TUTORIAL-2

Name - Anshika Rathi
Sec - CST - SPL2
Roll No. - 70

ans 1

```
void fun (int n) {
    int j=1, i=0;
    while(i<n){
        i+=j; j++; }
}
```

for $j=1$ $\quad$ $i=1;$
$j=2$ $\quad$ $i=1+2;$
$j=3$ $\quad$ $i=1+2+3;$
⋮ $\quad$ ⋮

(m levels)

for i

$\because 1+2+3+\ldots < n$

$\because 1+2+\ldots m < n$

$\therefore \dfrac{m(m+1)}{2} < n$

$m \approx \sqrt{n}$

$\therefore$ by summation method

$\Rightarrow \sum_{i=1}^{m} 1 \Rightarrow 1+1+\ldots \sqrt{n}$ times

$\therefore \boxed{T(n) = \sqrt{n}}$

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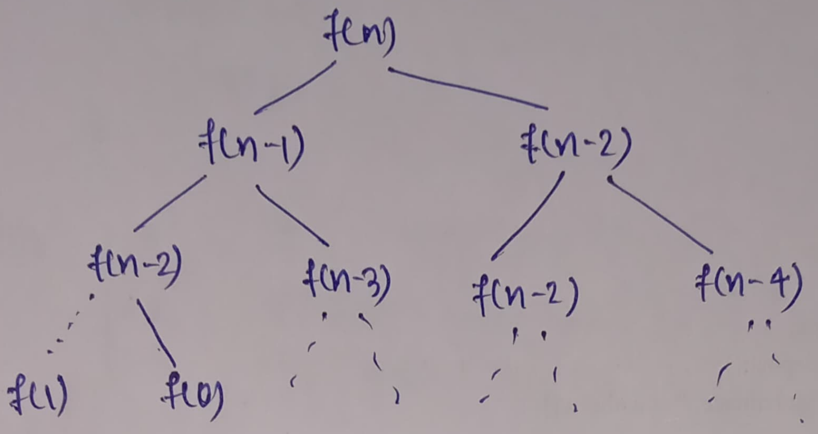
Ques2. For Fibonacci series -

$$f(n) = f(n-1) + f(n-2)$$

$$f(0) = 0 \quad f(1) = 1$$

By forming tree -



At every function call we get 2 function calls.

$\therefore$ For n levels -

we have, $2 \times 2 \times \ldots n$ times

$\therefore$ $\boxed{T(n) = 2^n}$

Maximum space -

considering recursive stack,

no. of calls maximum $= n$

For each call we have space complexity $O(1)$

$\therefore$ $\boxed{T(n) = O(n)}$

without considering recursive stack,

for each call we have time complexity $O(1)$

$\therefore$ $\boxed{T(n) = O(1)}$

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Que3.

(i) nlogn

Quick sort

```
void quicksort (int arr[], int low, int high)
{
    if (low < high)
    {
        int pi = ~~positi~~ partition (ar, low, high);
        quicksort (ar, low, pi-1);
        quicksort (ar, pi+1, high);
    }
}

int partition (int ar[], int low, int high)
{
    int pivot = arr[high]
    int i = (low-1);
    for (int j = low; j<=high-1; j++)
    {
        if (arr[i] <pivot)
        {
            i++;
            swap (&arr[i], &arr[j]);
        }
    }
    swap (&arr[i+1], &arr[high]);
    return (i+1);
}
```

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(2) n^3

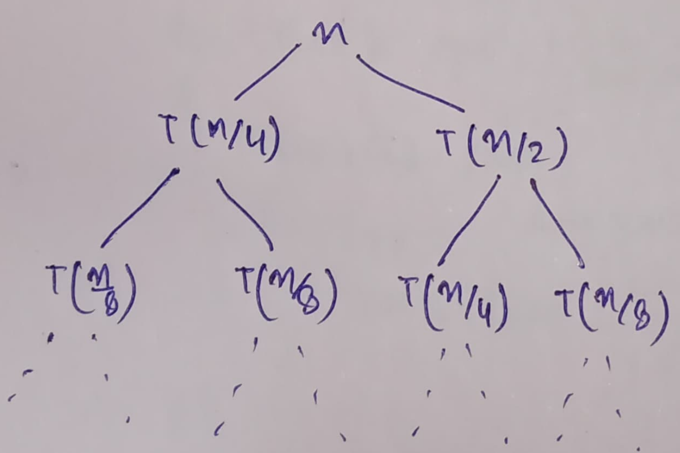
multiplication of two square matrix

```
for (i=0; i<r1 ; i++){
    for (j=0; j<c2 ; j++){
        for (k=0; k<c1; k++)
        {
            res[i][j] += a[i][k] * b[k][j];
        }
    }
}
```

