



CMPT 413: Computational Linguistics

ED3: Edit Distance and FSTs

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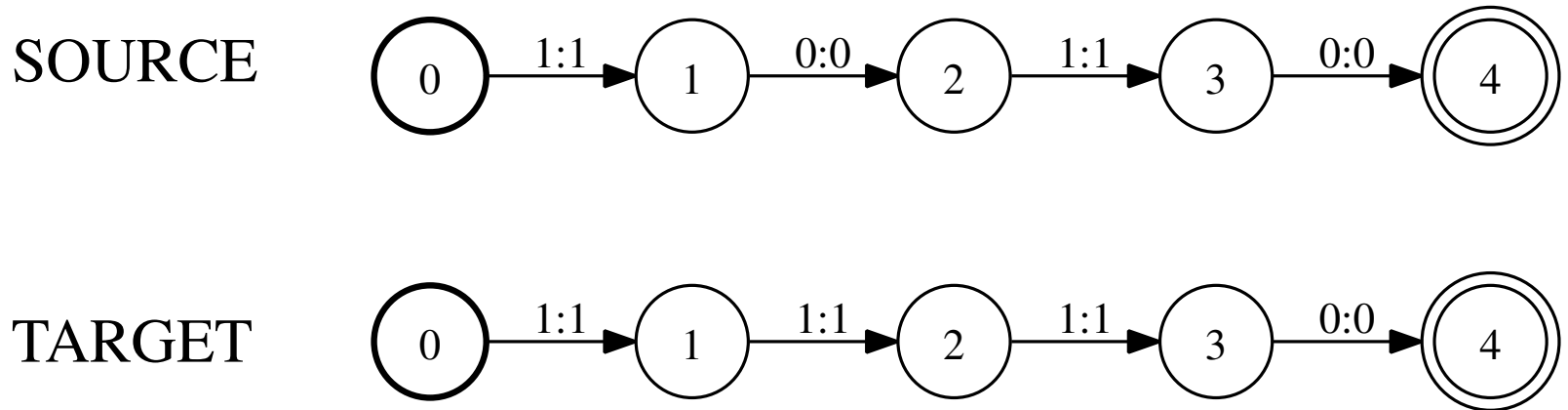
<http://www.cs.sfu.ca/~anoop>

Edit Distance and FSTs

- Algorithm using a Finite-state transducer:
 - Construct an Edit FST:
 - a transition $x:x$ gets zero cost for every x in the alphabet
 - a transition on $\epsilon:x$ (insertion) gets cost 1
 - $x:\epsilon$ (deletion) for any char x gets cost 1
 - optionally $x:y$ (substitution) gets a cost 1 for every x, y
 - Compose source FST, edit FST and target FST
 - Finding minimum cost edit distance == Finding the shortest path from start state to final state
 - Computes an error rate
 - For words this computes the word error rate
 - For letters the letter error rate

