CMPT 413
Computational Linguistics

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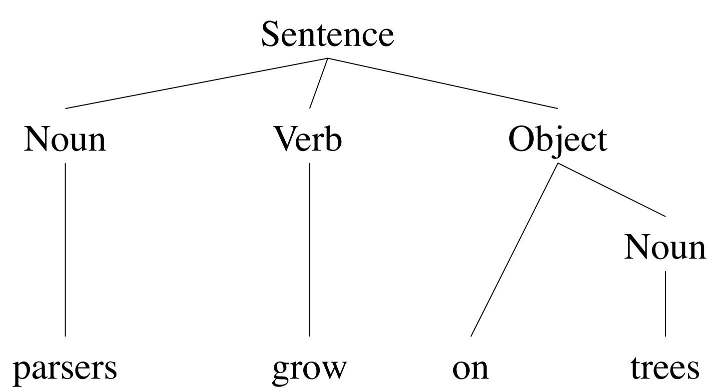
Context-free Grammars

- Set of rules by which valid sentences can be constructed.
- Example:
  - Sentence → Noun Verb Object
  - Noun → *trees* | *parsers*
  - Verb → *are* | *grow*
  - Object → *on* Noun | Adjective
  - Adjective → *slowly* | *interesting*
- What strings can Sentence *derive*?
- Syntax only – no semantic checking

Derivations of a CFG

- *parsers grow on trees*
- *parsers grow on* **Noun**
- *parsers grow* **Object**
- *parsers* **Verb Object**
- **Noun Verb Object**
- **Sentence**

Derivations and parse trees



Ambiguity

- An input is ambiguous with respect to a CFG if it can be derived with two different parse trees
- A parser needs a mechanical definition of ambiguity as it parses the input string
- Is a parser choice really ambiguous, i.e. does it lead to ambiguous parse trees? or not?
- We can formally define ambiguity in terms of the derivations possible in a CFG

Arithmetic Expressions

- $E \rightarrow E + E$
- $E \rightarrow E * E$
- $E \rightarrow (E)$
- $E \rightarrow - E$
- $E \rightarrow \mathbf{id}$

Leftmost derivations for **id + id * id**

$E \rightarrow E + E$
$E \rightarrow E * E$
$E \rightarrow (E)$
$E \rightarrow - E$
$E \rightarrow \mathbf{id}$

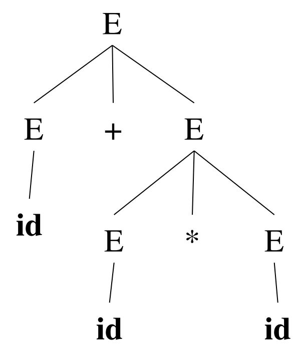
- $E \Rightarrow E + E$

$\Rightarrow \mathbf{id} + E$

$\Rightarrow \mathbf{id} + E * E$

$\Rightarrow \mathbf{id} + \mathbf{id} * E$

$\Rightarrow \mathbf{id} + \mathbf{id} * \mathbf{id}$



Leftmost derivations for **id + id * id**

$E \rightarrow E + E$
$E \rightarrow E * E$
$E \rightarrow (E)$
$E \rightarrow - E$
$E \rightarrow \mathbf{id}$

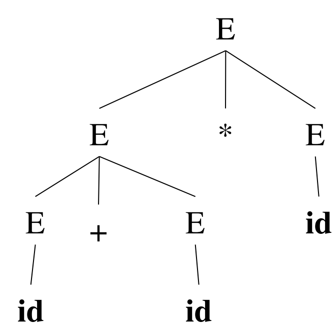
- $E \Rightarrow E * E$

$\Rightarrow E + E * E$

$\Rightarrow \mathbf{id} + E * E$

$\Rightarrow \mathbf{id} + \mathbf{id} * E$

$\Rightarrow \mathbf{id} + \mathbf{id} * \mathbf{id}$



Rightmost derivation for
id + id * id

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$E \rightarrow E * E$
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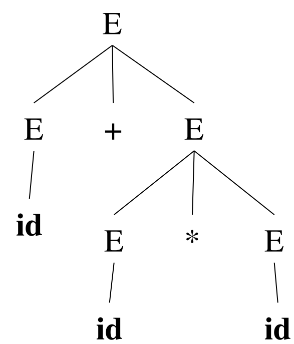