(3) $\log(\log n)$

```
for (i=2; i<n; i= i*i)
{
    count++;
}
```

Que 4. $T(n) = T(n/4) + T(n/2) + c \times n^2$



At level -

$0 \rightarrow cn^2$

$1 \rightarrow \frac{n^2}{4^2} + \frac{n^2}{2^2} = \frac{c5n^2}{16}$

$2 \rightarrow \frac{n^2}{8^2} + \frac{n^2}{16^2} + \frac{n^2}{4^2} + \frac{n^2}{8^2} = \left(\frac{5}{16}\right)^2 n^2$

⋮

max levels $= \frac{n}{2^k} = 1$

$\Rightarrow k = \log_2 n$

$\therefore T(n) = c\left(n^2 + (5/16)n^2 + (5/16)^2 n + \dots (5/16)^{\log n} n^2\right)$

$= cn^2\left[1 + (5/16) + (5/16)^2 + \dots (5/16)^{\log n}\right]$

$= cn^2 \times 1 \times \left(\frac{1-(5/16)^{\log n}}{1-5/16}\right)$

$= cn^2 \times \frac{11}{5}\left(1-(5/16)^{\log n}\right)$

$\therefore \boxed{T(n) = O(n^2 c)}$

ques:

```
int fun(int n) {
    for (i=1; i<=n; i++)
        for (j=1; j<n; j+=i)
            //O(1)
}
```

for i	j	
		$j = (n-1)/i$ times
1	1	
2	1+3+5	
3	1+4+7	
.	.	
.	.	
n	.	

$\sum_{i=1}^{n} \frac{(n-1)}{i}$

Sathi

$\therefore T(n) = \frac{(n-1)}{1} + \frac{(n-1)}{2} + \frac{(n-1)}{3} + \dots \frac{(n-1)}{n}$

$$T(n) = n\left[1+\frac{1}{2}+\frac{1}{3}+\cdots\frac{1}{n}\right] - 1\left[1+\frac{1}{2}+\frac{1}{3}+\cdots\frac{1}{n}\right]$$

$$= n\log n - \cancel{\log n}\ \log n$$

$$\therefore T(n) = O(n\log n)$$

Que6. `for (i=2; i<=n; i = pow(i,k))`

```
{
    O(1)
}
```

for i

2^1

2^k

2^{k^2}

2^{k^3}

$\vdots$

2^{k^m}

where, $2^{k^m} <= n$

$k^m = \log_2 n$

$m = \log_k \log_2 n$

$$\therefore \sum_{i=1}^{m} 1$$

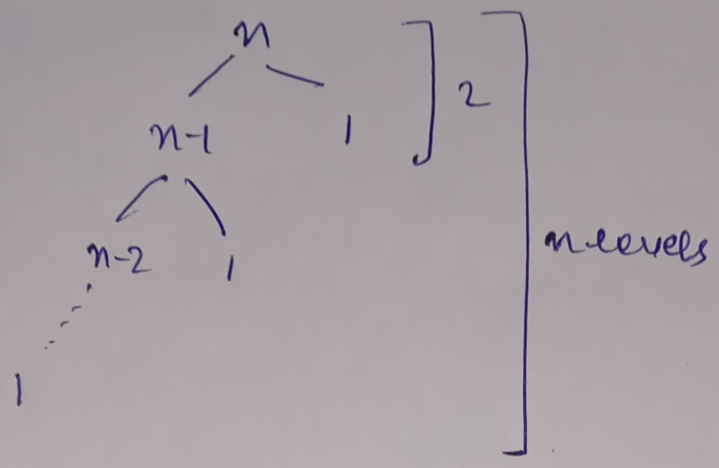
$\Rightarrow 1 + 1 + 1 + \ldots m$ times

$\Rightarrow \boxed{T(n) = O(\log_k \log n)}$

Aarti

Que7. Given algo divides array in 99% & 1% part

$\therefore T(n) = T(n-1) + O(1)$



'n' work is done at each level for merging.

$T(n) = (T(n-1) + T(n-2) + \dots T(1) + O(1)) \times n$

$= n \times n$

$\therefore \boxed{T(n) = O(n^2)}$

Lowest height = 2
Highest height = n

$\therefore$ $\boxed{\text{Difference} = n-2}$ $n > 1$

The given algo produces linear result.

Aarti

Ques 8. Considering for large values of 'n'

(a) $100 < \log\log n < \log n < (\log n)^2 < \sqrt{n} < n <$
$n\log n < \log(n!) < n^2 < 2^n < 4^n < 2^{2^n}$

(b) $1 < \log\log n < \sqrt{\log n} < \log n < \log 2n$
$< 2\log n < n < n\log n < 2n < 4n < \log(n!)$
$< n^2 < n! < 2^{2^n}$

(c) $96 < \log_8 n < \log 2n < 5n < n\log_6 n < n\log_2 n$
$< \log(n!) < 8n^2 < 7n^3 < n! < 8^{2n}$

Aarti