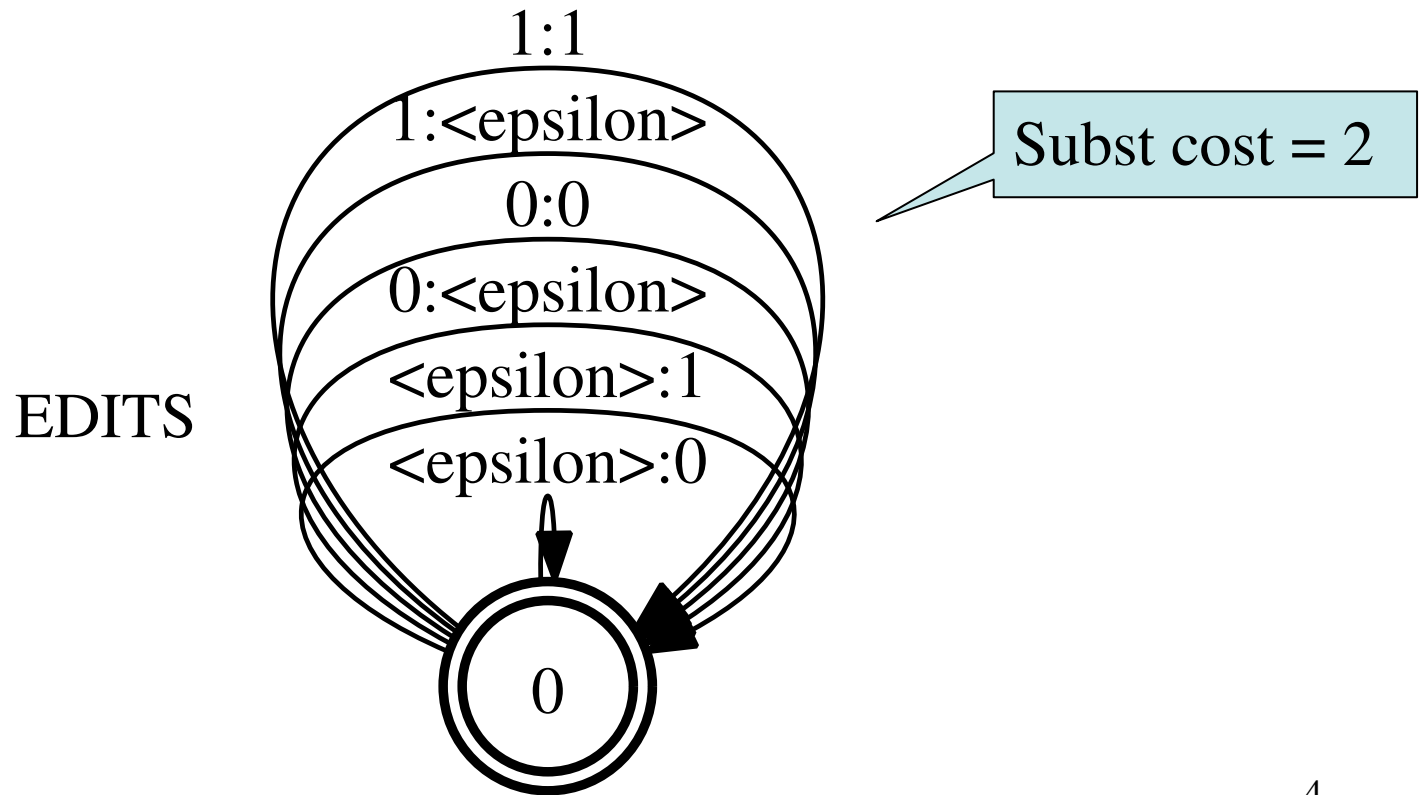
Edit Distance and FSTs

- Lets assume we want to edit source string 1010 into the target string 1110
- The alphabet is just 1 and 0