Rightmost derivation for
id + id * id

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$E \Rightarrow E * E$
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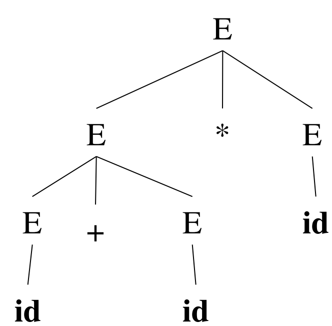


Ambiguity

- We can now define *ambiguity* for a context-free parser
- If a parser has a choice of two different leftmost derivations,
- or if a parser has a choice of two different rightmost derivations,
- for a particular input then that input is ambiguous

Parsing - Roadmap

- Parser is a decision procedure: builds a parse tree
- Top-down vs. bottom-up
- Recursive-descent with backtracking
- Bottom-up parsing (CKY)
- Shift-reduce parsing
- Combining top-down and bottom-up: Earley parsing

Top-Down vs. Bottom Up

Grammar: $S \rightarrow A\ B$ Input String: ccbca

$A \rightarrow c \mid \varepsilon$

$B \rightarrow cbB \mid ca$

Top-Down/leftmost		Bottom-Up/rightmost	
$S \Rightarrow AB$	$S \rightarrow AB$	$ccbca \Leftarrow Acbca$	$A \rightarrow c$
$\Rightarrow cB$	$A \rightarrow c$	$\Leftarrow AcbB$	$B \rightarrow ca$
$\Rightarrow ccbB$	$B \rightarrow cbB$	$\Leftarrow AB$	$B \rightarrow cbB$
$\Rightarrow ccbca$	$B \rightarrow ca$	$\Leftarrow S$	$S \rightarrow AB$

Ambiguity

- Grammar is ambiguous if more than one parse tree is possible for some sentences
- Examples in English:
  – Two sisters reunited after 18 years in checkout counter
- It is undecidable to check using an algorithm whether a grammar is ambiguous

Parsing CFGs

- Consider the problem of parsing with arbitrary CFGs
- For any input string, the parser has to produce a parse tree
- The simpler problem: print **yes** if the input string is generated by the grammar, print **no** otherwise
- This problem is called *recognition*

CKY Recognition Algorithm

- The Cocke-Kasami-Younger algorithm
- As we shall see it runs in time that is polynomial in the size of the input
- It takes space polynomial in the size of the input
- **Remarkable fact:** it can find all possible parse trees (exponentially many) in polynomial time

Chomsky Normal Form

- Before we can see how CKY works, we need to convert the input CFG into Chomsky Normal Form
- CNF is one of many grammar transformations that *preserve* the language
- CNF means that the input CFG G is converted to a new CFG G' in which all rules are of the form:

  $A \rightarrow B\ C$

  $A \rightarrow a$

Epsilon Removal

- First step, remove epsilon rules

  $A \rightarrow B\ C$

  $C \rightarrow \varepsilon \mid C\ D \mid a$

  $D \rightarrow b \quad B \rightarrow b$

- After ε-removal:

  $A \rightarrow B \mid B\ C\ D \mid B\ a \mid BC$

  $C \rightarrow D \mid C\ D\ D \mid a\ D \mid C\ D \mid a$

  $D \rightarrow b \quad B \rightarrow b$

Removal of Chain Rules

- Second step, remove chain rules

  A → B C | C D C

  C → D | a

  D → d   B → b
- After removal of chain rules:

  A → B a | B D | a D a | a D D | D D a | D D D

  D → d   B → b

Eliminate terminals from RHS

- Third step, remove terminals from the rhs of rules

  A → B a C d
- After removal of terminals from the rhs:

  A → B N_1 C N_2

  N_1 → a

  N_2 → d

Binarize RHS with Nonterminals

- Fourth step, convert the rhs of each rule to have two non-terminals

  $A \rightarrow B\ N_1\ C\ N_2$

  $N_1 \rightarrow a$

  $N_2 \rightarrow d$

- After converting to binary form:

  $A \rightarrow B\ N_3$ $\qquad N_1 \rightarrow a$

  $N_3 \rightarrow N_1\ N_4$ $\qquad N_2 \rightarrow d$

  $N_4 \rightarrow C\ N_2$

CKY algorithm

- We will consider the working of the algorithm on an example CFG and input string
- Example CFG:

  $S \rightarrow A\ X \mid Y\ B$

  $X \rightarrow A\ B \mid B\ A$ $\qquad Y \rightarrow B\ A$

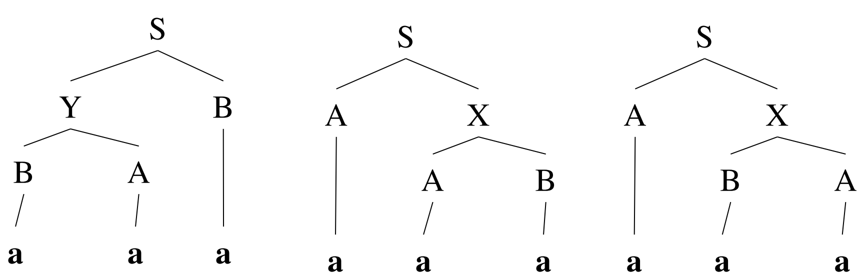
  $A \rightarrow a$ $\qquad B \rightarrow a$

- Example input string: *aaa*

CKY Algorithm

	0	1	2	3
0		A, B A → a B → a	X, Y X → A B	B A Y → B A	S $S \rightarrow A_{(0,1)}\ X_{(1,3)}$ $S \rightarrow Y_{(0,2)}\ B_{(2,3)}$
1			A, B A → a B → a	X, Y X → A B	B A Y → B A
2				A, B A → a B → a
		a	a	a

Parse trees



CKY Algorithm

Input string **input** of size *n*
Create a 2D table **chart** of size n^2
for i=0 **to** n-1
  chart[i][i+1] = A **if** there is a rule A → a and **input**[i]=a
for j=2 **to** N
  for i=j-2 **downto** 0
    for k=i+1 **to** j-1
      chart[i][j] = A **if** there is a rule A → B C **and chart** [i][k] = B **and chart**[k][j] = C
return *yes* **if chart**[0][n] has the start symbol
else return *no*

CKY algorithm summary

- Parsing arbitrary CFGs
- For the CKY algorithm, the time complexity is $O(|G|^2 n^3)$
- The space requirement is $O(n^2)$
- The CKY algorithm handles arbitrary ambiguous CFGs
- All ambiguous choices are stored in the chart
- For compilers we consider parsing algorithms for CFGs that do not handle ambiguous grammars

Parsing - Summary

- Parsing arbitrary CFGs: $O(n^3)$ time complexity
- Top-down vs. bottom-up
  - Recursive-descent parsing
  - Shift-reduce parsing
- Earley parsing
- Ambiguous grammars result in parser output with multiple parse trees for a single input string

Parsing - Additional Results

- $O(n^2)$ time complexity for linear grammars
  - All rules are of the form S → aSb or S → a
  - Reason for $O(n^2)$ bound is the linear grammar normal form: A → aB, A → Ba, A → B, A → a
- Left corner parsers
  - extension of top-down parsing to arbitrary CFGs
- Earley's parsing algorithm
  - $O(n^3)$ worst case time for arbitrary CFGs just like CKY
  - $O(n^2)$ worst case time for unambiguous CFGs
  - $O(n)$ for specific unambiguous grammars

(e.g. S → aSa | bSb | ε)

Non-CF Languages

$$L_1 = \{wcw \mid w \in (a|b)*\}$$

$$L_2 = \{a^n b^m c^n d^m \mid n \geq 1, m \geq 1\}$$

$$L_3 = \{a^n b^n c^n \mid n \geq 0\}$$

CF Languages

$$L_4 = \{wcw^R \mid w \in (a|b)*\}$$

$$S \rightarrow aSa \mid bSb \mid c$$

$$L_5 = \{a^n b^m c^m d^n \mid n \geq 1, m \geq 1\}$$

$$S \rightarrow aSd \mid aAd$$

$$A \rightarrow bAc \mid bc$$

Context-free languages and Pushdown Automata

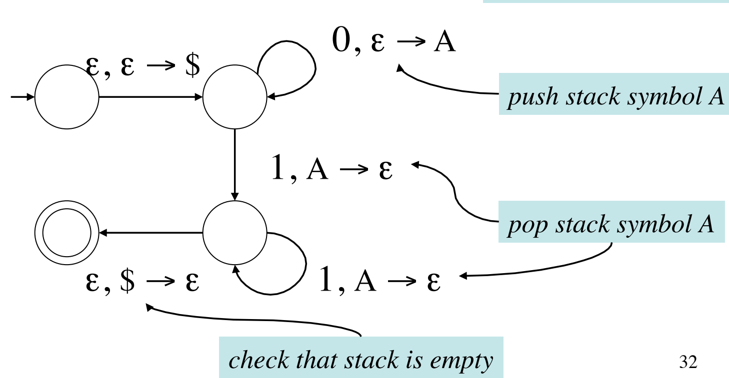
- Recall that for each regular language there was an equivalent finite-state automaton
- The FSA was used as a recognizer of the regular language
- For each context-free language there is also an automaton that recognizes it: called a **pushdown automaton (pda)**

