



# Sorting Algorithms

Design & Analysis

# \***1**\*A

# Sorting Problem

#### Sorting Problem

- Input: A sequence of n numbers(a1,a2,...an)
- Output: A permutation (reordering)(al',a2',...an') of the input sequence such that al'<=a2'<=...an

#### Efficiency and Passes

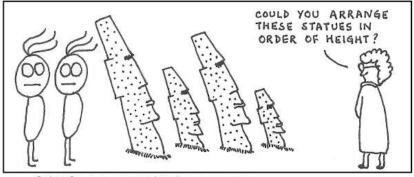
- Efficiency denotes how much time the algorithm takes to sort the elements
  - Efficiency of sorting algorithm is measured in terms of time complexity –big-Oh notations
  - $\rightarrow$  O(n<sup>2</sup>), O(nlogn)
- Passes -The phases in which elements are moving to acquire their proper position



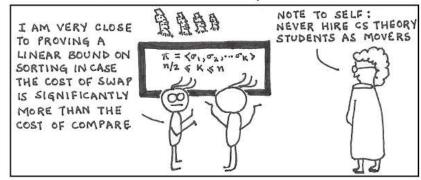


# Analysis of:

- Sorting Algorithms
  - ▶ Bubble Sort
  - Selection Sort
  - Insertion Sort



SEVERAL HOURS LATER ...

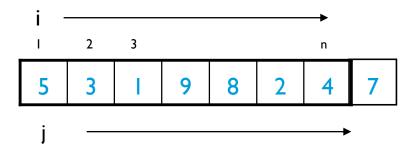






#### **Bubble Sort**

- Idea:
  - Repeatedly pass through the array
  - Swaps adjacent elements that are out of order



Easier to implement, but slower than Insertion sort

# Example







#### **Bubble Sort**

```
for(i=0;i<n-1;i++)
 for(j=0;j<n-i-1;j++)
  if (a[j]>a[j+1])
      temp=a[j];
      a[j]=a[j+1];
      a[j+1]=temp;
```

# \***1**\*A

#### Selection Sort

#### Idea:

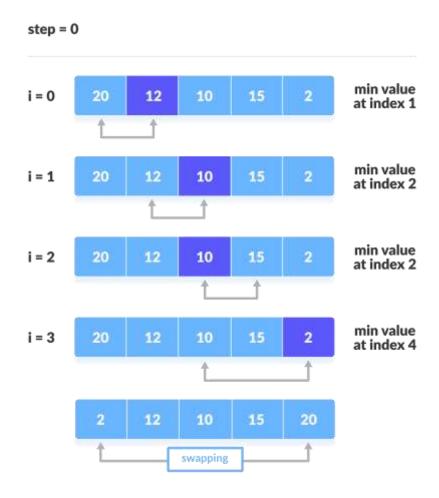
- Find the smallest element in the array
- Exchange it with the element in the first position
- Find the second smallest element and exchange it with the element in the second position
- Continue until the array is sorted

#### Disadvantage:

Running time depends only slightly on the amount of order in the file



# Example







5 3 4 1 2

Selection Sort



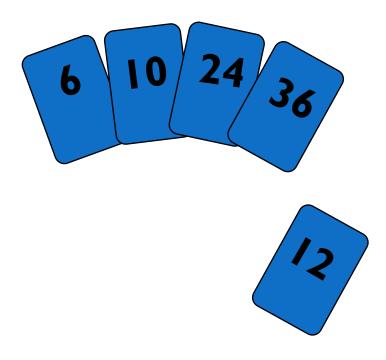
#### Selection Sort

```
for(i=0;i<n;i++)</pre>
         min=a[i];
          pos=i;
          for(j=i+1;j<n;j++)</pre>
                   if(a[j]<min)</pre>
                             min=a[j];
                             pos=j;
          if (pos!=i)
                   temp=a[i];
                   a[i]=a[pos];
                   a[pos]=temp;
```



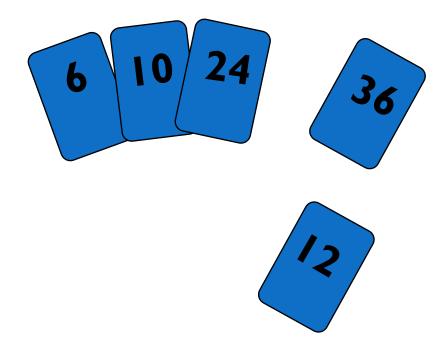
- Idea: like sorting a hand of playing cards
  - Start with an empty left hand and the cards facing down on the table.
  - Remove one card at a time from the table, and insert it into the correct position in the left hand
    - compare it with each of the cards already in the hand, from right to left
  - The cards held in the left hand are sorted
    - these cards were originally the top cards of the pile on the table