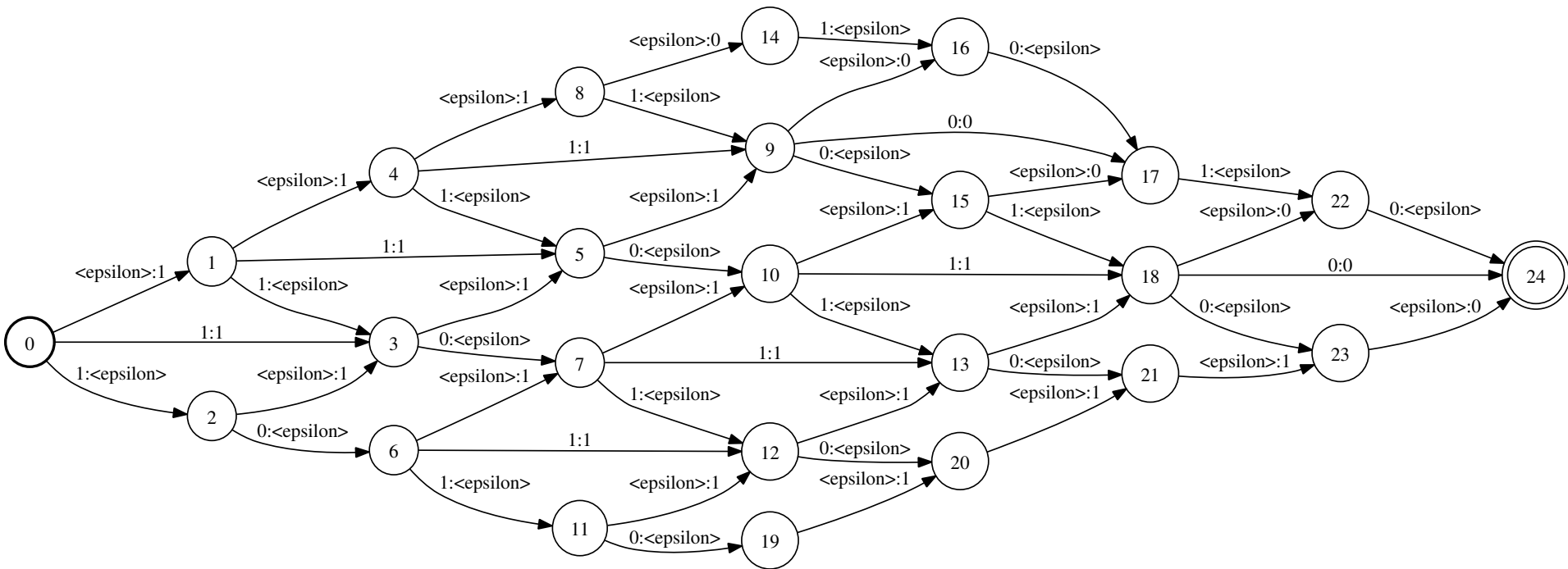
Edit Distance and FSTs

- Construct a FST that allows strings to be edited:
aka flower FST



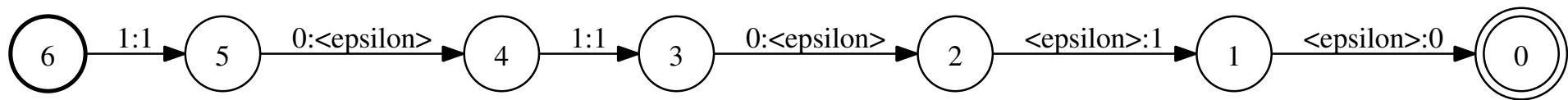
Edit Distance and FSTs

- Compose SOURCE and EDITS and TARGET



