d<f[k];d++)
{
for(int g=0;g<f[i];g++)
{
if(trail[i][g]==trail[k][d])
count=1;
}
if(count==0)
{
trail[i][f[i]-1]=trail[k][d];
}
}
}
}
}

```

```

        f[i]++;
    }
    count=0;
}

}
count=0;
}

for(i=0;i<n1;i++)
{
cout<<"Leading("<<nt[i]<<")= { ";
for(int j=0;j<l[i];j++)
{
    cout<<lead[i][j]<<" ";
}

cout<<" }\t\t\t\t";

cout<<"Trailing("<<nt[i]<<")= { ";
for(j=0;j<f[i];j++)
{
    cout<<trail[i][j]<<" ";
}

cout<<" }\n";
}

getch();
}

void leading(char x,int i)
{
char z;
for(int j=0;j<n[i];j++)
{
if(isupper(a[i][j][0]))
{
left[i]=nt[i];
right[i]=a[i][j][0];
if(!isupper(a[i][j][1]))
{
//cout<<a[i][j][1]<<" ";
lead[i][l[i]]=a[i][j][1];
l[i]++;
}
}
}

```

```
    }
}
else
{
    //cout<<a[i][j][0]<<" ";
lead[i][l[i]]=a[i][j][0];
l[i]++;
}
}
```

```
void trailing(char x,int i)
{
char z;
for(int j=0;j<n[i];j++)
{
if(isupper(b[i][j][0]))
{
if(!isupper(b[i][j][1]))
{
//cout<<a[i][j][1]<<" ";
trail[i][f[i]]=b[i][j][1];
f[i]++;
}
}
else
{
//cout<<a[i][j][0]<<" ";
trail[i][f[i]]=b[i][j][0];
f[i]++;
}
}
}
```

OUTPUT-

Enter the no. of production:5

E->TE'

$$E' \rightarrow +TE'/e$$

T->FT'

$$T' \rightarrow^* FT'/e$$

F->(E) /id

Leading

Leading (E)={,(,id}
Leading (E')={+,e}
Leading (T)={,(,id}
Leading (T')={*,e}
Leading (F)={,(,id}

Trailing
Trailing (E)={},\$}
Trailing (E')={} ,+,\$}
Trailing (T)={} ,+,\$}
Trailing (T')={} *,+,),\$}
F Trailing (F)={} *,+,),\$}

RESULT: Thus the Leading and Trailing was executed and verified Successfully

Computation of LR (0) items

AIM:To write a code for LR(0) Parser for Following Production:

E->E+T
T->T*T/F/F
F->(E)/char

ALGORITHM:

1. Initialize the stack with the start state.
2. Read an input symbol
3. while true do
 - 3.1 Using the top of the stack and the input symbol determine the next state.
 - 3.2 If the next state is a stack state
then
 - 3.2.1 stack the state
 - 3.2.2 get the next input symbol
 - 3.3 else if the next state is a reduce state
then
 - 3.3.1 output reduction number, k
 - 3.3.2 pop RHSk -1 states from the stack where RHSk is the right hand side of production k.
 - 3.3.3 set the next input symbol to the LHSk
 - 3.4 else if the next state is an accept state
then
 - 3.4.1 output valid sentence
 - 3.4.2 return
 - else
 - 3.4.3 output invalid sentence
 - 3.4.4 return

SOURCE CODE:

```
#include<string.h>
#include<conio.h>
#include<stdio.h>
```

```
intaxn[][][6][2]={
```



```

{
    if(btop>=0)
        b[btop--]='\0';
}

void display()
{
    for(i=0;i<=top;i++)
        printf("%d%c",a[i],b[i]);
}

void display1(char p[],int m)
{
    int l;
    printf("\t\t");
    for(l=m;p[l]!='\0';l++)
        printf("%c",p[l]);
    printf("\n");
}

void error()
{
    printf("\n\nSyntax Error");
}

void reduce(int p)
{
    intlen,k,ad;
    charsrc,*dest;
    switch(p)
    {
        case 1:dest="E+T";
            src='E';
            break;
        case 2:dest="T";
            src='E';
            break;
        case 3:dest="T*F";
            src='T';
            break;
        case 4:dest="F";
            src='T';
}

```

