

TUTORIAL - 3

Q1 Write linear search pseudocode to search an element in a sorted array with min-comparisons.

Ans1

```

Void linearSearch (int A[], int n, int key) {
    int flag = 0;
    for (int i = 0; i < n; i++)
        if (A[i] == key) {
            flag = 1;
            break;
        }
    if (flag == 0)
        cout << "Not found";
    else
        cout << "found";
}

```

Q2 Write pseudocode for iterative algorithm as been in iterative for ($i=1$ to $n-1$)

```

{
    t = A[i]; j = i - 1;
    while (j >= 0 & A[j] > t) {
        A[j+1] = A[j];
        j--;
    }
    A[j+1] = t;
}

```

Recursive

```

Void insertionSort (int arr[], int n) {
    if (n <= 1)
        return j;
}

```

Q3 $T(n) = 3T(n/2) + n^2$
 $n \log_2^3 = n^{1.5}$
 $n^{1.5} > n$

Ans 3

```

insertion sort (arr, n-1)
int last = arr[n-1], j = n-2;
while (j > 0 && arr[j] > last)
    arr[j+1] = arr[j]
    j--;
arr[j+1] = last;
}

```

Insertion sort is called online sorting as it takes an input element for iteration and produces a partial solution without requiring access to offline algorithm.

Q4 Complexity of all sorting algorithm has been discussed:

Algorithm	Best	Average	Worst
Bubble sort	$O(n)$	$O(n^2)$	$O(n^2)$
Selection sort	$O(n^2)$	$O(n^2)$	$O(n^2)$
Insertion sort	$O(n)$	$O(n^2)$	$O(n^2)$
Count sort	$O(n)$	$O(n)$	$O(n)$
Quick sort	$O(n \log n)$	$O(n \log n)$	$O(n^2)$
Merge sort	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
Heap sort	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
Template	stable instance		

Algorithm	Inplace	Stable	Online
Bubble	✓	✓	×
Selection	✓	×	×
Insertion	✓	✓	✓
Count	×	✓	×

Merge	X	✓	X
Quick	✓	X	X
Heap	✓	X	X

Q5 Write pseudo code for binary search the time space.

Ans Recursive =

```

int binary (int arr[], int l, int r, int Key)
{
    if (r >= l)
    {
        int mid = l + (r - l) / 2;
        if (arr[mid] == Key)
            return mid;
        if (arr[mid] > Key)
            return binary (arr, mid - 1, Key);
        else
            return (arr, mid + 1, r, Key);
    }
    return -1;
}

```

```

int binary search (int arr, int l, int r, int Key)
{
    while (l <= r) {
        int m = l + (r - l) / 2;
        if (arr[m] == Key);
            return m;
        if (arr[m] < Key)
            l = m + 1;
        else
            r = m - 1;
    }
    return -1;
}

```

Q7 Find two indices such that $A[i] + A[j] = k$ in minimum.

Ans7:

```
void Sum (int A[], int K, int n)
{
    Sort (A, A+n)
    int i=0; j=n-1;
    while (i < j)
        if (A[i] + A[j] == K)
            break;
        else if (A[i] + A[j] > K)
            j--;
    print (i, j);
}
```

Q8 Which sorting best for practical uses? Explain.

Ans8 For practical uses, ~~no~~ it would be best for very large data. Further, time complexity of merge sort is same in all cases, that is $O(n \log n)$.