

Separating Dependency from Constituency in a Tree Rewriting System

Anoop Sarkar

`anoop@linc.cis.upenn.edu`

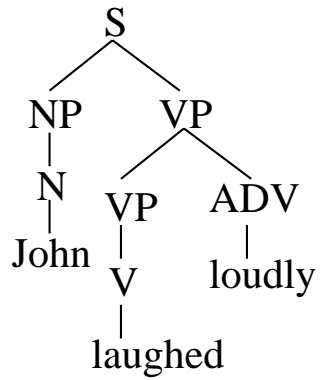
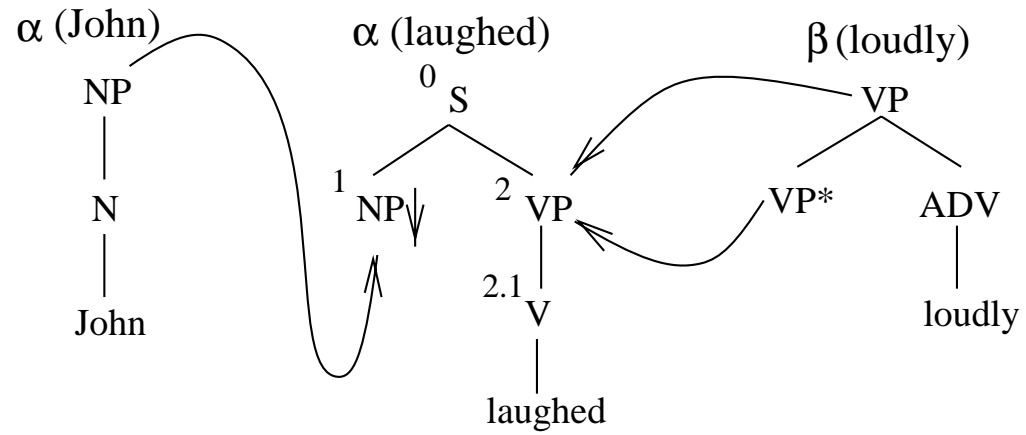
Dept. of Computer and Information Science
University of Pennsylvania

Motivation

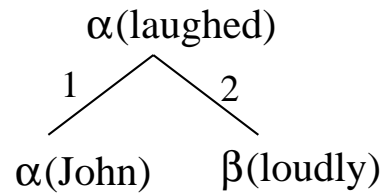
- TAG: reconcile locality and endemic treatment of long distance dependency with coordination
- most successful accounts exploit string adjacency
- CCG: function application interspersed with composition and type raising (Steedman 1996)
- despite being weakly equivalent the CCG analysis of coordination cannot be imported into TAG
- TAGs need a treatment which is structural and not string based.

TAGs

TAG G:

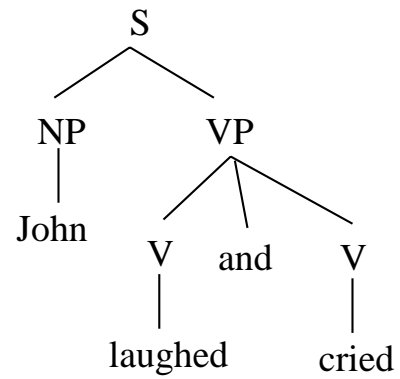
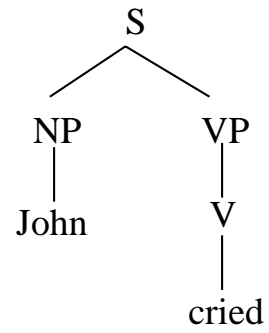
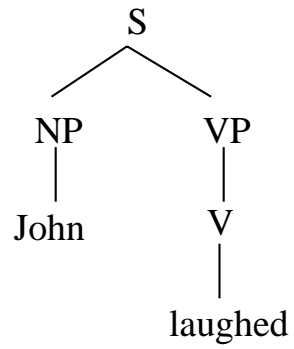