```

        break;
case 5:dest="(E)";
    src='F';
    break;
case 6:dest="i";
    src='F';
    break;
default:dest="\0";
    src='\0';
    break;
}
for(k=0;k<strlen(dest);k++)
{
    pop();
    popb();
}
pushb(src);
switch(src)
{
    case 'E': ad=0;
        break;
    case 'T': ad=1;
        break;
    case 'F': ad=2;
        break;
    default: ad=-1;
        break;
}
push(gotot[TOS()][ad]);
}
int main()
{
    int j,st,ic;
    char ip[20] = "\0", an;
    clrscr();
    printf("Enter any String :- ");
    gets(ip);
    push(0);
    display();
    printf("\t%s\n", ip);
}

```

```

for(j=0;ip[j]!='\0';)
{
    st=TOS();
    an=ip[j];
    if(an>='a'&an<='z')
        ic=0;
    else if(an=='+')
        ic=1;
    else if(an=='*')
        ic=2;
    else if(an=='(')
        ic=3;
    else if(an==')')
        ic=4;
    else if(an=='$')
        ic=5;
    else
    {
        error();
        break;
    }
    if(axn[st][ic][0]==100)
    {
        pushb(an);
        push(axn[st][ic][1]);
        display();
        j++;
        display1(ip,j);
    }
    if(axn[st][ic][0]==101)
    {
        reduce(axn[st][ic][1]);
        display();
        display1(ip,j);
    }
    if(axn[st][ic][1]==102)
    {
        printf("Given String is Accepted");
        break;
    }
}

```

```
    }  
    getch();  
    return 0;  
}
```

OUTPUT

Enter any String :- a+b*c

0	a+b*c
0a5	+b*c
0F3	+b*c
0T2	+b*c
0E1	+b*c
0E1+6	b*c
0E1+6b5	*c
0E1+6F3	*c
0E1+6T9	*c
0E1+6T9*7	c
0E1+6T9*7c5	

RESULT: Thus the LR(0) Program was executed and verified Successfully.

Construction of DAG

AIM: To write a C program to perform the

ALGORITHM:

RESULT: Thus the DAG was executed and verified Successfully.

EX.NO : 14

Intermediate Code Generation

AIM: To write a C program to implementation of code generation

ALGORITHM:

step 1: Start.

Step 2: Enter the three address codes.

Step 3: If the code constitutes only memory operands they are moved to the register and according to the operation the corresponding assembly code is generated.

Step 4: If the code constitutes immediate operands then the code will have a # symbol proceeding the number in code.

Step 5: If the operand or three address code involve pointers then the code generated will constitute pointer register. This content may be stored to other location or vice versa.

Step 6: Appropriate functions and other relevant display statements are executed.

Step 7: Stop.

SOURCE CODE:

```
#include<stdio.h>
#include<string.h>
voidpm();
voidplus();
voiddiv();
inti,ch,j,l,addr=100;
char ex[10],exp[10],exp1[10],exp2[10],id1[5],op[5],id2[5];
void main()
{
clrscr();
while(1)
{
printf("\n1.assignment\n2.arithmetic\n3.relational\n4.Exit\nEnter the choice:");
scanf("%d",&ch);
switch(ch)
{
case 1:
printf("\nEnter the expression with assignment operator:");
scanf("%s",exp);
l=strlen(exp);
exp2[0]='\0';
i=0;
while(exp[i]!='=')
{
i++;
}
strncat(exp2,exp,i);
strupr(exp);
exp1[0]='\0';
strncat(exp1,exp,l-(i+1));
```

```
strrev(exp1);
printf("Three address code:\n temp=%s\n s=temp\n",exp1,exp2);
break;
```

case 2:

```
printf("\nEnter the expression with arithmetic operator:");
scanf("%s",ex);
strcpy(exp,ex);
l=strlen(exp);
exp1[0]='\0';

for(i=0;i<l;i++)
{
if(exp[i]=='+||exp[i]=='-')
{
if(exp[i+2]=='/||exp[i+2]=='*')
{
pm();
break;
}
else
{
plus();
break;
}
}
else if(exp[i]=='/||exp[i]=='*)
{
div();
break;
}
}
break;
```

case 3:

```
printf("Enter the expression with relational operator");
scanf("%s%s%s",&id1,&op,&id2);
if(((strcmp(op,<)==0)||strcmp(op,>)==0)||strcmp(op,<")==0)||strcmp(op,>")==0)||strcmp(op,"==")==0)||strcmp(op,"!=")==0)
printf("Expression is error");
else
{
printf("\n%d\tif %s%s%sgoto %d",addr,id1,op,id2,addr+3);
addr++;
printf("\n%d\tT:=0",addr);
addr++;
```

```

printf("\n%d\t goto %d",addr,addr+2);
addr++;
printf("\n%d\t T:=1",addr);
}
break;
case 4:
exit(0);
}
}
}
void pm()
{
strrev(exp);
j=l-i-1;
strncat(exp1,exp,j);
strrev(exp1);
printf("Three address code:\ntemp=%s\ntemp1=%c%ctemp\n",exp1,exp[j+1],exp[j]);
}
void div()
{
strncat(exp1,exp,i+2);
printf("Three address code:\ntemp=%s\ntemp1=temp%c%c\n",exp1,exp[i+2],exp[i+3]);
}
void plus()
{
strncat(exp1,exp,i+2);
printf("Three address code:\ntemp=%s\ntemp1=temp%c%c\n",exp1,exp[i+2],exp[i+3]);
}

```

OUTPUT :

Example Generation of Three Address Project Output Result

1. assignment
2. arithmetic
3. relational
4. Exit

Enter the choice:1

Enter the expression with assignment operator:

a=b

Three address code:

temp=b

a=temp

1.assignment

2.arithmetic

3.relational

4.Exit

Enter the choice:2

Enter the expression with arithmetic operator:

a+b-c

Three address code:

temp=a+b

temp1=temp-c

1.assignment

2.arithmetic

3.relational

4.Exit

Enter the choice:2

Enter the expression with arithmetic operator:

a-b/c

Three address code:

temp=b/c

temp1=a-temp

1.assignment

2.arithmetic

3.relational

4.Exit

Enter the choice:2

Enter the expression with arithmetic operator:

a*b-c

Three address code:

temp=a*b

temp1=temp-c

1.assignment

2.arithmetic

3.relational

4.Exit

Enter the choice:2

Enter the expression with arithmetic operator:a/b*c

Three address code:

temp=a/b

temp1=temp*c

1.assignment

2.arithmetic

3.relational

4.Exit

Enter the choice:3

Enter the expression with relational operator

a

<=

b

100 if a<=b goto 103

101 T:=0

102 goto 104

103 T:=1

1.assignment

2.arithmetic

3.relational

4.Exit

Enter the choice:4

RESULT: Thus the Generation of Three Address was executed and verified Successfully.

EX.NO.:15

OPERATOR PRECEDENCE

AIM:

To write a C program to implement the concept of operator precedence.

ALGORITHM:

1. Start the program.
2. Include the required header files and start the declaration of main method.
3. Declare the required variable and define the function for pushing and popping the characters.
4. The operators are displayed in column and row wise and stored it in a queue.
5. Using a switch case find the values of the operators.

6. Display the precedence of the operator and generate the code for precedence of operator for the given expression.
7. Compile and execute the program for the output.
8. Stop the program

PROGRAM:

```
#include<stdio.h>
#include<conio.h>
#include<string.h>
#include<ctype.h>
#include<stdlib.h>
charstr[20],stk[20],pstk[20];
inttos=-1,flag=0,ptr=0,rm=-1,i,j;
char
q[9][9]={{'>', '>', '<', '<', '<', '<', '<', '<', '>'}, {'>', '>', '<', '<', '<', '<', '<', '<', '>'}, {'>', '>', '>', '>', '>', '>', '<', '<', '<'},
{'>', '>', '>', '>', '<', '<', '<', '<', '>'}, {'>', '>', '<', '<', '<', '<', '<', '<', '>'}, {'<', '<', '<', '<', '<', '<', '<', '<', '='},
{'>', '>', '>', '>', 'E', '>', 'E', '>'}, {'>', '>', '>', '>', 'E', '>', 'E', '>'}, {'<', '<', '<', '<', '<', 'E', '<', 'A'}};
char c[9]={'+', '-', '*', '/', '^', 'a', '(', ')', '$'};
voidpushin(char a)
{
tos++;
stk[tos]=a;
}
char popout()
{
char a;
a=stk[tos];
tos--;
return(a);
}
int find(char a)
{
switch(a)
{
case '+':return 0;
case '-':return 1;
case '*':return 2;
case '/':return 3;
case '^':return 4;
case '(':return 5;
case ')':return 6;
case 'a':return 7;
}}
```

```

case'$':return 8;
}
return-1;
}
void display(char a)
{
printf("\n SHIFT %c",a);
}
void display1(char a)
{
if(a!='(')
{
if(isalpha(a))
printf("\n REDUCE E--> %c",a);
else if(a==')')
printf("\n REDUCE E-->(E)");
else
printf("\n REDUCE E-->E %c E",a);
}}
intrel(char a,charb,char d)
{
if(isalpha (a))
a='a';
if(isalpha(b))
b='a';
if(q[find(a)][find(b)]==d)
return 1;
else
return 0;
}
void main()
{
clrscr();
printf("\n\n\t The productions used are:\n\t");
printf("E-->E*E/E+E/E^E/E*E/E-E\n\tE-->E/E \n\tE-->a/b/c/d/e.../z");
printf("\n\tEnter an expression that terminals with$:");
fflush(stdin);
i=-1;
while(str[i]!='$')
{
i++;
scanf("%c",&str[i]);
}
for(j=0;j<i;j++)
{
if((str[j]=='('||(str[j]==')')||(str[j+1]=='('||(str[j+1]=='))'))

```

```

{
}
else
if(((isalpha(str[j])==0)&&(isalpha(str[j+1])==0))||(isalpha(str[j])!=0)&&(isalpha(str[j+1])!=0)))
{
printf("ERROR");
getch();
exit(0);
}
if(((isalpha(str[0]))!=0)|(str[0]=='(')&&(((isalpha(str[i-1]))!=0)|(str[i-1]==')')))
{
pushin('$');
printf("\n\n\n\t+\t*\t^\\ta\t(\t)\t$\n\n");for(i=0;i<9;i++)
{
printf("%c",c[i]);for(j=0;j<9;j++)
printf("\t%c",q[i][j]);printf("\n");}getch();while(1){if(str[ptr]=='$'
&&stk[tos]=='$'){printf("\n\n\t ACCEPT!");break;}else
if(rel(stk[tos],str[ptr], '<')||rel(stk[tos],str[ptr], '==')){display(str[ptr]);pushin(str[ptr]);ptr++;}el
se
if(rel(stk[tos],str[ptr], '>')){do{rm++;pstk[rm]=popout();display1(pstk[rm]);}while(!rel(stk[to
s],pstk[rm], '<'));}
else{printf("\n\n\t NOT ACCEPTED!!!!!!");}
getch();
exit(1);
}
getch();
}
else {
printf("ERROR");
getch();
}
}

```

OUTPUT:

The productions used are:

E-->E*E/E+E/E^E/E*E/E-E

$E \rightarrow E/E$

E-->a/b/c/d/e.../z

Enter an expression that terminals with \$: a+b*c\$

ENTER THE INFORMATION THAT APPLIES WITH YOURS

+ - * / ^ a () \$

+ >><<<<○○○
- >><<<<○○○

```
*    >>>><<>><>
/    >>>><<>><>
^    >><<<>><>
a    <<<<<    =    <    E
(    >>>>    E    >    E    >
)    >>>>    E    >    E    >
$    <<<<<    E    <    A
```

SHIFT a

REDUCE E--> a

SHIFT +

SHIFT b

REDUCE E--> b

SHIFT *

SHIFT c

REDUCE E--> c

REDUCE E-->E * E

REDUCE E-->E + E

ACCEPT!

RESULT:

Thus the C program implementation for operator precedence is executed and verified.