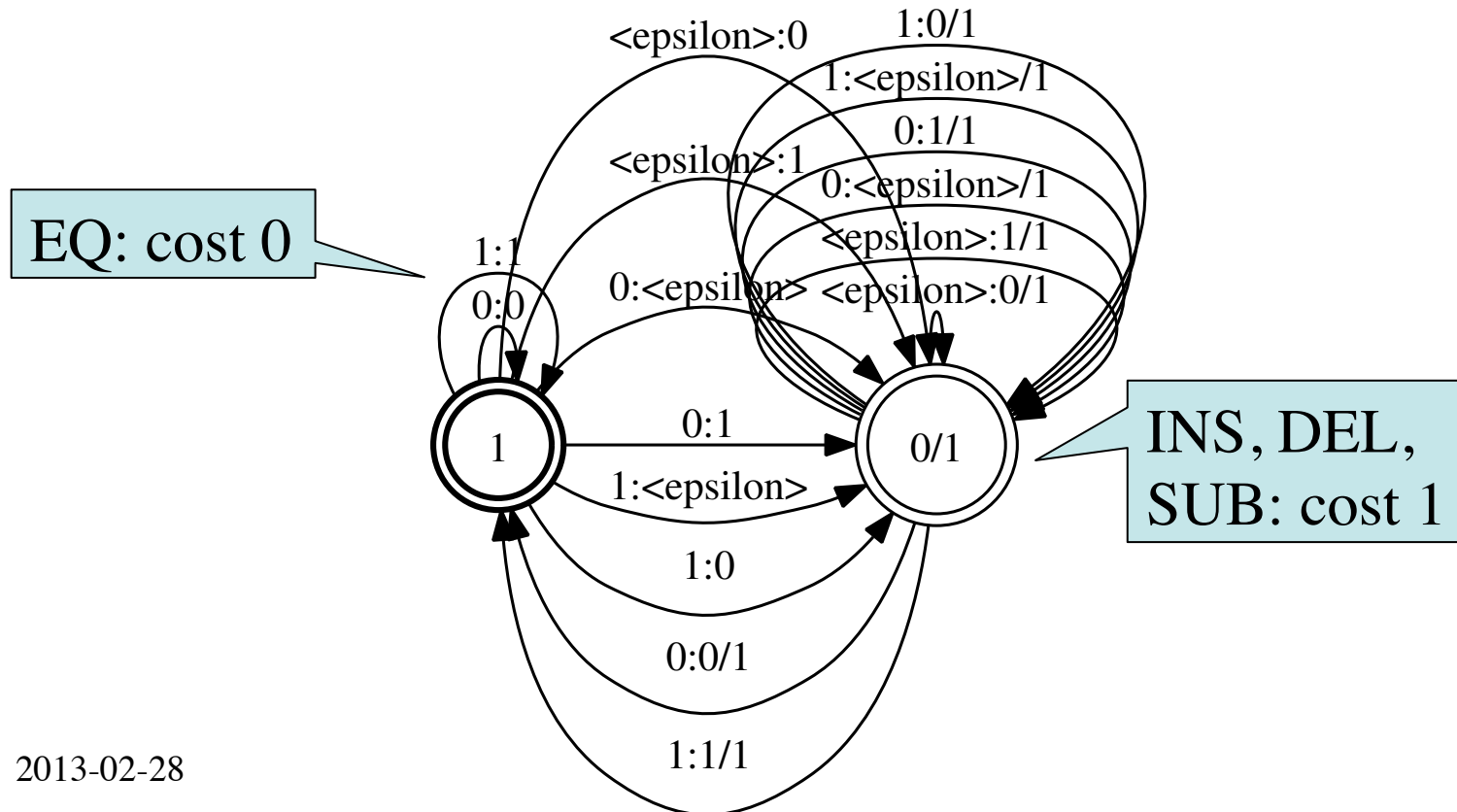
Edit Distance and FSTs

- The shortest path is the minimum edit FST from SOURCE (1010) to TARGET (1110)



Edit Distance and FSTs

However, if we want a substitution cost of 1 (instead of 2) then we have to create a different transducer.

