

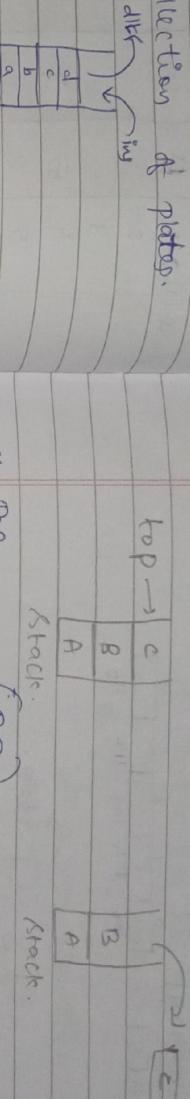
## 03 = STACK

\* Linear dg.

\* LIFO

\* ins & del at top

Eg → collection of plates.



\* Basic operations →

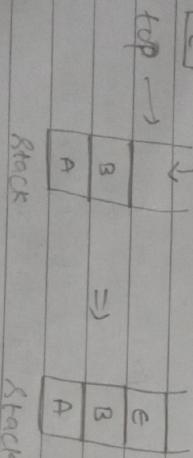
- 1) push → ins
- 2) pop → del.

\* other () →

- peek() → get the top data element of stack without removing it.
- isFull() → Check if stack is full. (ins)
- isEmpty() → Check if the stack is empty. (del)

\* Push operation =

process of putting a new element onto stack.



\* Program for push =

```
void push(int data){  
    if (!isfull()) {  
        stack[top] = data;  
        top++;  
    } else {  
        printf("Stack full");  
    }  
}
```

\* When you are ins. an element / push  
→ increment. address.

\* pop operation ⇒ pop → decrement address.

→ 3rd element 'C' → ins address ⇒ top = 4 (at point)  
so top=top+1 ⇒ 4+1=5.  
∴ stack[4]=data  
⇒ stack[5]='C';



- \* Stack can be maintained in a program by using a array of elements & pointer variables Top & max.
- \* Top points containing the location of the top elements of the stack & max gives the max no. of elements in stack.
- \* Stack is empty if top = -1 or Top = null.

$\rightarrow$  (A) for push

PUSH(STACK, TOP, MAX, ITEM)

1. If TOP = MAX then print overflow.

$\hookrightarrow$  return

2. Set TOP = TOP + 1

3. Set STACK[TOP] = ITEM

4. Return [TOP]

$\rightarrow$  (A) for pop

POP(STACK, TOP, ITEM)

1. If TOP = -1, then print underflow

$\hookrightarrow$  return

2. Set ITEM = STACK[TOP]

3. Set TOP = TOP - 1

4. Return [STOP].

$\Rightarrow$  Representation of stack as L.S =

[Top]

$\rightarrow$  [24]  $\rightarrow$  [8]  $\rightarrow$  [9]  $\rightarrow$  [10 null]

$\rightarrow$  Any action of stack =

\* Back tracking  $\Rightarrow$  is an algorithmic technique

- \* Disadv of using an array to represent a stack is that the array must be specified to have a fixed size.
- \* If the size of the array can't be determined in advance, we can use L.S representation of stack.
- \* Head pointer of L.S is used as the TOP of stack.

for solving prob by trying to build a soln incrementally.

- \* polish notation
- a recursion
- & evaluating exp.

$\Rightarrow$  polish notation =  $\left. \begin{array}{l} ab \\ a+b \\ ab \\ a/b \end{array} \right\}$  exp.

\* Algebraic exp can be written using 3 represnt bt equivalent notations namely infix, postfix & prefix notations.

3) infix notation =

operator symbol is placed b/w its 2 operands in most arithmetic (o).  
 eg  $\rightarrow (A+B)$ .

4) prefix notation = (polish notation)

notation in which operator is placed before its 2 operands.

eg  $\rightarrow \underline{\underline{A+B}}$   $\Rightarrow$   $\begin{matrix} +AB \\ \downarrow \\ \text{exp} \end{matrix}$

5) postfix notation = (reverse polish notation)

operator placed after the operands.

eg  $\rightarrow \underline{\underline{A+B}}$   $\Rightarrow$   $\begin{matrix} AB+ \\ \downarrow \\ \text{exp} \end{matrix}$

$\Rightarrow$  Recursion using stack =

\* () that calls itself  $\rightarrow$  recursive () if thus techniques  $\rightarrow$  recursion.

Eg  $\rightarrow$  int main() {  
 ~~  
 fact(n);  
 ~~  
 }

{ void fact(n){  
 ~~  
 fact(n-1);  
 ~~  
 }