Derived Tree



Derivation Tree

Coordination in a TAG: Problems

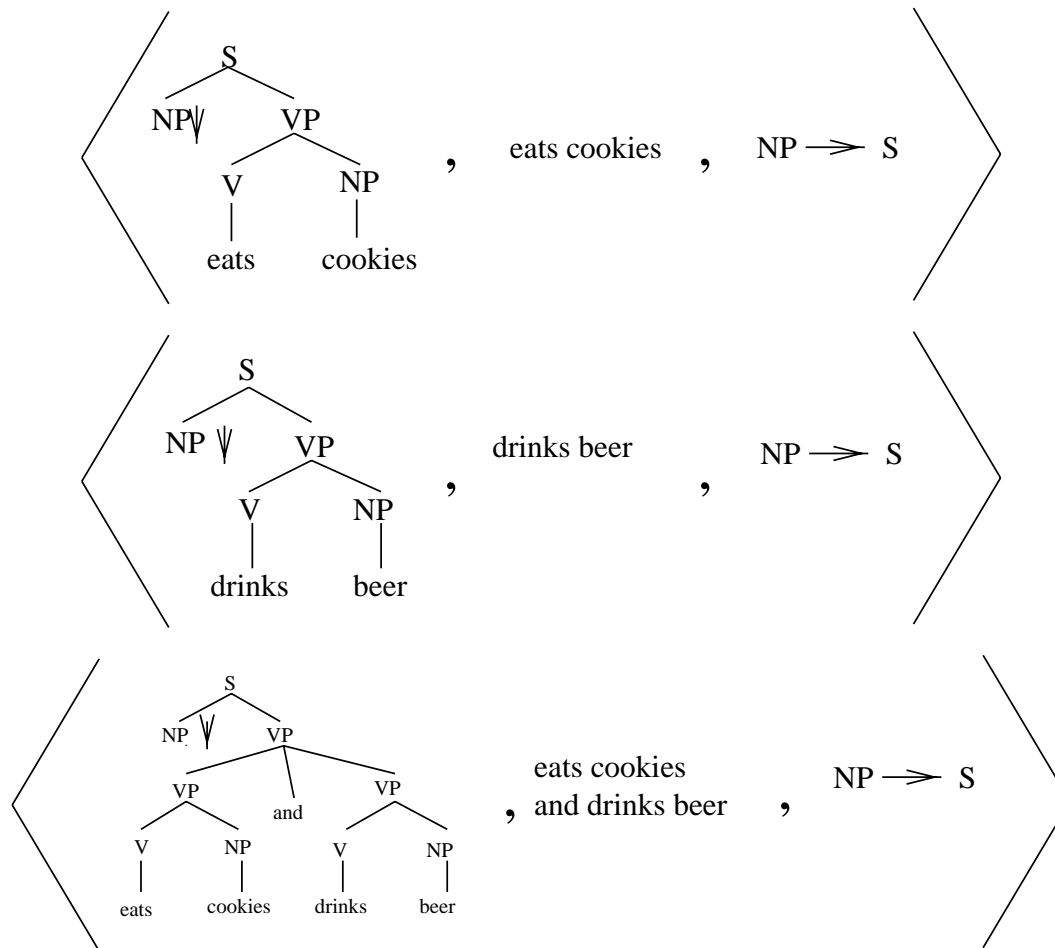


John laughed and cried

Motivation

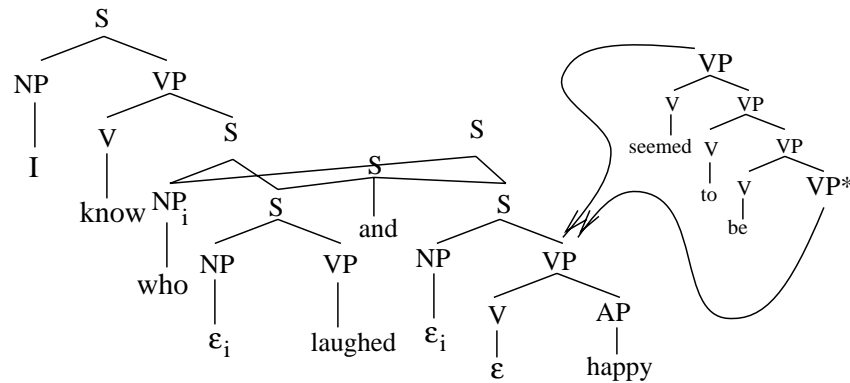
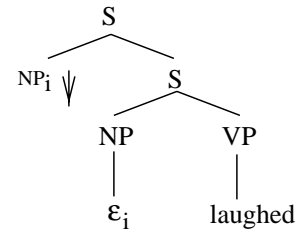
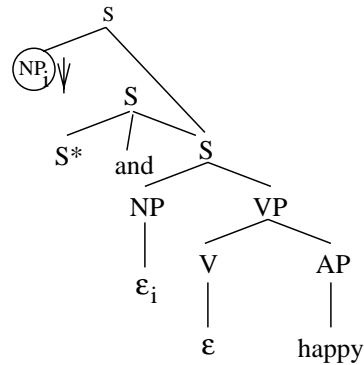
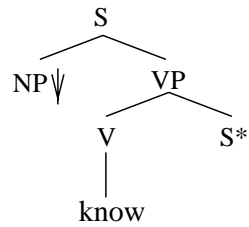
- TAG conflates the notion of dependency and constituency
- is it possible to separate constituent structure from dependency ?
- exploit notions from parallel rewriting systems to handle coordination (Engelfriet, Rozenberg and Slutzki, 1980)
- practical underpinnings: a parser for coordination
- ambiguity (multiple derivations) in non-coordinate sentences
- previous approaches in TAG give unsatisfactory consequences (Joshi and Schabes, 1991; Sarkar and Joshi, 1996)
- avoid structure merging in the parse or unrooted elementary trees