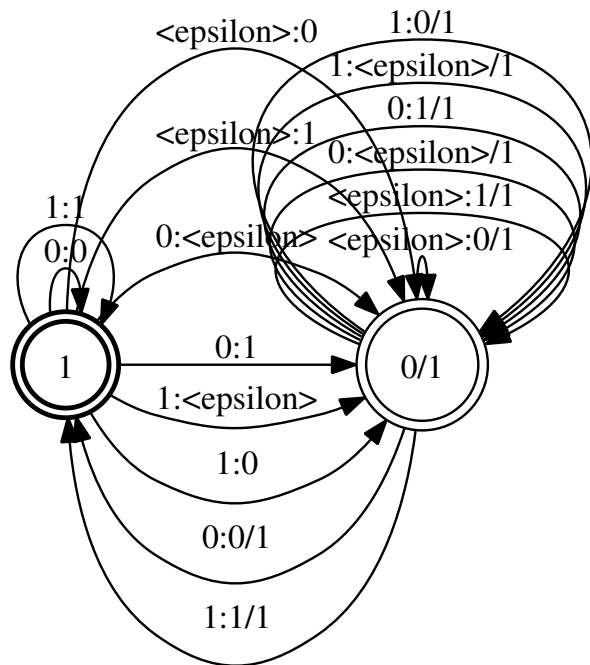


Edit distance and FSTs

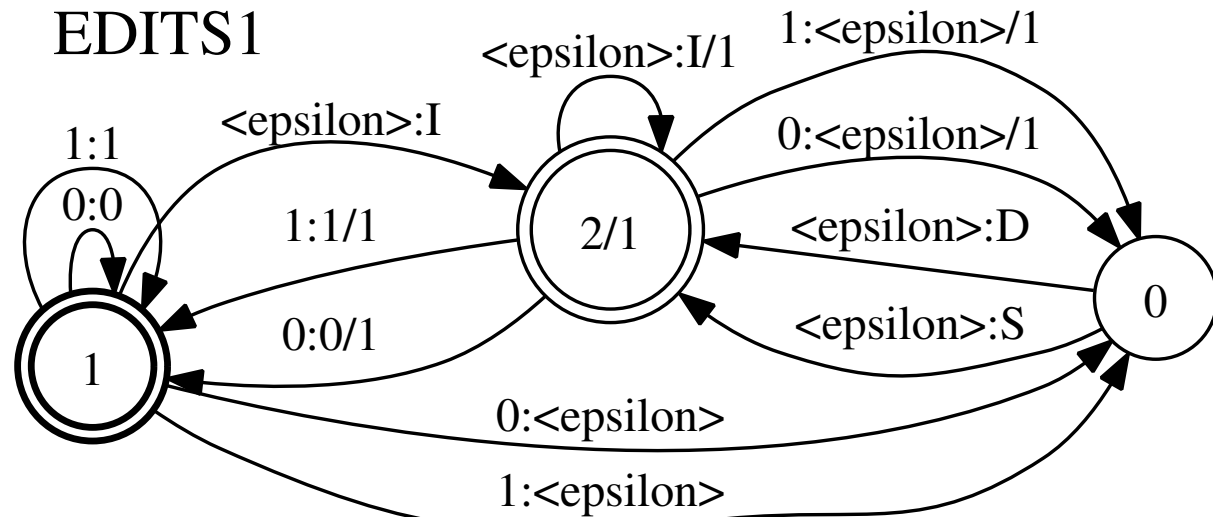
- One problem is scaling to larger character sets
- For 95 ascii symbols the Levenshtein edit transducer will have 9215 transitions
- For 10,000 words the Edit FST needs 100,020,000 (100M) transitions
- Number of transitions = $(V+1)^2-1$
- Solution: De-compose the Edit FST

has $(V+1)^2-1$
transitions

EDITS =
EDITS1 \circ EDITS2

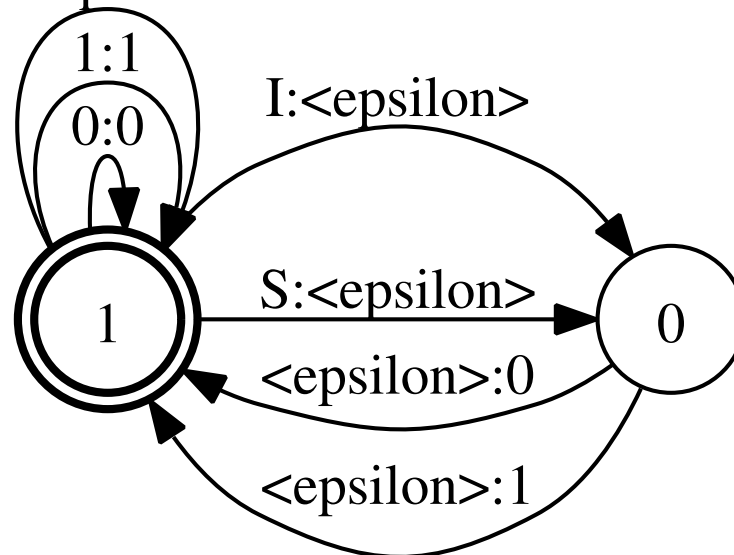


EDITS1



has $\approx 4V$
transitions

D:<epsilon>



EDITS2

has $\approx 4V$
transitions



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