

A decorative graphic on the left side of the slide consisting of two overlapping parallelograms. The front one is blue and the back one is light green. They are positioned diagonally, with the blue one partially covering the green one.

# ProgTeam Spring Week 4

Strings I: Hashing



# Strings

- List of characters, usually a-z, A-Z, 0-9
  - Different from normal array as “alphabet” is fairly small
  - Allows for different approaches to problems than just a normal array



# Hashing

- Pro: “Jackknife” for string problems; probably 90% of problems can be solved by throwing hashing at it
- Con: Slower than many other algorithms; has large constant factor
  - Can require some manipulation to figure out which strings you’re comparing



# Hashing

- “Unique” integer returned for a string
- Common hash: for primes  $p$  and  $m$  ( $p < m$ )
  - $(s[0] * p^0 + s[1] * p^1 + \dots + s[n] * p^n) \bmod m$
  - Important!-  $p > \text{size of alphabet}$  (otherwise we have collisions)
    - With large enough  $p$ , doesn't require small alphabet



# Hashing

- With  $m = 10^9 + 7$ , odds of collision are  $\approx 10^{-9}$
- With 100,000 strings, odds of *any* two colliding are 99.95%
  - (see also: Birthday paradox)
- Solution: hash with two different powers
  - Odds of collision are  $\approx 10^{-18}$
  - With 100,000 strings, about 1 in 100 million
- Example:  $p_1 = 131$ ,  $p_2 = 499$ 
  - (Both greater than char max; fine for strings)



# Rolling Hash

- $(s[0] * p^0 + s[1] * p^1 + \dots + s[n] * p^n) \bmod m$
- Notice:
  - $\text{hash}(\text{"abcba"}) = (\text{hash}(\text{"bcba"}) - 'a') / p$
  - With a prime 'm', we can calculate the "modulo inverse"  $p^{-1}$
- Now we can get a hash of any substring using prefix sums!
- $\text{Hash}(L,R) = (\text{Sum}[R] - \text{Sum}[L - 1]) * \text{modinv}(p^L)$
- Examples of this on the Github