

Semantics for **DYNJIT**^{source} \rightarrow **DYNJIT**^{target}

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CoreCPY can be always converted to **DYNJIT^{source}** in a straightforward approach because, **the stack usage of the bytecode generated by valid Python source code is finite.**

Regarding the perspective of partial evaluation, we know **CoreCPY** is a flowchart language with a finite stack(a stack S_N has a size N).

By assigning the abstract value **Dyn** s to a named variable s , and assigning the abstract value **Slot** i to a datum stored in the i -th element from the bottom of stack(BOS) if initialized, we will have finite configurations:

$$\mathbf{Configurations} \subset \mathbf{Labels} \times [\mathbf{Dyn} \ s_1, \mathbf{Dyn} \ s_2, \dots, \mathbf{Dyn} \ s_k]^k \times \mathcal{P}(\{\mathbf{Slot} \ 1, \dots, \mathbf{Slot} \ N\})$$

where \mathcal{P} means getting the power set, k is the number of named variables not matter what it is(cell, free, etc.) and, N , as we've mentioned above, is the size of the stack, which is finite for any given code.

Note that

$$\left| \mathbf{Labels} \times [\mathbf{Dyn} \ s_1, \mathbf{Dyn} \ s_2, \dots, \mathbf{Dyn} \ s_k]^k \times \mathcal{P}(\{\mathbf{Slot} \ 1, \dots, \mathbf{Slot} \ N\}) \right|$$

is simply finite, this leads to finite **Configurations**, and also, finite basic blocks when translating the use of stack to unnamed registers.

DYNJIT^{source} and DYNJIT^{target}

DYNJIT^{source} Syntax

$$\begin{aligned}
 \langle \text{repr} \rangle &::= S \text{ constant} \\
 &\quad | D \text{ var} \\
 \langle \text{instr} \rangle &::= \langle \text{var} \rangle = \text{call } \langle \text{repr} \rangle (\langle \text{repr} \rangle^*) \\
 &\quad | \langle \text{var} \rangle = \langle \text{repr} \rangle \\
 &\quad | \text{return } \langle \text{repr} \rangle \\
 &\quad | \text{goto label} \\
 &\quad | \text{if } \langle \text{repr} \rangle \text{ goto label goto label} \\
 \langle \text{move} \rangle &::= \langle \text{var} \rangle \leftarrow \langle \text{var} \rangle \\
 \langle \Phi \rangle &::= \text{label} : \langle \text{move} \rangle^* \\
 \langle \text{basicblock} \rangle &::= \text{label label} : \Phi [\langle \Phi \rangle^*] \langle \text{instr} \rangle + \\
 \langle \text{entryblock} \rangle &::= \text{label entry} : \Phi [\langle \Phi \rangle^*] \langle \text{instr} \rangle + \\
 \langle \text{basicblocks} \rangle &::= \langle \text{entryblock} \rangle \langle \text{basicblock} \rangle^*
 \end{aligned}$$

DYNJIT^{target} Syntax

$$\begin{aligned}
 \langle \text{absvalue} \rangle &::= (\langle \text{repr} \rangle, \langle \text{type} \rangle) \\
 \langle \text{expr} \rangle &::= \langle \text{absvalue} \rangle \\
 &\quad | \text{call } \langle \text{expr} \rangle (\langle \text{expr} \rangle^*) \\
 \langle \text{stmt} \rangle &::= \langle \text{var} \rangle = \langle \text{expr} \rangle \\
 &\quad | \text{goto label} \\
 &\quad | \text{return } \langle \text{expr} \rangle \\
 &\quad | \text{if } \langle \text{expr} \rangle \text{ goto label goto label} \\
 &\quad | \text{do } \langle \text{expr} \rangle \\
 &\quad | \text{label label} \\
 &\quad | \{ \langle \text{stmts} \rangle \} \\
 &\quad | \text{switch } \langle \text{expr} \rangle \langle \text{case} \rangle^* \\
 \langle \text{case} \rangle &::= | \langle \text{type} \rangle \rightarrow \langle \text{stmt} \rangle \\
 \langle \text{stmts} \rangle &::= \langle \text{stmt} \rangle^*
 \end{aligned}$$

Types

```
type fptr = int
type meth = int
type t =
| TopT
| BottomT
| NoneT
| FPTrT of fptr
| MethT of meth
| NomT of string * (string, t) map
| TupleT of t list
| TypeT of t list
| CellT of t list
| UnionT of t list
| IntrinT of intrinsic
```

Intrinsics and Constants

```
type intrinsic =  
| IsInstanceOf  
| TypeOf  
| BuildTuple  
| Upcast  
| Downcast
```

```
type const =  
| NoneC  
| UndefC  
| TypeL   of t  
| FPtrC   of fptr  
| MethC   of meth  
| IntC     of int  
| FloatC   of float  
| StrC     of string  
| TupleC   of const list  
| IntrinsicC of intrinsic
```