Structure Merging (Joshi and Schabes 91)



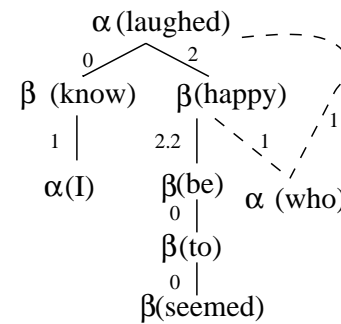
Coordination as Adjunction

Elementary trees



I know who laughed and seemed to be happy

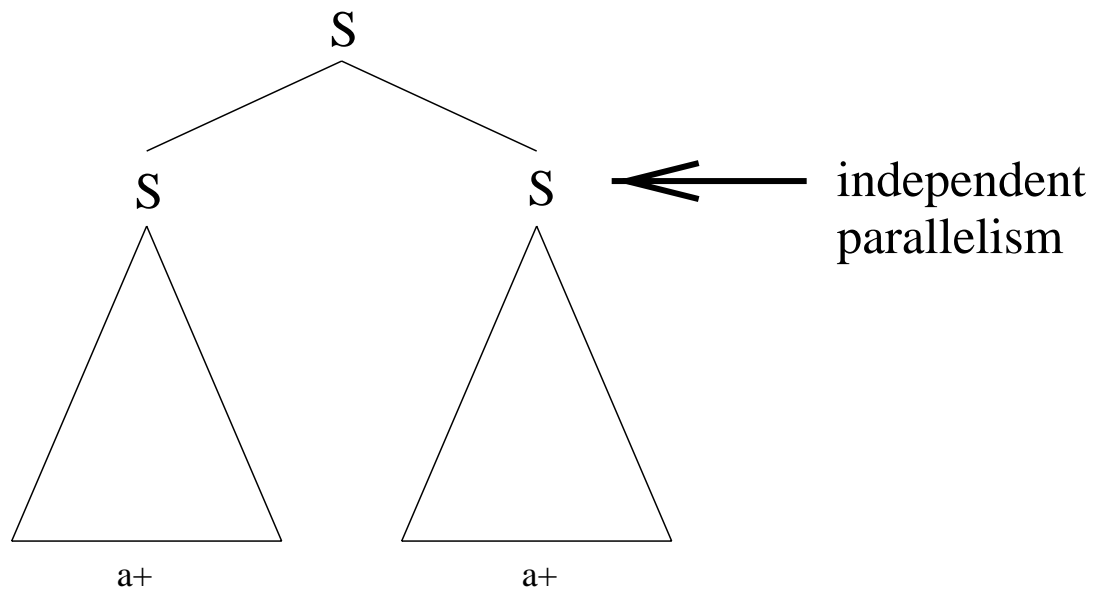
Derived structure



Derivation structure

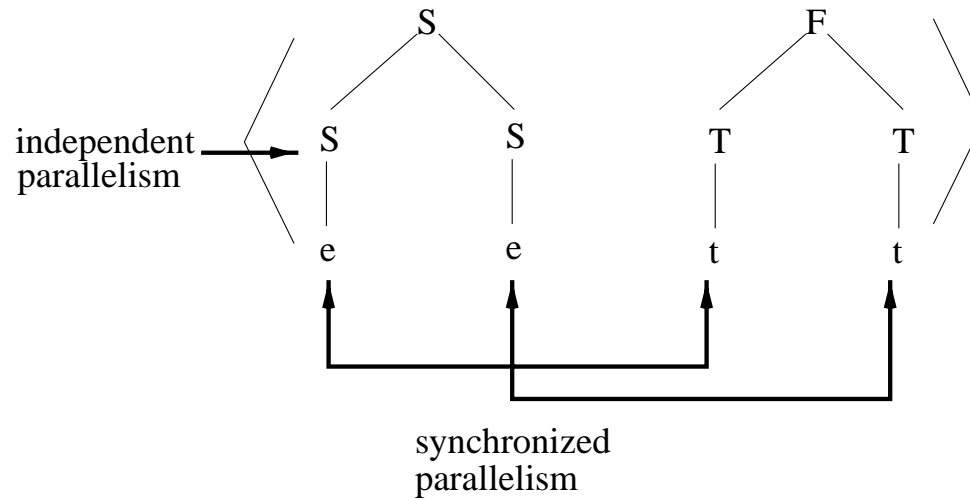
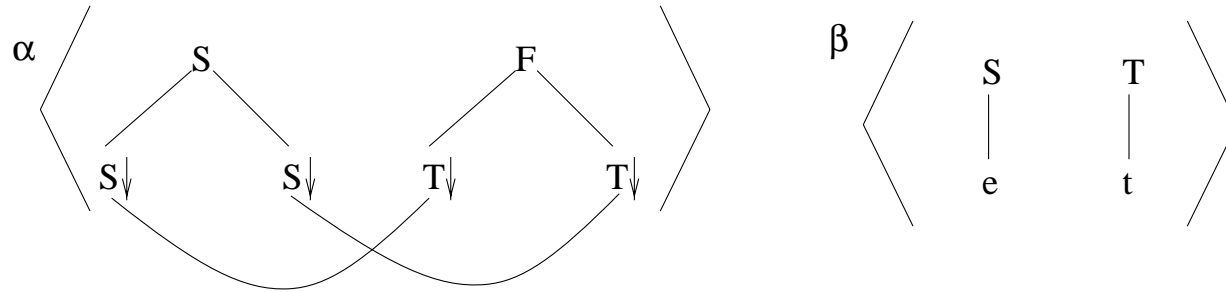
Independent Parallelism – CFGs