Pushdown Automata

- PDA has
  - an alphabet (terminals) and
  - stack symbols (like non-terminals),
  - a finite-state automaton, and
  - stack

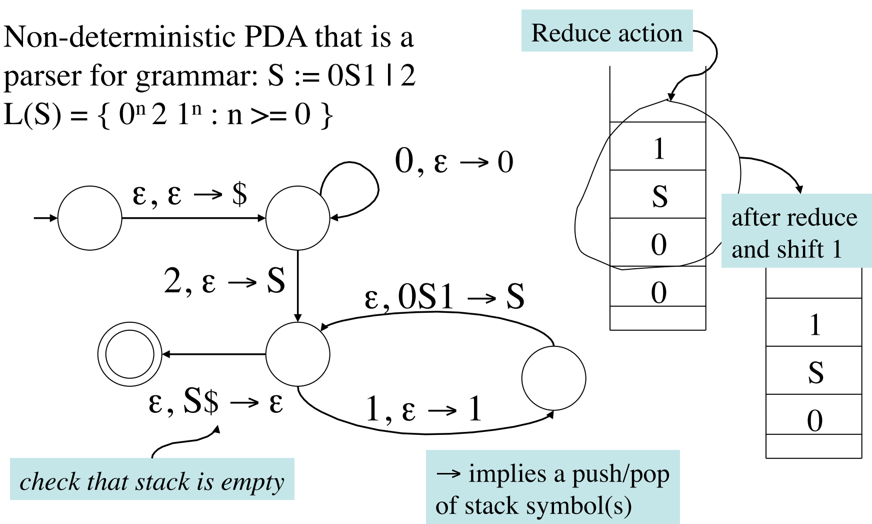
e.g. PDA for language
$L = \{ 0^n 1^n : n >= 0 \}$

→ implies a push/pop of stack symbol(s)



Shift-reduce parser as a pda

Non-deterministic PDA that is a parser for grammar: S := 0S1 | 2
L(S) = { $0^n 2 1^n$: n >= 0 }



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Context-free languages and Pushdown Automata

- Similar to FSAs there are non-deterministic pda and deterministic pda
- Unlike in the case of FSAs we cannot always convert a npda to a dpda
- The construction of a pda will provide us with the algorithm for parsing (take in strings and provide the parse tree)

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CKY algorithm for PCFGs

- We will consider the working of the algorithm on an example PCFG and input string
- Example PCFG:

  $S \rightarrow A\ X\ (0.3) \mid Y\ B\ (0.7)$

  $X \rightarrow A\ B\ (0.1) \mid B\ A\ (0.9) \qquad Y \rightarrow B\ A\ (1.0)$

  $A \rightarrow a\ (1.0) \quad B \rightarrow a\ (1.0)$
- Example input string: *aaa*

CKY Algorithm

Max(0.1, 0.9)

0.3 * 0.9 = 0.27
Max(0.27, 0.7)

	0	1	2	3
0		A 1.0, B 1.0 $A \rightarrow a_{1.0}$ $B \rightarrow a_{1.0}$	X 0.9, Y 1.0 $X \rightarrow A\ B_{0.1} \mid B\ A_{0.9}$ $Y \rightarrow B\ A_{1.0}$	S 0.7 $S \rightarrow A_{(0,1)}\ X_{(1,3)}\ _{0.3}$ $S \rightarrow Y_{(0,2)}\ B_{(2,3)}\ _{0.7}$
1			A 1.0, B 1.0 $A \rightarrow a_{1.0}$ $B \rightarrow a_{1.0}$	X 0.9, Y 1.0 $X \rightarrow A\ B_{0.1} \mid B\ A_{0.9}$ $Y \rightarrow B\ A_{1.0}$
2				A 1.0, B 1.0 $A \rightarrow a_{1.0}$ $B \rightarrow a_{1.0}$
		a	a	a

Parse trees

PCFG is consistent:
0.7 + 0.27 + 0.03 = 1.0