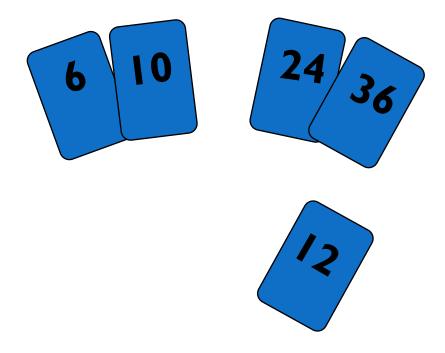


To insert 12, we need to make room for it by moving first 36 and then 24.

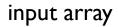






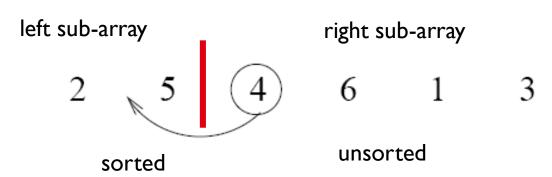




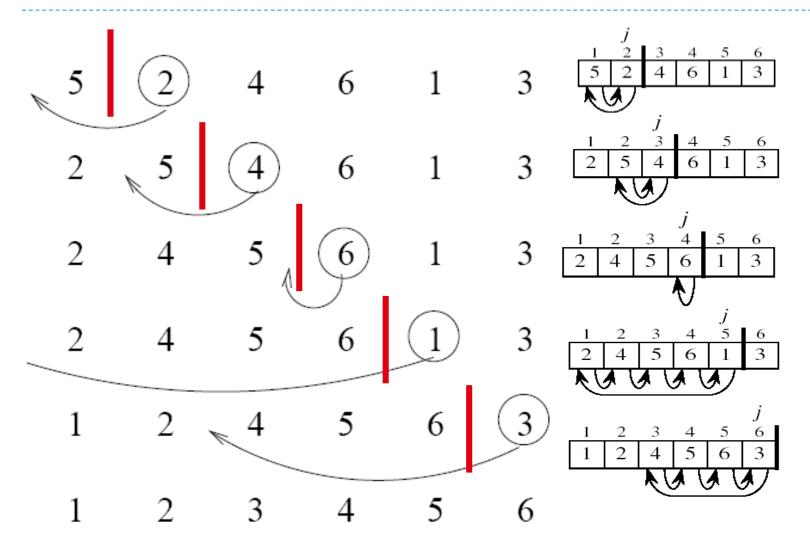


5 2 4 6 I 3

at each iteration, the array is divided in two sub-arrays:









#### **INSERTION-SORT**

for 
$$j \leftarrow 2$$
 to  $n$ 

do key  $\leftarrow A[j]$ 

| lnsert  $A[j]$  into the sorted sequence  $A[1..j-1]$ 

i  $\leftarrow j-1$ 

while  $i > 0$  and  $A[i] > key$ 

do  $A[i+1] \leftarrow A[i]$ 
 $i \leftarrow i-1$ 
 $A[i+1] \leftarrow key$ 

Insertion sort – sorts the elements in place



# Analysis of Insertion Sort

# INSERTION-SORT(A) for $j \leftarrow 2$ to n **do** key $\leftarrow$ A[j] ▷Insert A[j] into the sorted sequence A[l..j-l] i ← j - I while i > 0 and A[i] > key do $A[i + 1] \leftarrow A[i]$ $i \leftarrow i - I$ $A[i + 1] \leftarrow key$



# Best Case Analysis

▶ The array is already sorted

- "while i > 0 and A[i] > key"
- ►  $A[i] \le \text{key upon the first time the } \text{while loop test is run (when } i = j-1)$

# Worst Case Analysis



- The array is in reverse sorted order "while i > 0 and A[i] > key"
  - Always A[i] > key in while loop test
  - ▶ Have to compare key with all elements to the left of the j-th position  $\Rightarrow$  compare with j-1 elements



# Insertion Sort - Summary

- Advantages
  - ▶ Good running time for "almost sorted" arrays  $\Theta(n)$
- Disadvantages
  - $\bullet$   $\Theta(n^2)$  running time in worst and average case

Thank you!