CFG G :
 $S \rightarrow S S$
 $S \rightarrow a$



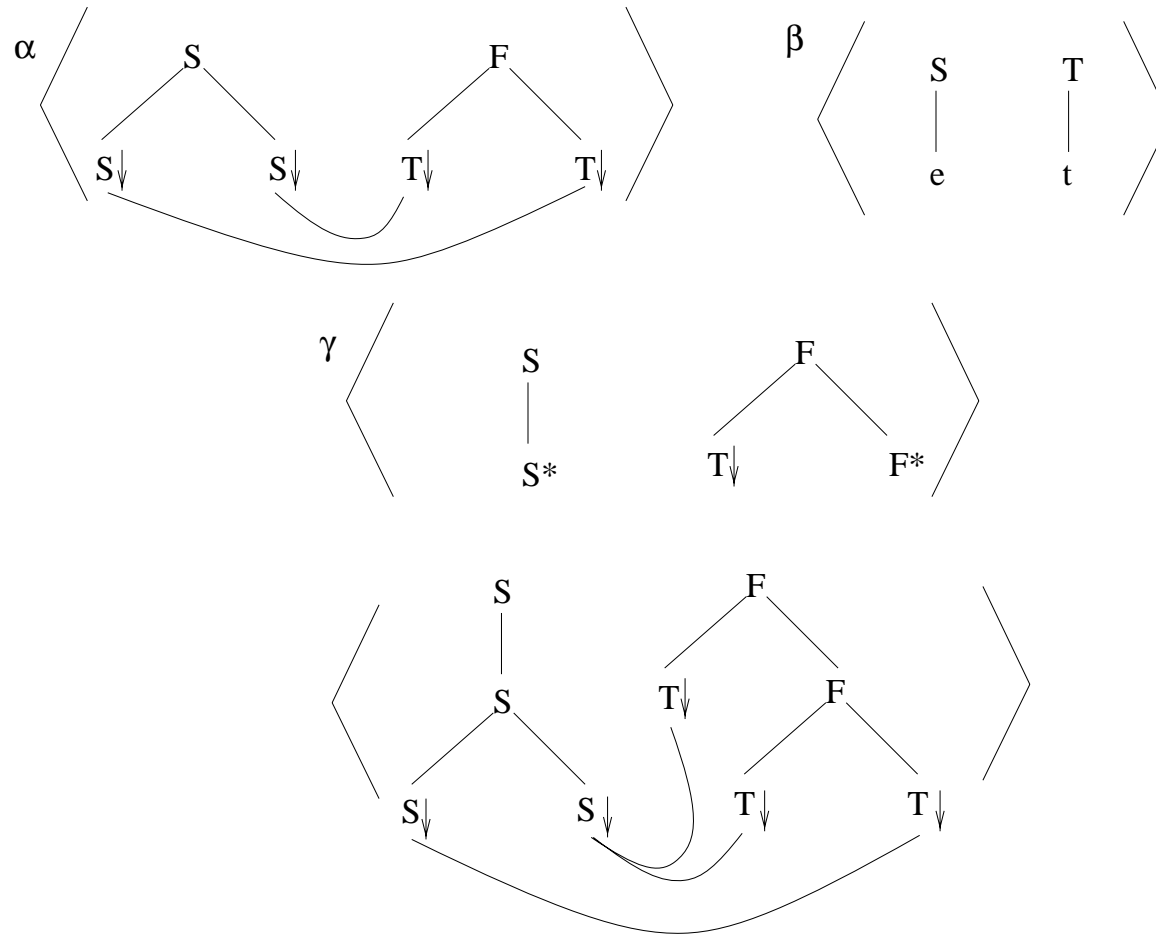
Synchronized Parallelism – STAGs

STAG S:

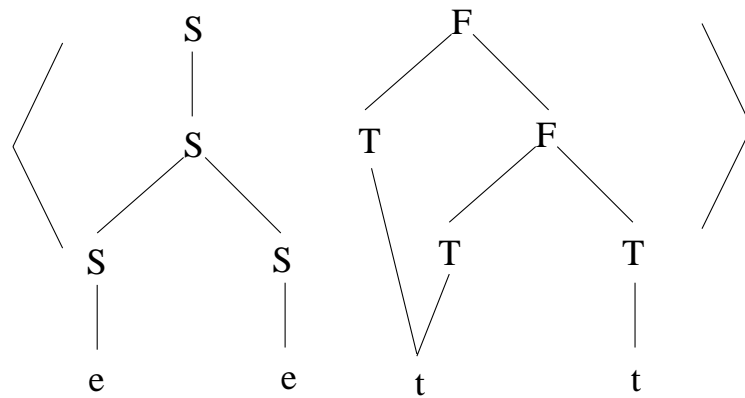
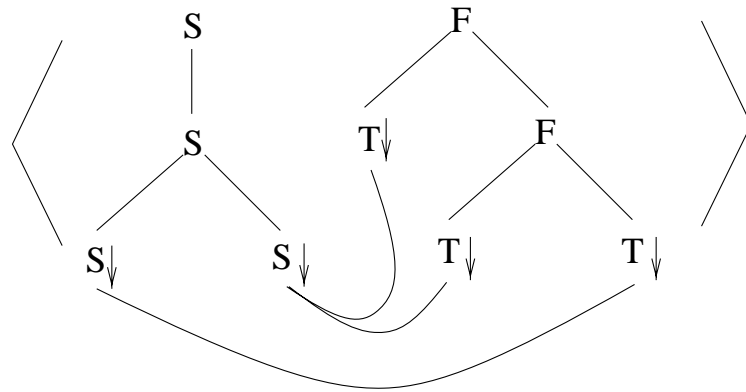


Non-local 'sharing' using Synchronized Parallelism

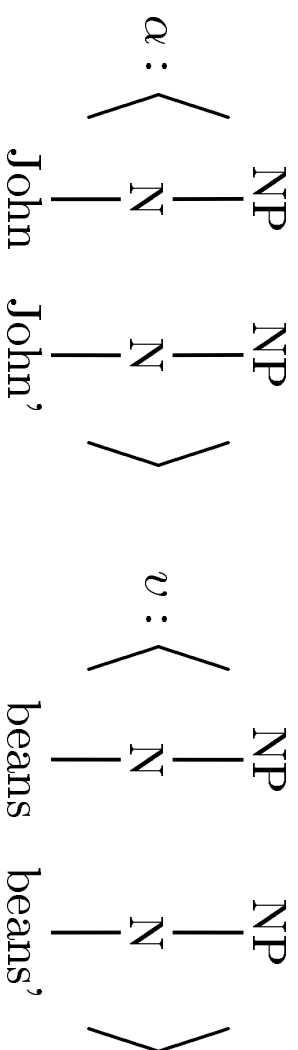
STAG S':



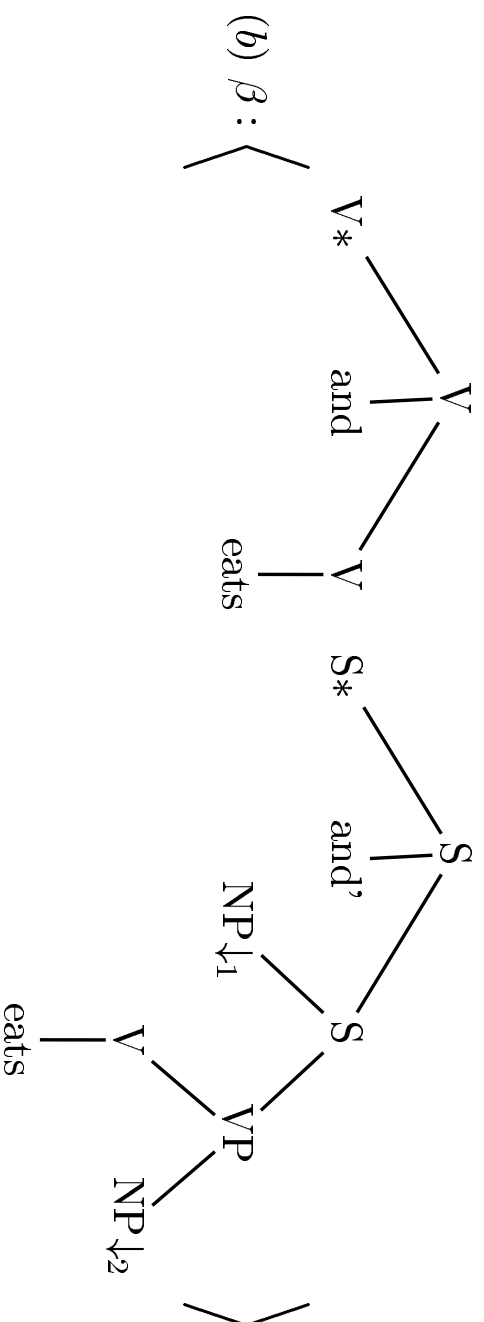
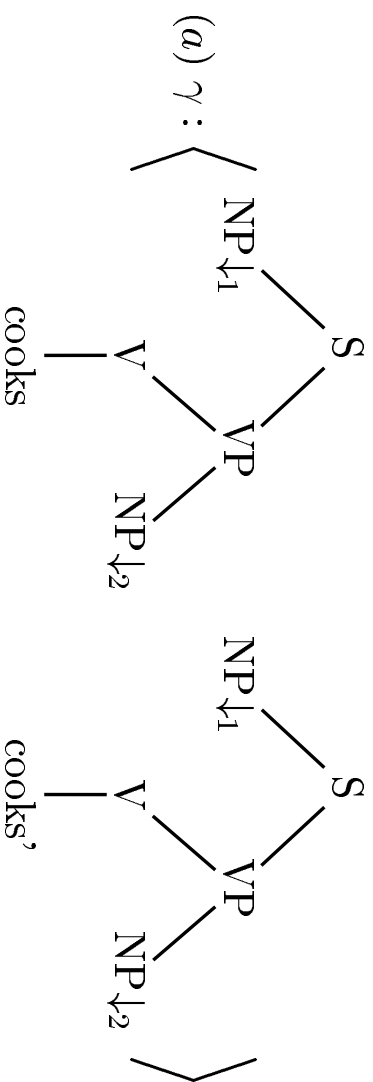
Non-local 'sharing' using Synchronized Parallelism



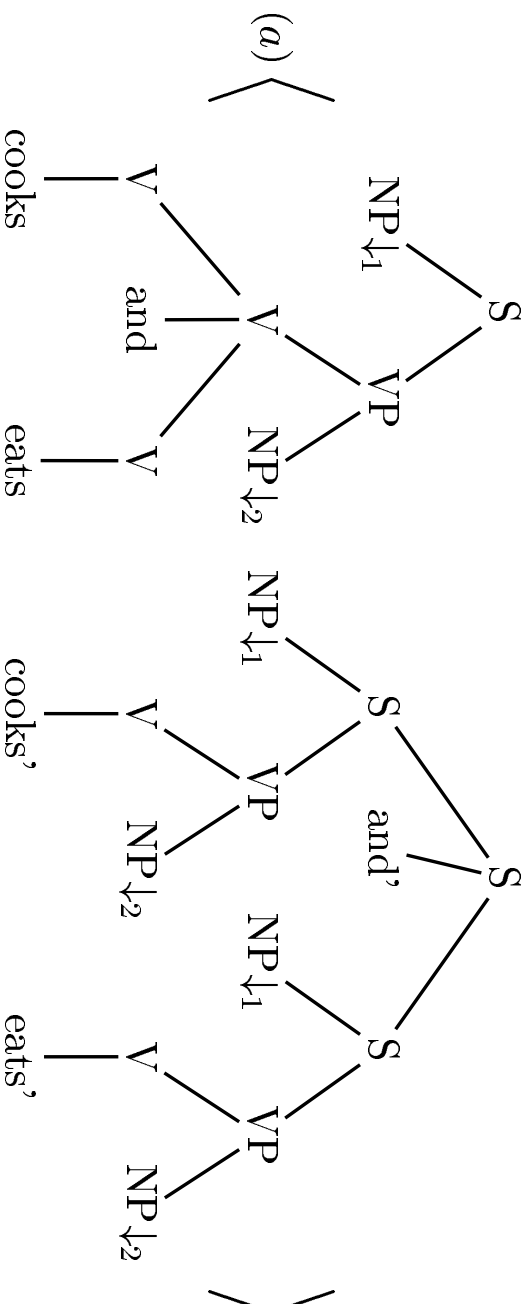
LSTAG: by example



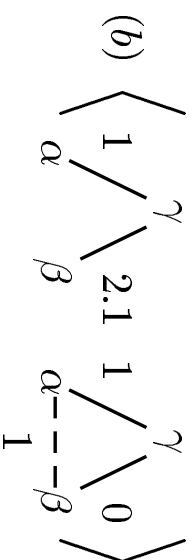
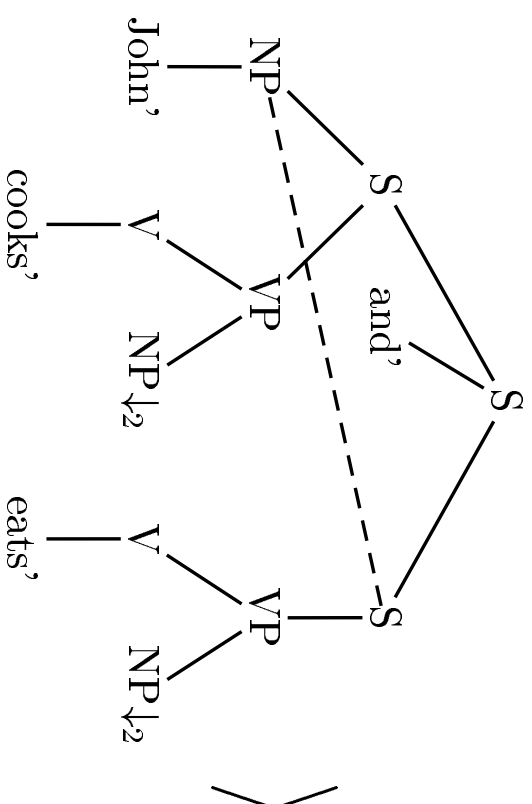
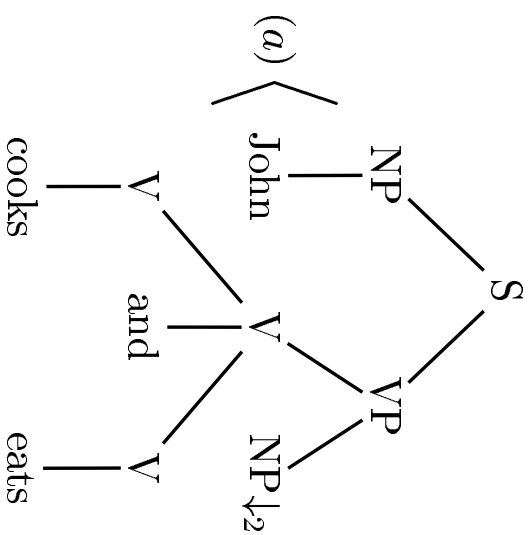
LSTAG: by example



LSTAG: by example



LSTAG: by example



LSTAG: formal definition

An LSTAG G is defined as a 4-tuple $\langle G_L, G_R, \Delta, \Phi \rangle$ where

- G_L, G_R are standard TAGs,
- Δ and Φ are disjoint sets of sets of links and
- for each pair $\gamma = \langle \gamma_L, \gamma_R \rangle$, where $\gamma_L \in G_L$ and $\gamma_R \in G_R$,
- $\delta_\gamma \in \Delta$ is a subset of *links* in γ and
- $\phi_{\gamma_R} \in \Phi$ is a distinguished subset of *links*
- for each link $\smile \in \phi_{\gamma_R}$, $\eta \smile \eta$, where η is a node address in γ_R .
- δ_R and ϕ_{γ_R} have some canonical order \prec .

LSTAG: formal definition

- adjunction (similarly substitution) of $\langle \beta_L, \beta_R \rangle$ into γ is given by

$$\langle \gamma'_L, \gamma'_R \rangle = \langle \gamma_L[a_L, \beta_L], \gamma_R[a_R, \beta_R] \rangle$$

- $\langle \gamma'_L, \gamma'_R \rangle$ is the new derived structure with new set of links $\delta_\gamma \sqcup \phi_{\beta_R}$.

- *for all* $\curvearrowright_{\gamma_i} \in \delta_\gamma, \curvearrowright_{\beta_i} \in \phi_{\beta_R}, 1 \leq i \leq n$

$$\delta_\gamma \sqcup \phi_{\beta_R} \stackrel{def}{=} (\curvearrowright_{\gamma_1} \sqcup \curvearrowright_{\beta_1}) \cup \dots \cup (\curvearrowright_{\gamma_n} \sqcup \curvearrowright_{\beta_n})$$

where

$$(\curvearrowright_{\gamma_1} \curvearrowright \curvearrowright_{\gamma_2}), \dots, (\curvearrowright_{\gamma_{n-1}} \curvearrowright \curvearrowright_{\gamma_n})$$

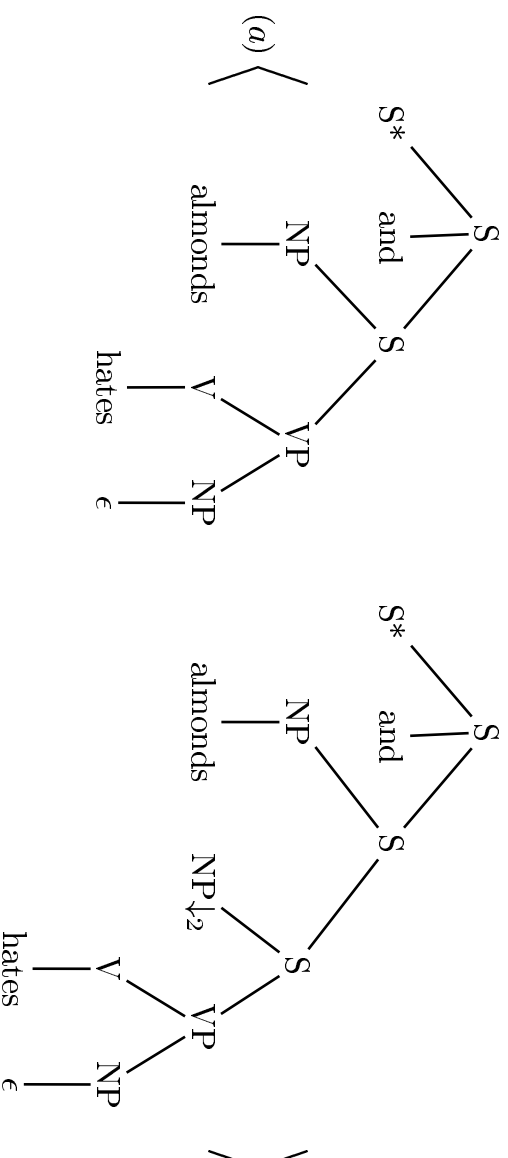
and

$$(\curvearrowright_{\beta_{R_1}} \curvearrowright \curvearrowright_{\beta_{R_2}}), \dots, (\curvearrowright_{\beta_{R_{n-1}}} \curvearrowright \curvearrowright_{\beta_{R_n}})$$

- $\curvearrowright_i \sqcup \curvearrowright_j$ is a set of links defined as follows. If $a_{L_i} \curvearrowright_i a_{R_i}$ and $a_{R_j} \curvearrowright_j a_{R_j}$, then

$$\curvearrowright_i \sqcup \curvearrowright_j \stackrel{def}{=} \{a_{L_i} \curvearrowright a_{R_i}\} \cup \{a_{L_i} \curvearrowright a_{R_j}\}$$

Restrictions



***Peanuts John likes and almonds hates.**

- these restrictions are not part of the formalism
- should be treated as well formedness conditions on elementary structures in the grammar
- they apply to NPs and not to Vs in English, for instance

Summary

- a structural approach towards coordination
- one which distinguishes linguistic dependency from the notion of constituency
- avoids structure merging in the parse and avoids unrooted elementary trees which were shortcomings of previous approaches
- extensions are formally better understood by using techniques from parallel rewriting systems
- the linguistic analyses presented in (Joshi and Schabes, 1991; Sarkar and Joshi, 1996) are easily adopted

Questions you can ask

- why are previous TAG approaches to coordination unsatisfactory ?
- does the approach cover more interesting cases of coordination ?
- can you handle gapping ?
- is this implemented ?

Practical underpinnings

- Xtag homepage:
<http://www.cis.upenn.edu/~xtag>